

## 2. Arrays and Structures

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# Agenda

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- Arrays
- Dynamically Allocated Arrays
- Structures
- (Polynomials)
- (Sparse Matrices)
- Representation of Multidimensional Arrays
- Ordered Lists → Not in text
- Strings

# What is an Array?

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- Conventional definition of array
  - “A consecutive set of memory location”



< Array A of size n >

- Easy, but it's based on perspective of implementation  
→ Let's try to find deeper understanding about array
- “What is **the essence** of the array?”  
→ Let's think about an alternative definition as ADT.
  - Separate implementation details from definition of array.

# Array As an ADT

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- Abstract Data Type *Array*
  - Objects: a set of pairs  $\langle \text{index}, \text{value} \rangle$ 
    - Each value of index is mapped with a value from the set item
    - *Index*: finite ordered set of one or more dim.
      - Ex)  $\{0, \dots, n-1\}$ : 1D index
      - $\{(0,0), (0,1), \dots, (n-1, m-1)\}$  : 2D index

“type” of  
elements



# Operations of Array

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- Create
  - Creates an array given its size
- Retrieve
  - Retrieve an element given its index
- Store
  - Store an element given its index and content
- Destroy
  - Destroy the array

# Array As an ADT

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- Abstract Data Type *Array*

- Operations

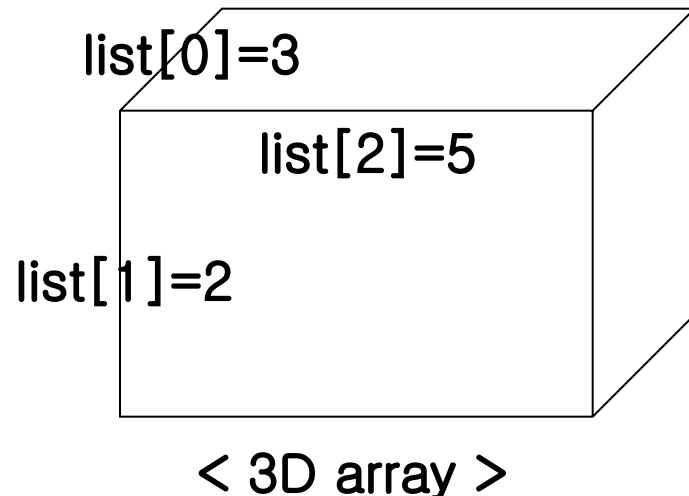
For  $A \in \text{Array}$ ,  $i \in \text{index}$ ,  $x \in \text{item}$  and  $j, \text{size} \in \text{integer}$

- **Array Create(j, list)**

$::=$  return an *Array* of  $j$  dims, where list is a  $j$ -tuple whose  $k$ -th element is the size of  $k$ -th dim. Items are undefined.

$j = 3$  (dim.)  
**int A[3][2][5];**

list[0]: 3 (depth)
list[1]: 2 (height)
list[2]: 5 (width)



# Array As an ADT

---

- Abstract Data Type *Array*

- Operations (cont.)

For  $A \in \text{Array}$ ,  $i \in \text{index}$ ,  $x \in \text{item}$  and  $j, \text{size} \in \text{integer}$

- Item Retrieve( $A, i$ )

$::=$  if( $i \in \text{index}$ ), return item associated with  $i$  in Array  $A$   
else return error

**value = A[i];**

- Array Store( $A, i, x$ )

$::=$  if( $i \in \text{index}$ ), return an array that is identical to Array  $A$  except the new pair  $\langle i, x \rangle$  has been inserted  
else return error.

