

Practical Attacks on Real World Crypto Implementations

Juraj Somorovsky

Ruhr-Universität Bochum, Hackmanit GmbH

Recent years revealed many crypto attacks...

- ESORICS 2004, Ba 2011 BEAS'I' of SSL to Chosen Plaintext Attack
- Eurocrypt 2002, Vaudena 2013 Lucky13 Induced by CBC Padding—Applications to SSL, II CL., II
- Crypto 1998, Bleich 2012 XML Encryption hertext Attacks Against Protocols based on 2016 DROWN andard PKCS #1



Standards updated

Countermeasures defined

What could go wrong in RWC implementations?



Overview

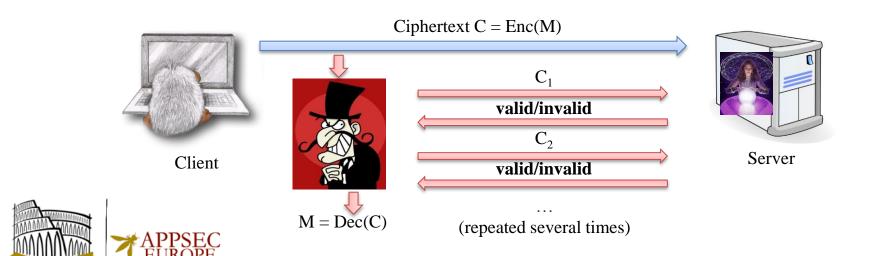
- 1. Bleichenbacher's Attack
 - XML Encryption
 - TLS
- 2. Invalid Curve Attack
 - · TLS
- 3. Padding Oracle Attack
 - TLS
- 4. TLS-Attacker





RSA-PKCS#1 v1.5

- Used to encrypt symmetric keys
- Vulnerable to an adaptive chosen-ciphertext attack



RSA-PKCS#1 v1.5: Countermeasures

- 1. Use RSA-OAEP (PKCS#1 v2)
- 2. Apply specific countermeasure

```
generate random
decrypt ciphertext: m = dec(c)
if ( padding correct )
   proceed with m
else
   proceed with random
```



Overview

- 1. Bleichenbacher's Attack
 - XML Encryption
 - TLS
- 2. Invalid Curve Attack
 - · TLS
- 3. Padding Oracle Attack
 - TLS
- 4. TLS-Attacker





RSA PKCS#1 v1.5 in XML Encryption

Hybrid encryption:

```
k = Dec_pkcs(priv,C1)
m = Dec_aes128(k,C2)
```

1 Dec_pkcs



Dec_aes128



Attack Countermeasure

Hybrid encryption:

```
k = Dec_pkcs(priv,C1)
m = Dec_aes128(k,C2)
```

Dec_pkcs



Dec_aes128



Case Apache WSS4J

Hybrid encryption:

```
k = Dec_pkcs(priv,C1)
m = Dec_aes128(k,C2)
```

Dec_pkcs



2 Dec_aes128



Case Apache WSS4J

Hybrid encryption:

```
k = Dec_pkcs(priv,C1)
m = Dec_aes128(k,C2)
```

Dec_pkcs

k Random:
128 B



Case Apache WSS4J

- Original bug much more complicated
- CVE-2015-0226
- Dennis Kupser, Christian Mainka, Jörg Schwenk, Juraj Somorovsky: How to Break XML Encryption – Automatically (WOOT'15)
- Found automatically using WS-Attacker
- https://github.com/RUB-NDS/WS-Attacker



Overview

- 1. Bleichenbacher's Attack
 - XML Encryption
- TLS
 - 2. Invalid Curve Attack
 - · TLS
 - 3. Padding Oracle Attack
 - TLS
 - 4. TLS-Attacker





How About TLS?

- Christopher Meyer, Juraj Somorovsky, Jörg Schwenk, Eugen Weiss, Sebastian Schinzel, Erik Tews: Revisiting SSL/TLS Implementations: New Bleichenbacher Side Channels and Attacks. USENIX Security 2014
- Practical attacks on JSSE Bouncy Castle, Cavium Accelerator
- Bug in OpenSSL



Case JSSE

- No direct TLS error messages
- Uses PKCS#1 unpadding function:

