

# 算法基础 hmv1

Q1 1. 正确

$$F_m(n) = \sum_{i=1}^m f_i(n) = \cancel{m} m O(g(n))$$

对于  $\forall m$ , 满足  $F_m(n) = m O(g(n)) = O(g(n))$

2. 错误

$$\text{令 } f_1(n) = n = g(n) = O(n) = (1)n$$

$$F(n) = \sum_{i=1}^n f_i(n) = n^2 \neq O(n) \quad \text{则错}$$

3. 正确

要对  $\{f_1, \dots\}$  满足  $\forall i \in \mathbb{N}^+, f_i(n) = o(f_{i+1}(n)) = o(g(n))$

$$\text{由 } f_1(n) = o(g(n))$$

~~构造~~ 构造  $f_2(n), \dots$  使  $\forall i, f_i(n) = o(g(n))$

可构造  $f_i(n) = \sqrt[k]{g(n)}$   $k > 1$  且时满足

$$\text{又使 } f_i(n) = o(f_{i+1}(n))$$

$$\text{构造 } f_{i+1}(n) = \sqrt[k]{f_i(n)} \quad k > 1$$

$$\text{则若 } f_{i+1}(n) = \sqrt[k]{f_i(n) g(n)} \quad \text{时成立}$$

即存在。 正确。

4. 错误

$$f(n) + g(n) = \Omega(h(n))$$

使  $f(n)$  和  $g(n)$  可以在特定  $n$  值下取到 0, 但满足  $\Omega(h(n))$  即可。



$$\begin{aligned}
 \text{不妨 } f(n) + g(n) &= n \left( \sin(n\pi + \frac{\pi}{4}) + 2\sqrt{2} \right) \\
 &= n \left( \frac{\sqrt{2}}{2} \sin n\pi + \frac{\sqrt{2}}{2} \cos n\pi + \frac{\sqrt{2}}{2} \right) \\
 &= \frac{\sqrt{2}n}{2} (\sin n\pi + 1) + \frac{\sqrt{2}n}{2} (\cos n\pi + 1)
 \end{aligned}$$

$$\begin{aligned}
 \text{令 } \Omega(h(n)) &= O(n) \\
 \text{则 } f(n) + g(n) &= \Omega(h(n))
 \end{aligned}$$

但可以取到  $n$  使  $f(n), g(n) \notin \Omega(h(n))$

5.





$$Q2 \quad \{ A(0, n) = n+1$$

$$n \geq 0$$

$$A(m+1, 0) = A(m, 1) \quad m \geq 0$$

$$A(m+1, n+1) = A(m, A(m+1, n)) \quad m, n \geq 0$$

1) 递归. 由  $A(0, n) = n+1 \geq 1$ .

一般情况下  $A(m, n) \geq 1$ .

当  $m=0$   $A(0, n) = n+1$

当  $m \neq 0$   $A(m, n)$  进入递归

假设  $m=k$  递归可以从  $A(0, n)$  终止.

当  $m=k+1$  时  $A(k+1, n)$

①  $n=0$  时  $A(k+1, 0) = A(k, 1)$ , 由假设可以终止.

$n \neq 0$  时 ~~假设~~ 假设  $n=t$  时可以终止.

当  $n=t+1$  时  $A(k+1, t+1)$

$= A(k, A(k+1, t))$

由假设  $A(k+1, t)$  终止, 设为  $i$

$= A(k, i)$  由假设终止.

则终止。



(2) ~~证  $A(m+1, n) > A(m, n)$~~

~~对  $A(m, n)$ , 当  $m=0$   $A(0, n) = n+1$~~

~~当  $m=k$~~

~~对  $A(m+1, n)$  当  $n=0$   $A(m+1, 0) = A(m, 1)$~~

先证  $A(m, n+1) > A(m, n)$

对  $A(m, n+1)$  当  $m=0$   $A(0, n+1) = n+2 > n+1 = A(0, n)$

当  $m=k+1, k \geq 0$

$$A(k+1, n+1) = A(k, A(k+1, n))$$

下证  $A(m, n) > m+n$

当  $m=0$   $A(0, n) = n+1 > n$

假设  $m=k$  时成立  $A(k, n) > k+n$

$m=k+1$  时  $A(k+1, n)$  当  $n=0$   $A(k+1, 0) = A(k, 1) > k+1$

假设  $n=t$  成立  $A(k+1, t) > k+1+t$

$n=t+1$  时  $A(k+1, t+1) = A(k, A(k+1, t))$

~~$> k + A(k+1, t)$~~

~~$= k + A(k+1, t) > 2k+1$~~

$> k + A(k+1, t) > k + k+1+t = 2k+t+1$

$> k+t+2$

综上  $A(m, n) > m+n$





继续证明:  $A(k+1, n+1) = A(k, A(k+1, n)) > k + A(k+1, n)$   
 $> A(k+1, n)$  ~~(k=0时不取等, k=0时~~  
 ~~$A(k+1, n)$~~  证.  $A(m, n+1) > A(m, n)$

2.  $A(m+1, n) > A(m, n)$

当  $n=0$   $A(m+1, 0) > A(m, 1) > A(m, 0)$  (由首证)

