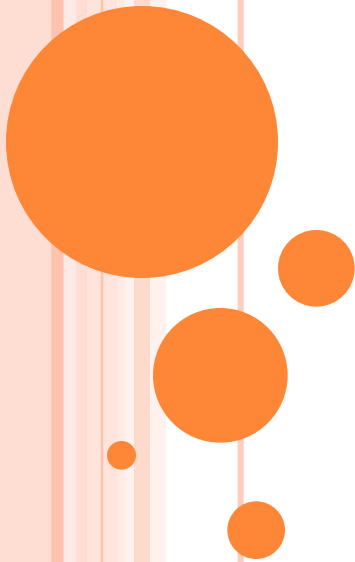


CSE 4112

LEXICAL ANALYSIS

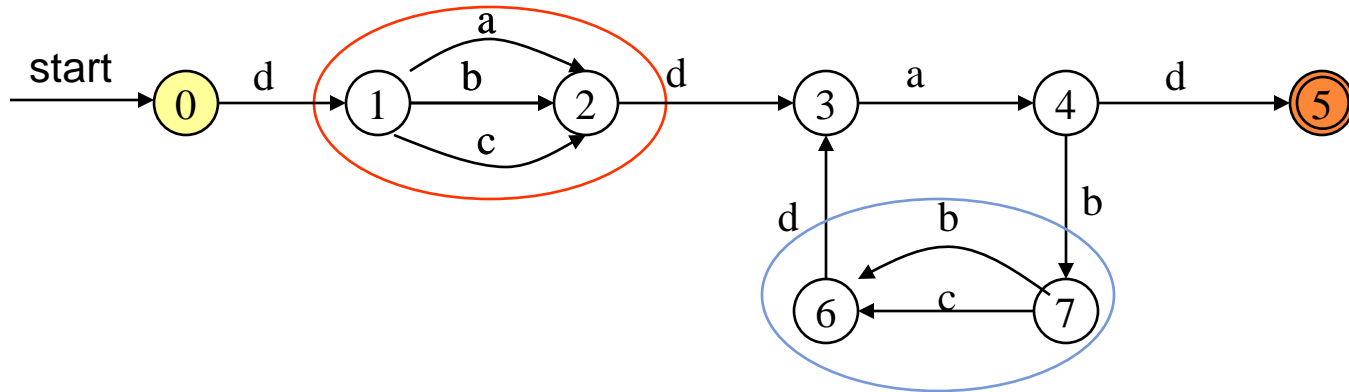
Lecture 03



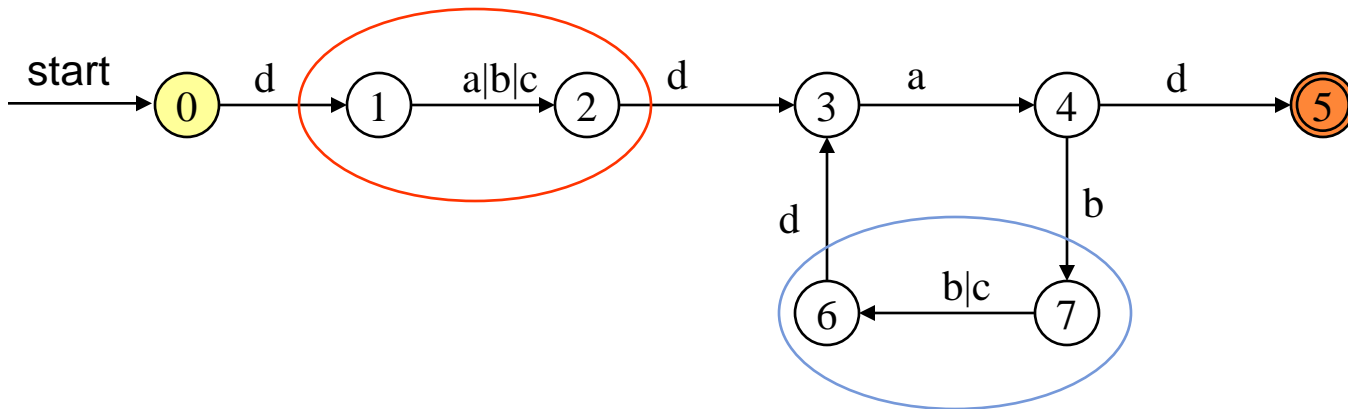
CONVERTING DFAS TO RES

1. Combine serial links by concatenation
2. Combine parallel links by alternation
3. Remove self-loops by Kleene closure
4. Select a node (other than initial or final) for removal. Replace it with a set of equivalent links whose path expressions correspond to the in and out links
5. Repeat steps 1-4 until the graph consists of a single link between the entry and exit nodes.

