Consider the following languages for **DFA**, Consider ∑={a,b}.

A1: {w| w has at least two ‘a’}

A2: {w| w has at most two ‘b’}

Now construct a **DFA** which will recognize the language, A= A1 U A2, using the formal procedure of Closure.

[**Consider the states of the first machine are a1, a2, a3. And the states of the second machine are b1, b2, b3, b4]