

Modified Grammar Group 6

1. **<program>** ==> <otherFunctions> <mainFunction>

FIRST: {TK_MAIN, TK_FUNID}

FOLLOW: {\$}

2. <mainFunction>==> TK_MAIN <stmts> TK_END

FIRST: {TK_MAIN}

FOLLOW: {\$}

3. <otherFunctions>==> <function><otherFunctions> | ∈

FIRST: {TK_FUNID, ∈}

FOLLOW: {TK_MAIN}

4. <function>==>TK_FUNID <input_par> <output_par> TK_SEM <stmts>
TK_END

FIRST: {TK_FUNID}

FOLLOW: {TK_FUNID, TK_MAIN}

5. <input_par>==>TK_INPUT TK_PARAMETER TK_LIST TK_SQL
<parameter_list> TK_SQR

FIRST: {TK_INPUT}

FOLLOW: {TK_OUTPUT, TK_SEM}

6. $\langle \text{output_par} \rangle \implies \text{TK_OUTPUT TK_PARAMETER TK_LIST TK_SQL}$
 $\langle \text{parameter_list} \rangle \text{ TK_SQR} \mid \epsilon$

FIRST: {TK_OUTPUT, ϵ }

FOLLOW: {TK_SEM}

7. $\langle \text{parameter_list} \rangle \implies \langle \text{dataType} \rangle \text{ TK_ID } \langle \text{remaining_list} \rangle$

FIRST: {TK_INT, TK_REAL, TK_RECORD, TK_UNION, TK_RUID}

FOLLOW: {TK_SQR, }

8. $\langle \text{dataType} \rangle \implies \langle \text{primitiveDatatype} \rangle \mid \langle \text{constructedDatatype} \rangle$

FIRST: {TK_INT, TK_REAL, TK_RECORD, TK_UNION, TK_RUID}

FOLLOW: {TK_ID, TK_COLON}

9. $\langle \text{primitiveDatatype} \rangle \implies \text{TK_INT} \mid \text{TK_REAL}$

FIRST: {TK_INT, TK_REAL}

FOLLOW: {TK_ID, TK_COLON}

10. $\langle \text{constructedDatatype} \rangle \implies \text{TK_RECORD TK_RUID} \mid \text{TK_UNION TK_RUID} \mid \text{TK_RUID}$

FIRST: {TK_RECORD, TK_UNION, TK_RUID}

FOLLOW: {TK_ID, TK_COLON}

FIRST(TK_RECORD) = {TK_RECORD}, FIRST(TK_UNION) = {TK_UNION},
FIRST(TK_RUID) = {TK_RUID}.

This shows that FIRST(TK_RECORD), FIRST(TK_UNION), FIRST(TK_RUID) are all disjoint.

This implies that this rule is in LL(1).

11. $\langle \text{remaining_list} \rangle \Rightarrow \text{TK_COMMA } \langle \text{parameter_list} \rangle \mid \epsilon$

FIRST: {TK_COMMA, ϵ }

FOLLOW: {TK_SQR}

12. $\langle \text{stmts} \rangle \Rightarrow \langle \text{typeDefinitions} \rangle \langle \text{declarations} \rangle \langle \text{otherStmts} \rangle \langle \text{returnStmt} \rangle$

FIRST: {TK_TYPE}

FOLLOW: {TK_END}

13. $\langle \text{typeDefinitions} \rangle \Rightarrow \langle \text{typeDefinition} \rangle \langle \text{typeDefinitions} \rangle \mid \langle \text{definetypestmt} \rangle \langle \text{typeDefinitions} \rangle \mid \epsilon$

FIRST: {TK_RECORD, TK_UNION, TK_DEFINETYPE}

FOLLOW: {TK_TYPE}

$\text{FIRST}(\langle \text{typeDefinition} \rangle) = \{\text{TK_RECORD}, \text{TK_UNION}\},$