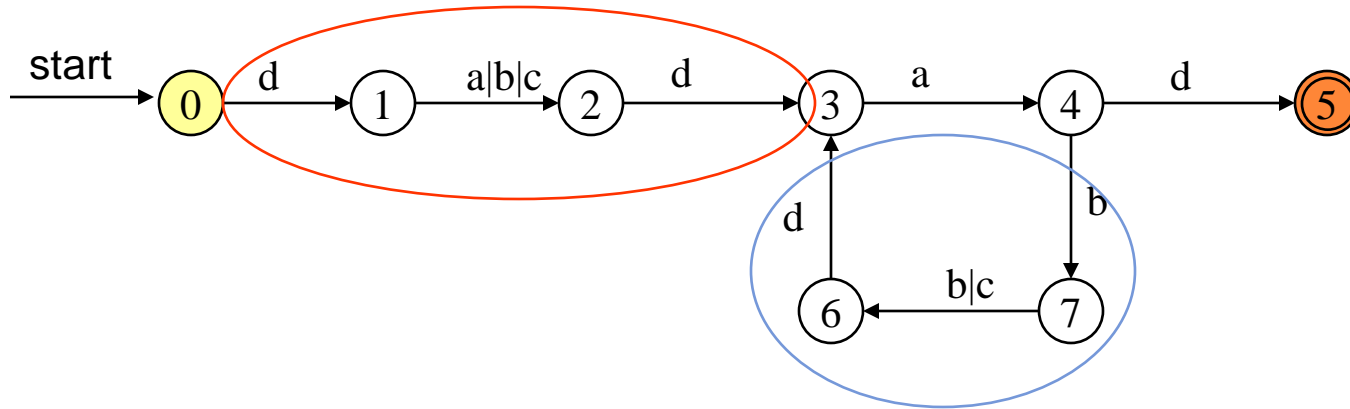
EXAMPLE



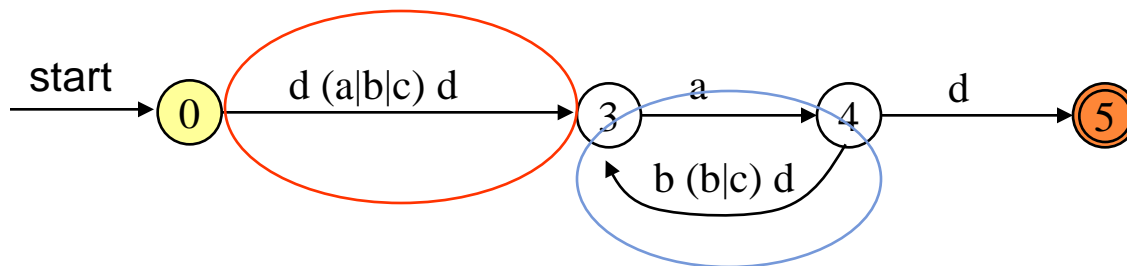
parallel edges become alternation



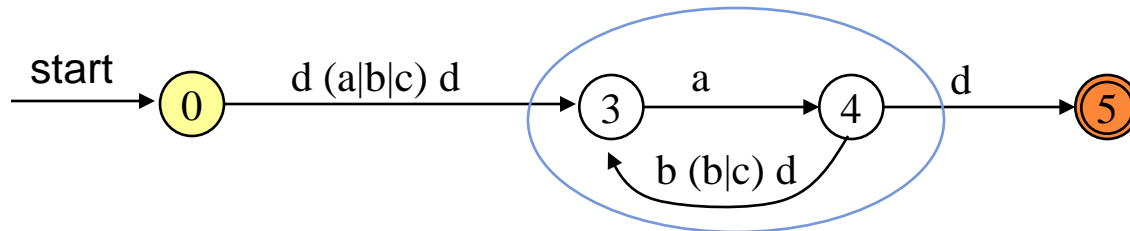
EXAMPLE



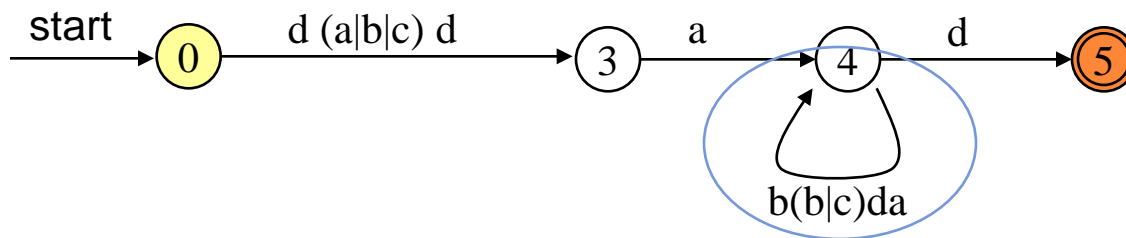
serial edges become concatenation



EXAMPLE



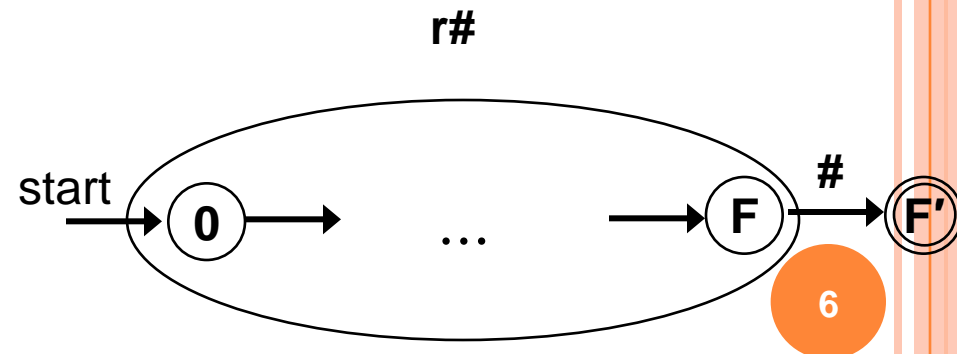
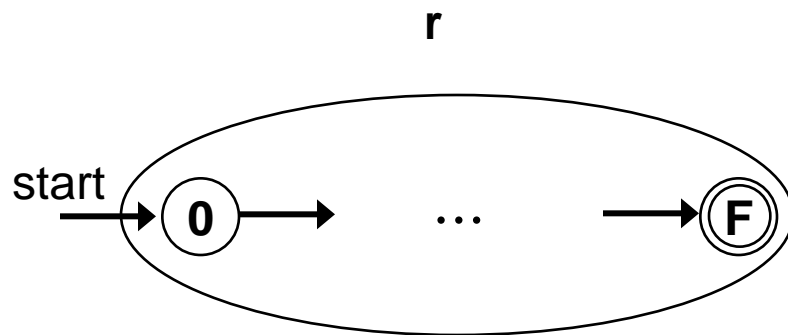
Find paths that can be “shortened”



REGULAR EXPRESSION TO DFA

○ Important States of NFA

- If it has a non- ϵ out-transition
- $move(s,a)$ is non-empty if s is important
- Accepting states are not important states
 - Adding a unique marker $\#$ after the RE r (i.e. $r\#$) we can make the accepting states important
 - Now a state with a transition on $\#$ will be accepting state

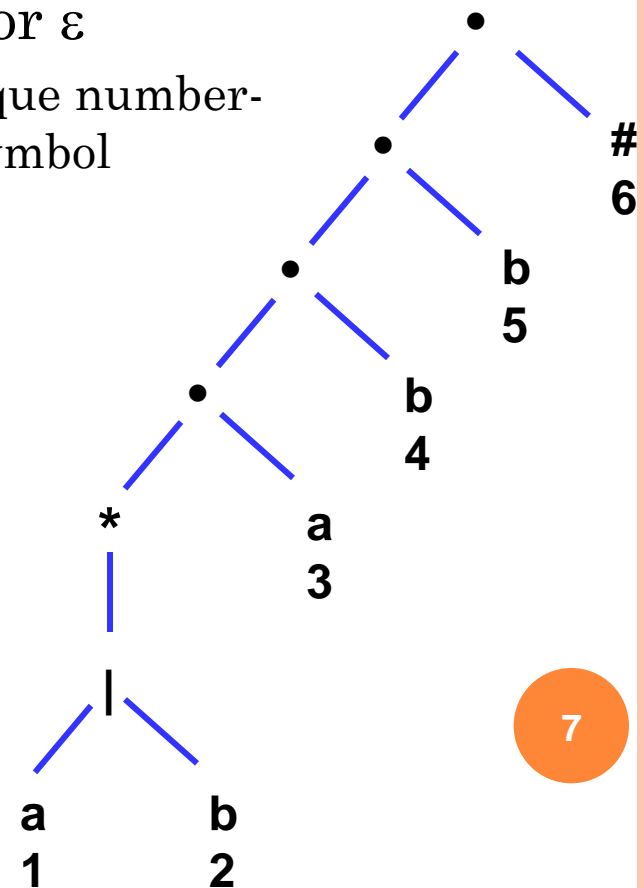


SYNTAX TREE

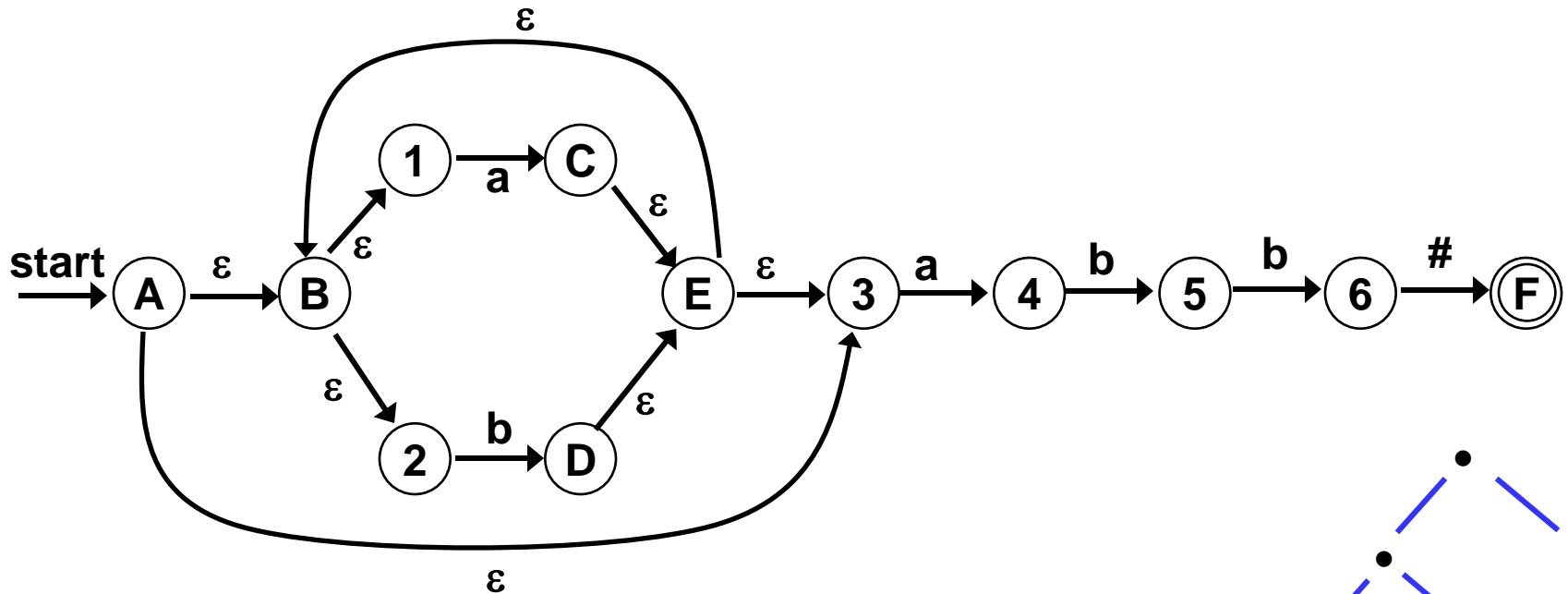
- Augmented RE ($r\#$) can be represented by a syntax tree