当  $n=t+1$   $A(m+1, t+1) = A(m, A(m+1, t))$   
 $> A(m, m+1+t) \geq A(m, t+1)$

综上  $A(m+1, n) > A(m, n)$

3)  $d(x) = w(x)$

由  $A(0, n) = n+1 \geq 1$ , ~~令  $x \geq 1$~~  取  $x_1 < x_2$

$A(d(x_1), d(x_2)) \leq x_1 < x_2$

$d(x_1) \leq d(x_2)$

则  $d(x)$  单调



Q3.1) 1.  $T(n) = 3T(\frac{n}{2}) + n^2$   $T(n) = \Theta(n^2)$

2.  $T(n) = 4T(\frac{n}{2}) + n^2$   $T(n) = \Theta(n^2 \log n)$

3.  $T(n) = T(\frac{n}{2}) + 2^n$   $T(n) = \Theta(2^n)$

4.  $T(n) = 2^n T(\frac{n}{2}) + n^n$   $2^n \neq \text{const}$

5.  $T(n) = 16T(\frac{n}{4}) + n$   $T(n) = \Theta(n^2)$

6.  $T(n) = 2T(\frac{n}{2}) + n \lg n$   $T(n) = \Theta(n \lg^2 n)$

7.  $T(n) = 2T(\frac{n}{2}) + \frac{n}{\lg n}$

$\log_2^2 = 1$  对  $\forall \varepsilon > 0$

$$\lim_{n \rightarrow +\infty} \frac{\frac{n}{\lg n}}{n^{1-\varepsilon}} = \lim_{n \rightarrow +\infty} \frac{n^\varepsilon}{\lg n} = +\infty$$

$$\lim_{n \rightarrow +\infty} \frac{\frac{n}{\lg n}}{n^{1+\varepsilon}} = \lim_{n \rightarrow +\infty} \frac{1}{n \lg n} = 0$$

$$\lim_{n \rightarrow +\infty} \frac{\frac{n}{\lg n}}{n \lg^k n} = 0 \quad k > 0$$

$\nexists f(n) \notin O(n^{1-\varepsilon})$

$\nexists \Omega(n^{1+\varepsilon})$

$\nexists \Theta(n \lg^k n)$

8.  $T(n) = 2T(\frac{n}{4}) + n^{0.5}$   $T(n) = \Theta(n^{0.5})$

9.  $T(n) = 0.5T(\frac{n}{2}) + \frac{1}{n}$   $0.5 < 1$





$$10. T(n) = 16T(\frac{n}{4}) + n! \quad T(n) = \Theta(n!)$$

$$11. T(n) = \sqrt{2}T(\frac{n}{2}) + \lg n \quad T(n) = \Theta(n^{\frac{1}{2}})$$

$$12. T(n) = 3T(\frac{n}{2}) + n \quad T(n) = \Theta(n^{\log_2 3})$$

$$13. T(n) = 2T(\frac{n}{2}) + \sqrt{n} \quad T(n) = \Theta(n)$$

$$14. T(n) = 4T(\frac{n}{2}) + cn \quad T(n) = \Theta(n^3)$$

$$15. T(n) = 3T(\frac{n}{4}) + n \lg n \quad T(n) = \Theta(n \lg n)$$

$$16. T(n) = 3T(\frac{n}{3}) + \frac{n}{2} \quad T(n) = \Theta(n)$$

$$17. T(n) = 6T(\frac{n}{3}) + n^2 \lg n \quad T(n) = \Theta(n^2 \lg n)$$

$$18. T(n) = 4T(\frac{n}{2}) + \frac{n}{\lg n} \quad T(n) = \Theta(n^2)$$

$$19. T(n) = 64T(\frac{n}{8}) - n^2 \lg n \quad -n^2 \lg n < 0$$

$$20. T(n) = 7T(\frac{n}{3}) + n^2 \quad T(n) = \Theta(n^2)$$

$$21. T(n) = 4T(\frac{n}{2}) + \lg n \quad T(n) = \Theta(n^2)$$

$$22. T(n) = T(\frac{n}{2}) + n(2 - \cos n) \quad T(n) = \Theta(n)$$

$$23. T(n) = 2T(\frac{n}{2}) + n \lg n \lg \lg n \quad \text{同 7}$$

$$24. T(n) = 2T(\frac{n}{2}) + n(\lg \lg n)^2 \quad \text{同 7}$$

$$25. T(n) = 2T(\frac{n}{2}) + \frac{n(\lg \lg n)^5}{\lg n \lg \lg n} \quad \text{同 7}$$



$$12) 7. T(n) = 2T\left(\frac{n}{2}\right) + \frac{n}{\lg n}$$

$$e = (1, -1), m = 1$$

$$T(n) = \frac{n}{\lg n} \lg n \lg \lg n = n \lg \lg n$$

$$23. T(n) = 2T\left(\frac{n}{2}\right) + n \lg n \lg \lg n$$

$$e = (1, 1, 1) \quad m = 2$$

$$T(n) = n \lg n \lg \lg n \lg n = n \lg^2 n \lg \lg n$$

$$24. T(n) = 2T\left(\frac{n}{2}\right) + n (\lg \lg n)^r \quad r \neq 0$$

$$e = (1, 0, r) \quad m = 2$$

$$T(n) = n (\lg \lg n)^r \lg n \quad r \neq 0$$

$$25. T(n) = 2T\left(\frac{n}{2}\right) + \frac{n (\lg \lg \lg n)^s}{\lg n \lg \lg n} \quad s \neq 0$$

$$e = (1, -1, -1, s) \quad m = 3$$

$$T(n) = \begin{cases} \frac{n (\lg \lg \lg n)^s}{\lg n \lg \lg n} \lg n \lg \lg n \lg n \lg \lg n \lg \lg n & s > -1 \\ \frac{n (\lg \lg \lg n)^s}{\lg n \lg \lg n} \lg n \lg \lg n \lg \lg n \lg \lg n \lg \lg n & s = -1 \\ n & s < -1 \end{cases}$$