$\text{FIRST}(\langle \text{definetypestmt} \rangle) = \{\text{TK_DEFINETYPE}\}, \text{FIRST}(\epsilon) = \{\epsilon\},$

$\text{FOLLOW}(\langle \text{typeDefinitions} \rangle) = \{\text{TK_TYPE}\}.$

This shows that $\text{FIRST}(\langle \text{typeDefinition} \rangle)$, $\text{FIRST}(\langle \text{definetypestmt} \rangle)$, $\text{FIRST}(\epsilon)$ are all disjoint, $\text{FOLLOW}(\langle \text{typeDefinitions} \rangle) \cap \text{FIRST}(\langle \text{typeDefinition} \rangle) = \emptyset$ and $\text{FOLLOW}(\langle \text{typeDefinitions} \rangle) \cap \text{FIRST}(\langle \text{typeDefinition} \rangle) = \emptyset$.

This implies that this rule is in LL(1).

14. $\langle \text{typeDefinition} \rangle \Rightarrow \text{TK_RECORD TK_RUID } \langle \text{fieldDefinitions} \rangle \text{TK_ENDRECORD}$

FIRST: {TK_RECORD, TK_UNION}

FOLLOW: {TK_RECORD, TK_UNION, TK_DEFINETYPE, TK_TYPE}

15. $\langle \text{typeDefinition} \rangle \Rightarrow \text{TK_UNION TK_RUID } \langle \text{fieldDefinitions} \rangle \text{ TK_ENDUNION}$

FIRST: {TK_RECORD, TK_UNION}

FOLLOW: {TK_RECORD, TK_UNION, TK_DEFINETYPE, TK_TYPE}

16. $\langle \text{fieldDefinitions} \rangle \Rightarrow \langle \text{fieldDefinition} \rangle \langle \text{fieldDefinition} \rangle \langle \text{moreFields} \rangle$

FIRST: {TK_TYPE, TK_RECORD, TK_UNION}

FOLLOW: {TK_ENDRECORD, TK_ENDUNION}

17. $\langle \text{fieldDefinition} \rangle \Rightarrow \text{TK_TYPE } \langle \text{dataType} \rangle \text{TK_COLON TK_FIELDID TK_SEM } | \langle \text{typeDefinition} \rangle \langle \text{fieldDefinition} \rangle$

FIRST: {TK_TYPE, TK_RECORD, TK_UNION}

FOLLOW: {TK_TYPE, TK_RECORD, TK_UNION, TK_ENDRECORD, TK_ENDUNION}

$\text{FIRST}(\text{TK_TYPE}) = \{\text{TK_TYPE}\}$, $\text{FIRST}(\langle \text{typeDefinition} \rangle) = \{\text{TK_RECORD}, \text{TK_UNION}\}$.

This shows that $\text{FIRST}(\langle \text{typeDefinition} \rangle)$ and $\text{FIRST}(\text{TK_TYPE})$ are disjoint.

This implies that this rule is in LL(1).

18. $\langle \text{moreFields} \rangle \Rightarrow \langle \text{fieldDefinition} \rangle \langle \text{moreFields} \rangle | \epsilon$

FIRST: {TK_TYPE, TK_RECORD, TK_UNION, ϵ }

FOLLOW: {TK_ENDRECORD, TK_ENDUNION}

19. $\langle \text{declarations} \rangle \Rightarrow \langle \text{declaration} \rangle \langle \text{declarations} \rangle | \epsilon$

FIRST: {TK_TYPE, ϵ }

FOLLOW: {TK_ID, TK_WHILE, TK_IF, TK_READ, TK_WRITE, TK_SQL, TK_RETURN, TK_CALL}

20. <declaration>====> TK_TYPE <dataType> TK_COLON TK_ID <global_or_not> TK_SEM

FIRST: {TK_TYPE}

FOLLOW: {TK_TYPE, TK_ID, TK_WHILE, TK_IF, TK_READ, TK_WRITE, TK_SQL, TK_RETURN, TK_CALL}

FIRST(TK_TYPE) = {TK_TYPE}.

