

# The Definition of Snail Programming Language

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## 1 Snail の構文定義

EBNF 記法を用いて Snail の具象構文を以下に示す.

$$\begin{aligned} \text{toplevel} ::= & \text{let } [\text{rec}] \text{ var } \{ \text{var } [ : \langle \text{type} \rangle ] \} : \langle \text{type} \rangle = \langle \text{term} \rangle \{ \langle \text{mutual-recursion-top-let} \rangle \} \\ & | \text{typedef cons } \{ \text{var} \} = [ | ] \{ \langle \text{type-dec} \rangle | \} \langle \text{type-dec} \rangle \{ \langle \text{mutual-recursion-type} \rangle \} \end{aligned}$$
$$\text{mutual-recursion-type} ::= \text{and cons } \{ \text{var} \} = [ | ] \{ \langle \text{type-dec} \rangle | \} \langle \text{type-dec} \rangle$$
$$\text{mutual-recursion-top-let} ::= \text{and var } \{ \text{var } [ : \langle \text{type} \rangle ] \} : \langle \text{type} \rangle = \langle \text{term} \rangle$$
$$\text{type-dec} ::= \text{cons } [\text{of } \langle \text{type} \rangle]$$
$$\begin{aligned} \text{type} ::= & \langle \text{type} \rangle \rightarrow \langle \text{type} \rangle \\ & | ! ' [ \langle \text{expmod} \rangle ' ] ' \{ ' \langle \text{type} \rangle ' \} ' \\ & | \langle \text{simple-type} \rangle \\ & | \langle \text{type} \rangle \langle \text{simple-type} \rangle \end{aligned}$$
$$\begin{aligned} \text{expmod} ::= & \text{int} \\ & | \infty \end{aligned}$$
$$\begin{aligned} \text{simple-type} ::= & ' ( ' \langle \text{type} \rangle ' ) ' \\ & | \text{var} \\ & | \text{cons} \\ & | () \end{aligned}$$
$$\begin{aligned} \text{pattern} ::= & \langle \text{simple-pattern} \rangle \\ & | \langle \text{pattern} \rangle \langle \text{simple-pattern} \rangle \\ & | \langle \text{simple-pattern} \rangle \text{binop} \langle \text{simple-pattern} \rangle \end{aligned}$$

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simple-pattern ::= '(' <pattern> ')'
                |  var
                |  cons '[' <simple-pattern> ']'
                |  []
                |  -

```

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mutual-recursion-let ::= and var {var [ : <type> ]} : <type> = <term>

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term ::= <simple-term>
        |  <term> <simple-term>
        |  let [rec] var {var [ : <type> ]} : <type> = <term> {<mutual-recursion-let>} in <term>
        |  fun {var [ : <type> ]} → <term>
        |  match <term> with [ | ] {<pattern> → <term> | } <pattern> → <term>
        |  if <term> then <term> else <term>

```

```

simple-term ::= '(' <term> [ : <type> ] ')'
                |  ! <term>
                |  int
                |  float
                |  string
                |  bool
                |  var
                |  cons [<simple-term>]
                |  ()
                |  []
                |  list

```

終端記号の意味を以下のように定義する.

- *var* 先頭が小文字で始まる文字列.
- *cons* 先頭が大文字で始まる文字列.
- *list* 組み込みリストの構文糖衣, [1,2,3] など.
- *string* 文字列リテラル.
- *int* 整数リテラル.
- *float* 小数リテラル.
- *bool* 真偽値リテラル.
- その他 予約語.

## 2 Snail の型システム

Snail は次のような型付け規則を持つ.

$\frac{}{\vdash \text{int} : \text{Int}}$	(INT)	$\frac{\Gamma, x : [A]_r \vdash e : B}{\Gamma \vdash \text{fun } (!x : !_r A) \rightarrow e : !_r A \multimap B}$ (FUN-EXP)
$\frac{}{\vdash \text{float} : \text{Float}}$	(FLOAT)	$\frac{\Gamma \vdash e : A \multimap B \quad \Delta \vdash e' : A}{\Gamma + \Delta \vdash e \ e' : B}$ (APP)
$\frac{}{\vdash \text{string} : \text{String}}$	(STRING)	$\frac{\Gamma \vdash e : \text{Bool} \quad \Delta \vdash e_1 : A \quad \Delta \vdash e_2 : A}{\Gamma + \Delta \vdash \text{if } e \text{ then } e_1 \text{ else } e_2 : A}$ (IF)
$\frac{}{\vdash \text{bool} : \text{Bool}}$	(BOOL)	$\frac{\Gamma \vdash e : A \quad \Delta, x : A \vdash e' : B}{\Gamma + \Delta \vdash \text{let } x = e \text{ in } e' : B}$ (LET)
$\frac{}{x : A \vdash x : A}$	(ID)	$\frac{\Gamma \vdash e : !_r A \quad \Delta, x : [A]_r \vdash e' : B}{\Gamma + \Delta \vdash \text{let } !x = e \text{ in } e' : B}$ (LET-EXP)
$\frac{\Gamma, x : A \vdash e : B}{\Gamma, x : [A]_1 \vdash e : B}$	(DER)	$\frac{[\Gamma], x : [A]_p \vdash e : A \quad \Delta, x : [A]_\infty \vdash e' : B}{\infty * [\Gamma] + \Delta \vdash \text{let rec } x = e \text{ in } e' : B}$ (LET-REC)
$\frac{[\Gamma] \vdash e : B}{r * [\Gamma] \vdash !e : !_r B}$	(PR)	
$\frac{\Gamma, x : A \vdash e : B}{\Gamma \vdash \text{fun } x \rightarrow e : A \multimap B}$	(FUN)	$\frac{\Delta \vdash e : B \quad \Gamma <: \Delta}{\Gamma, \Theta \vdash e : B}$ (SUB)