**A[i] = x;** is equivalent to “ $A = \text{Store}(A, i, x);$ ”

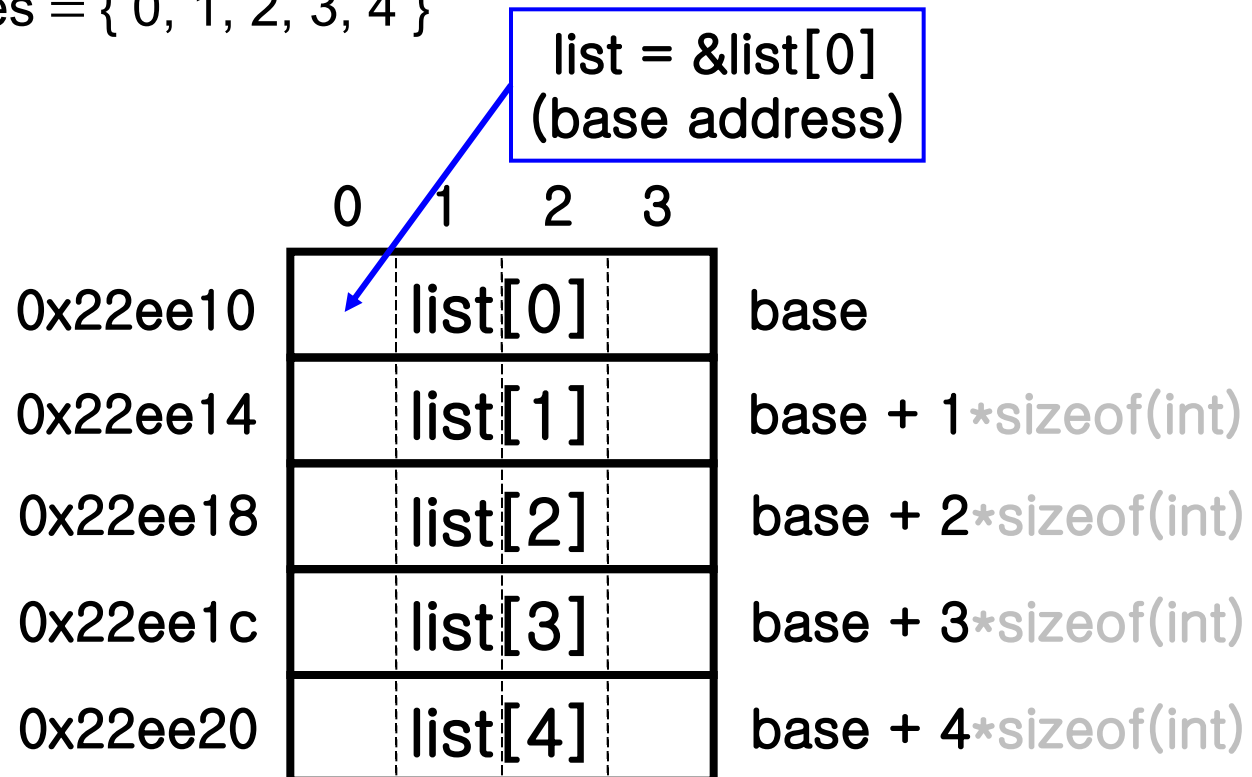
# Array in C Language

- Can be viewed as an implementation of ADT array

Ex) `int list[5];` // integer array containing 5 elements

`int *plist[5];` // pointer array containing 5 elements

- Indices = { 0, 1, 2, 3, 4 }





# Array in C Language

---

- Addition/subtraction operation on pointer implies multiplication by sizeof(type)
  - for `int *p`,  
if `p == 0x22ee10`,  
then `p+i == 0x22ee10 + i * sizeof(int)`
  - `A[i] == *(A+i);` // In C, array is implemented by pointer

cf. How can we access  $p + i$  (in bytes)?,

```
int *pp = (int *)((char *)p + i);
```

Note! **`(char *)p + i == 0x22ee10 + i`**, because `sizeof(char) = 1`

# Array and Pointer

---

- size information of array is contained in array.

```
void function(int array_arg[])
{
    int noEntry = sizeof(array_arg) / sizeof(int);    //we don't know the size
    printf("(%d, %d)\n", sizeof(array_arg), noEntry);
}

void main()
{
    int array[5];
    int noEntry = sizeof(array) / sizeof(int);

    printf("(%d, %d)\n", sizeof(array), noEntry);

    function (array);
}
```

**practice**

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- (Sparse Matrices)
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# One-Dimensional Arrays

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- Statically allocated array vs. dynamically allocated array
  - `int list[100];`
  - `int *list = (int*)malloc(100 * sizeof(int));`
    - Later list should be deallocated by “`free(list);`”
  - After allocation, its use is almost the same as the statically allocated array.  
Ex) `for(i = 0; i < 100; i++)`  
    `list[i] = i;`
- Why dynamically allocated array?
  - **The size of array is decided at runtime.**

*calloc*  
allocates bytes  
and initializes them to 0

# Allocating 1D Array

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Ex) Reading a series of numbers from the user

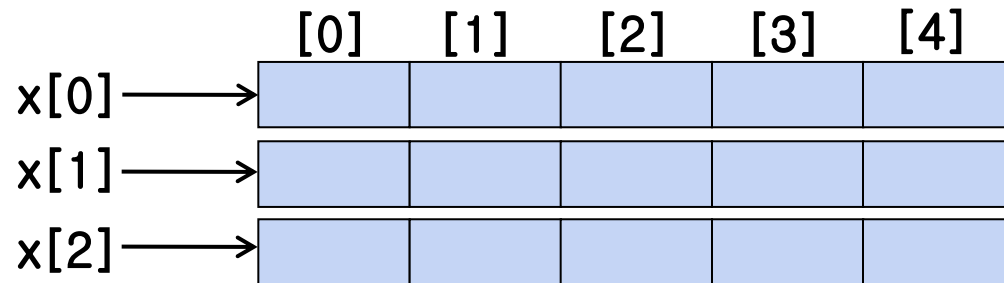
```
int n = 0, *list = NULL;
printf("Enter the number of integers to read: ");
scanf("%d", &n);
if (n < 1) {
    printf( "Improper value of n. \n");
    exit(-1);
}

list = malloc(n * sizeof(int));
if(list == NULL){
    printf( "Failed to allocate memory.\n");
    exit(-1);
}
```

# Two-Dimensional Arrays

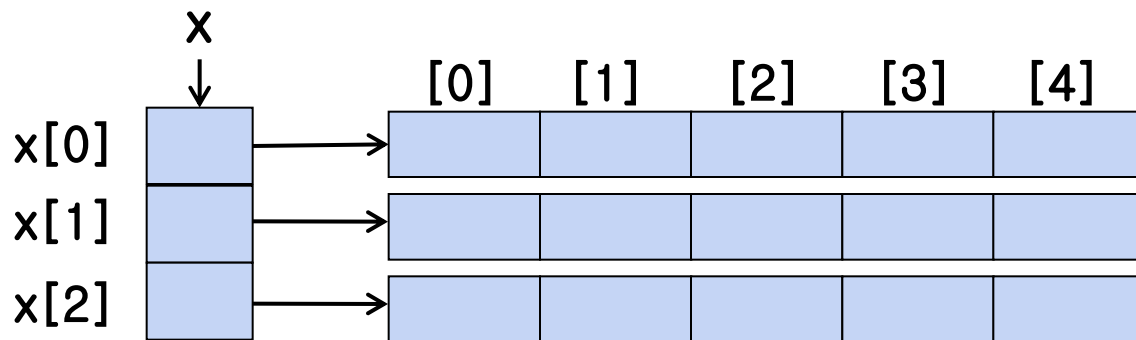
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- 2D is represented as a 1D array, where each element is a 1D array  
Ex) `int x[3][5];`



