

# Sparsity Methods for Systems and Control

## What is Sparsity?

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1 Redundant Dictionary

2 Underdetermined Systems

3 The  $\ell^0$  Norm

4 Exhaustive Search

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2 Underdetermined Systems

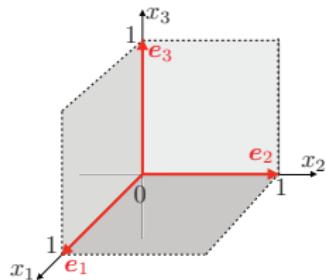
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# Standard basis for $\mathbb{R}^3$

- Standard basis  $\{e_1, e_2, e_3\}$ :

$$e_1 = \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}, \quad e_2 = \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix}, \quad e_3 = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix}.$$



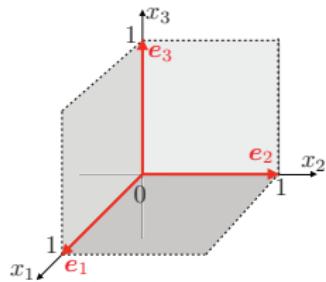
- Any vector  $y \in \mathbb{R}^3$  can be represented as

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# General basis for $\mathbb{R}^3$

- Any three **linearly independent** vectors  $\phi_1, \phi_2$ , and  $\phi_3$  in  $\mathbb{R}^3$  form a basis for  $\mathbb{R}^3$ .
- Any vector  $y \in \mathbb{R}^3$  can be represented as

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# Redundant basis

- Three linearly independent vectors:

$$\phi_1 = e_1 + e_2 = \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix}, \quad \phi_2 = e_2 + e_3 = \begin{bmatrix} 0 \\ 1 \\ 1 \end{bmatrix}, \quad \phi_3 = e_3 + e_1 = \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix}.$$

- Set of 6 vectors (**redundant basis**)

$$\{e_1, e_2, e_3, \phi_1, \phi_2, \phi_3\}$$

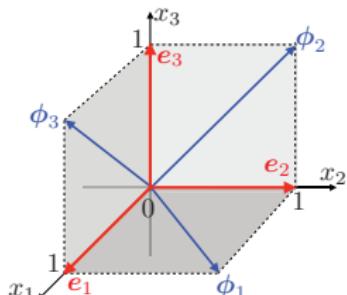
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# Redundant basis

- For a vector  $y \in \mathbb{R}^3$ , we want a signal representation (**redundant representation**):

$$y = \sum_{i=1}^3 \alpha_i e_i + \sum_{i=1}^3 \beta_i \phi_i.$$

- There are **infinitely many** solutions for  $\alpha_i$  and  $\beta_i$  ( $i = 1, 2, 3$ ).

$$(\alpha_1, \alpha_2, \alpha_3, \beta_1, \beta_2, \beta_3) = (y_1, y_2, y_3, 0, 0, 0),$$

$$(\alpha_1, \alpha_2, \alpha_3, \beta_1, \beta_2, \beta_3) = (-y_3, -y_1, -y_2, y_1, y_2, y_3).$$

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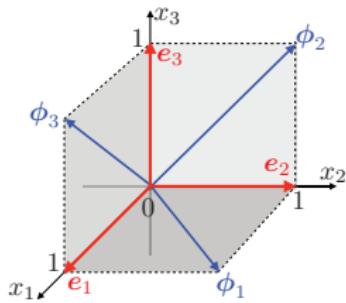
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# Sparse representation

- A vector  $y = (1, 1, 1)^\top$  on the plane spanned by  $e_1$  and  $\phi_2$ .
- A coefficient set is obtained as

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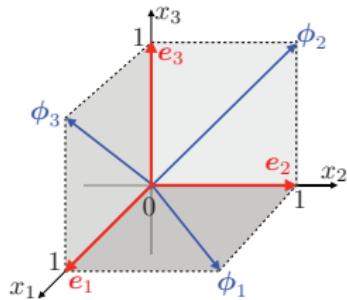


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# Redundant dictionary

- How do you explain this picture by using words in a small dictionary that does not have the word "elephant"?