- Leaves contain: Alphabet symbols or ϵ
 - Each non- ϵ leaf is associated with a unique number-*position* of the leaf and *position* of the symbol
- Internal nodes contain: Operators
 - cat-node*, *or-node* or *star-node*

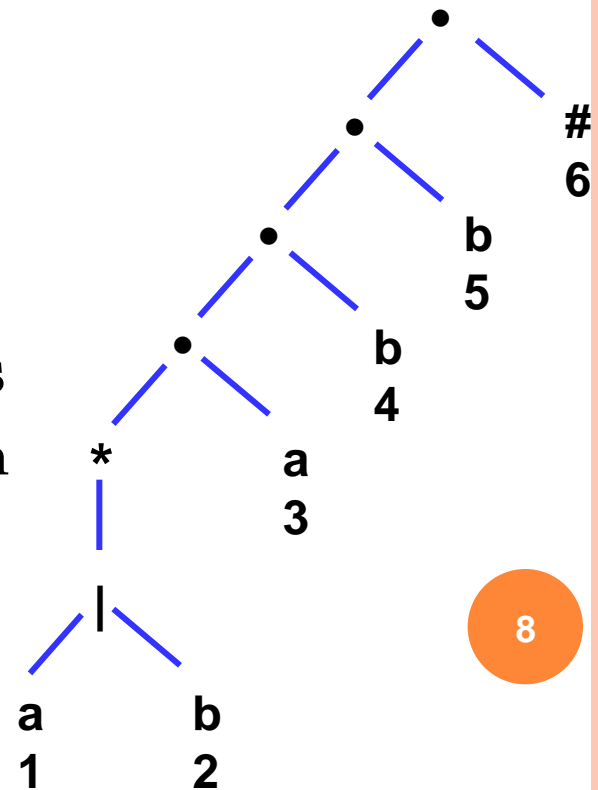
- Syntax tree for $r\# = (a \mid b)^*abb\#$



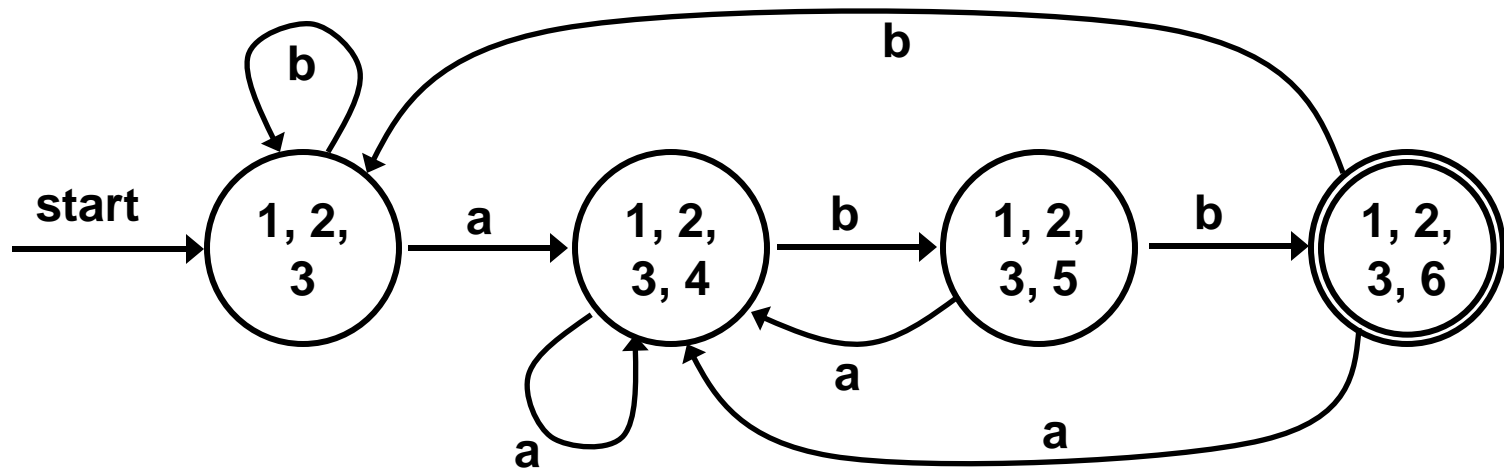
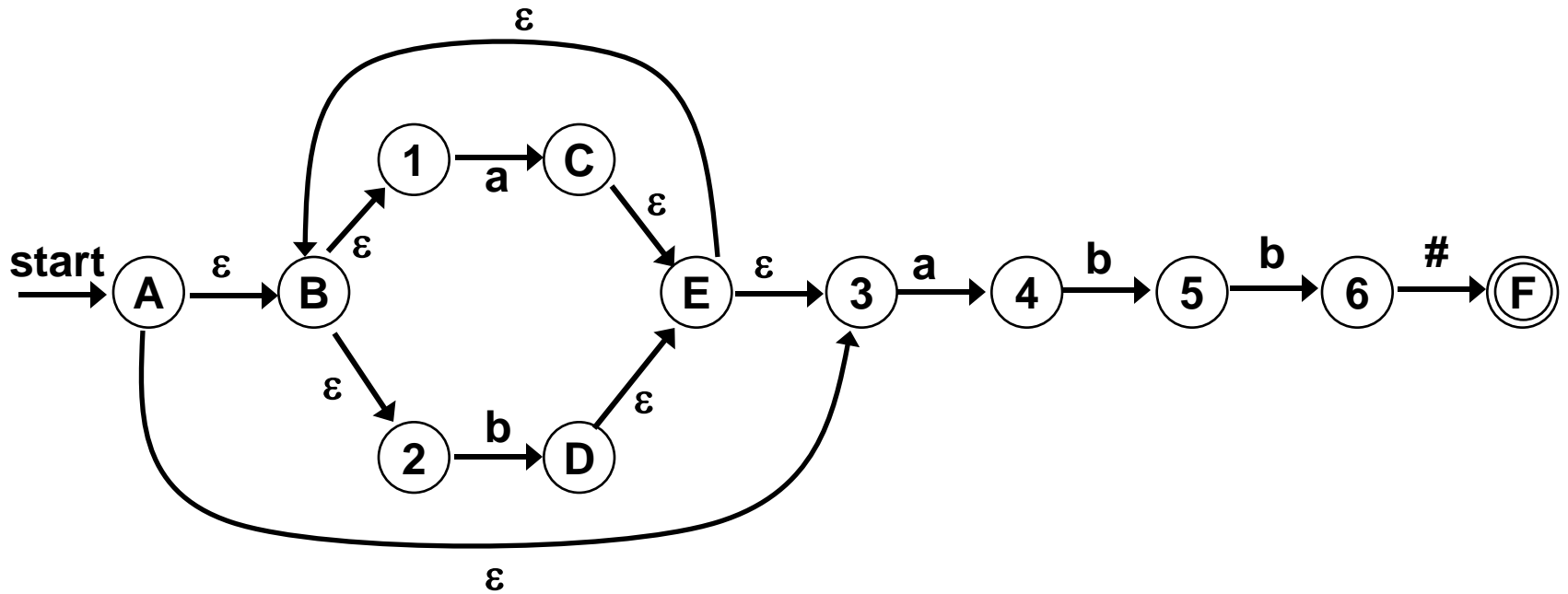
NFA FOR $(A \mid B)^*ABB\#$



- Lettered states are non-important states
- Number states are important states
 - Numbers correspond to the number in syntax tree