As there is just one production this rule is already in LL(1).

21. <global_or_not>====>TK_COLON TK_GLOBAL| ϵ

FIRST: {TK_COLON, ϵ }

FOLLOW: {TK_SEM}

FIRST(TK_COLON) = {TK_COLON}, FIRST(ϵ) = { ϵ },
FOLLOW(<global_or_not>) = {TK_SEM}.

This shows that FIRST(<TK_COLON>) and FIRST(ϵ) are disjoint, and
FOLLOW(<global_or_not>) \cap FIRST(TK_COLON) = ϕ .

This implies that this rule is in LL(1).

22. <otherStmts>====> <stmt><otherStmts> | ϵ

FIRST: {TK_ID, TK_WHILE, TK_IF, TK_READ, TK_WRITE, TK_SQL, ϵ }

FOLLOW: {TK_RETURN, TK_ENDIF, TK_ELSE, TK_ENDWHILE}

23. <stmt>====> <assignmentStmt> | <iterativeStmt>|<conditionalStmt>|<ioStmt>|
<funCallStmt>

FIRST: {TK_ID, TK_WHILE, TK_IF, TK_READ, TK_WRITE, TK_SQL, ϵ }

FOLLOW: {TK_ID, TK_WHILE, TK_IF, TK_READ, TK_WRITE, TK_SQL, TK_ENDWHILE, TK_ELSE, TK_ENDIF, TK_RETURN, TK_CALL }

24. $\langle \text{assignmentStmt} \rangle \implies \langle \text{singleOrRecId} \rangle \text{ TK_ASSIGNOP}$
 $\langle \text{arithmeticExpression} \rangle \text{ TK_SEM}$

FIRST: {TK_ID}

FOLLOW: {TK_ID, TK_WHILE, TK_IF, TK_READ, TK_WRITE, TK_SQL, TK_ENDWHILE, TK_ELSE, TK_ENDIF, TK_RETURN, TK_CALL }

25. $\langle \text{singleOrRecId} \rangle \implies \text{TK_ID} \langle \text{nestedDot} \rangle$

FIRST: {TK_ID}

FOLLOW: {TK_ASSIGNOP, TK_CL, TK_MUL, TK_DIV, TK_ID, TK_NUM, TK_RNUM, TK_OP}

FIRST(TK_ID) = {TK_ID}.

As there is just one production this rule is already in LL(1).

26. $\langle \text{nestedDot} \rangle \implies \text{TK_DOT} \text{ TK_FIELDID} \mid \epsilon$

FIRST: {TK_DOT, ϵ }

FOLLOW: {TK_ASSIGNOP, TK_CL, TK_MUL, TK_DIV, TK_ID, TK_NUM, TK_RNUM, TK_OP}

FIRST(TK_DOT) = {TK_DOT}, FIRST(ϵ) = { ϵ }, FOLLOW($\langle \text{nestedDot} \rangle$) = {TK_ASSIGNOP, TK_CL, TK_MUL, TK_DIV, TK_ID, TK_NUM, TK_RNUM, TK_OP}.

This shows that FIRST(TK_DOT) and FIRST(ϵ) are disjoint, and FOLLOW($\langle \text{nestedDot} \rangle$) \cap FIRST(TK_DOT) = ϕ .

This implies that this rule is in LL(1).

27. $\langle \text{funCallStmt} \rangle \Rightarrow \langle \text{outputParameters} \rangle \text{ TK_CALL TK_FUNID TK_WITH TK_PARAMETERS } \langle \text{inputParameters} \rangle \text{ TK_SEM}$

FIRST: {TK_SQL, TK_CALL}

FOLLOW: {TK_ID, TK_WHILE, TK_IF, TK_READ, TK_WRITE, TK_SQL, TK_ENDWHILE, TK_ELSE, TK_ENDIF, TK_RETURN, TK_CALL },

FIRST($\langle \text{outputParameters} \rangle$) = {TK_SQL, ϵ }

As there is just one production this rule is already in LL(1).