# Redundant dictionary

- A set of vectors  $\{\phi_1, \phi_2, \dots, \phi_n\}$  in  $\mathbb{R}^m$ .
- If  $m < n$  and  $m$  vectors in this set are linearly independent, then this is called a **redundant dictionary**.
- The elements  $\phi_1, \phi_2, \dots, \phi_n$  in the dictionary is called **atoms** (not “words”).
- For a vector  $y \in \mathbb{R}^m$ , we find coefficients  $\alpha_1, \alpha_2, \dots, \alpha_n$  such that

$$y = \sum_{i=1}^n \alpha_i \phi_i.$$

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# Sparse representation

- Define a matrix  $\Phi$  and a vector  $x$  as

$$\Phi \triangleq [\phi_1 \quad \phi_2 \quad \dots \quad \phi_n] \in \mathbb{R}^{m \times n}, \quad x \triangleq \begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_n \end{bmatrix} \in \mathbb{R}^n.$$

- Then the relation

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is compactly rewritten as

$$y = \Phi x.$$

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## Problem (Sparse Representation)

Given a vector  $\mathbf{y} \in \mathbb{R}^m$  and a dictionary matrix  $\Phi \in \mathbb{R}^{m \times n}$  with  $m < n$ .  
Find the simplest (i.e. sparsest) representation of  $\mathbf{y}$  that satisfies

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# Linear equations

- linear equations with unknowns  $x_1$ ,  $x_2$ , and  $x_3$ :

$$x_1 + x_2 + x_3 = 3$$

$$x_1 - x_3 = 0$$

- There are infinitely many solutions
- All solutions

$$x_1 = t, \quad x_2 = -2t + 3, \quad x_3 = t,$$

where  $t \in \mathbb{R}$  is a parameter.

- Such a system of equations is called an **underdetermined system**.
- How can we specify one **unique** vector among these?

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# Minimum solution

- Let us consider a detective, like Edogawa Conan<sup>1</sup>, who solve this problem.
- The two proofs (equations) are insufficient and he should seek one more **independent** proof.
- If he gets one more proof saying *the criminal is the smallest one among the suspects*.
- The  $\ell^2$ -norm

$$\begin{aligned}\|x\|_2^2 &= x_1^2 + x_2^2 + x_3^2 \\ &= t^2 + (-2t + 3)^2 + t^2 \\ &= 6(t - 1)^2 + 3.\end{aligned}$$

is minimized by  $t = 1$ .

- The unique solution is  $(x_1, x_2, x_3) = (1, 1, 1)$ .

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# Linear equations in matrix form

- Linear equations in a matrix form:

$$\Phi x = y.$$

- $\Phi$  is an  $m \times n$  matrix where  $m < n$  (we call this a **fat matrix**).
- Assume  $\Phi$  has **full row rank**, that is,

$$\text{rank}(\Phi) = m.$$

- For any vector  $y \in \mathbb{R}^m$ , there exists at least one solution  $x$  that satisfies  $\Phi x = y$ .
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## Norm

A norm in  $\mathbb{R}^n$  should satisfy

- ① For any vector  $x \in \mathbb{R}^n$  and any number  $\alpha \in \mathbb{R}$ ,  $\|\alpha x\| = |\alpha| \|x\|$ .
- ② For any  $x, y \in \mathbb{R}^n$ ,  $\|x + y\| \leq \|x\| + \|y\|$ .
- ③  $\|x\| = 0 \iff x = 0$ .

- The  $\ell^2$  norm (or Euclidean norm)

$$\|x\|_2 \triangleq \sqrt{x_1^2 + x_2^2 + \dots + x_n^2}.$$

- The  $\ell^1$  norm

$$\|x\|_1 \triangleq |x_1| + |x_2| + \dots + |x_n|.$$

- The  $\ell^\infty$  norm (or maximum norm)