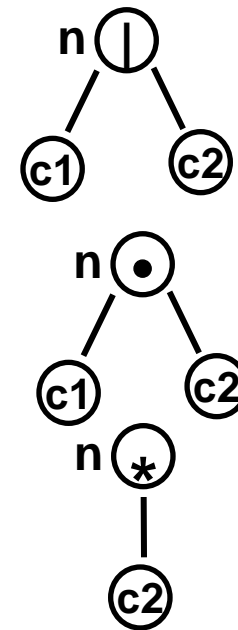


DFA FOR $(A \mid B)^*ABB\#$



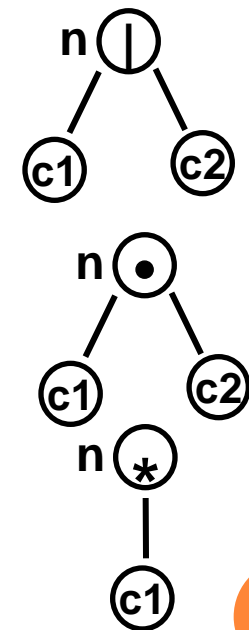
TERMINOLOGY

- Nullable:
 - Nodes that are the root of some sub-expression that generate empty string
- If n is a leaf labeled by ε then
 - **nullable (n) = true**
- If n is a leaf labeled with position i
 - **nullable (n) = false**
- If n is an or-node ($|$) with children $c1$ and $c2$
 - **nullable (n) = nullable($c1$) or nullable ($c2$)**
- If n is an cat-node (\bullet) with children $c1$ and $c2$
 - **nullable (n) = nullable($c1$) and nullable ($c2$)**
- If n is an star-node ($*$) with children $c1$
 - **nullable (n) = true**



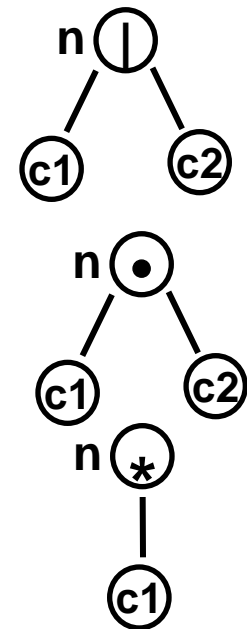
TERMINOLOGY

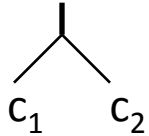
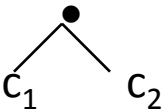

- **Firstpos(n):**
 - Set of positions that can match the first symbol of a string generated by the sub-expression rooted at n
- If n is a leaf labeled by ϵ then
 - **firstpos (n) = \emptyset**
- If n is a leaf labeled with position i
 - **firstpos (n) = {i}**
- If n is an or-node (|) with children c1 and c2
 - **firstpos (n) = firstpos(c1) \cup firstpos (c2)**
- If n is a cat-node (\bullet) with children c1 and c2
 - **firstpos(n) = If nullable (c1) then firstpos(c1) \cup firstpos (c2) else firstpos(c1)**
- If n is an star-node (*) with children c1
 - **firstpos (n) = firstpos(c1)**



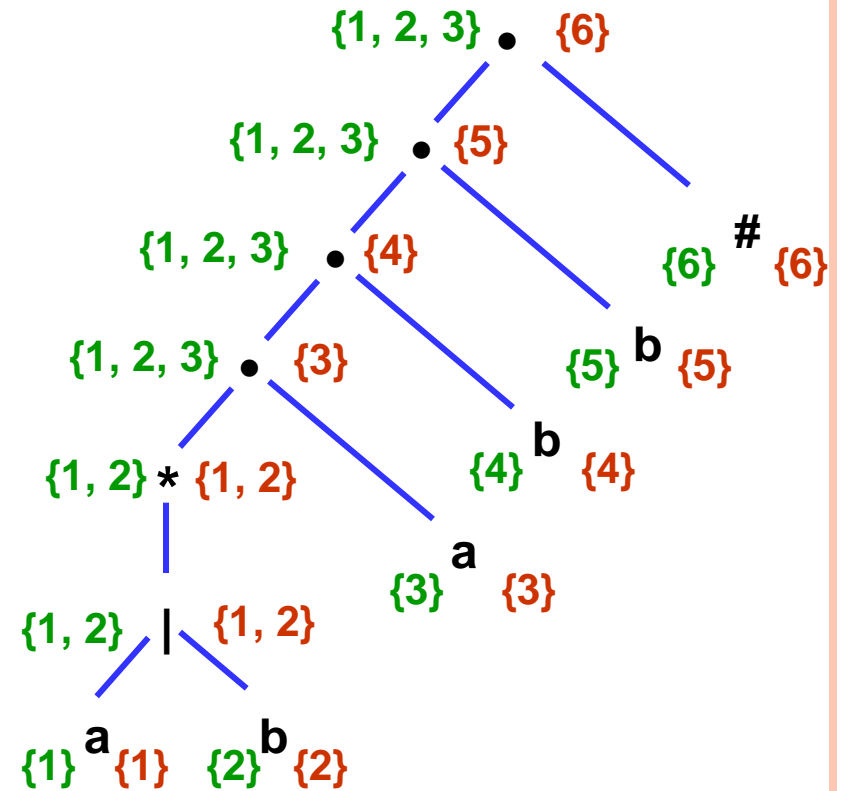
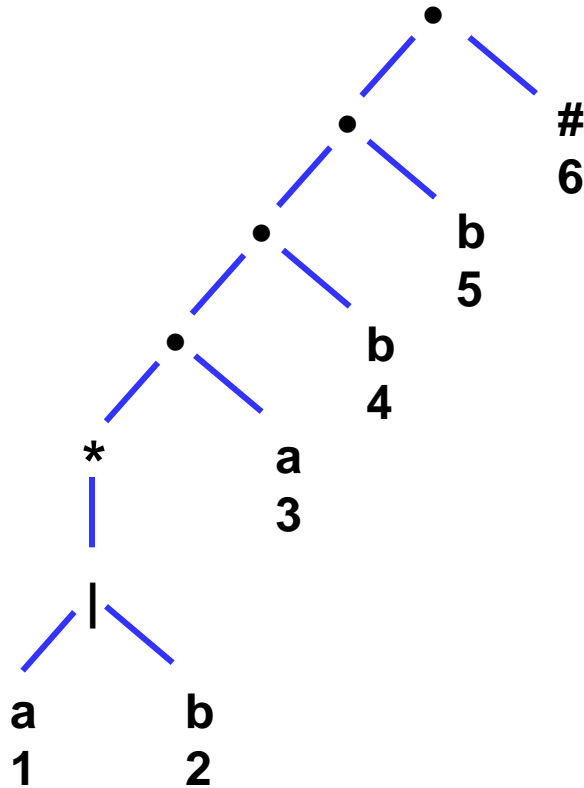
TERMINOLOGY

- Lastpos(n):
 - Set of positions that can match the last symbol of a string generated by the sub-expression rooted at n
- If n is a leaf labeled by ϵ then
 - **lastpos (n) = \emptyset**
- If n is a leaf labeled with position i
 - **lastpos (n) = {i}**
- If n is an or-node (|) with children c1 and c2
 - **lastpos (n) = lastpos(c1) \cup lastpos (c2)**
- If n is an cat-node (\bullet) with children c1 and c2
 - **lastpos(n) = If nullable (c2) then lastpos(c1) \cup lastpos (c2)**
 - **else lastpos(c2)**
- If n is an star-node (*) with children c1
 - **lastpos (n) = lastpos(c1)**



n	nullable(n)	firstpos(n)	lastpos(n)
leaf labeled ε	true	Φ	Φ
leaf labeled with position i	false	$\{i\}$	$\{i\}$
	$\text{nullable}(c_1) \text{ or } \text{nullable}(c_2)$	$\text{firstpos}(c_1) \cup \text{firstpos}(c_2)$	$\text{lastpos}(c_1) \cup \text{lastpos}(c_2)$
	$\text{nullable}(c_1) \text{ and } \text{nullable}(c_2)$	if ($\text{nullable}(c_1)$) $\text{firstpos}(c_1) \cup \text{firstpos}(c_2)$ else $\text{firstpos}(c_1)$	if ($\text{nullable}(c_2)$) $\text{lastpos}(c_1) \cup \text{lastpos}(c_2)$ else $\text{lastpos}(c_2)$
	true	firstpos(c_1)	lastpos(c_1)