- Dynamically allocated 2D array

4 memory blocks



# Allocating 2D Arrays

---

- Function to allocate 2D array

```
int ** make2dArray(int rows, int cols)
{ /* create a two dimensional rows X cols array */
    int **x = NULL, i = 0;

    x = malloc(rows * sizeof (*x)); /* get memory for row pointers */

    for (i = 0; i < rows; i++)      /* get memory for each row */
        x[i] = malloc(cols * sizeof (**x));

    return x;
}
```

**practice**

- Usage

```
int **myArray = NULL;
myArray = make2dArray(3,5);
myArray[1][2] = 6;
```

# Deallocating 2D Arrays

---

- Function to deallocate 2D array

```
void free2dArray(int **array, int rows)
// deallocates a two dimensional array
{
    int i = 0;

    for (i = 0; i < rows; i++) /* free memory for each row */
        free(array[i]);
    free(array);                /* free memory for row pointers */

    return;
}
```

**practice**

- Usage

```
int **myArray = NULL;
myArray = make2dArray(3,5);
myArray[1][2] = 6;
...
free2dArray(myArray, 5);
```



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# Structures

---

- **Struct:** collection of data items, each item may differ in type

```
struct tPerson {  
    char name[10];    // name in char array  
    int age;          // age in integer  
    float salary;     // salary of person in float  
} person1;  
  
struct tPerson person2;  
    // equivalent to the previous declaration of "person1"
```

Diagram annotations:

- A blue callout box labeled "tag" points to the `struct tPerson` definition.
- A blue callout box labeled "instance name" points to the `person1` instance declaration.

- Access to member: '.' (member operator), "->" operator  
Ex) `person1.name`, `person1.age`, `person1.salary`

# Why Structures ?

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- Structures explicitly represent the relation among attributes.

Ex) Representation of 20 persons

- Using arrays

```
char name[20][10];  
int age[20];  
float salary[20];
```

→ Difficult to understand  
relation between attributes

- Using structure

```
struct tPerson {  
    char name[10];  
    int age;  
    float salary;  
} person[20];
```

→ Easy to understand

**name[k], age[k], salary[k]**

**vs.**

**person[k]**

# Structures

---

- Struct can be defined as a type by *typedef*

Ex) **typedef** struct {    // tag was omitted  
      // members

} human\_being;

type name

**typedef** struct tPerson human\_being;

type name

tag

**typedef** struct tPerson{

      // members

} human\_being ;

tag

type name

# To make instances

---

```
struct tPerson {  
    char name[10];  
    int age;  
} person1;
```

```
struct tPerson {  
    char name[10];  
    int age;  
};  
struct tPerson person1;
```

```
typedef struct {  
    char name[10];  
    int age;  
} Tperson;  
Tperson person1;
```

```
typedef struct tPerson Tperson;  
struct tPerson {  
    char name[10];  
    int age;  
};  
Tperson person1;
```

# Structures

---

- Comparison of structure variables is **NOT defined** in C language

Ex) struct Person person1, person2;

printf("person1 == person2 = %d\n", person1 == person2);

**// compile error occurs**

**cf) Comparison of array is defined as comparison of pointer**

- **Assignment of structure is provided in ANSI C**

Ex) struct Person person1, person2;

...

**person2 = person1;      // all members are copied**

- But not available in old-style C

**cf) Assignment of array is NOT allowed.**

# Comparison of Structure Variables

---

- A function to compare two instances of human\_being structure

```
int humansEqual(human_being person1, human_being person2)
{
    if(strcmp(person1.name, person2.name))
        return FALSE;
    if(person1.age != person2.age)
        return FALSE;
    if(person1.salary != person2.salary)
        return FALSE;

    return TRUE;
}
```

# Structures as Parameters

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- Passing struct as parameter

```
struct MyStruct {  
    char *name;  
    int age;  
    // ...  
};
```

```
int main()  
{  
    struct MyStruct a;  
    Print(a);  
    PrintPtr(&a);  
}
```

```
void Print(struct MyStruct a)  
{  
    // assignment of structure  
    printf("name = %s\n", a.name);  
    ...  
}
```

```
void PrintPtr(struct MyStruct *p)  
{  
    printf("name = %s\n", p->name);  
    ...  
}
```

→ more efficient  
Why?



# Embedded Structures

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- Embedded structure

```
typedef struct {  
    int month;  
    int day;  
    int year;  
} date;
```

```
typedef struct {  
    char name[10];  
    ...  
    date dob; // date of birth  
} human_being;
```

```
Ex) human_being john;  
    john.dob.year = 1971;  
    john.dob.month = 4;  
    john.dob.day = 15;
```

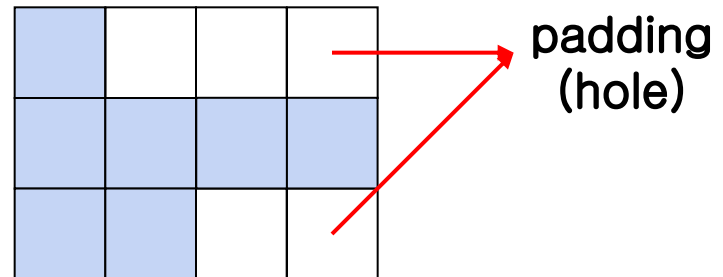
# Structures

---

- Internal representation of structure
  - Padding (or hole) can be included for memory alignment
    - To improve access speed
    - $\text{sizeof}(\text{struct}) \neq \sum \text{sizeof}(\text{members})$
  - Memory alignment often improves efficiency on many CPU's.