$$\|x\|_\infty \triangleq \max\{|x_1|, |x_2|, \dots, |x_n|\}$$

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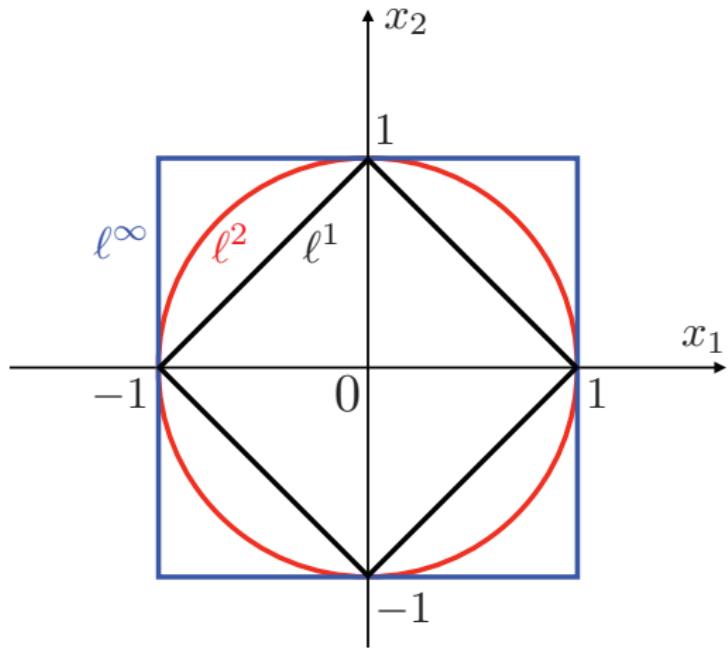
$$\|x\|_1 \triangleq |x_1| + |x_2| + \dots + |x_n|.$$

- The  $\ell^\infty$  norm (or maximum norm)

$$\|x\|_\infty \triangleq \max\{|x_1|, |x_2|, \dots, |x_n|\}$$

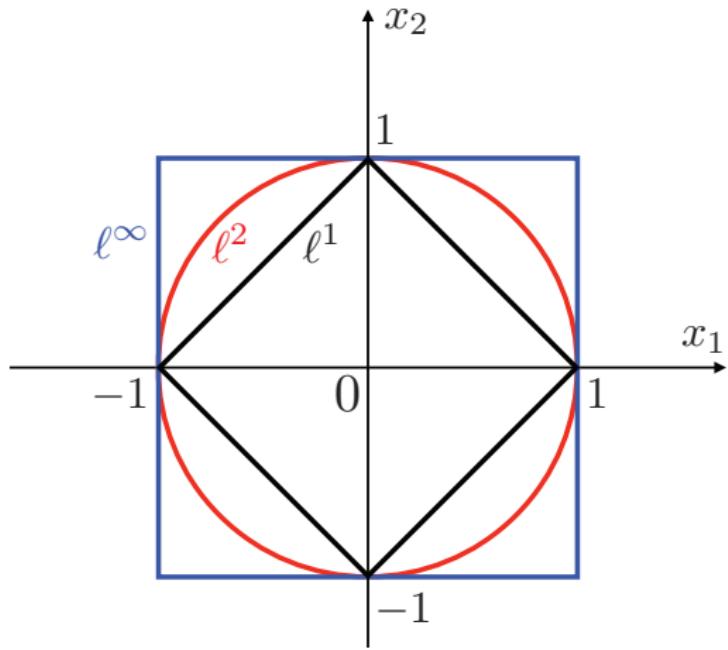
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Contour curves ( $\|x\|_p = 1$ ) of  $\ell^1, \ell^2, \ell^\infty$  norms.



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## $\ell^0$ norm

- Consider a vector  $x = [x_1, x_2, \dots, x_n]^\top \in \mathbb{R}^n$ .
- The  $\ell^0$  norm of  $x$  is defined by

$$\|x\|_0 \triangleq |\text{supp}(x)|,$$

where

- $\text{supp}$  is the support of  $x$ , namely,

$$\text{supp}(x) \triangleq \{i \in \{1, 2, \dots, n\} : x_i \neq 0\},$$

- $\ell^0$  norm is the number of non-zero elements.