FIRSTPOS AND LASTPOS EXAMPLE



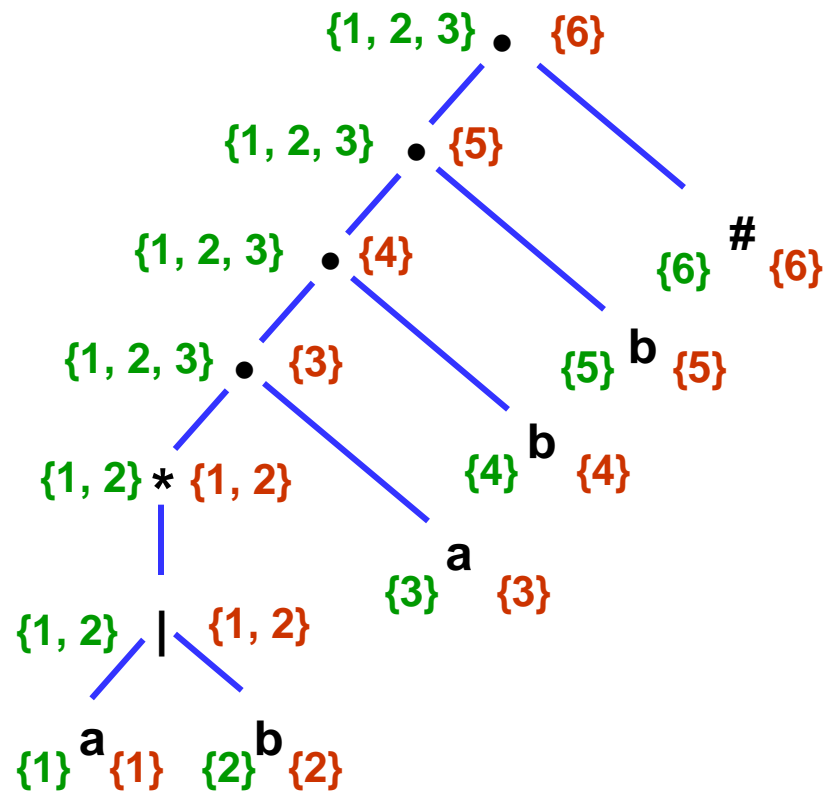
TERMINOLOGY

- Followpos(i):
 - Tells what positions can follow position i in the syntax tree
- Rule 1:

If n is a cat-node with left child $c1$ and right child $c2$ and i is a position in lastpos($c1$), then all positions in firstpos($c2$) are in followpos(i)
- Rule 2:

If n is a star node, and i is a position in lastpos(n), then all positions in firstpos(n) are in followpos(i)
- After computing firstpos and lastpos for each node followpos of each position can be computed by making depth-first traversal of the syntax tree

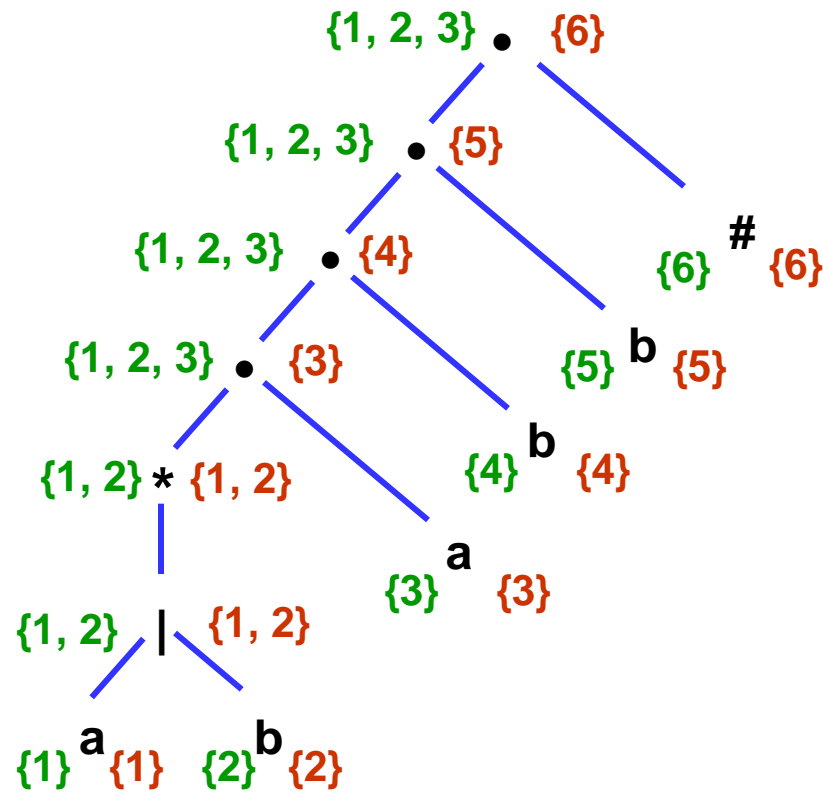
FOLLOWPOS EXAMPLE



Node	followpos
1	{1,2,
2	{1,2,
3	{
4	{
5	{
6	{

- At star-node:
 - $lastpos(*) = \{1,2\}$ and $firstpos(*) = \{1,2\}$
 - According to Rule 2:
 - $followpos\{1\} = \{1,2\}$
 - $followpos\{2\} = \{1,2\}$

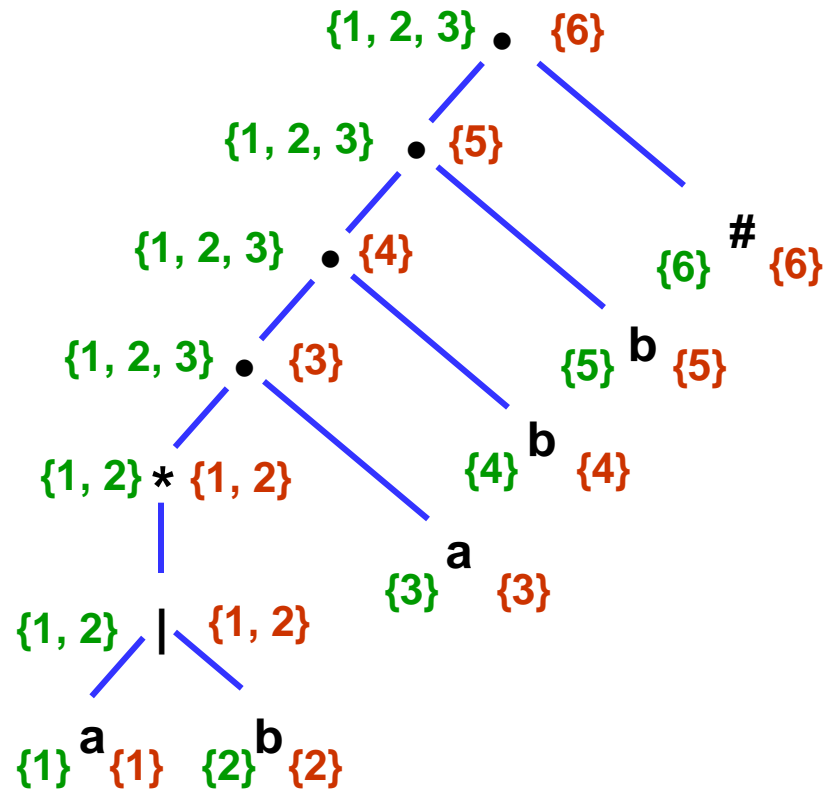
FOLLOWPOS EXAMPLE



Node	Followpos
1	{1,2,3
2	{1,2,3
3	{
4	{
5	{
6	{

- At cat-node above the star-node, '*' is left child and 'a' is right child
 - $lastpos(*) = \{1,2\}$ and $firstpos(a) = \{3\}$
 - According to Rule 1:
 - $followpos\{1\} = \{3\}$
 - $followpos\{2\} = \{3\}$