28. $\langle \text{outputParameters} \rangle \Rightarrow \text{TK_SQL } \langle \text{idList} \rangle \text{ TK_SQR TK_ASSIGNOP } | \epsilon$

FIRST: {TK_SQL, ϵ }

FOLLOW: {TK_CALL}

29. $\langle \text{inputParameters} \rangle \Rightarrow \text{TK_SQL } \langle \text{idList} \rangle \text{ TK_SQR}$

FIRST: {TK_SQL}

FOLLOW: {TK_SEM}

30. $\langle \text{iterativeStmt} \rangle \Rightarrow \text{TK_WHILE TK_OP } \langle \text{booleanExpression} \rangle \text{ TK_CL } \langle \text{stmt} \rangle \langle \text{otherStmts} \rangle \text{ TK_ENDWHILE}$

FIRST: {TK_WHILE}

FOLLOW: {TK_ID, TK_WHILE, TK_IF, TK_READ, TK_WRITE, TK_SQL, TK_ENDWHILE, TK_ELSE, TK_ENDIF, TK_RETURN, TK_CALL }

31. $\langle \text{conditionalStmt} \rangle \Rightarrow \text{TK_IF TK_OP } \langle \text{booleanExpression} \rangle \text{ TK_CL TK_THEN } \langle \text{stmt} \rangle \langle \text{otherStmts} \rangle \langle \text{optionalElse} \rangle \text{ TK_ENDIF}$

FIRST: {TK_IF}

FOLLOW: {TK_ID, TK_WHILE, TK_IF, TK_READ, TK_WRITE, TK_SQL, TK_ENDWHILE, TK_ELSE, TK_ENDIF, TK_RETURN, TK_CALL }

FIRST(TK_IF) = {TK_IF}.

As there is just one production this rule is already in LL(1).

32. <optionalElse>====> TK_ELSE <stmt><otherStmts> | ϵ

FIRST: {TK_ELSE, ϵ }

FOLLOW: {TK_ENDIF}

FIRST(TK_ELSE) = {TK_MUL, TK_DIV}, FIRST(ϵ) = { ϵ },
FOLLOW(<optionalElse>) = {TK_ENDIF}.

This shows that FIRST(TK_ELSE) and FIRST(ϵ) are disjoint, and
FOLLOW(<optionalElse>) \cap FIRST(TK_ELSE) = ϕ .

This implies that this rule is in LL(1).

33. <ioStmt>====>TK_READ TK_OP <var> TK_CL TK_SEM | TK_WRITE TK_OP
<var> TK_CL TK_SEM

FIRST: {TK_READ, TK_WRITE}

FOLLOW: {TK_ID, TK_WHILE, TK_IF, TK_READ, TK_WRITE, TK_SQL, TK_ENDWHILE, TK_ELSE, TK_ENDIF, TK_RETURN, TK_CALL }

34. <arithmeticExpression> =====> <term><N1>

FIRST: {TK_ID, TK_NUM, TK_RNUM, TK_OP}

FOLLOW: {TK_CL, TK_SEM}

As there is just one production this rule is already in LL(1).

35. $\langle N1 \rangle \implies \langle \text{operator1} \rangle \langle \text{term} \rangle \langle N1 \rangle \mid \epsilon$

FIRST: {TK_PLUS, TK_MINUS, ϵ }

FOLLOW: {TK_CL, TK_SEM}

FIRST($\langle \text{operator1} \rangle$) = {TK_PLUS, TK_MINUS}, FIRST(ϵ) = { ϵ },
FOLLOW($\langle N1 \rangle$) = {TK_CL, TK_SEM}.

