

# Homework 4 - Design and Analysis of Algorithms - Naiara Alonso Montes

## Problem 1

### Original problem

$$\begin{aligned}
 &\min 2x_1 - 3x_2 + 4x_3 - x_4 \\
 &s. t. \quad x_1 + 2x_2 - x_4 \leq 3 \\
 &\quad \quad 2x_1 + x_2 - 3x_3 \geq -1 \\
 &\quad \quad x_1 - x_2 + 3x_3 + 3x_4 = 2 \\
 &\quad \quad \quad x_1 \geq 0 \\
 &\quad \quad \quad x_4 \geq 0
 \end{aligned}$$

### Process

1. From *min* to *max*

$$\begin{aligned}
 &\max -2x_1 + 3x_2 - 4x_3 + x_4 \\
 &s. t. \quad x_1 + 2x_2 - x_4 \leq 3 \\
 &\quad \quad 2x_1 + x_2 - 3x_3 \geq -1 \\
 &\quad \quad x_1 - x_2 + 3x_3 + 3x_4 = 2 \\
 &\quad \quad \quad x_1 \geq 0 \\
 &\quad \quad \quad x_4 \geq 0
 \end{aligned}$$

2. From equality to inequality

$$\begin{aligned}
 &\max -2x_1 + 3x_2 - 4x_3 + x_4 \\
 &s. t. \quad x_1 + 2x_2 - x_4 \leq 3 \\
 &\quad \quad 2x_1 + x_2 - 3x_3 \geq -1 \\
 &\quad \quad x_1 - x_2 + 3x_3 + 3x_4 \leq 2 \\
 &\quad \quad x_1 - x_2 + 3x_3 + 3x_4 \geq 2 \\
 &\quad \quad \quad x_1 \geq 0 \\
 &\quad \quad \quad x_4 \geq 0
 \end{aligned}$$

3. From *greater than* to *lower than*

$$\begin{aligned}
 &\max -2x_1 + 3x_2 - 4x_3 + x_4 \\
 &s. t. \quad x_1 + 2x_2 - x_4 \leq 3 \\
 &\quad \quad -2x_1 - x_2 + 3x_3 \leq 1 \\
 &\quad \quad x_1 - x_2 + 3x_3 + 3x_4 \leq 2 \\
 &\quad \quad -x_1 + x_2 - 3x_3 - 3x_4 \leq 2 \\
 &\quad \quad \quad x_1 \geq 0 \\
 &\quad \quad \quad x_4 \geq 0
 \end{aligned}$$

4. From *unbounded* to *bounded*

$$x_1; x_2 = z_2 - z_3; x_4 = z_4 - z_5; x_4 = z_6$$

$$\max -2z_1 + 3z_2 - 3z_3 - 4z_4 + 4z_5 + z_6$$

$$\begin{aligned} s. t. \quad & z_1 + 2z_2 - 2z_3 - z_6 \leq 3 \\ & -2z_1 - z_2 + z_3 + 3z_4 - 3z_5 \leq 1 \\ & z_1 - z_2 + z_3 + 3z_4 - 3z_5 + 3z_6 \leq 2 \\ & -z_1 + z_2 - z_3 - 3z_4 + 3z_5 - 3z_6 \leq 2 \\ & z_1 \geq 0 \\ & z_2 \geq 0 \\ & z_3 \geq 0 \\ & z_4 \geq 0 \\ & z_5 \geq 0 \\ & z_6 \geq 0 \end{aligned}$$

## 5. In matrix representation

$$b^T = (-2 \quad 3 \quad -3 \quad -4 \quad 4 \quad 1) \quad (1)$$

$$b^T = (3 \quad 1 \quad 2 \quad -2) \quad (2)$$

$$A = \begin{pmatrix} 1 & 2 & -2 & 0 & 0 & -1 \\ -2 & -1 & 1 & 3 & -3 & 0 \\ 1 & -1 & 1 & 3 & -3 & 3 \\ -1 & 1 & -1 & -3 & 3 & -3 \end{pmatrix} \quad (3)$$