- The  $\ell^0$  norm counts the number of non-zero elements in  $x$ .
- The  $\ell^0$  norm *does not* satisfy the first property in the definition of norm, and it is sometimes called the  $\ell^0$  pseudo-norm.

$$\|2x\|_0 = \|x\|_0 \neq 2\|x\|_0.$$

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# $\ell^0$ optimization problem

Now the problem of sparse representation is formulated as follows:

## $\ell^0$ optimization problem

Given a vector  $y \in \mathbb{R}^m$  and a full-row-rank matrix  $\Phi \in \mathbb{R}^{m \times n}$  with  $m < n$ . Find the optimizer  $x^*$  of the optimization problem:

$$\underset{x \in \mathbb{R}^n}{\text{minimize}} \quad \|x\|_0 \quad \text{subject to} \quad y = \Phi x.$$

This problem is called the  $\ell^0$  optimization.

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1 Redundant Dictionary

2 Underdetermined Systems

3 The  $\ell^0$  Norm

4 Exhaustive Search

# How to solve it?

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- We can try an **exhaustive search** for this optimization.

# Exhaustive search: example

- Find the minimum- $\ell^0$  solution  $(x_1, x_2, x_3)$  that satisfies

$$x_1 + x_2 + x_3 = 3$$

$$x_1 - x_3 = 0$$

- First, try  $(x_1, x_2, x_3) = (0, 0, 0)$ . This is not a solution.
- Second, try  $(x_1, 0, 0)$ ,  $(0, x_2, 0)$ , and  $(0, 0, x_3)$ .

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# Exhaustive Search (step 1)

- If  $y = \mathbf{0}$ , then output  $x^* = \mathbf{0}$  as the optimal solution and quit.
- Otherwise, proceed to the next step.

## Exhaustive Search (step 2)

- Find a vector  $x$  with  $\|x\|_0 = 1$  that satisfies the equation  $y = \Phi x$ .  
That is, set

$$x_1 \triangleq \begin{bmatrix} x_1 \\ 0 \\ \vdots \\ 0 \end{bmatrix}, \quad x_2 \triangleq \begin{bmatrix} 0 \\ x_2 \\ 0 \\ \vdots \\ 0 \end{bmatrix}, \dots, \quad x_n \triangleq \begin{bmatrix} 0 \\ \vdots \\ 0 \\ x_n \end{bmatrix}$$

and search  $x_i \in \mathbb{R}$  ( $i = 1, 2, \dots, n$ ) that satisfies

$$y = \Phi x_i = x_i \phi_i.$$

The algorithm starts with  $x_1$  and checks if  $y = \Phi x_1$  holds. If it does, then  $x_1$  is the solution. If not, then the algorithm moves to  $x_2$  and repeats the process.

After all  $x_i$  have been checked, the algorithm returns to the previous step and continues until a solution is found.

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## Exhaustive Search (step 3)

- Find a vector  $x$  with  $\|x\|_0 = 2$  that satisfies the equation  $y = \Phi x$ .  
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$$x_{1,2} \triangleq \begin{bmatrix} x_1 \\ x_2 \\ 0 \\ \vdots \\ 0 \end{bmatrix}, \quad x_{1,3} \triangleq \begin{bmatrix} x_1 \\ 0 \\ x_3 \\ 0 \\ \vdots \\ 0 \end{bmatrix}, \dots, \quad x_{n-1,n} \triangleq \begin{bmatrix} 0 \\ \vdots \\ 0 \\ x_{n-1} \\ x_n \end{bmatrix}$$

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# Exhaustive Search (step $k$ )

- Do similar procedures for  $\|x\|_0 = k, k = 3, 4, \dots, m$ .

# Is exhaustive search useful?

- It is easily implemented.
- The computation time to find a solution grows exponentially with problem size  $m$ .
- Suppose  $m = 100$ .

Exhaustive search is a simple algorithm that checks every possible combination of variables to find the best one.

It is called "exhaustive" because it checks every possible combination of variables.

The number of combinations is  $2^m$ , where  $m$  is the number of variables.

For  $m = 100$ , there are approximately  $10^{30}$  combinations.

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- Sparsity of a vector is measured by its  $\ell^0$  norm.
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