This shows that FIRST($\langle \text{operator1} \rangle$) and FIRST(ϵ) are disjoint, and
FOLLOW($\langle N1 \rangle$) \cap FIRST($\langle \text{operator1} \rangle$) = ϕ .

This implies that this rule is in LL(1).

36. $\langle \text{term} \rangle \implies \langle \text{factor} \rangle \langle N2 \rangle$

FIRST: {TK_ID, TK_NUM, TK_RNUM, TK_OP}

FOLLOW: {TK_CL, TK_SEM, TK_PLUS, TK_MINUS}

As there is just one production this rule is already in LL(1).

37. $\langle \text{factor} \rangle \implies \langle \text{var} \rangle \mid \text{TK_OP } \langle \text{arithmeticExpression} \rangle \text{TK_CL}$

FIRST: {TK_ID, TK_NUM, TK_RNUM, TK_OP}

FOLLOW: {TK_MUL, TK_DIV, TK_CL, TK_PLUS, TK_MINUS, TK_SEM}

FIRST($\langle \text{var} \rangle$) = {TK_ID, TK_NUM, TK_RNUM}, FIRST(TK_OP) = {TK_OP}.

This shows that FIRST($\langle \text{var} \rangle$) and FIRST(TK_OP) are disjoint

This implies that this rule is in LL(1).

38. $\langle N2 \rangle \implies \langle \text{operator2} \rangle \langle \text{factor} \rangle \langle N2 \rangle \mid \epsilon$

FIRST: {TK_MUL, TK_DIV, ϵ }

FOLLOW: {TK_CL, TK_SEM, TK_PLUS, TK_MINUS}

$\text{FIRST}(\langle \text{operator2} \rangle) = \{\text{TK_MUL}, \text{TK_DIV}\}$, $\text{FIRST}(\epsilon) = \{\epsilon\}$, $\text{FOLLOW}(\langle \text{N2} \rangle) = \{\text{TK_CL}, \text{TK_SEM}, \text{TK_PLUS}, \text{TK_MINUS}\}$.

This shows that $\text{FIRST}(\langle \text{operator2} \rangle)$ and $\text{FIRST}(\epsilon)$ are disjoint, and $\text{FOLLOW}(\langle \text{N2} \rangle) \cap \text{FIRST}(\langle \text{operator2} \rangle) = \phi$.

This implies that this rule is in $\text{LL}(1)$.

39. $\langle \text{operator1} \rangle \implies \text{TK_PLUS} \mid \text{TK_MINUS}$

FIRST: $\{\text{TK_PLUS}, \text{TK_MINUS}\}$

FOLLOW: $\{\text{TK_OP}, \text{TK_ID}, \text{TK_NUM}, \text{TK_RNUM}\}$

$\text{FIRST}(\text{TK_PLUS}) = \{\text{TK_PLUS}\}$ and $\text{FIRST}(\text{TK_MINUS}) = \{\text{TK_MINUS}\}$.

This shows that $\text{FIRST}(\text{TK_PLUS})$ and $\text{FIRST}(\text{TK_MINUS})$ are disjoint.

This implies that this rule is in $\text{LL}(1)$.

40. $\langle \text{operator2} \rangle \implies \text{TK_MUL} \mid \text{TK_DIV}$

FIRST: $\{\text{TK_MUL}, \text{TK_DIV}\}$

FOLLOW: $\{\text{TK_OP}, \text{TK_ID}, \text{TK_NUM}, \text{TK_RNUM}\}$

$\text{FIRST}(\text{TK_MUL}) = \{\text{TK_MUL}\}$ and $\text{FIRST}(\text{TK_DIV}) = \{\text{TK_DIV}\}$.

This shows that $\text{FIRST}(\text{TK_MUL})$ and $\text{FIRST}(\text{TK_DIV})$ are disjoint.

This implies that this rule is in $\text{LL}(1)$.