## 6. Dual problem

$$\min 3y_1 + y_2 + 2y_3 - 2y_4$$

$$\begin{aligned} s. t. \quad & y_1 - 2y_2 + y_2 - y_4 \geq -2 \\ & 2y_1 - y_2 - y_3 + y_4 \geq 3 \\ & -2y_1 + y_2 + y_3 - y_4 \geq -3 \\ & 3y_2 + 3y_3 - 3y_4 \geq -4 \\ & -3y_2 - 3y_3 + 3y_4 \geq 4 \\ & -y_1 + 3y_3 - 3y_4 \geq 1 \\ & y_1 \geq 0 \\ & y_2 \geq 0 \\ & y_3 \geq 0 \\ & y_4 \geq 0 \end{aligned}$$

## Problem 2

(a)

$$\begin{aligned}
& \max f_{s,a} + f_{s,b} - f_{c,s} \\
s.t. \quad & -f_{s,a} + f_{a,c} + f_{a,t} \leq 0 \\
& f_{s,a} - f_{a,c} - f_{a,t} \leq 0 \\
& -f_{s,b} + f_{b,c} + f_{b,t} \leq 0 \\
& f_{s,b} - f_{b,c} - f_{b,t} \leq 0 \\
& -f_{a,c} - f_{b,c} + f_{c,s} + f_{c,t} \leq 0 \\
& f_{a,c} + f_{b,c} - f_{c,s} - f_{c,t} \leq 0 \\
& f_{s,a} \leq 8 \\
& f_{s,b} \leq 7 \\
& f_{a,c} \leq 4 \\
& f_{b,c} \leq 3 \\
& f_{c,s} \leq 2 \\
& f_{c,t} \leq 2 \\
& f_{a,t} \leq 3 \\
& f_{b,t} \leq 5 \\
& f_{s,a} \geq 0 \\
& f_{s,b} \geq 0 \\
& f_{a,c} \geq 0 \\
& f_{b,c} \geq 0 \\
& f_{c,s} \geq 0 \\
& f_{c,t} \geq 0 \\
& f_{a,t} \geq 0 \\
& f_{b,t} \geq 0
\end{aligned}$$

$$c^T = (1 \quad 1 \quad 0 \quad 0 \quad -1 \quad 0 \quad 0 \quad 0) \quad (4)$$

$$b^T = (0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 8 \quad 7 \quad 4 \quad 3 \quad 2 \quad 2 \quad 3 \quad 5) \quad (5)$$

$$A = \begin{pmatrix}
-1 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \\
1 & 0 & -1 & 0 & 0 & 0 & -1 & 0 \\
0 & -1 & 0 & 1 & 0 & 0 & 0 & 1 \\
0 & 1 & 0 & -1 & 0 & 0 & 0 & -1 \\
0 & 0 & -1 & -1 & 1 & 1 & 0 & 0 \\
0 & 0 & 1 & 1 & -1 & -1 & 0 & 0 \\
1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 1
\end{pmatrix} \quad (6)$$

(b)

$$\min 8y_7 + 7y_8 + 4y_9 + 3y_{10} + 2y_{11} + 2y_{12} + 3y_{13} + 5y_{14}$$

$$\begin{aligned} s. t. \quad & -y_1 + y_2 + y_7 \geq 1 \\ & -y_3 + y_4 + y_8 \geq 1 \\ & y_1 - y_2 - y_5 + y_6 + y_9 \geq 0 \\ & y_3 - y_4 - y_5 + y_6 + y_{10} \geq 0 \\ & y_5 - y_6 + y_{11} \geq -1 \\ & y_5 - y_6 + y_{12} \geq 0 \\ & y_1 - y_2 + y_{13} \geq 0 \\ & y_3 - y_4 + y_{14} \geq 0 \\ & y_1 \geq 0 \\ & y_2 \geq 0 \\ & y_3 \geq 0 \\ & y_4 \geq 0 \\ & y_5 \geq 0 \\ & y_6 \geq 0 \\ & y_7 \geq 0 \\ & y_8 \geq 0 \\ & y_9 \geq 0 \\ & y_{10} \geq 0 \\ & y_{11} \geq 0 \\ & y_{12} \geq 0 \\ & y_{13} \geq 0 \\ & y_{14} \geq 0 \end{aligned}$$

(c)

The dual of the general form is the LP for the min cut problem. Using paths we can create the general form. In this LP problem,  $E$  is the set of edges and  $V$  is the set of vertex.