Ex)

```
struct MyStruct {  
    char ch;  
    int i;  
    short s;  
}
```

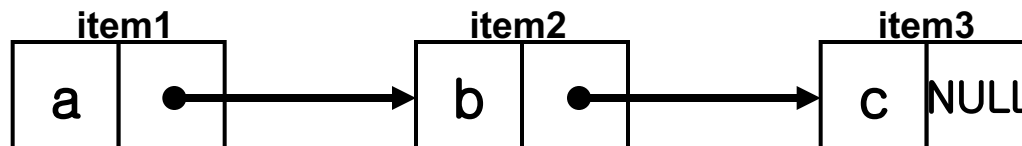


# Structures

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- Self-referential structures

```
typedef struct tList {  
    char data;  
    struct tList *link;           // pointer to the same type  
} List;  
List item1, item2, item3;  
item1.data = 'a';    item2.data = 'b';    item3.data = 'c';  
item1.link = &item2;  
item2.link = &item3;  
item3.link = NULL;
```



# Structures

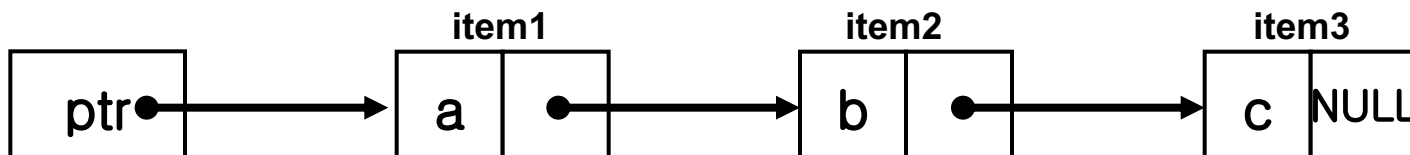
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- Self-referential structures

```
List *ptr;
```

```
ptr = &item1;
```

```
while (ptr){  
    printf("%c\n", ptr->data);  
    ptr = ptr->link;  
}
```



# Contacts with structure

---

- read and parse lines from a file.
  - each line of files is in the following format.
    - name; birthday; email; phone\_number
    - ex) henry choi; 20190303; hchoi@handong.edu; 010-1234-5678
- add the data into array of structure named 'Contact'

**practice**

# Agenda

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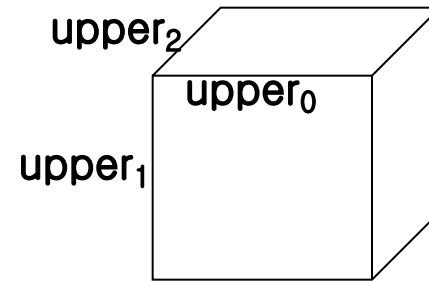
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# Multidimensional Arrays

- Multidimensional array

$A[\text{upper}_0][\text{upper}_1][\text{upper}_2]$

- Total size =  $\prod \text{upper}_i$



- Internal representation: multidimensional array are serialized into 1D memory

Ex) 2D array  $A[m][n]$

$$\begin{array}{c}
 \text{m} \left[ \begin{array}{c} \text{n} \\ \hline \begin{bmatrix} A[0][0], A[0][1], \dots, A[0][n-1] \\ A[1][0], A[1][1], \dots, A[1][n-1] \\ \dots, A[i][j], \dots \\ A[m-1][0], A[m-1][1], \dots, A[m-1][n-1] \end{bmatrix} \end{array} \right. \\
 \text{< 2D array >}
 \end{array}$$

$A[0][0], A[0][1], \dots,$   
 $A[0][n-1], A[1][0],$   
 $A[1][1], \dots, A[1][n-1], \dots,$   
 $A[m-1][0], A[m-1][1], \dots,$   
 $A[m-1][n-1]$

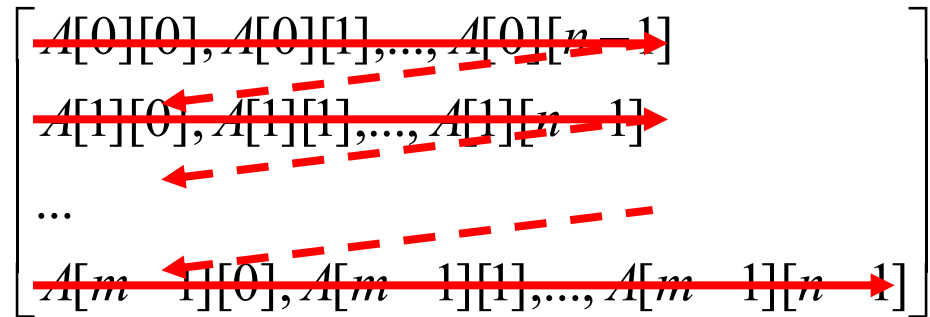
< Memory >

# Multidimensional Arrays

- Representation of 2D array

- Row-major format**

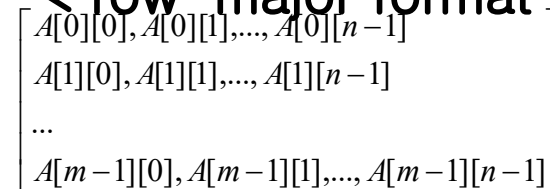
- $A[i][j] \equiv *((\text{int}^*)A + i * n + j)$



< row-major format >

- Column-major format (less popular)

