

# Context-Aware Scheduling Algorithm in Smart Home System

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**Abstract:** The main idea of pervasive computing is to make computing exist everywhere in the physical world. The smart home system is an important realisation of pervasive computing whose aim is to provide system users with an intelligent life experience. The key technique used to realise this is context awareness. Contexts in the living space can provide large amounts of information regarding users' behaviours and habits. Together with an information system, it can automatically execute many common operations of applications, instead of users, and can make the applications "smart". However, since contexts in the environment are diverse and sensitive, it is difficult to choose the ones that are most useful to the users' current activity. A proper scheduling strategy should first consider the users' demand. This paper proposes a context-aware scheduling algorithm that is based on correlation, with the purpose of improving the utilization rate of context collections. Experiments show that with the priority based on correlation in low-level contexts, the scheduling of reasoning tasks can reduce the cost of transmission.

**Key words:** pervasive computing; smart-home; reasoning scheduling; correlation

## I. INTRODUCTION

The idea of pervasive computing is to bring about the existence of computers anywhere in our physical world without sensation of users,

users are the centre. Every computing unit will cater to the users' habit and make autonomous interactions with users' activities. In that case, people in the pervasive environment can concentrate on their working or life enjoying without spending energy on operating the machines. Furthermore, pervasive computing also will provide users customised information services by perceiving and predicting users' demand. Pervasive computing infrastructure in the home should be able to sense the context in which specific situations take place and adapt to them according to its location of use, the people and objects that are around, and changes of those entities over time, how to make an effective use of the context information is an important and difficult topic. Context scheduling plays a vital important role in this field. The main purpose of context scheduling in smart home environment is to make a sequence of tasks performing. With the sequence, applications in the system can provide proper services to the users with effective using of context information collected from the environment. The context scheduling algorithm in smart home system should consider not only the user preference but also the cost of context transmission. In scheduling theory, a real-time system comprises a set of real-time tasks; each task consists of an infinite or finite stream of jobs. The task set can be scheduled by a number of policies including fixed priority or dynamic priority algorithms [1-3]. The common dynamic priority scheduling algo-

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This paper proposes a priority scheduling algorithm based on context correlation to improve the average utilization of low-level context in the pervasive computing environment, such as smart home system. Experimental results show that the proposed algorithm can reduce the times of cache replacement and the times of data transmission.

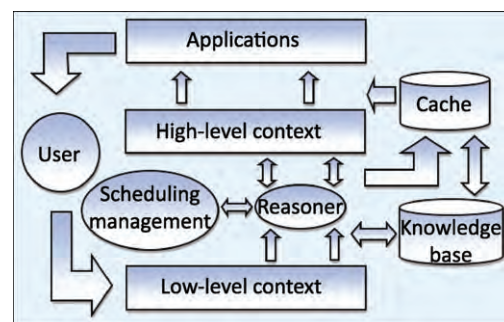
rithms do not consider context-aware of pervasive environment. This paper focuses the scheduling problem on the communication process between low-level context and the context reasoning module, and proposes a context priority scheduling algorithm based on correlation with the purpose to solve the problem of behaviour disorder of system caused by the conflict in context reasoning, and to improve the utilization rate of context collections in the smart home environment.

The rest of this paper is organised as follows. In Section II we review introduce related works and discuss a context-aware model in smart home. In Section III we present an improved context-aware scheduling algorithm based on correlation, and illustrate the algorithm on an example. Section IV provides experimental support for our improved context-aware scheduling algorithm. Finally, some conclusions are drawn in Section V.

## II. CONTEXT-AWARE IN SMART HOME SYSTEM

### 2.1 Related work

Many applications of pervasive computing have to access some related contexts in order to provide the right services at the right time and the right places. However, there are some challenges for applications which need to store a lot of information accessed in the future previously, and to expend the information of communication channel timely in smart home. These challenges make the design of smart home applications much more difficult than other applications. The smart home system should effectively integrate communication and computing networks among the previously separated electronic devices in home and incorporate core tenets of pervasive computing. Many researchers have focused on context-aware smart applications. Georgia Tech's Aware Home experimented with many technologies for capturing low-level sensor data about human activities but did not focus on automated activity recognition. The MIT House project developed an instrumented condominium



**Fig.1** Model of the context scheduling in smart home system

for studying activity recognition in a naturalistic environment. This space is designed as a living laboratory from which experiments in activity recognition can be conducted in a naturalistic manner [4-9]. A context-aware smart home always utilises contextual information of the involved entities to adapt its behaviour and provide context-aware services. Context-aware in smart home system should have the ability to forecast user behaviours actually. For example, whether a homeowner is at work or on vacation, the smart home will alert him to what is going on, and security systems can be built to provide an immense amount of help in an emergency. And the cost to realise the system must be reduced to the level that user can accept.

