

Graphs

Instructor: Dr. Sunho Lim (Ph.D., Assistant Professor)

Lecture 13

sunho.lim@ttu.edu

Adapted partially from Data Structures and Algorithms in C++, Adam Drozdek, 4th Edition, Cengage Learning; and Algorithms and Data Structures, Douglas Wilhelm Harder, Mmath

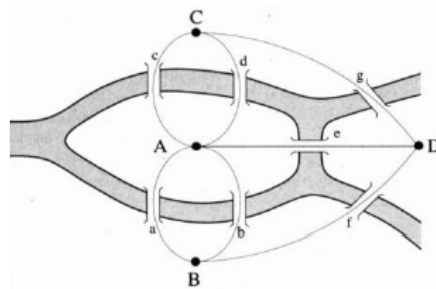
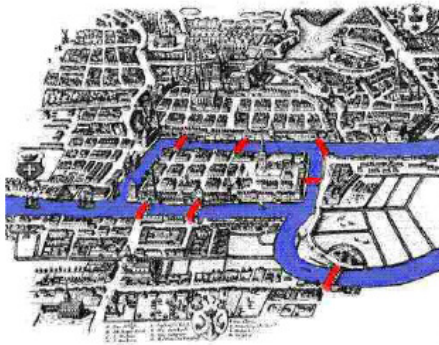
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Introduction

- Graph theory
 - start with Euler who was asked to find a nice path across the seven Königsberg bridges
 - the (Eulerian) path should cross over each of the **seven bridges exactly once**



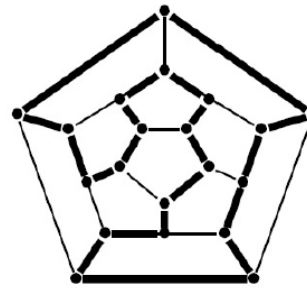
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Introduction (cont.)

- Another early bird,
 - Sir William Rowan Hamilton (1805-1865)
 - In 1859, developed a toy based on finding a path **visiting all cities** in a graph **exactly once** and sold it to a toy maker in Dublin
 - never was a big success...



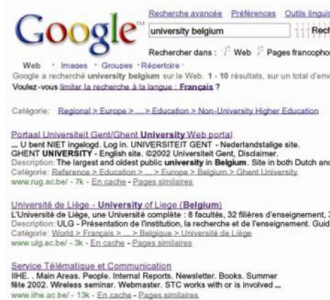
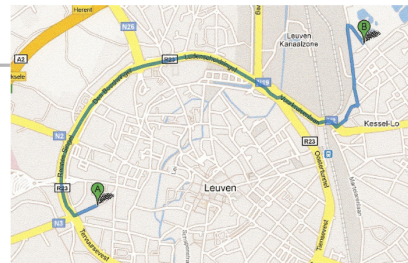
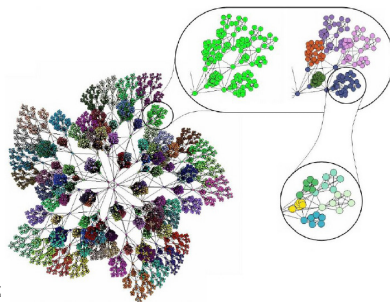
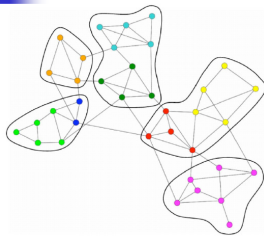
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Introduction (cont.)

- Now used in everywhere...
 - large scale problem, ranking, shortest path, etc.



CS.



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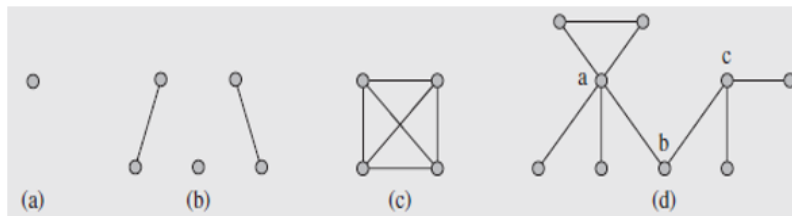
Introduction (cont.)

- Trees
 - quite flexible, but inherent limitation -- only express hierarchical structures
- Graphs
 - a collection of nodes and the connections between them
 - **generalize** a **tree**



Terminologies

- A **simple graph**
 - $G = (V, E)$ consists of a (finite) set denoted by V , and a collection E , of **unordered pairs** $\{u, v\}$ of distinct elements from V
 - V , called a **vertex** or a **point** or a **node**
 - E , called an **edge** or a **line** or a **link**
 - The number of vertices, $|V|$, and edges, $|E|$



Terminologies (cont.)