41. $\langle \text{booleanExpression} \rangle \implies \text{TK_OP} \langle \text{booleanExpression} \rangle \text{TK_CL} \mid \text{TK_OP} \langle \text{booleanExpression} \rangle \text{TK_CL}$

FIRST: $\{\text{TK_OP}\}$

FOLLOW: $\{\text{TK_CL}\}$

42. $\langle \text{booleanExpression} \rangle \implies \langle \text{var} \rangle \langle \text{relationalOp} \rangle \langle \text{var} \rangle$

FIRST: {TK_ID, TK_NUM, TK_RNUM}

FOLLOW: {TK_CL}

43. $\langle \text{booleanExpression} \rangle \implies \text{TK_NOT TK_OP} \langle \text{booleanExpression} \rangle \text{TK_CL}$

FIRST: {TK_NOT}

FOLLOW: {TK_CL}

FIRST(TK_NOT) = {TK_NOT}, FIRST($\langle \text{var} \rangle$) = {TK_ID, TK_NUM, TK_RNUM}
and FIRST(TK_OP) = {TK_OP}.

This shows that FIRST(TK_NOT), FIRST($\langle \text{var} \rangle$) and FIRST(TK_OP) are all disjoint.

This implies that this rule is in LL(1).

44. $\langle \text{var} \rangle \implies \langle \text{singleOrRecId} \rangle \mid \text{TK_NUM} \mid \text{TK_RNUM}$

FIRST: {TK_ID, TK_NUM, TK_RNUM}

FOLLOW: {TK_MUL, TK_DIV, TK_CL, TK_PLUS, TK_MINUS, TK_SEM, TK_LT, TK_LE, TK_EQ, TK_GT, TK_GE, TK_NE }

FIRST($\langle \text{singleOrRecId} \rangle$) = {TK_ID}, FIRST(TK_NUM) = {TK_NUM},
FIRST(TK_RNUM) = {TK_RNUM}, FOLLOW: {TK_MUL, TK_DIV, TK_CL, TK_PLUS, TK_MINUS, TK_SEM, TK_LT, TK_LE, TK_EQ, TK_GT, TK_GE, TK_NE }.

This shows that FIRST($\langle \text{singleOrRecId} \rangle$), FIRST(TK_NUM), FIRST(TK_RNUM) are all disjoint, FOLLOW($\langle \text{var} \rangle$) \cap FIRST(TK_NUM) = ϕ and FOLLOW($\langle \text{var} \rangle$) \cap FIRST(TK_RNUM) = ϕ .

This implies that this rule is in LL(1).

45. <logicalOp>====>TK_AND | TK_OR

FIRST: {TK_AND, TK_OR}

FOLLOW: {TK_OP}

46. <relationalOp>====> TK_LT | TK_LE | TK_EQ |TK_GT | TK_GE | TK_NE

FIRST: {TK_LT, TK_LE, TK_EQ, TK_GT, TK_GE, TK_NE}

FOLLOW: {TK_ID, TK_NUM, TK_RNUM}

47. <returnStmt>====>TK_RETURN <optionalReturn> TK_SEM

FIRST: {TK_RETURN}

FOLLOW: {TK_END}

48. <optionalReturn>====>TK_SQL <idList> TK_SQR | ∈

FIRST: {TK_SQL, ∈}

FOLLOW: {TK_SEM}

49. <idList>====> TK_ID <more_ids>

FIRST: {TK_ID}

FOLLOW: {TK_SQR}

50. <more_ids>====> TK_COMMA <idList> | ∈

FIRST: {TK_COMMA, ∈}

FOLLOW: {TK_SQR}

51. <definetypstmt>==>TK_DEFINETYPE <A> TK_RUID TK_AS TK_RUID

FIRST: {TK_DEFINETYPE}

FOLLOW: {TK_RECORD, TK_UNION}

52. <A>==>TK_RECORD | TK_UNION

FIRST: {TK_RECORD, TK_UNION}

FOLLOW: {TK_RUID}