Some researchers have more focused on context-aware information abstraction and pre-treatment rather than context scheduling. Ref. [4] presented a distributed system model for context-aware mobile computing, which can collect information from sensors, reason from known database, and thus adopt corresponding activities. A. Ranganathan et al. proposed a context model based on first order predicate calculus. They use first order model to express complex rules, which involves contexts. This knowledge expression enables automated inductive and deductive reasoning to be easily done on contextual information [8]. Ref. [9] clarified the notion of context and laid out foundations for the design and development of context-aware applications. In this paper, the context-aware scheduling algorithm based on correlation in smart home environment also

inherits this method.

## 2.2 An interactive system model of context-aware in smart home system

In general, the context is a meaning of the broad concept that can be considered context is affecting the behaviour of a group of computing entities implicit input. Context is typically the location, identity, and state of people, groups, and computational and physical. objects. And the context in home environment can be divided into two categories: physical context, including the position, temperature, humidity, noise, etc.; complex context, including the human behaviour, user's facial expression and so on. Ref. [10] provided a basic model of context model which describe the implementation process of context collection in a common intelligent environment as shown Figure 1. To enhance user's experience, the context must be given different priorities taken from our previous work [6]. The higher priority means that the context is more sensitive to the user and easily to bring about user's special behaviour. Users' habits would be gathered in the historical records by contrasting the data access. Then the priority of context would be modified in order to make sure that context has high priority is the one that sensitive to the users. The communicating channels also should be given. By collecting feedback information from users directly, some personalised adjusts would be given to the priority of context.

## 2.3 Preference of user experience

In smart home environment, man is the centre of the whole system. Services are customised by the actions and preference of the user. Without user's behaviour and preference, the pervasive computing environment will not provide service with high-level personalization. The sensors cannot easily collect those indirect-contexts directly; however, we can learn the user preference by backtracking the historical records.

Different from the application layer, user

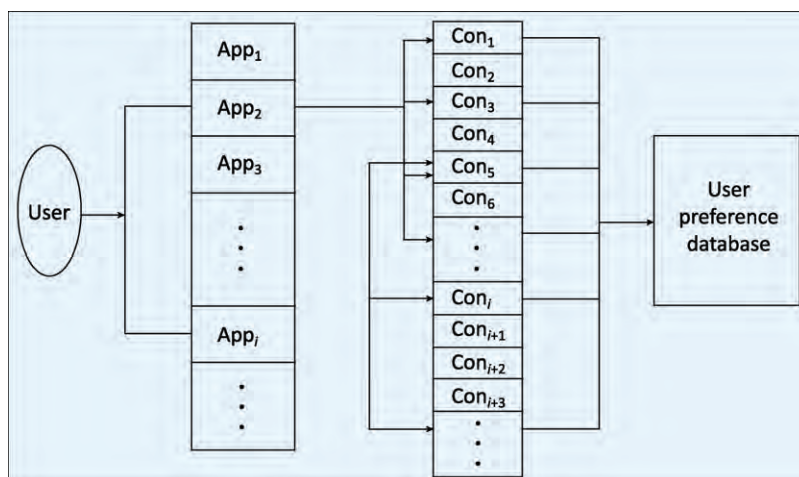


Fig.2 Obtain the user preference

will not have directly interaction with the reasoning and scheduling module. We can only collect user preference by tracking the application access. The more appearance of one context the more sensitive it will be to the user in the system.

## 2.4 Illustrative use case

Air-conditioner, TV set, lighting system, water heater, cleaning robot, security system and other intelligent devices are placed in the smart home system. At 8:00 a.m., Miss WANG leaves home and goes to work. System judges no one in the house then opens the security system. Cleaning robot begins to sweep the floor, and other devices turn off to keep safe. At 7:00 p.m., when Miss WANG comes back home, the system perceives that there are user activities in the house. Then it will turn on the lighting system where user being. The air-conditioner works, and the TV set begin to broadcast the program that she preferring. At 11:00 p.m., she goes to sleep. System turns on security system and turns off the TV set. Air-conditioner keeps on working but switch to "sleep" mode.