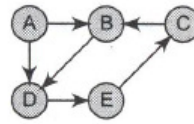


Figure 13.2 Directed graph

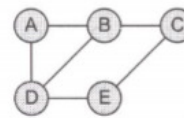
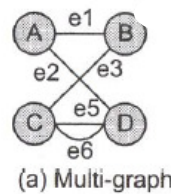


Figure 13.1 Undirected graph

- A **directed graph, digraph**
 - $G = (V, E)$, $(v_i, v_j) \neq (v_j, v_i)$
 - In a simple graph (undirected graph), $(v_i, v_j) = (v_j, v_i)$
- A **multigraph**
 - two vertices can be joined by **multiple edges**
- A **pseudograph**
 - a multigraph allowing for **loops**
 - a vertex can be joined with itself by an edge



(a) Multi-graph



pseudograph

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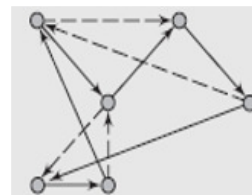
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Terminologies (cont.)

- A **path**
 - a sequence of edges, $\text{edge}(v_1, v_2)$, $\text{edge}(v_2, v_3)$, ..., $\text{edge}(v_{n-1}, v_n)$
 - denoted as path $v_1, v_2, v_3, \dots, v_{n-1}, v_n$
 - if $v_1 = v_n$ and **no edge is repeated**,
 - **circuit**
 - If the vertices in a circuit are different,
 - **cycle**



circuit in a digraph



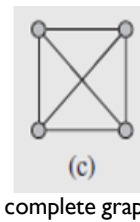
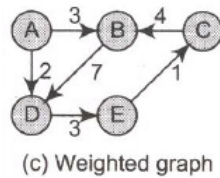
cycle in the digraph

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Terminologies (cont.)

- A **weighted graph**
 - an assigned number (e.g., weight, cost, distance, length, etc.) on each edge
- A **complete graph**
 - **exactly one edge** between each pair of distinct vertices



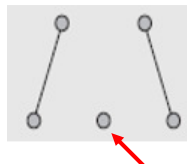
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Terminologies (cont.)

- A **subgraph** G' of graph $G = (V, E)$,
 - $G' = (V', E')$, where $V' \subseteq V$ and $E' \subseteq E$
- V_i and V_j are **adjacent**,
 - if the edge (V_i, V_j) is in E
 - such an edge is called **incident** with the vertices V_i and V_j
- The **degree** of a vertex v ,
 - $\deg(v)$, the **number of edges** incident with v
 - if $\deg(v) = 0$, v is an **isolated vertex**



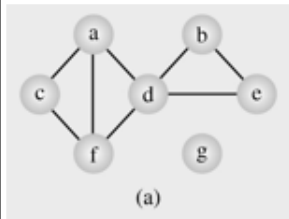
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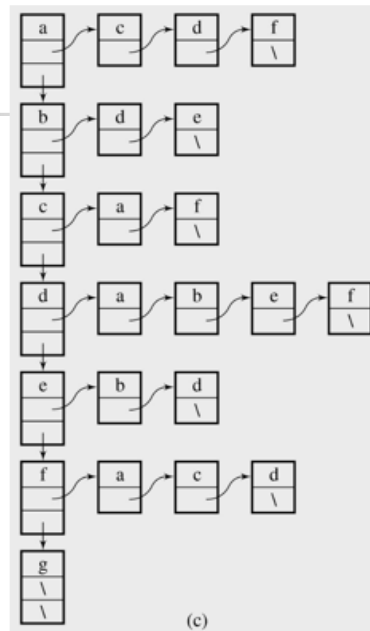
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Graph Representation

- An **adjacency list**
 - specify all vertices adjacent to each vertex of the graph



a	c	d	f
b	d	e	f
c	a	f	
d	a	b	e
e	b	d	
f	a	c	d
g			



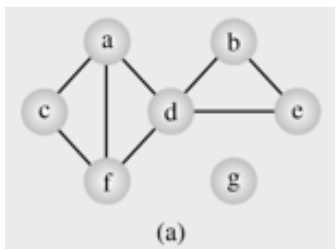
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Graph Representation (cont.)

- An **adjacency matrix**
 - a $|V| \times |V|$ **binary matrix** where,

$$a_{ij} = \begin{cases} 1 & \text{if there exists an edge } (v_i, v_j) \\ 0 & \text{otherwise} \end{cases}$$



	a	b	c	d	e	f	g
a	0	0	1	1	0	1	0
b	0	0	0	1	1	0	0
c	1	0	0	0	0	1	0
d	1	1	0	0	1	1	0
e	0	1	0	1	0	0	0
f	1	0	1	1	0	0	0
g	0	0	0	0	0	0	0

(d)

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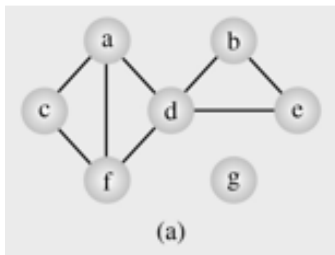


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Graph Representation (cont.)

- An **incidence matrix**
 - a $|V| \times |E|$ **binary matrix** where,
 - $a_{ij} = 1$ if **edge e_j** is **incident** with **vertex v_i** , otherwise 0



	ac	ad	af	bd	be	cf	de	df
a	1	1	1	0	0	0	0	0
b	0	0	0	1	1	0	0	0
c	1	0	0	0	0	1	0	0
d	0	1	0	1	0	0	1	1
e	0	0	0	0	1	0	1	0
f	0	0	1	0	0	1	0	1
g	0	0	0	0	0	0	0	0

(c)

