

# Test 2 Study Guide

- Master Theorem

- $T(n) = 2T(\frac{n}{2}) + 1$ 
  - $a = 2, b = 2, d = 0$
  - $T(n) \in \Theta(n)$
- $T(n) = 2T(\frac{n}{2}) + n$ 
  - $a = 2, b = 2, d = 1$
  - $T(n) \in \Theta(n \log n)$
- $T(n) = 7T(\frac{n}{2}) + n^2$ 
  - $a = 7, b = 2, d = 2$
  - $T(n) \in \Theta(n^{\log_2 7})$
- $T(n) = 2T(\frac{n}{3}) + n$ 
  - $a = 2, b = 3, d = 1$
  - $T(n) \in \Theta(n)$
- $T(n) = 8T(\frac{n}{2}) + n^2$ 
  - $a = 8, b = 2, d = 2$
  - $T(n) \in \Theta(n^3)$

- Algorithms

- Given an integer array of size  $n$ , find the maximum value of the sum of elements of a subarray where no two elements of the subarray are adjacent to each other. Design a  $\Theta(n)$  algorithm using a decrease and conquer approach. Your algorithm just needs to find the maximum value of the sum. It does not need to identify the elements of the subarray.
  - Size 1 is just the element, size 2 is just the max between the two elements
  - If we know how to solve size  $k$ , how do we solve size  $k+1$ ?
    - Three cases
      - If the new element is not part of the solution. the solution for  $k$  and  $k+1$  will be the same
      - If it is included in the solution, the  $k$ th element cannot be. So the optimal solution will be the optimal solution for  $k-1$  plus the element at  $k+1$
      - If we have negative numbers,  $A[k+1]$  might be the max just by itself

- The total solution then is  $\max\{v_1, v_0 + A[k + 1], A[k + 1]\}$
- Alg MaxNonAdjacent ( $A[0 \dots n-1]$ )
  - $v_0 = A[0]$
  - $v_1 = \max(v_0, A[1])$
  - for  $i = 2$  to  $n-1$ 
    - $v = \max\{v_1, v_0 + A[i], A[i]\}$
  - return  $v$
- Design an algorithm to compute the integer part of the square root of a non-negative integer  $n$ . For example, for input  $n = 5$ , the output is 2. You can only use the four basic algebraic operations (+, −, ×, /) and the floor() operation, i.e., rounded down. Your algorithm should have a time complexity of  $O(\log n)$ .
  - alg sqrt( $n$ )
    - $l = 0$
    - $r = n$
    - while true
      - $x = \text{floor}(\frac{l+r}{2})$
      - if  $x*x \leq n$ 
        - if  $(x+1)*(x+1) \geq n$ 
          - return  $x$
        - else
          - $left = x+1$
      - else
        - $r = x$
  - Given an array  $A[0 \dots n - 1]$  of  $n$  distinctive integers, reorder it such that  $A[0] < A[1] > A[2] < A[3] \dots$ . • Design an  $O(n \log n)$  time and  $\Theta(1)$  space algorithm. Partial credits are given to solutions that produce a correct order but do not meet the time or space complexity requirements. • (3 Bonus Points) Design a  $\Theta(n)$  time and  $\Theta(1)$  space algorithm.
    - Presort then swap elements 1 and 2, 3 and 4, and so on, that will achieve  $n \log n$
    - For linear
      - Alg WiggleSort( $A$ )
        - $up = \text{true}$
        - for  $i = 0$  to  $A.\text{len} - 2$ 
          - if  $up$

- if  $A[i] > A[i+1]$ 
  - $\text{swap}(A[i], A[i+1])$
- $\text{up} = \text{false}$
- else
  - if  $A[i] < A[i+1]$ 
    - $\text{swap}(A[i], A[i+1])$
  - $\text{up} = \text{true}$
- OR just use quickselect to find median, then do l r l r l r...