In these common scenarios [11], smart home system controls intelligent devices with proper applications. These applications learn different scenes from contexts information in the environment. The lighting system mostly has no need to work during the daytime. But in the evening, it becomes the most important

device for user. It is better to let the robot do cleaning work when no one is in the house and let the air-conditioner work when user back home. Some context information affects other one's importance. Such as time, luminance and user activities, we should give these contexts high priority in order to realise the dynamic regulation of the priority of intelligent devices in different scenes. Task scheduling strategy should consider these relationships between different contexts and modify execution sequence with the environment information changing.

### III. PRIORITY SCHEDULING ALGORITHM OF CONTEXT-AWARE

#### 3.1 Algorithm process and basic definition

Having presented context information types in smart home system, we now discuss how applications can effectively use context information. Context reasoning is the essential part of context-aware computing. And task scheduling is the important constituent part of any system. Also, context reasoning needs an effective strategy to manage reasoning tasks. The more important task has the higher priority, and if they have the same priority, whether it is disruptive becomes the primary basis of scheduling. The main steps of priority calculation are shown as Algorithm 1.

**Algorithm 1** The main steps of priority calculation

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```

1: Priority contextPriority (context  $c_j$ ) {
2:   while (task queue is not empty)
3:      $T = \text{head of the task queue}$ 
4:     while (queue  $T$  is not empty)
5:       if (queue  $T$  has correlation with  $c_j$ )
6:         //if task  $T$  need context  $c_j$  to reason, it
6:         //has correlation with  $c_j$ 
6:         put  $T$  into queue  $T_r$ 
6:          $T = T \rightarrow \text{next}$ 
7:       end if
8:     end while
9:   if (queue  $T_r$  is not empty)
10:    while (queue  $T_r$  is not empty)
11:       $t = t + T_{e_j}$ 
11:       $n = n + 1$ 
11:      //calculate the whole execution time
11:      //and number of tasks having correlation with  $c_i$ 
11:       $T_r = T_r \rightarrow \text{next}$ 

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12:   end while
13:    $T_{e_{ave}} = t/n$ 
14: end if
15:  $\text{trans}_i = T_{t_j} / T_{t_{ave}}$ 
16:  $P_j = (r_j / T_{e_{ave}}) * (n/N) * \text{trans}_j + C$ 
17: return  $P_i$ 
18: end while
19: }

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To evaluate the priority of reasoning task, some attributes will be considered:

First, the direct correlation of low-level contexts among tasks is useful. In pervasive computing environment, transferring cost probably restrict the efficiency of system processing. Choose the task having the most same low-level contexts with current task as the next task to be scheduling will reduce the cache replacement times.

Second, low-level contexts with long refresh period have high reusability. Give high priority to these contexts also can reduce cache replacement times.

Third, times a low-level context appearing in tasks queue can show the importance of the context to the system.

In this paper, we give basic symbols definition as follows:

- 1) Define  $c_j$  as the low-level context,  $j$  is the number of the context;
- 2) Define  $t_i$  as the query task,  $i$  is the number of it;
- 3) Define  $r_j$  as the refreshing period of  $c_j$ . That means  $c_j$  will lose efficacy when exceeding  $r_j$ , and need to be collected again;
- 4) Define  $T_{t_j}$  as the transmission cost of once refreshing  $c_j$ , and  $T_{t_{ave}}$  as the average transmission cost of all the contexts;
- 5) Define  $N$  as the number of all tasks in the query queue, and  $n_j$  as the number of tasks which having correlation with low-level context  $c_j$ ;
- 6) Define  $T_{e_i}$  as the execution time of  $t_i$ , and  $T_{e_{ave}}$  as the average execution time of  $t_i$ ;

#### 3.2 Priority of the contexts

In the smart home environment, contexts have different updating ratios. Compared with the updating frequently ones, contexts with long refreshing period are more valuable to be stored.

red in the cache. Long efficient time allows more applications access contexts in one refreshing period. But, if not paid much attention by the system, a context even with long effective duration should not have a high priority. With the purpose of increasing the low-level context utilization ratio and reducing the cost of communications between sensors and reasoner, we give high priority to the context that has high access rate during one refreshing time. Besides, refreshing context with higher transmission cost will make more time cost. So we give high priority to the context having high transmission cost. The priority of context will be defined with three main qualifications: the refreshing period, the transmission cost of context and the times accessed by applications.