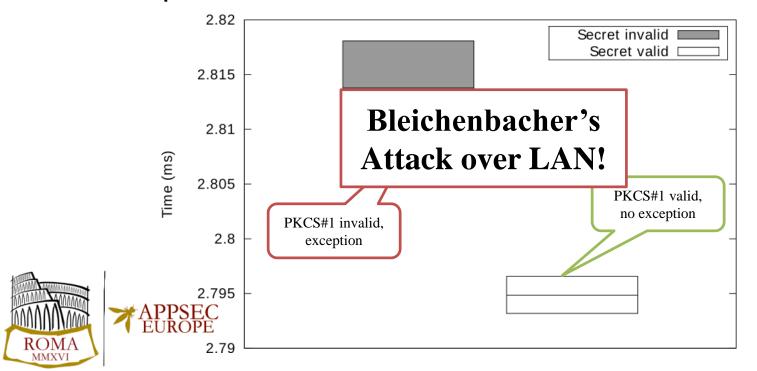
```
private byte [] unpadV15 (byte[] padded) {
   if (PKCS valid) {
      return unpadded text;
   } else {
      throw new BadPaddingException();
   }
}
```

Caught, random generated...what's wrong?



Case JSSE (CVE-2014-411)

Exception consumes about 20 microseconds!



Overview

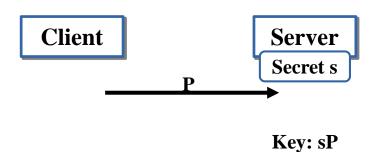
- 1. Bleichenbacher's Attack
 - XML Encryption
 - TLS
- 2. Invalid Curve Attack
 - · TLS
- 3. Padding Oracle Attack
 - TLS
- 4. TLS-Attacker

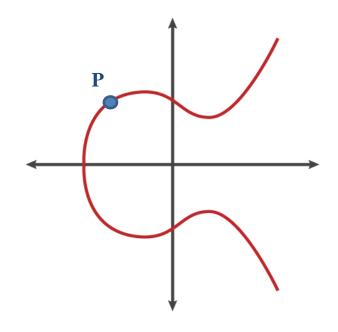




Elliptic Curve

- Set of points over a finite field
- Used e.g. for key exchange







Invalid Curve Attack

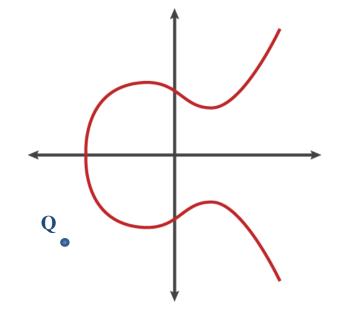
- Crypto 2000: Biehl, Meyer, Müller
- Attacker sends an invalid point of small order (e.g. 5)



Attacker computes:



 $s_1 = s \mod 5$



Invalid Curve Attack

- Choose points of small co-prime order (5, 7, 11, ...)
- Send to the server
- Compute:

```
s_1 = s \mod 5
```

 $s_2 = s \mod 7$

 $s_3 = s \mod 11$

 $s_4 = s \mod 13$

Compute s with CRT



Overview

- 1. Bleichenbacher's Attack
 - XML Encryption
 - TLS
- 2. Invalid Curve Attack



- 3. Padding Oracle Attack
 - TLS
- 4. TLS-Attacker





Practical Attacks?

- Tibor Jager, Jörg Schwenk, Juraj Somorovsky: Practical Invalid Curve Attacks on TLS-ECDH. ESORICS 2015
- Analyzed 8 libraries
- 2 vulnerable
 - Bouncy Castle: 3300 TLS queries
 - Oracle JSSE: 17000 TLS queries
- Further vulnerability found in a Hardware Security module



Impact

- Attacks extract server private keys
- Java servers using EC certificates vulnerable
 - For example Apache Tomcat



Demo



Overview

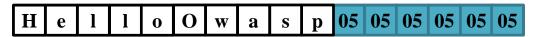
- 1. Bleichenbacher's Attack
 - XML Encryption
 - TLS
- 2. Invalid Curve Attack
 - · TLS
- 3. Padding Oracle Attack
 - TLS
- 4. TLS-Attacker





