Nano: Interpreters

CSE 130 Week ???

The plan

- 1. Interpreters
- 2. Hw4 concepts

The plan

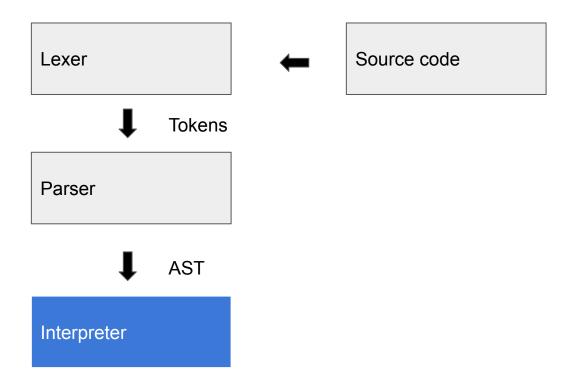
- 1. Interpreters
- 2. Hw4 concepts
 - a. Environments
 - b. Closures
 - c. Native ops

Interpreters

Interpreters

An interpreter executes another program in some language without compilation

The big picture



A boring interpreter

```
Expr ::= IntLit Int | Add Expr Expr -- n | e1 + e2
eval :: Expr -> Int
eval e = ??
```

How do we implement eval?

A boring interpreter

```
Expr ::= IntLit Int | Add Expr Expr -- n | e1 + e2
eval :: Expr -> Int
eval (IntLit x) = x
eval (Add e1 e2) = (eval e1) + (eval e2)
```

Nano is not so simple...

```
data Expr
  = FInt Int
    EBool Bool
                             -- []
    ENil
  | EVar Id
                             -- X
    EBin Binop Expr Expr
    EIf Expr Expr Expr
                             -- If e1 then e2 else e3
    ELet Id Expr Expr
                             -- let x = e1 in e2
  | EApp Expr Expr
                             -- e1 e2
    ELam Id Expr
                             -- \x. e
```

What does it mean to "evaluate" an Expr?

```
eval :: ??
eval = ??
```

What does it mean to "evaluate" an Expr?

```
eval :: Env -> Expr -> ??
eval = ??
```

Our output type needs to be able to represent any possible result -- a boolean, a list, etc...

What does it mean to "evaluate" an Expr?

```
eval :: Env -> Expr -> Value
eval = ??
data Value
 = VTnt Tnt
  | VBool Bool
                            -- will discuss later
  | VClos Env Id Expr
                            -- []
   VNil
  | VCons Value Value
                        -- X:XS
  | VPrim (Value -> Value) -- will discuss later
```

Environments

How should we evaluate this:

let
$$x = 5$$
 in $x + x$

Environments

How should we evaluate this:

let
$$x = 5$$
 in $x + x$

We need to know the value of "x" while evaluating "x + x"

```
type Env = [(Id, Value)]
```

Environments

```
eval :: Env -> Expr -> Value eval env e = ??
```

You might need to update the environment when recursively evaluating subexpressions

Closures

Why closures?

```
let x = 1
in let foo = \n -> x + n
in let x = 2
in foo x
```

How should we evaluate this?

Why closures?

```
let x = 1
in let foo = \n -> x + n -- x = 1
in let x = 2
in foo x
```

Why closures?

```
let x = 1

in let foo = n -> x + n -- x = 1

in let x = 2 -- x = 2

in foo x -- x = 2, foo x = 1 + 2 = 3
```

When create a closures?

```
let x = 1
in let foo = \n -> x + n -- VClos (x + n) "n" [(x, 1)]
    in let x = 2
        in foo x
data Value
                       -- (VClos e "x" env) is
                       -- A function with argument "x"
  | VClos Env Id Expr
                       -- and body e that was defined
                       -- in an environment env
```

```
data Binop = ... | Cons

data ENil = ... | ENil

data Value = ... | VNil | VCons Value Value
```

```
data Binop = ... | Cons

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data Value = ... | VNil | VCons Value Value
```

Now, add support for "head" and "tail"...

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How might we do this?

We need to be able to define primitive function -- sort of a standard library

One function constructor:

VClos Env Id Expr

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One function constructor:

VClos Env Id Expr

This won't let us define "head" or "tail"!

The Expr type doesn't allow us to pattern match on the list -- no way to represent these functions in our Expr language

```
data Value = ... | VPrim (Value -> Value)
```

Now you can implement normal Haskell functions over Values and use them

Function application

```
let f = \x -> x + 1
in f 3
```

Function application

```
let f = \langle x - \rangle x + 1 - - [("f", VClos[]"x"("x" + 1))]
in f 3
```

How does one evaluate (f 3)?

Function application

```
let f = \langle x - \rangle x + 1 - - [("f", VClos [] "x" ("x" + 1))]
in f 3
```

Function application

```
let factorial = \n ->
    if n <= 0
        then 1
        else n * (factorial (n - 1))
in factorial 3</pre>
```