Given a context  $c_j$ ,  $r_j$  is its refreshing period. Define its sustainable freshness  $F_j$  as follows:

$$F_j = \frac{r_j}{Te_{ave}} \quad (1)$$

where  $F$  indicates the efficiency of access of context  $c_j$ . Within one refreshing period, the more tasks which related to  $c_j$  can be executed the more valuable for  $c_j$  to be stored in the cache. We give the expression of  $Te_{ave}$  as follows:

$$Te_{ave} = \frac{\sum_{i=1}^n t_i}{n} \quad (2)$$

Another important parameter to calculate the priority of  $c_j$  is the visit possibility. Define  $A_r$  as the percentage of tasks having correlation with context  $c_j$  among all the tasks.  $A_r$  predicts that how many tasks in the task queue will query the value of  $c_j$ :

$$A_r = \frac{n}{N} \quad (3)$$

Define  $trans_j$  as the indication of transmission cost of  $c_j$ :

$$trans_j = \frac{Tt_j}{Tt_{ave}} \quad (4)$$

Higher value of  $trans_j$  means that if  $c_j$  lose efficacy, higher cost to refresh it should be paid.

With the three basic parameters above, this paper gives the definition of the priority  $P_j$  to

the context  $c_j$  as follows:

$$P_j = F_j \times A_r \times trans_j + C \quad (5)$$

For any low-level context  $c_j$ , the tasks having short execution time will have higher implementation rate within one refreshing period than the ones having long execution time. So we give the former higher priority to be reasoned. But if the refreshing period of the context is very short, the reasoning result will not be valid for long. So we reduce the priority of contexts in this kind. In addition, if the context has long time to be ignored by the applications, this means that we have no need to give high priority to it. Another parameter  $C$  in this expression means the user preference. For different living habits, users also have their own concern. If the users care about some context in the environment (such as temperature) specially, they can give the high priority to it.

### 3.3 Description of task and partition of task set

This paper use an acyclic directed graph  $G(V, E)$  to represent a group of tasks which should be scheduled.  $V$  is the nonempty finite set of nodes.  $E$  is the set of directed lines in different nodes. The set of nodes  $V=\{task_i\}$  ( $i=1, 2, \dots, n$ ) means a group of tasks which should be scheduled. The set of lines  $E=\{e_{ij}\}$  means the precedence relationship in different tasks.  $e_{ij}$  shows that  $task_i$  should finish before  $task_j$  starts. If  $e_{ij}$  exist in the directed graph,  $t_i$  is the pre-order task of  $t_j$ , and  $t_j$  is the successor task of  $t_i$ .

This paper gives the description of a task with four-tuple:  $task_i$  (set queue number, contexts matrix, execution time, priority). In considering the precedence relationship in different tasks, the set queue number of task is given as follows:

- 1) Choose the tasks in  $G(V, E)$  where in-degree is 0 and put them into the set queue  $TQ_1$ ;
- 2) Decreased the out-degree of all the successor tasks in  $TQ_1$  by 1, then choose the ones that in-degree is 0 and put them into the set queue  $TQ_2$ ;
- 3) Repeat 1)-2) until all tasks in  $G(V, E)$  are in the corresponding set queues  $TQ_i$  ( $i=1, 2, \dots, n$ ).

**Table I** The priority of low-level contexts in assumption

Low-level contexts	$c_1$	$c_2$	$c_3$	$c_4$	$c_5$	$c_6$	$c_7$	$c_8$
Priority	6	3	4	4	1	2	5	3

**Table II** The relationship table of context-aware elements

	$c_1$	$c_2$	$c_3$	$c_4$	$c_5$	$c_6$	$c_7$	$c_8$
$t_1$	1	1	0	0	1	0	0	1
$t_2$	1	0	1	0	1	0	0	1
$t_3$	0	1	0	0	1	1	1	0
$t_4$	0	0	1	1	1	1	0	0
$t_5$	0	0	1	0	0	0	1	0
$t_6$	1	0	0	0	1	1	1	0

The lower set queue number means higher priority of the task queue. For reducing the “BUMP” phenomenon [12] in the scheduling process, the task queue with higher priority will be operated prior at the ones with lower priority.