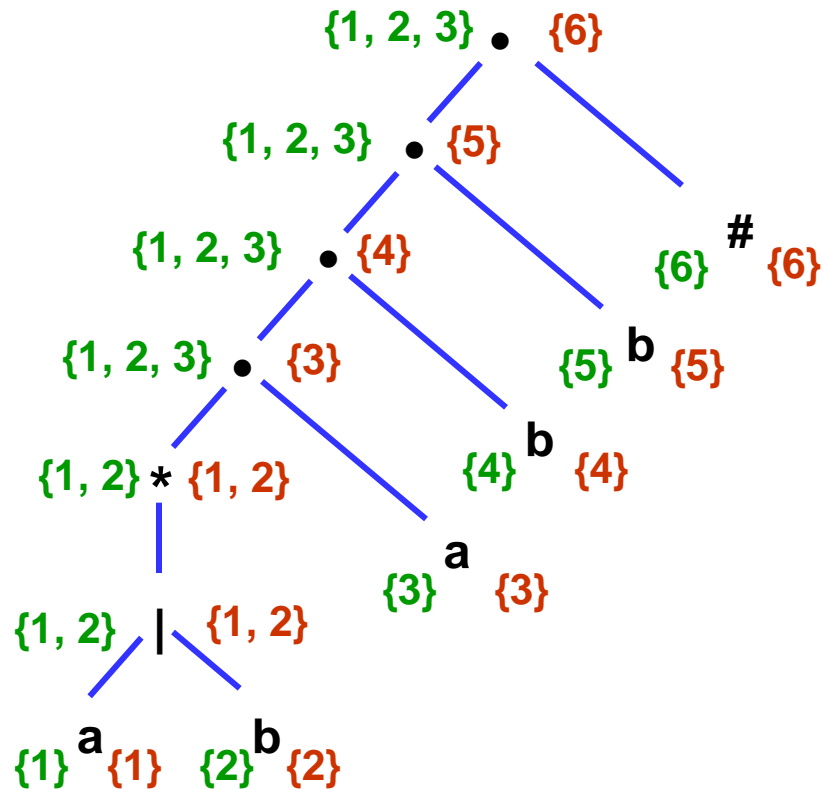
FOLLOWPOS EXAMPLE



Node	Followpos
1	{1,2,3
2	{1,2,3
3	{4
4	{5
5	{6
6	{

- At next cat-node '•' is left child and 'b' is right child
 - $lastpos(\bullet) = \{3\}$ and $firstpos(b) = \{4\}$
 - According to Rule 1:
 - $followpos\{3\} = \{4\}$
- Similarly, $followpos\{4\} = \{5\}$ and $followpos\{5\} = \{6\}$

FOLLOWPOS EXAMPLE

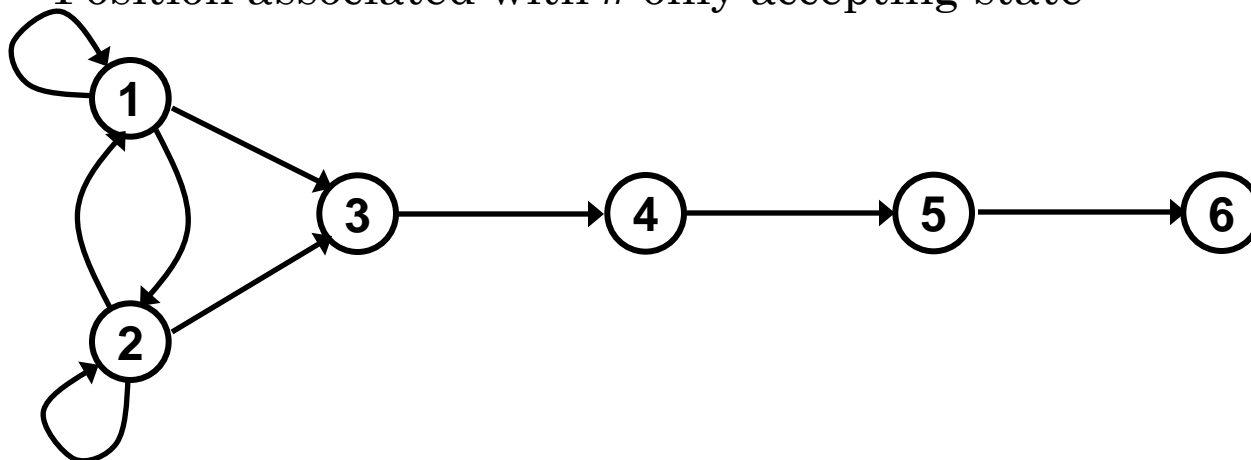


Node	followpos
1	$\{1,2,3\}$
2	$\{1,2,3\}$
3	$\{4\}$
4	$\{5\}$
5	$\{6\}$
6	-

FOLLOWPOS GRAPH

- A node for each position
- Edge from node i to node j if $j \in \text{followpos}\{i\}$
- *followpos* graph becomes equivalent NFA without ϵ -transition if
 - All positions in *firstpos* of root become start state
 - Label edge $\{i,j\}$ by the symbol at position j
 - Position associated with $\#$ only accepting state

Node	followpos
1	{1,2,3}
2	{1,2,3}
3	{4}
4	{5}
5	{6}
6	-



CONSTRUCTION OF DFA FROM RE

- Input: A regular expression r
- Output: A DFA D that recognizes $L(r)$
- Method:
 1. Construct syntax tree ST for augmented RE $r\#$
 2. Construct the functions `nullable`, `firstpos`, `lastpos` and `followpos` for ST
 3. Construct $Dstates$: set of states of D
 $Dtrans$: transition table for D

CONSTRUCTION OF DFA FROM RE

○ Algorithm

Initially, the only unmarked state in **Dstates** is *firstpos*(**root**)
while there is an unmarked state **T** in **Dstates** do begin

 Mark **T**;

 For each input symbol **a** do begin

 Let **U** be the set of positions that are in followpos(**p**) for
 some position **p** in **T** such that the symbol at position
 p is **a**

 If **U** is not empty and is not in **Dstates** then

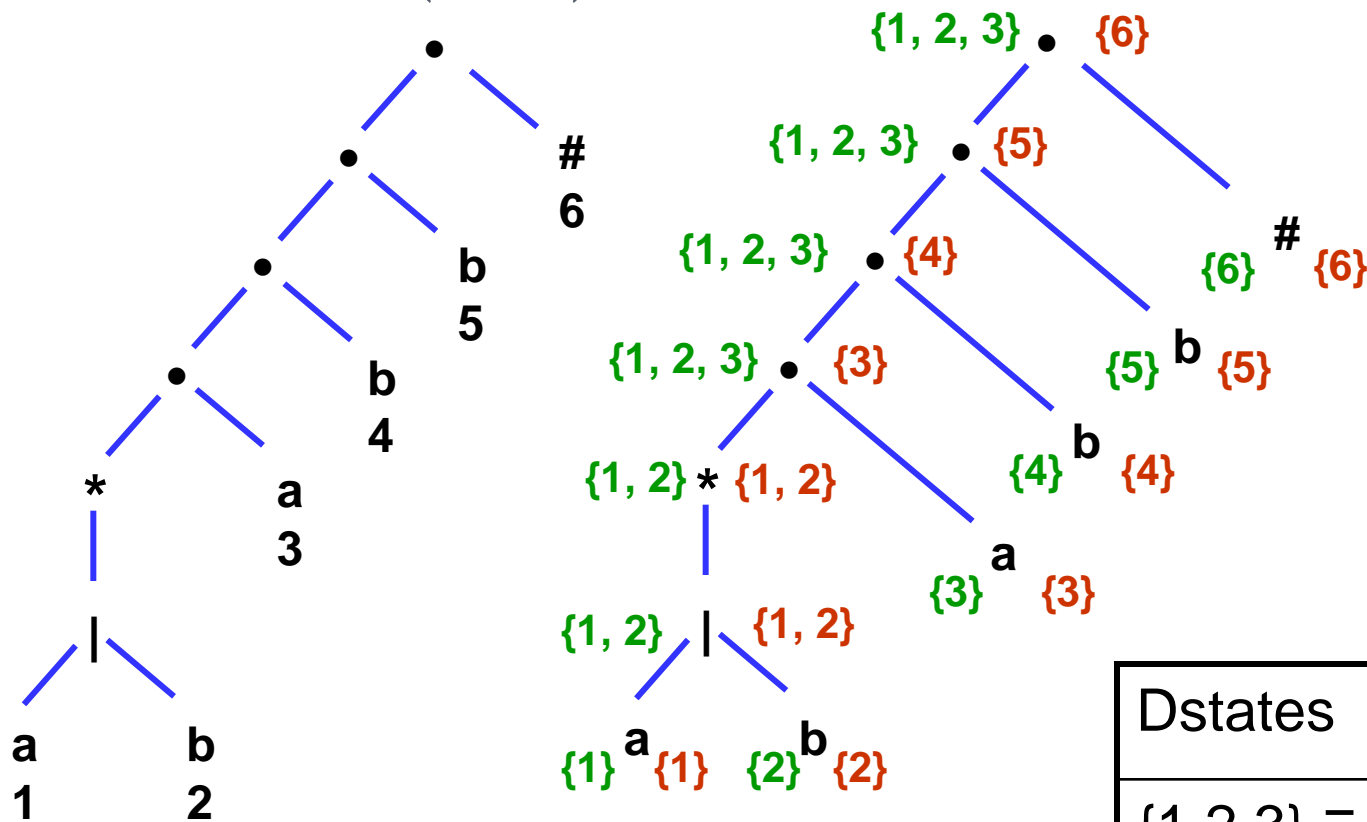
 Add **U** as an unmarked states to **Dstates**