- $A[i][j] \equiv *((\text{int}^*)A + i + j * m)$



< column-major format >



# Example

---

- Test program

```
const int rows = 3 , cols = 5;  
int array2D[rows][cols];  
int *array1D = (int *)array2D;  
int i = 0, j = 0, n = 0;
```

```
for(i = 0; i < rows; i++)  
    for(j = 0; j < cols; j++)  
        array2D[i][j] = n++;
```

```
for(i = 0; i < rows; i++)  
    for(j = 0; j < cols; j++)  
        printf("array2D[%d][%d](%p) = %d, array1D[%d](%p) = %d\n",  
            i, j, &array2D[i][j], array2D[i][j],  
            i*cols+j, &array1D[i*cols+j], array1D[i*cols+j]);
```

0	1	2	3	4
5	6	7	8	9
10	11	12	13	14

# Example

---

- Result

```
array2D[0][0] (0x22ede0) = 0, array1D[0] (0x22ede0) = 0
array2D[0][1] (0x22ede4) = 1, array1D[1] (0x22ede4) = 1
array2D[0][2] (0x22ede8) = 2, array1D[2] (0x22ede8) = 2
array2D[0][3] (0x22edec) = 3, array1D[3] (0x22edec) = 3
array2D[0][4] (0x22edf0) = 4, array1D[4] (0x22edf0) = 4
array2D[1][0] (0x22edf4) = 5, array1D[5] (0x22edf4) = 5
array2D[1][1] (0x22edf8) = 6, array1D[6] (0x22edf8) = 6
array2D[1][2] (0x22edfc) = 7, array1D[7] (0x22edfc) = 7
array2D[1][3] (0x22ee00) = 8, array1D[8] (0x22ee00) = 8
array2D[1][4] (0x22ee04) = 9, array1D[9] (0x22ee04) = 9
array2D[2][0] (0x22ee08) = 10, array1D[10] (0x22ee08) = 10
array2D[2][1] (0x22ee0c) = 11, array1D[11] (0x22ee0c) = 11
array2D[2][2] (0x22ee10) = 12, array1D[12] (0x22ee10) = 12
array2D[2][3] (0x22ee14) = 13, array1D[13] (0x22ee14) = 13
array2D[2][4] (0x22ee18) = 14, array1D[14] (0x22ee18) = 14
```

# Multidimensional Arrays

---

- Generalized row-major format
  - `int A[upper0] [upper1]...[uppern-1]`
    - $A[i_0] [i_1] \dots [i_{n-1}] = *((\text{int}^*)A + i_0 * \text{upper}_1 * \text{upper}_2 \dots \text{upper}_{n-1}$   
 $+ i_1 * \text{upper}_2 \dots \text{upper}_{n-1}$   
 $+ i_2 * \text{upper}_3 \dots \text{upper}_{n-1}$   
 $+ i_{n-2} * \text{upper}_{n-1}$   
 $+ i_{n-1})$

$$= *(A + \sum_{j=0}^{n-1} i_j a_j) \quad \begin{cases} a_j = \prod_{k=j+1}^{n-1} \text{upper}_k, \text{ for } 0 \leq j < n-1 \\ a_{n-1} = 1 \end{cases}$$

# Efficient Access to 2D Array [sup]

---

- Efficient way

```
int matrix[numRows][numColumns];
for ( row = 0; row < numRows; row++ ) {
    for ( column = 0; column < numColumns; column++ ){
        matrix[ row ][ column ] = 0;
    }
}
```

**assuming row major format**

- Inefficient way

```
for ( column = 0; column < numColumns; column++ ){
    for ( row = 0; row < numRows; row++ ) {
        matrix[ row ][ column ] = 0;
    }
}
```

# Efficient Access to 2D Array [sup]

---

- 1D implementation

```
int matrix1D[numRows*numColumns];  
for (i = 0; i < numRows*numColumns; i++ )  
    matrix1D[i] = 0;
```

**practice  
check the execution time**

→ About 11% faster than 2D array initialization in C++

→ About 47% faster than 2D array initialization in Java

**High-dimensional arrays are usually more expensive than low-dimensional array**

- More optimized version (?): pointer operation

```
int *pLimit = matrix1D + numRows*numColumns;  
int *p;  
for(p = matrix1D; p < pLimit; p++)  
    *p = 0;
```

→ In some environment, this optimization provides little improvement in efficiency.

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# Ordered List [Aho, Hopcroft, Ullman]

---

- **Ordered list**: a sequence of zero or more elements of a given type

$$a_0, a_1, a_2, \dots, a_{n-1}$$

- n: length of list
- Important property: all elements are **linearly ordered** according to its 'position'
  - Implementation of 'position' can vary

Ex) Student attendance roll (ordered by student number)

- Implementation: **array**, linked list, cursor, ...

# Ordered List

---

- Operations on ordered list

For  $L \in \text{list}$ ,  $x \in \text{element}$ ,  $p \in \text{position}$

- Insert( $L, x, p$ )** ::= insert  $x$  at position  $p$  in list  $L$   
Ex) **Insert( $\langle a, b, c, \underline{d}, e \rangle, n, \underline{3}$ )**  $\rightarrow L = \langle a, b, c, \underline{n}, d, e \rangle$
- Delete( $L, p$ )** ::= delete element at position  $p$  on list  $L$   
Ex) **Delete( $\langle a, b, \underline{c}, d, e \rangle, \underline{2}$ )**  $\rightarrow L = \langle a, b, d, e \rangle$
- MakeNull( $L$ )** ::= make  $L$  an empty list
- Locate( $L, x$ )** ::= return position of  $x$  on list  $L$   
if  $x$  does not exist on  $L$ , return *error*
- Retrieve( $L, p$ )** ::= return element at position  $p$  on list  $L$