### 3.4 Scheduling algorithm based on correlation

Each reasoning task has correlation with several low-level contexts. In this scheduling algorithm, priority of the task will inherit from the contexts it related to. That means a reasoning task may relate to different low-level contexts, the priority of it can be defined from low-level contexts that used for reasoning [12]. For example, assume a task  $t_{\text{light}}$  is for controlling the light in the restroom. During the daytime, it may have a low priority because the low-level context “light” will not be requested many times; but at night, “light” became the important context for system users. Its priority became higher than in the daytime. By inheriting, the priority of  $t_{\text{light}}$  also rises.

Suppose some low-level contexts in the smart home environment and their priorities are given in Table I.

We established the relation based on two-dimensional correlation table according to the different corpus  $C$ . In Table II, the horizontal direction contains different low-level contexts, and the vertical direction contains the tasks that reasoning from low-level context to high-level context. As the name suggests, correlation is used to characterise the relationship between

each reasoning task. Given a visit to low-level contexts corpus  $C=(c_1, c_2, c_3, \dots, c_n)$ , and reasoning task queue  $TQ=(t_1, t_2, t_3, \dots, t_n)$ . If  $t_i$  has relationship to  $c_j$ , we mark  $t_i \times c_j=1$  at the intersection of the line correspond to  $t_i$  and the column correspond to  $c_j$  in the relationship table; if not, we mark  $t_i \times c_j=0$  in the same place. Given a task queue of high-level context query request, and each of them is related to the sequences of low-level context as follows:  $t_1(c_1, c_2, c_5, c_8)$ ,  $t_2(c_1, c_3, c_5, c_8)$ ,  $t_3(c_2, c_5, c_6, c_7)$ ,  $t_4(c_3, c_4, c_5, c_6)$ ,  $t_5(c_3, c_7)$ ,  $t_6(c_1, c_5, c_6, c_7)$ .

The relationship graph of corresponding to correlation is shown in Table II.

Priority scheduling algorithm will be built based on relations and the priority of the various elements. From the graph, we know the two-dimensional relation table quantifies the correlation between the tasks. Firstly, group these contextual elements. Secondly, start scheduling management. Algorithm 2 describes the main steps of the scheduling algorithm based on correlation.

The most important step of the algorithm is to calculate the priority  $Pt_i$  of task  $t_i$ . In this paper, the priority of task will be generated with 5 attributes:

1)  $Pc_i$  is the priority which indicates the correlation between task  $t_i$  and the current task. Add together the priority of all contexts that both associating with  $t_i$  and the current task, and then divide the number of low-level contexts  $N_{\text{con}}$  which having been added together. This is the basic priority of task  $t_i$ :

$$Pc_i = \frac{\sum P_j}{N_{\text{con}}} \quad (6)$$

2) Considered that the same task would have different importance at different time, the priority of it should be fixed automatically. As was talked above,  $t_{\text{light}}$  is much more important to the user during the night than in the daytime. And the application controlling the light is sure to run much frequently at night. So, when a task  $t_i$  has been executed recently, we give it a higher priority  $Pf_i$ .

The more times it has been executed, the higher priority it has.



$$Pf_i = 1 - a^{\frac{m_i}{M-m_i}} \quad (7)$$

In Eq. (7),  $M$  indicates the number of recent tasks having been executed;  $m_i$  is the times that  $t_i$  has been executed in recent  $M$  executed tasks;  $a \in (0, 1)$  is the parameter.

As  $m_i$  increasing, the growth rate of  $Pf_i$  reduces until approaching 0 (when  $m_i$  equal to  $M$ ). That means  $t_i$  has completely held the processor.