Dtrans[**T**,**a**]=**U**

 End

end

DFA FOR $(A \mid B)^*ABB\#$



Node	followpos
1	{1,2,3}
2	{1,2,3}
3	{4}
4	{5}
5	{6}
6	-

firstpos{root} = {1,2,3} \equiv A (unmarked)

For the input symbol **a**, positions are 1, 3

\therefore followpos(1) \cup followpos{3}

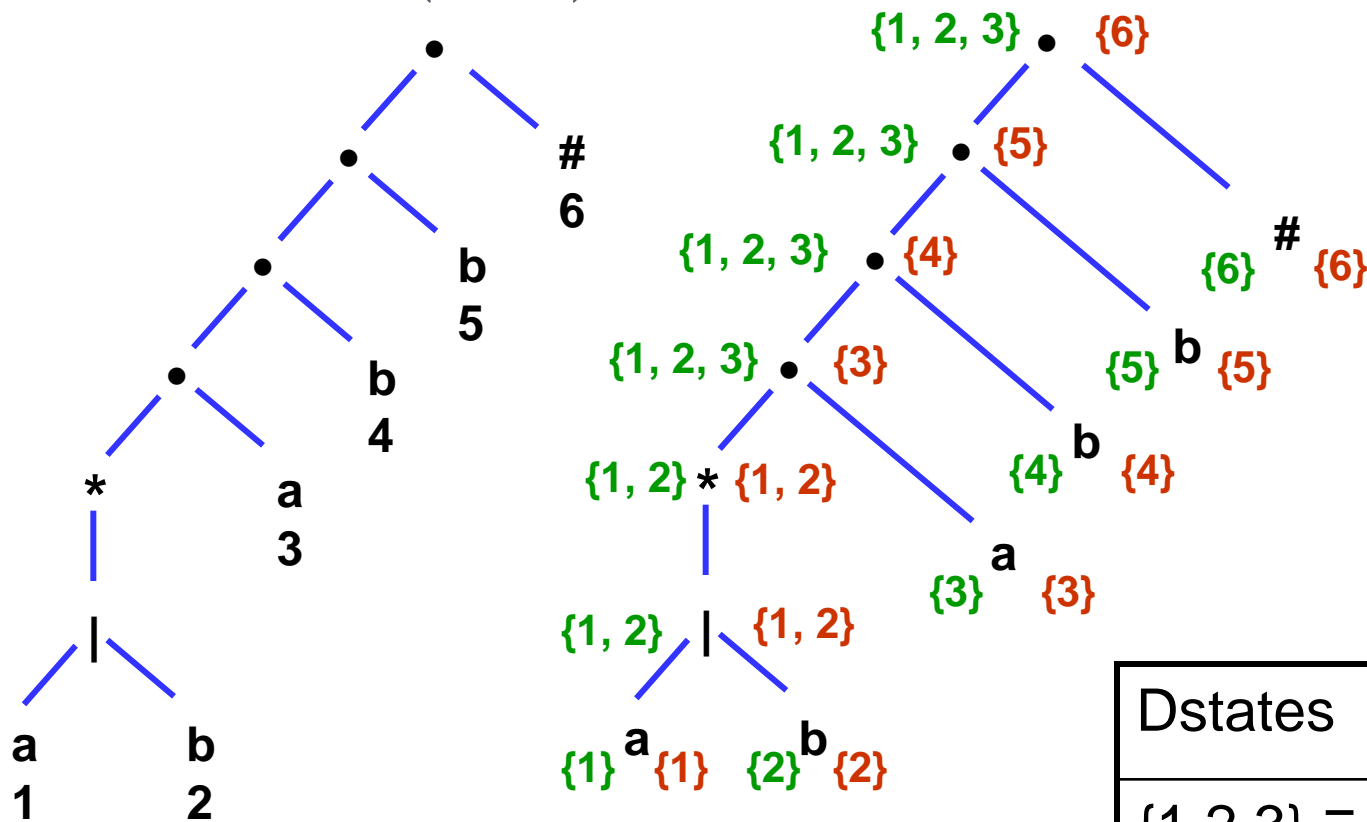
= {1,2,3,4} \equiv B

For the input symbol **b**, positions are 2

\therefore followpos(2) = {1,2,3,} \equiv A

Dstates	a	b
{1,2,3} \equiv A	B	A
{1,2,3,4} \equiv B		

DFA FOR $(A \mid B)^*ABB\#$



Node	followpos
1	{1,2,3}
2	{1,2,3}
3	{4}
4	{5}
5	{6}
6	-

$\{1,2,3,4\} \equiv B$ (unmarked)

For the input symbol **a**, positions are 1, 3

$\therefore \text{followpos}(1) \cup \text{followpos}\{3\}$

$= \{1,2,3,4\} \equiv B$

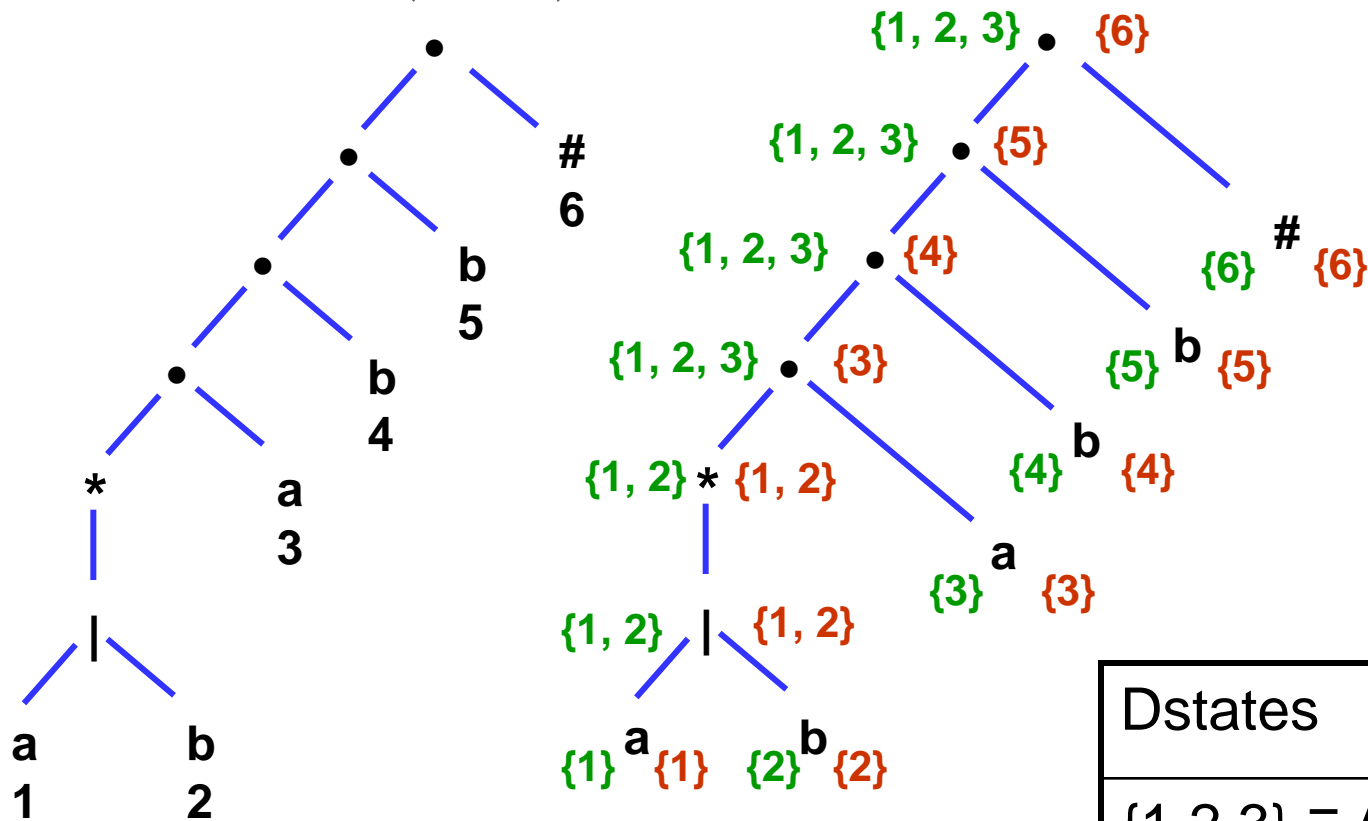
For the input symbol **b**, positions are 2, 4