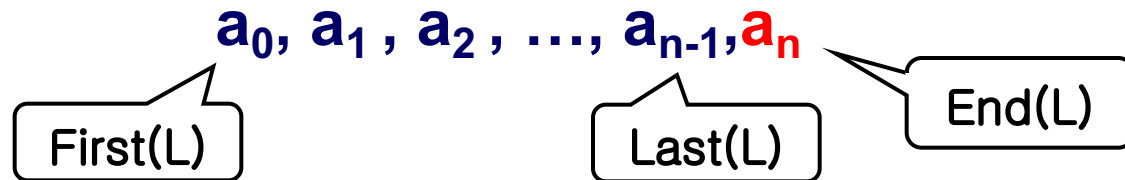
# Ordered List

---

- Operations on ordered list

For  $L \in \text{list}$ ,  $x \in \text{element}$ ,  $p \in \text{position}$

- Next/Previous( $L, p$ )** ::= return position of following/preceding position  $p$  on list  $L$
- First( $L$ )/Last( $L$ )** ::= return position of first/last element
- End( $L$ )** ::= return position of next to last element
  - End( $L$ ) is just for boundary condition, and should not be used to access the element

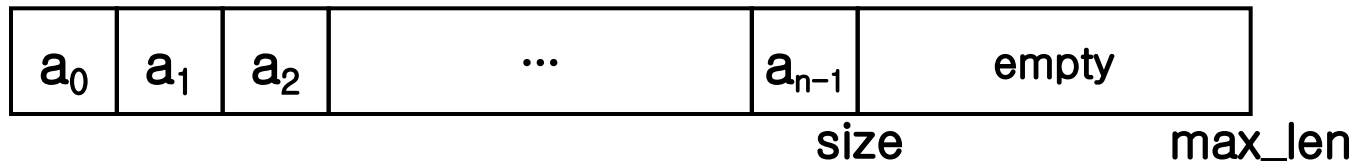


```
for(p = First(L); p < End(L); p = Next(L, p)){  
    ...  
}
```

# Array Implementation of List

---

- List elements are stored in contiguous cells of array
  - Position is represented by integer



- Data representation
  - Array **elements** to store elements, whose size is **max\_len**
  - Integer **size** to store # of elements
- Operations
  - Implementations of most operations are straightforward.
  - For Insert/Delete, elements should be shifted.

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# String

---

- **String**: concatenation of characters
  - $S = s_0s_1s_2\dots s_{n-1}$ 
    - $s_i$ : characters from the character set of the programming language

Ex) `char str[6] = "HELLO";` // in C  
// `str = 'H' + 'E' + 'L' + 'L' + 'O' + '\0'`

**Let's define String by ADT**

# String ADT

---

- Abstract Data Type *String*
  - Objects: a finite sequence of zero or more characters
  - Functions

$s, t \in \text{String}, i, j, m \in \text{non-negative integers}$

    - **String Null( $m$ )** := return a string whose maximum length is  $m$  but initially set to null string("");
    - **int Compare( $s, t$ )** := if  $s$  equals  $t$  return 0  
else if  $s$  precedes  $t$  return -1  
else return +1
    - **Boolean IsNull( $s$ )** :=  $s$  is a null string return TRUE  
else return FALSE
    - **int Length( $s$ )** := return # of characters

# String ADT

---

- Abstract Data Type *String*
  - Functions (cont.)

$s, t \in \text{String}, i, j, m \in \text{non-negative integers}$

    - **String Concat( $s, t$ )** := if  $t$  is not a null string  
return a string whose elements are  
those of  $s$  followed by those of  $t$   
else return  $s$   
Ex) Concat("abc", "123") == "abc123"
    - **String Substr( $s, i, m$ )** := if( $i > 0 \ \&\& \ (i+m-1) < \text{Length}(s)$ )  
return string containing char's of  $s$  at  
positions  $i, i+1, \dots, i+m-1$   
Ex) Substr("ABCDEF~~G~~", 3, 2) = "DE"

# String in C Language

---

- String is represented by character array terminated with null character ('\0') in C.

Ex) `char s[100] = "dog";`

0	1	2	3	99
d	o	g	\0	...

`char s[] = "dog";`

d	o	g	\0
---	---	---	----

# String Functions in C Language

---

Declared in *string.h*

- strcat, strncat : string concatenation
- strcmp, strncmp : string comparison
  - strcmp("ABC", "ABC") → 0
  - strcmp("ABC", "AB**B**") → 1
  - strcmp("ABC", "AB**D**") → -1
  - strncmp("Hello", "Hello, World", 5) → 0
- strcpy, strncpy : string copy
- strlen : string length



# String Functions in C Language

---

- strchr, strrchr, strstr: find character or substring  
char s[] = "Hello, World";
  - strchr(s, 'l') → s + 2      // search from left
  - strrchr(s, 'l') → s + 10      // search from right
- strtok(s, delim) : return token delimited by delim
  - strtok(s, ",") → "Hello"
- strspn(s, spanset), strcspn, strpbrk :  
scan s for characters in (not in) spanset  
char s[] = "Hello, World";
  - strspn(s, "Hle") → 4      // length of "Hell"
  - strcspn(s, "WXYZ") → 7      // length of "Hello, " right before 'W'
  - strpbrk(s, "WXYZ") → s + 7 // like strcspn, except pointer.

