

Exam Review 1

Tuesday, September 28, 2021 11:01 AM

Pg 190

- Kruskal & Prim's
- Knapsack & Huffman's
- Can you trace an algorithm?
- Can you write pseudocode?
- Practice techniques (ALL)

ex1: $\log \log n$ & $\log^2 n$

$$\lim_{n \rightarrow \infty} \left[\frac{(\log n)^2}{\log \log n} \right]'$$

$$\lim_{n \rightarrow \infty} \left[\frac{2(\log n)(1/n)}{(\frac{1}{\log n})(1/n)} \right] = \lim_{n \rightarrow \infty} \frac{2(\log n)}{n} \cdot (n \log n)$$

$$= \lim_{n \rightarrow \infty} 2(\log n) \cdot \log n = 2(\log n)^2$$

$$\lim_{n \rightarrow \infty} = \frac{2 \log^2 n}{c} = \infty \quad \therefore \log \log n \in O(\log^2 n)$$

ex2: Show that $\log \log n \in o(\log^2 n)$
using the definition

$$\log \log n < c \cdot (\log^2 n)$$

$$\text{Let } c=2 \quad N=2$$

$$0 < 2$$

$$\text{ex 3: } n = 2^{\log n} \rightarrow n^{\log^2} = n$$

$$n^{1/\log n} = 2^{\log n (1/\log n)}$$

$$n^{1/\log n} = 2$$

$$\lim_{n \rightarrow \infty} \frac{g(n)}{f(n)} = \begin{cases} c & \text{implies } g(n) \in \Theta(f(n)) \text{ if } c > 0 \\ 0 & \text{implies } g(n) \in o(f(n)) \\ \infty & \text{implies } f(n) \in o(g(n)) \end{cases}$$

Definition

A **binary search tree** is a binary tree of items (ordinarily called keys), that come from an ordered set, such that

1. Each node contains one key.
2. The keys in the left subtree of a given node are less than or equal to the key in that node.
3. The keys in the right subtree of a given node are greater than or equal to the key in that node.

22. Group the following functions by complexity category.

$n \ln n$ $(\lg n)^2$ $5n^2 + 7n$ $n^{5/2}$
 $n!$ $2^{n!}$ 4^n n^n $n^n + \ln n$
 $5^{\lg n}$ $\lg(n!)$ $(\lg n)!$ \sqrt{n} e^n $8n + 12$ $10^n + n^{20}$

$$n \ln n \in O(n \log n)$$

$$(\lg n)^2 \in O(\log^2 n)$$

3) Show that $\log(n!) \in \Omega(n \log n)$

$$\log(n!) = \log(n) + \log(n-1) + \dots + \log(n/2)$$

$$\geq \log(n) + \log(n-1) + \dots + \log(n/2) \leftarrow n/2 \text{ terms}$$

$$\geq \log\left(\frac{n}{2}\right) + \log\left(\frac{n}{2}\right) + \dots + \log\left(\frac{n}{2}\right)$$

$$= \frac{n}{2} \log \frac{n}{2}$$

$$= \frac{n}{2} (\log n - \log 2)$$

$$\log(n!) = \frac{n}{2} (\log n - 1)$$

$$\frac{1}{2} (n \log n - n)$$

$$\left(\frac{3}{2}\right)^n \text{ vs } 2^n$$