- Definition of variables

$$y_e = \{\{u, v\}, \dots, \{n, m\}\} \in E$$

$$x_u = \{v_a, v_b, \dots, v_x\} \in V$$

$$x_v = \{v_a, v_b, \dots, v_x\} \in V$$

$$x_u \neq x_v$$

$P$  is the set of paths from  $s$  to  $t$

In this case,  $y_e$  will be equal to 1 if the edge is in a given cut, else if 0.

$$\begin{aligned} \min \quad & \sum_{\{x_u, x_v\}} c(\{x_u, x_v\}) \cdot y_e(\{x_u, x_v\}) \\ s. t. \quad & \sum_{\{x_u, x_v\}} y_e(\{x_u, x_v\}) \geq 1 \text{ for all paths in } P \text{ from } s \text{ to } t \\ & y_e(\{x_u, x_v\}) \in [0, 1] \end{aligned}$$

The result of the LP problem will provide the value of the cut.

(d)

By the constraint from the LP problem of the min-cut provided in the previous exercise, if an edge in a path belongs to the a given cut, the result of the constraint must be at least equal to 1, as all edges  $y_e$  were assigned with value 1.

## Problem 3

(a) and (b)

$x_1 = \text{type I}; x_2 = \text{type II}$

$$\max 100x_1 + 120x_2$$

$$s. t. \quad 4x_1 + 6x_2 \leq 90$$

$$3x_1 + 3x_2 \leq 60$$

$$x_1 \geq 0$$

$$x_2 \geq 0$$

I choose to maximize  $x_1$

$$x_1 = 0; x_2 = 0$$

$$y_1 = ?; y_2 = x_2$$

$$y_1 = 60 - 3x_1 + 3y_2$$

$$x_1 = -\frac{1}{3}y_1 + 20 - y_2; x_2 = y_2$$

$$100 \cdot \left(-\frac{1}{3}y_1 + 20 - y_2\right) + 120y_2 = -\frac{100}{3}y_1 + 2000 - 100y_2 + 120y_2 = -\frac{100}{3}y_1 +$$

$$-\frac{4}{3}y_1 + 80 - 4y_2 + 6y_2 \leq 90 = -\frac{4}{3}y_1 + 2y_2 \leq 10$$

$$\max_{y_1, y_2} \quad 1 - \frac{100}{3}y_1 + 20y_2 + 2000$$

$$s. t. \quad -\frac{4}{3}y_1 + 2y_2 \leq 10$$

$$-y_1 + 6y_2 \leq 0$$

$$-\frac{1}{3}y_1 - y_2 \geq -20$$

$$y_2 \geq 0$$

I choose to maximize  $y_2$

$$y_1 = 0; y_2 = 20$$

$$z_1 = y_1; z_2 = 10 + \frac{4}{3}y_1 - 2y_2$$

$$y_1 = z_1; y_2 = 5 + \frac{2}{3}z_1 - \frac{1}{2}z_2$$

$$-\frac{100}{3}z_1 + 20 \cdot \left(5 + \frac{2}{3}z_1 - \frac{1}{2}z_2\right) + 2000 = -20z_1 - 10z_2 + 100$$

It is not possible to maximize anymore

$$z_1 = 1; z_2 = 0$$

$$y_1 = z_1 = 0; y_2 = 5 + \frac{2}{3} \cdot 0 - \frac{1}{2} \cdot 0 = 5$$

$$x_1 = -\frac{1}{3} \cdot 0 + 20 - 5 = 15; x_2 = y_2 = 5$$

## Solution

$x_1 = 15$  and  $x_2 = 5$ . The constraints are met.

## (c) and (d)

The dual representation is:

$$\begin{aligned} \min & 90y_1 + 60y_2 \\ \text{s. t. } & 4y_1 + 3y_2 \geq 100 \\ & 6y_1 + 3y_2 \geq 120 \\ & y_1 \geq 0 \\ & y_2 \geq 0 \end{aligned}$$

It can be solved using a system of 2 linear equations, and then checking if the constraints are met, in this case the result of the min objective function must be equal to the result of the max function.

$$\begin{aligned} 4y_1 + 3y_2 &= 100 \\ 6y_1 + 3y_2 &= 120 \\ y_2 &= 40 - 2y_1 \\ 4y_1 + 3(40 - 2y_1) &= 100 \\ y_1 &= 10 \\ y_2 &= 40 - 2(10) \\ y_2 &= 20 \end{aligned}$$

Both min and max results match.

In [ ]: `!jupyter nbconvert --to html HW4.ipynb`