$\therefore \text{followpos}(2) \cup \text{followpos}\{4\}$

$= \{1,2,3,5\} \equiv C$

Dstates	a	b
$\{1,2,3\} \equiv A$	B	A
$\{1,2,3,4\} \equiv B$	B	C
$\{1,2,3,5\} \equiv C$		

DFA FOR $(A \mid B)^*ABB\#$



Node	followpos
1	{1,2,3}
2	{1,2,3}
3	{4}
4	{5}
5	{6}
6	-

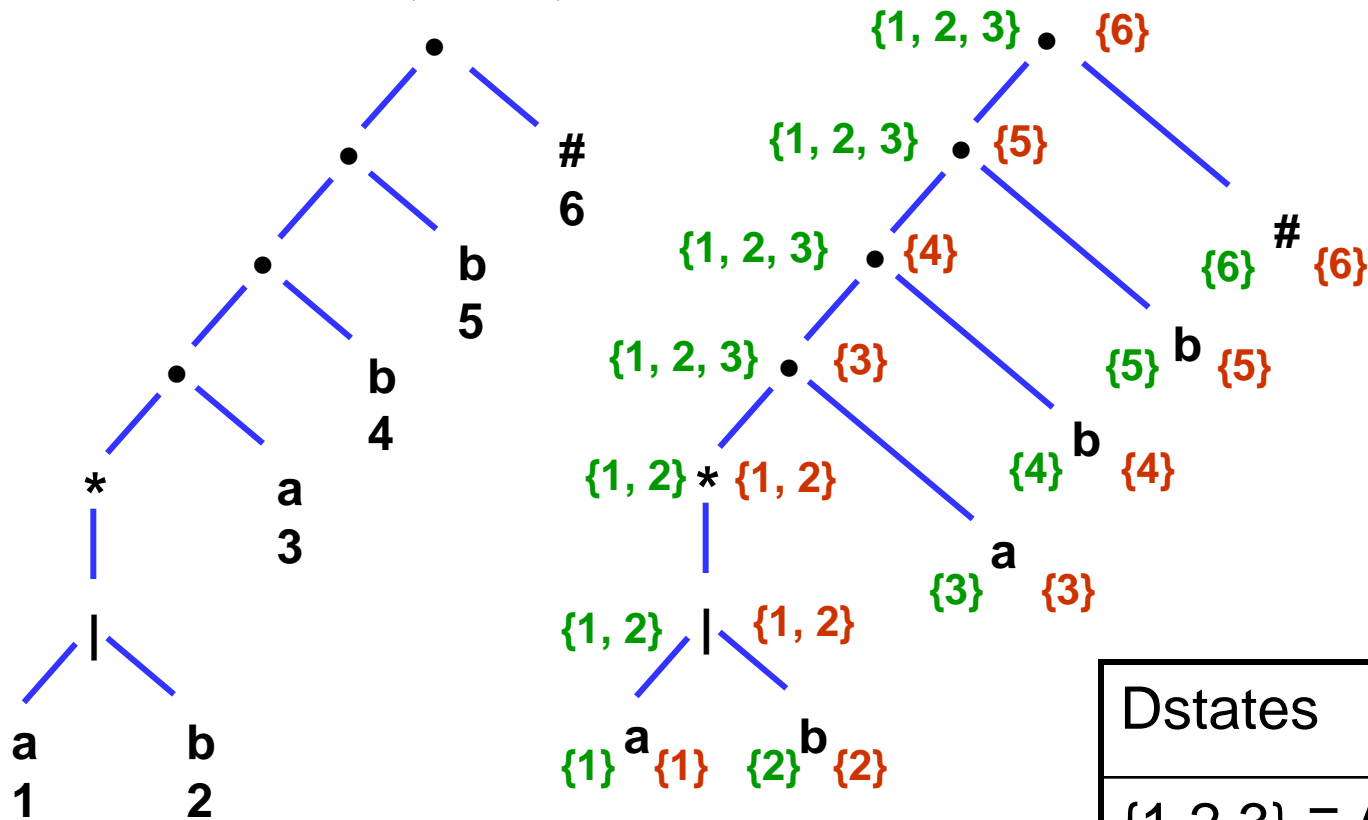
$\{1,2,3,5\} \equiv C$ (unmarked)

For the input symbol **a**, positions are 1, 3
 $\therefore \text{followpos}(1) \cup \text{followpos}\{3\}$
 $= \{1,2,3,4\} \equiv B$

For the input symbol **b**, positions are 2, 5
 $\therefore \text{followpos}(2) \cup \text{followpos}\{5\}$
 $= \{1,2,3,6\} \equiv D$

Dstates	a	b
$\{1,2,3\} \equiv A$	B	A
$\{1,2,3,4\} \equiv B$	B	C
$\{1,2,3,5\} \equiv C$	B	D
$\{1,2,3,6\} \equiv D$		

DFA FOR $(A \mid B)^*ABB\#$



Node	followpos
1	$\{1, 2, 3\}$
2	$\{1, 2, 3\}$
3	$\{4\}$
4	$\{5\}$
5	$\{6\}$
6	-

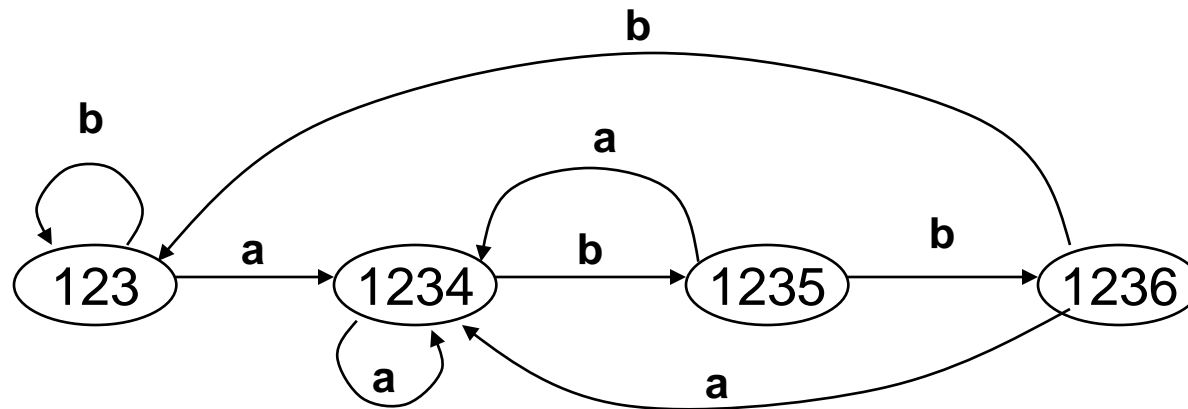
$\{1, 2, 3, 6\} \equiv D$ (unmarked)

For the input symbol **a**, positions are 1, 3
 $\therefore \text{followpos}(1) \cup \text{followpos}(3)$
 $= \{1, 2, 3, 4\} \equiv B$

For the input symbol **b**, positions are 2
 $\therefore \text{followpos}(2)$
 $= \{1, 2, 3\} \equiv A$

Dstates	a	b
$\{1, 2, 3\} \equiv A$	B	A
$\{1, 2, 3, 4\} \equiv B$	B	C
$\{1, 2, 3, 5\} \equiv C$	B	D
$\{1, 2, 3, 6\} \equiv D$	B	A

DFA FOR $(A \mid B)^*ABB\#$



ASSIGNMENT 01

- Uploaded here:
https://piazza.com/university_of_dhaka/other/cse4102/re_sources
- Last date of submission:
23/01/2017
- Hand Written Assignment in A4 paper
 - Cover page including:
your name, roll, assignment no, date, etc.

Thank You