#1.) Indicate whether 
$$A \in O(B)$$
,  $A \in \Omega(B)$ , or  $A \in \Theta(B)$ 

A B O 
$$\Re$$
  $\Theta$ 

1)  $\sqrt{\frac{i}{09}n}$  1 Yes Yes Yes

2)  $\sqrt{\log n}$  loglogn NO Yes NO

3)  $T(n)$   $\frac{1}{5}n$  Yes Yes Yes

$$g(n) \in \mathcal{O}(f(n)) \longrightarrow g(n) \leq c \cdot f(n)$$

$$g(n) \in \mathcal{N}(f(n)) \longrightarrow g(n) \geq c \cdot f(n)$$

$$g(n) \in \mathcal{O}(f(n)) \longrightarrow d \cdot f(n) \leq g(n) \leq c \cdot f(n)$$

$$\frac{1}{1} N_{102} = O(1) \Rightarrow N_{102} = 1 \cdot C$$

$$(\log n)^{1/2} \leq c \cdot \log(\log n)$$

3) 
$$T(n) \in O(\frac{1}{5}n)$$
  $\Rightarrow 5 + 2 \leq (\frac{1}{5}n)$ 

$$T(n) = \begin{cases} 5T(n/5) + 2 & n > 1 \\ 1 & n = 1 \end{cases}$$

$$T(n) = Q \cdot T(%) + (\cdot n^{k})$$

$$= \Theta(v)$$

Use definitions to Prove: 
$$N + 3n^2 \in O(n^2)$$

Upper Bound 
$$N+3n^2 \leq C \cdot n^2 \rightarrow \frac{N+3n^2}{C} \leq n^2$$
Let  $N=1$ 

$$\frac{(1)+3(1)^2}{1000}$$
  $\leq 1$   $\sqrt{1000}$   $\leq 1$ 

| 
$$| (1) + 3(1)^2 \ge (1) \cdot (1)^2$$
 |  $| (1) + 3(1)^2 \ge (1) \cdot (1)^2$  |  $| (1) + 3(1)^2 \ge (1) \cdot (1)^2$  |

$$n+3n^2 \in \Theta(n^2)$$

Prove: 
$$\sum_{i=1}^{n} i = \frac{n(n+i)}{2}$$

$$\sum_{i=1}^{n} i = 1 + 2 + 3 + 4 \dots n = \frac{n(n+1)}{2}$$

I. H: Assume that 
$$\sum_{i=1}^{k} i = \frac{K(K+1)}{Z}$$

$$I.S: W.T.S: \sum_{i=1}^{k+1} i = \frac{(k+1)(k+2)}{2}$$

$$\sum_{i=1}^{K+1} i = \sum_{j=1}^{K} i + K+1$$

$$= (k+1) \left(\frac{\kappa}{2} + \frac{2}{2}\right) \left[-\frac{(\kappa+1)(\kappa+2)}{2}\right]$$

What are six steps one must do When designing an algorithm?

- 1. Idea/intuition
- 2. Pseudocode
- 3, code
- 4. Experimental run time
- 5. Prove run time
- 6. Prove correctness