# strtok

---

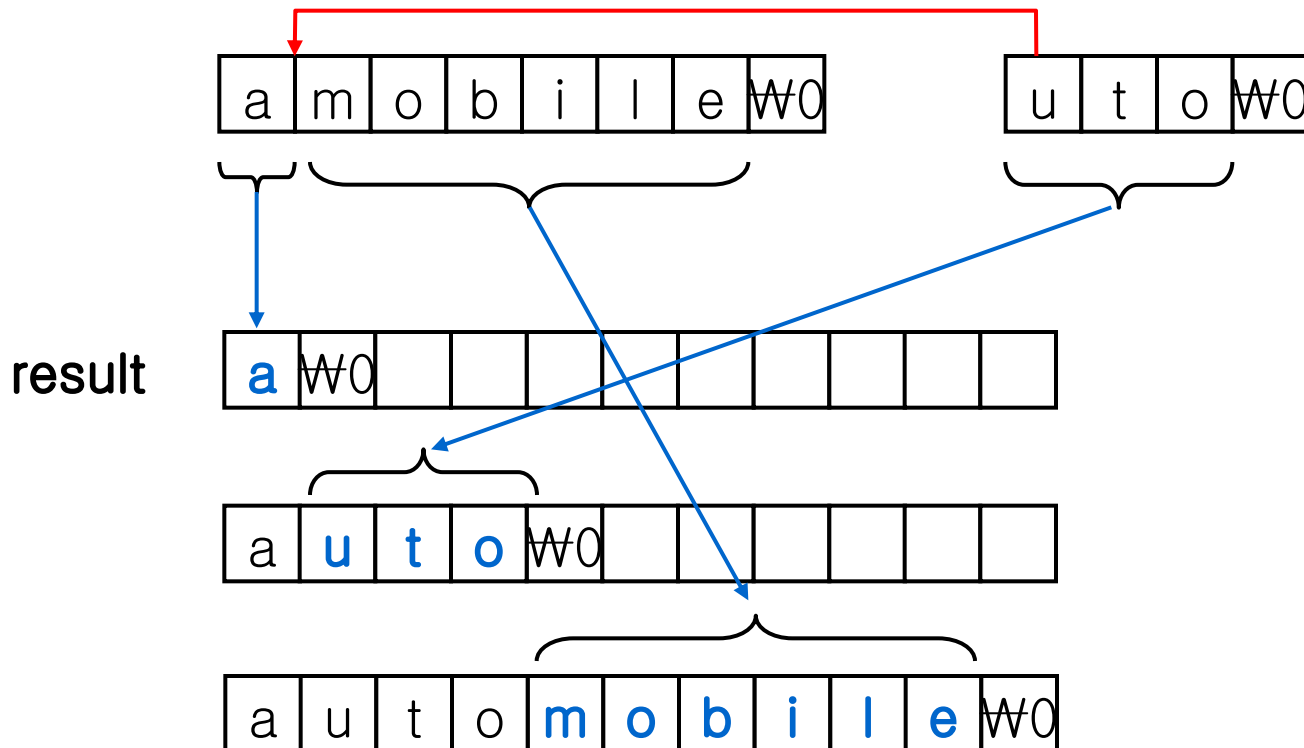
```
char array[] = "I love you, HGU";  
char *ptr = strtok(array, " ");  
  
while (ptr) {  
    printf("%s\n", ptr);  
    ptr = strtok(NULL, " ");  
}
```

**output:**

**I  
love  
you,  
HGU**

# String Insertion

Example) Given string1 = “amobile” and string2 = “uto”, insert string2 into position 1 of string1 to make “**auto**mobile”



# String Copy Functions

---

- `strcpy(dest, src)`: copy *src* to *dest*

Ex)

```
char src[] = "ABCDE";
```

```
char dest[100] = "0123456789";
```

```
strcpy(dest, src);           // dest = "ABCDE"
```

- `strncpy(dest, src, n)`: copy no more than *n* characters of *src* to *dest*.

If *n* is equal or less than length of *src*,

**'\0' is NOT appended to the result automatically.**

Ex)

```
strncpy(dest, src, 3) // dest = "ABC3456789";
```

# String Concatenation Functions

---

- `strcat(dest, src)`: append string *src* to the end of *dest*  
Ex)  
`char src[] = “, World”;`  
`char dest[100] = “Hello”;`  
`strcat(dest, src);`                      `// dest == “Hello, World”;`
- `strncat(dest, src, n)`: append no more than *n* characters of *src* to the end of *dest*. **‘\0’ is always appended.**  
Ex)  
`strncat(dest, src, 3);` `// dest == “Hello, W”`

# String Insertion

---

```
void strnins(char *s, char *t, int i)
{
    char temp[100];
    //char* temp = (char*)malloc(strlen(s)+strlen(t)+1);

    if(strlen(s) == 0)
        strcpy(s, t);
    else if(strlen(t)) {
        strncpy(temp, s, i); // \0 is not attached
        temp[i] = 0;         // don't forget
        strcat(temp, t);
        strcat(temp, s+i);
        strcpy(s, temp);
    }
    //free(temp);
}
```

practice

# String Pattern Matching

---

- Find a pattern from string
  - `strstr(string, pattern)` in standard C library
    - if(pattern exists in string) return position of pattern in string
    - else return NULL