AES

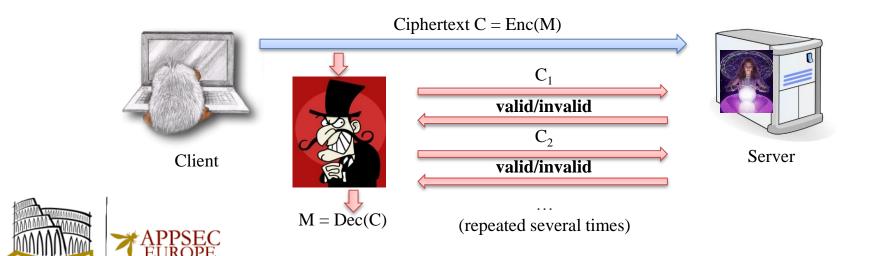
- Block cipher
- Cipher block chaining mode (CBC) of operation
- Last block has to be padded
 - Message: HelloOwasp
 - Padding size: 16 10 = 6





AES-CBC

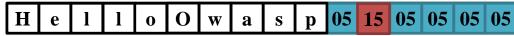
- Vulnerable to an adaptive chosen-ciphertext attack
- Padding oracle attack

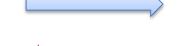


Padding Oracle in AES-CBC

Flipping bits in AES-CBC possible









- Countermeasure: Authenticated encryption
 - Messages cannot be modified



Overview

- 1. Bleichenbacher's Attack
 - XML Encryption
 - TLS
- 2. Invalid Curve Attack
 - · TLS
- 3. Padding Oracle Attack



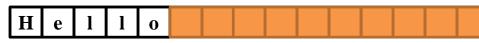
4. TLS-Attacker



AES-CBC in TLS

pad mac

- MAC-Pad-Encrypt
- Example:
 - Two blocks
 - Message: Hello
 - MAC size: 20 bytes
 - Padding size: 32 5 20 = 7







06 06 06 06 06 06 06

AES-CBC in TLS

pad mac

- Challenge: not to reveal padding validity
- Always:
 - Padding validation
 - MAC validation
- Same error message









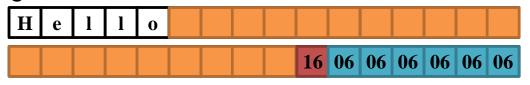
Botan Vulnerability

```
cbc decrypt record(record, record len);
verify padding(record, record len);
if (padding bad)
      pad size = 0;
size t mac pad size = mac size + pad size;
if(record len < mac pad size)</pre>
      throw TLS Exception(Alert::DECODING ERROR);
u16bit plaintext length = record len - mac pad size;
verify mac(plaintext, plaintext length, mac size);
 f(mac_bad || padding_bad)
       throw TLS Exception(Alert::BAD RECORD MAC);
```

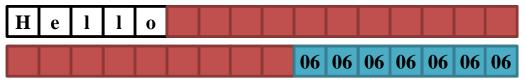
Botan Vulnerability



Bad padding: BAD_RECORD_MAC



Bad MAC: BAD_RECORD_MAC



Special case: Decoding_Error







Constant Time Processing

- Timing side-channel can introduce a padding oracle
- Hard to exploit
- Patches applied after Lucky13

Let's analyze some of them ...



