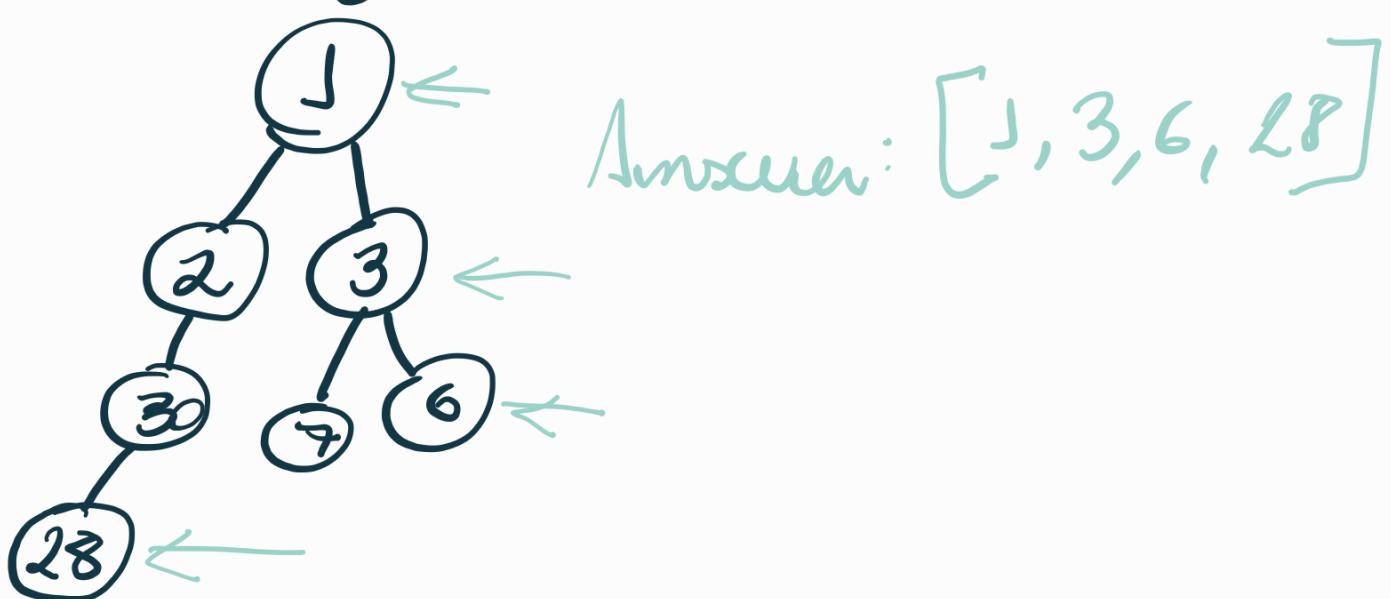
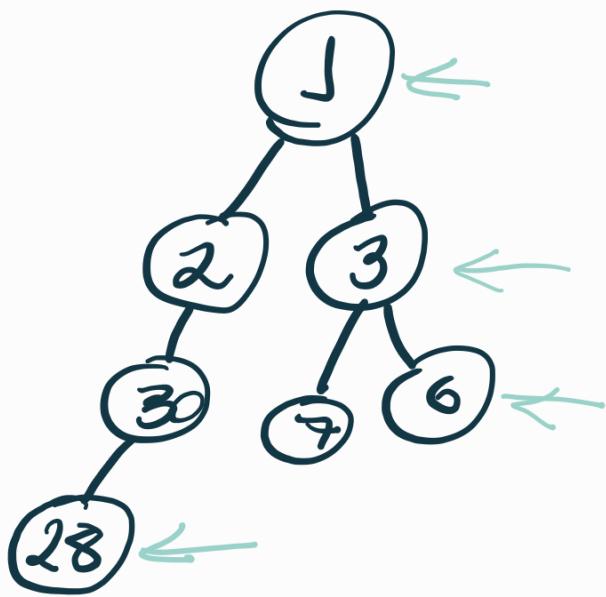


- o) given the root of a BT  
1, list only the values of  
the nodes that are on  
the outer right side of it  
e.g.:



obs: final list has to be ordered  
top to bottom.

- o) Root also counts

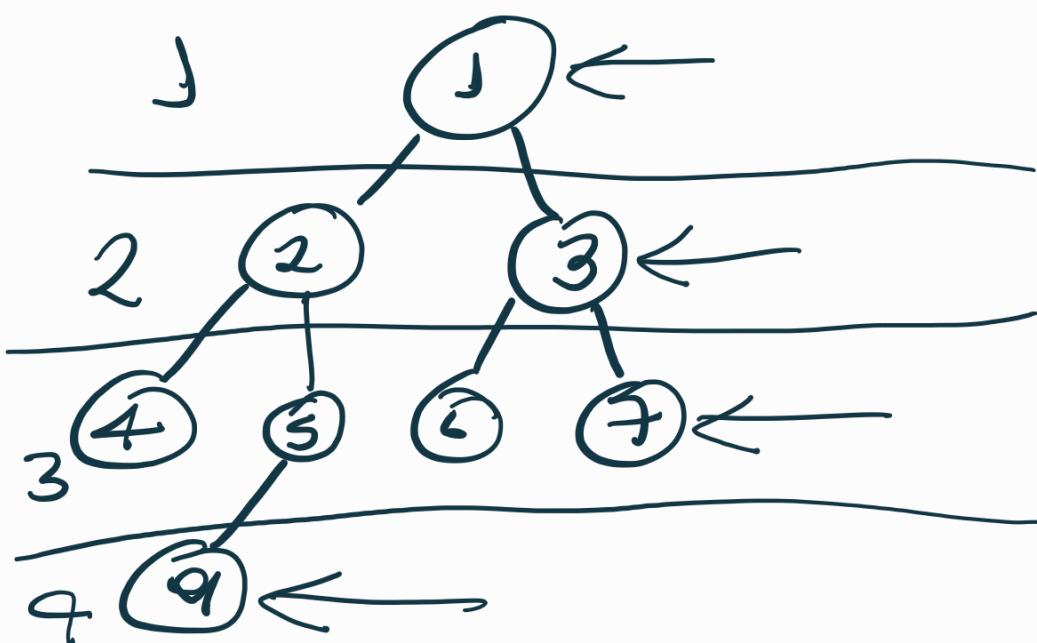


→ As right-most items matter the most, focus them first before checking left

→ Process right sides and add them to the list

→ Every recursion should know the number of its current layer.

↳ So as you do a deep dive pass it as an incremental parameter.



→ each level can have only one valid value on the result list.

→ given that layer 0 (root) will always add one element to the result on the following recursion

↳ if the current size of my result array equals my current layer, then i know the current layer isn't added an value yet.

↳ and thus has to be added.

→ for this to work root (source) has to be handled as layer 0 instead of layer 1.

→ with this approach, if left side of root has more layers than right, then they will all be accounted for.

↳

↳ As result will only account for layers that haven't had values accounted until its traversal.



Alternate Solution:



→ I could also think of it in terms of layers (levels)

1. → each layer has an right-most node

→ Create an storage  $\langle \mathcal{X}, \mathcal{V} \rangle$

1. Run through the whole right side and add values to its resp.-  
pective layers.

→ Obs: if a layer already has

an value in storage then  
You can't replace it

→ After visiting the entire  
right side, do visit left  
side and add the extra  
layers that were not  
present on the right

b, finally after having  
traversed it all

c, iterate the  $\langle K, V \rangle$  items  
in storage and create

an list from it.

L, then between that list.