Ex) `s = "example";`  
`p = strstr(s, "amp") → s + 2;`  
`index = p - s;`

# Simple Implementation

---

- Exhaustive matching

string

a	b	a	b	b	a	a	b	a	a	W0
---	---	---	---	---	---	---	---	---	---	----

pattern

a	a	b	W0
---	---	---	----

a	a	b	W0
---	---	---	----

a	a	b	W0
---	---	---	----

⋮

a	a	b	W0
---	---	---	----

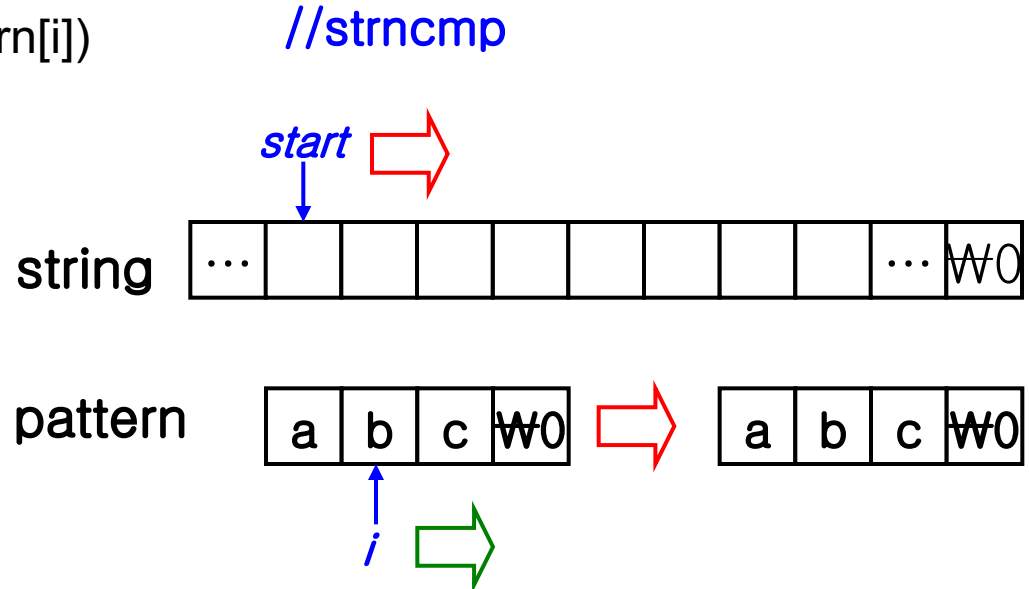
Found!



# Simple Implementation

```
int pattern_matching(char *string, char *pattern)
{
    int start, i;
    int lens = strlen(string);
    int lenp = strlen(pattern);

    for(start = 0; start + lenp <= lens; start++){
        for(i = 0; i < lenp; i++){
            if(string[start+i] != pattern[i])
                break;
        }
        if(i == lenp) //found
            return start;
    }
    return -1; // not found
}
```

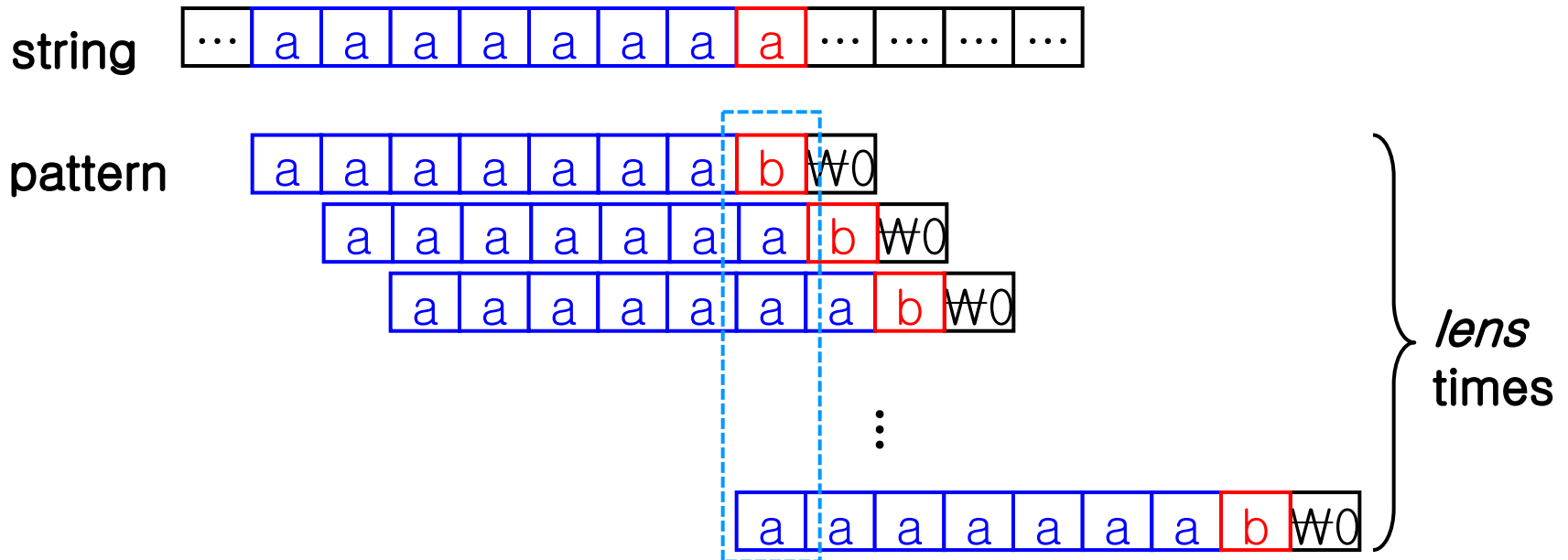


# Simple Implementation

- Complexity of pattern\_matching:  $O(\text{lens} * \text{lenp})$

Ex) string = "aa...a", pattern = "a...ab"

Each 'a's in *string* is compared with most of 'a's in *pattern*



For an efficient algorithm, check the KMP (Knuth, Morris, Pratt) algorithm

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**questions or comments?**  
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