Constant Time Validation

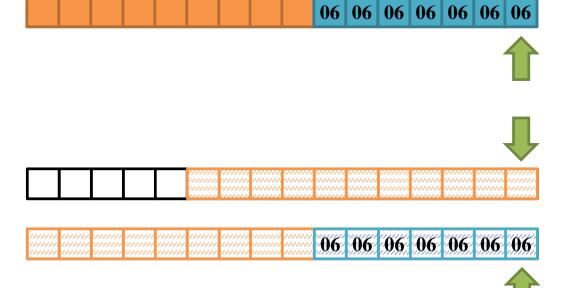






Mask data





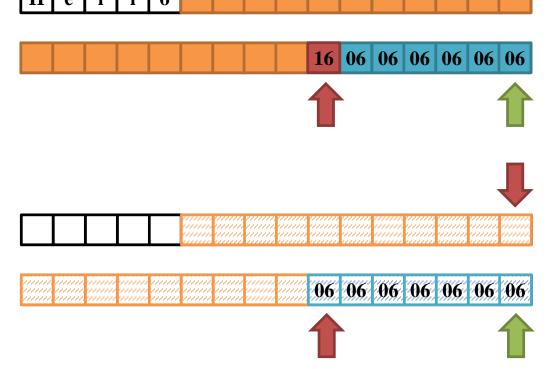
Constant Time Validation





Mask data

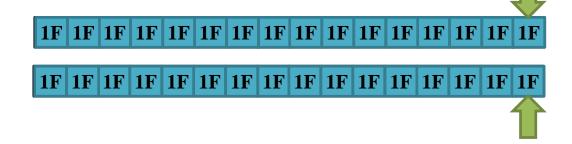




OpenSSL Vulnerability



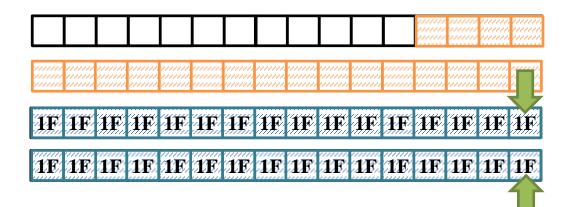
Decrypted data





Mask data



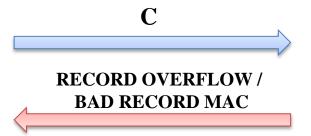


OpenSSL Vulnerability

- CVE-2016-2107
- Leads to a different server response









- Exploitable in BEAST scenarios
 - Decryption of 16 bytes possible



http://web-in-security.blogspot.it/2016/05/curious-padding-oracle-in-openssl-cve.html

MatrixSSL Vulnerability

- Tried to fix Lucky13
- Introduced a buffer overread
- Fixed in 3.8.3
- https://github.com/matrixssl/matrixssl/blob/master/CHANGES.md



Overview

- 1. Bleichenbacher's Attack
 - XML Encryption
 - TLS
- 2. Invalid Curve Attack
 - · TLS
- 3. Padding Oracle Attack
 - TLS
- 4. TLS-Attacker





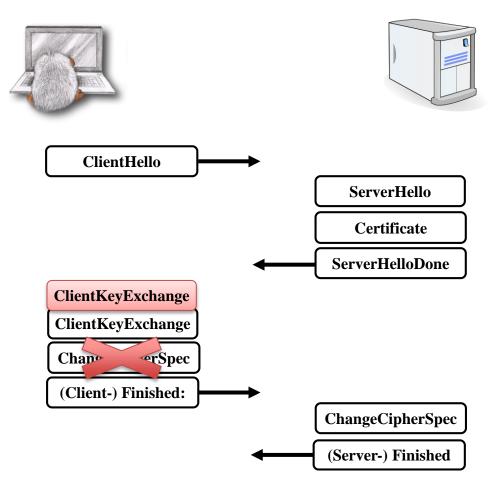
Recent Attacks on TLS

- Not only crypto attacks ...
- Buffer overflows / overreads
 - Heartbleed
- Invalid protocol flows
 - FREAK
 - Early CCS
- Tool for flexible protocol executions needed



TLS-Attacker

- Flexible
 - Protocol flow definition
 - Variable modifications





TLS-Attacker

```
WorkflowExecutor workflowExecutor =
configHandler.initializeWorkflowExecutor(transportHandler, tlsContext);
List<ProtocolMessage> protocolMessages = new LinkedList<>();
protocolMessages.add(new ClientHelloMessage(ConnectionEnd.CLIENT));
protocolMessages.add(new ServerHelloMessage(ConnectionEnd.SERVER));
protocolMessages.add(new CertificateMessage(ConnectionEnd.SERVER));
protocolMessages.add(new ServerHelloDoneMessage(ConnectionEnd.SERVER));
protocolMessages.add(new RSAClientKeyExchangeMessage(ConnectionEnd.CLIENT));
protocolMessages.add(new RSAClientKeyExchangeMessage(ConnectionEnd.CLIENT));
protocolMessages.add(new ChangeCipherSpecMessage(ConnectionEnd.CLIENT));
protocolMessages.add(new FinishedMessage(ConnectionEnd.CLIENT));
protocolMessages.add(new ChangeCipherSpecMessage(ConnectionEnd.SERVER));
protocolMessages.add(new FinishedMessage(ConnectionEnd.SERVER));
workflowExecutor.executeWorkflow():
```

TLS-Attacker

- https://github.com/RUB-NDS/TLS-Attacker
- Crypto attacks
- Fuzzing
- Protocol suite framework
- Written in Java



Demo

Conclusions

- Old attacks relevant for RWC implementations
- Old algorithms in the newest standards
 - RSA PKCS#1 v1.5 (attack: 1998)

2008: TLS 1.2

2013: XML Encryption 1.1

2015: JSON Web Encryption

New add-hoc countermeasures can introduce new flaws



Conclusions

- For standard designers:
 - Remove old crypto
- For developers:
 - Analyze possible side-channels, best practices
 - Check point is on curve
- For pentesters:
 - More tools / analyses of crypto applications needed
 - TLS-Attacker