**Algorithm 2** The main steps of the scheduling algorithm

```

1: Schedule taskSchedule() {
2:   while (task queue is not empty)
3:     Con = head of contexts queue of current task
4:     while (queue Con is not empty)
5:       search Con in task queue
6:       if (task has correlation with Con)
7:         put task into queue  $T_{can}$ 
8:       end if
9:     end while
//pick up the tasks having correlation with current task
10:    while ( $T_{can}$  is not empty)
11:      Concan = head of contexts queue of candidate task
12:      while (Concan is not empty)
13:        if (Concan is in Con)
14:           $T_{can} \rightarrow Pt += Con_{can}$ 
15:           $N_{con}++$ 
16:        end if
17:      end while // calculate the Pc
18:       $m = \text{count (recently executed times of } T_{can})$ 
19:       $T_{can} \rightarrow Pt += (1 - a^{M-m}) * C_1$ 
//calculate the Pf
20:       $T_{can} \rightarrow Pt += \text{count (wait time of } T_{can}) * C_2$ 
//calculate the  $Pw$ 
21:       $d = \text{count (times } T_{can} \text{ has been dropped)}$ 
22:       $T_{can} \rightarrow Pt += e^d * C_3$  //calculate the  $Pd$ 
23:       $T_{can} \rightarrow Pt += T_{can} \rightarrow E$ 
24:    end while
//calculate the priority of tasks

25:    if ( $T_{can}$  is not empty)
26:      sort  $T_{can}$  and put the highest one into TASK
27:    return (TASK)
28:  end if
29: end while
30: }
```

3) To solve the problem of “starvation”, we give task  $t_i$  an addition priority  $Pw_i$  increasing with its waiting time:

$$Pw_i = \text{wait}_i \quad (8)$$

$\text{wait}_i$  is the indication of the time that task  $t_i$  has been waiting for execution. The longer  $t_i$  has been waiting, the higher the value of  $\text{wait}_i$  is. They bear the linear relation.

4) If task  $t_i$  has cannot be executed within its deadline, the task will be dropped. When the applications apply for the same reason task again, priority compensation will be given to the task. Define  $\text{Drop}_i$  as the times task  $t_i$  has been dropped. Then we give the  $Pd_i$  as:

$$Pd_i = e^{\text{Drop}_i} \quad (9)$$

As with the previous rule, this priority compensation  $Pd_i$  is in the purpose of solving the “starvation” problem. Using the exponential function,  $Pd_i$  will increase fast as the  $\text{Drop}_i$  increasing. In this way, after a less times of dropping,  $t_i$  will have a high priority with  $Pd_i$ .

5) Another thing must be considered is the “emergency” (such as alarm, doorbell and phone ring). We give these emergencies a special priority parameter  $E$ .

We give  $E$  a high value to ensure that the emergency can have the highest priority.

With these 5 attributes, the reasoning task  $t_i$  will generate its priority  $Pt_i$  for scheduling as follow:

$$Pt_i = Pc_i + C_1 \times Pf_i + C_2 \times Pw_i + C_3 \times Pd_i + E \quad (10)$$

In Eq. (10),  $C_1$ ,  $C_2$ ,  $C_3$  are the weight factors. They indicate the weight of each addition priority contribute to the task priority.

## IV. EXPERIMENTAL ANALYSIS

### 4.1 Basic experiment scene

In this section, we present experimental evaluation of the scheduling algorithm described above.

This experiment scene is based on the smart space model [13]. In the smart space, various context-aware applications work in cooperation. In this paper, we assume that 50 contexts concerned by the user are distributed in the space, 30 potential task requested by system applications might occur in random. Different from the model [14], in this experiment we

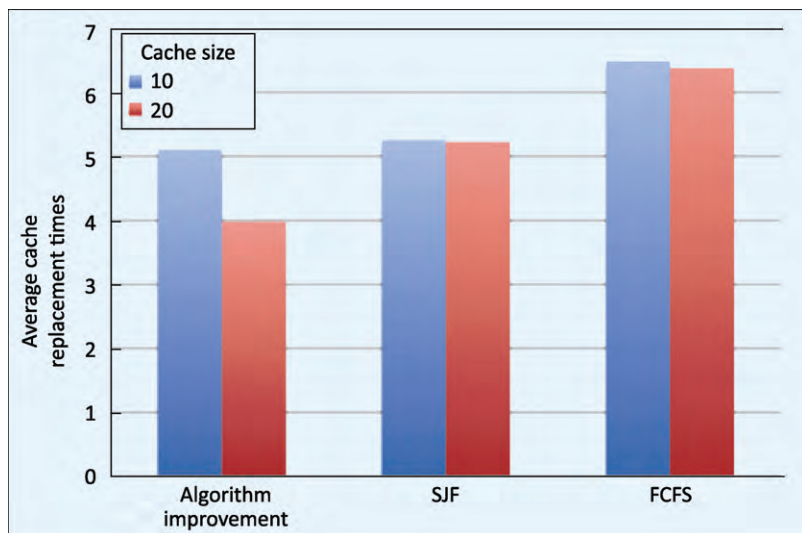


Fig.3 Performance comparison with different cache sizes

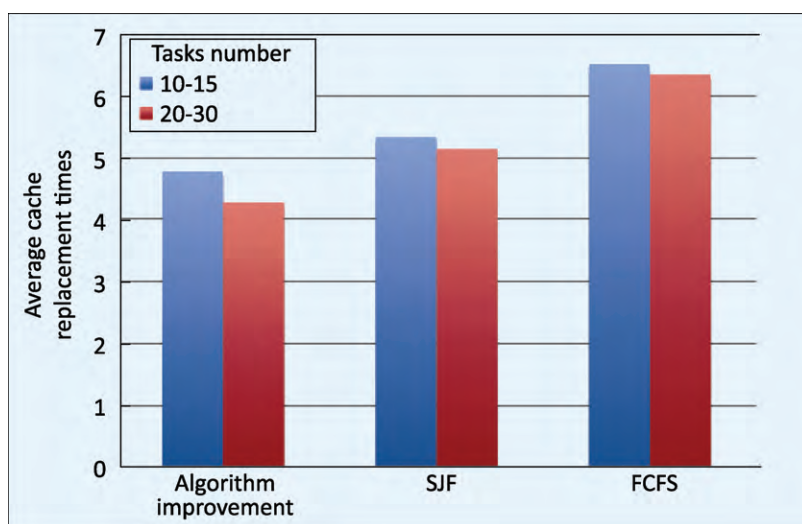


Fig.4 Performance comparison with different quantity of tasks

Table III The range of parameters in the experiment

Parameters	Range
Contexts number related to every task	5-10
Refresh period of every context	30-300 s
Transmission cost of every context	0-3 s

mainly concern with the transmission between reasoned and context sensors. We will compare this algorithm improvement with two classic real-time scheduling algorithms: FCFS and SJF in the main index that data transmission times between reason unit and receiver unit. Parameters related to the system will be generated randomly according to normal distribution.

## 4.2 Comparison of performance with different algorithms

Each time of experiments, the FCFS, SJF and the algorithm improvement will schedule a tasks queue separately. In order to simplify the experiment, we assume that new task will not be inserted into the queue in every scheduling process. Every low-level context tasks request for will be stored in the cache. If the cache is full, replace the context of lowest priority with the new one.

The simulation demonstrates that when the cache size is small, the performance of the algorithm improvement is nearly the same to the SJF and FCFS. As shown in the Table IV, in average the algorithm improvement consumes 2%-5% less cache replacement times than that of the SJF and 14%-20% less than that of FCFS. The main reason with small memory size cannot afford enough long time to store each context, and the improvement will not work completely. In this situation, SJF performs better than FCFS, because SJF is advanced in cache replacement by calculating the priority of the context.

Figure 3 shows that algorithm improvement performing better in large cache environment. When cache volume change from 10 to 20, the replacement times of cache reduced by 22%-24% with algorithm improvement, while the other two algorithms have little improvement. The reason is that the performance of SJF and FCFS do not mainly rely on the cache hits.

Another main analysis is the effect of the quantity of the tasks in the request queue. We assume that all tasks will be executed before deadline, and get the performances of these algorithms showed in Figure 4. As the tasks number in the queue increasing, the algorithm improvement reduces the replacement times of cache.

The experiment results indicate that this scheduling algorithm improvement performance is better than the classic real-time scheduling algorithms in large cache and mass tasks system. In this environment, the algorithm improvement can obviously reduce the replace-



ment times of cache. In other word, context in the smart space can be more effective within its refreshing period with the lower transmissions cost.

## V. CONCLUSION AND FUTURE WORK

This paper discussed a context-based model of smart home system. Context reasoning needs an effective strategy to manage reasoning tasks in pervasive environment. Based on context correlation, we proposed a context priority-scheduling algorithm in order to improve the average utilization of low-level context in the pervasive computing environment. Simulation experiments show that this algorithm can reduce the times of cache replacement and the times of data transmission. In order to further reduce the cost of context transmission, our next research work is to optimise the algorithm of cache replacement and make coordination with the scheduling algorithm proposed in this paper.

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