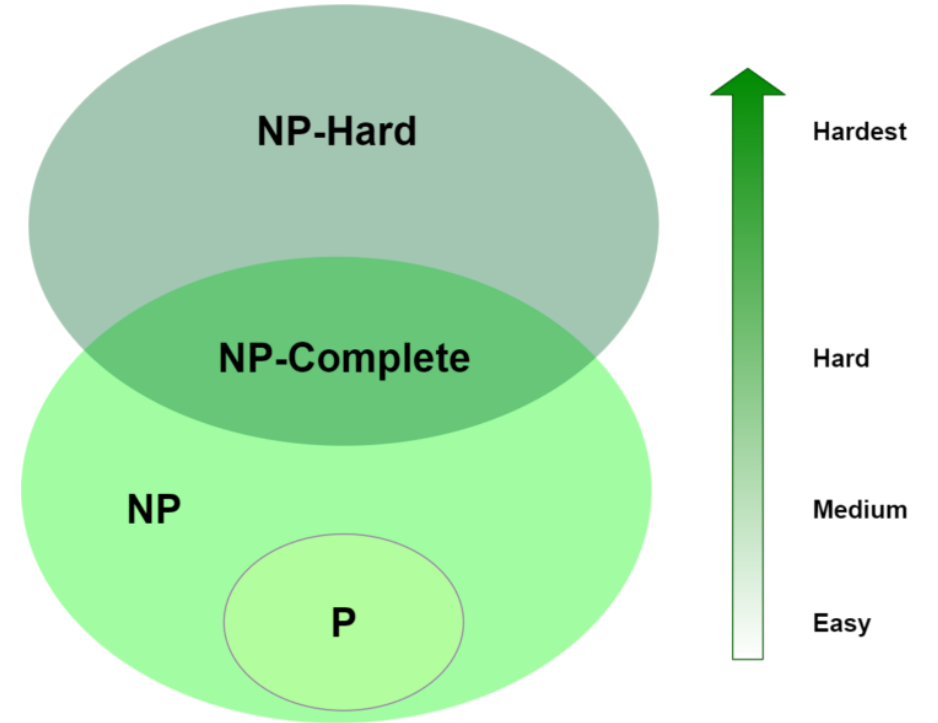


# DSE 2256 DESIGN & ANALYSIS OF ALGORITHMS

## Lecture 40

### Limitations of Algorithm Power

P, NP, NP-Complete, NP-Hard Problems



# Key Terminologies and Definitions

1. **Decision Problems:** A problem that can be posed as a **yes-no** question of the input values.

## 2. **Deterministic and Non-Deterministic Algorithms**

- In **deterministic algorithm**, for a given particular input, the computer will always produce the same output.
- In **non-deterministic algorithm**, for the same input, the computer may produce different output in different runs.

### **Formal Definition of a Non-Deterministic Algorithm:**

A non-deterministic algorithm is a **two-stage procedure** that takes as its input an **instance  $I$**  of a **decision problem** and does the following:

- **Nondeterministic ("guessing") stage:** An arbitrary **string  $S$**  is generated that can be thought of as a candidate solution to the given **instance  $I$** .
- **Deterministic ("verification") stage:** A deterministic algorithm takes both  **$I$  and  $S$**  as its input and outputs yes if  $S$  represents a solution to instance  $I$ .

# Key Terminologies and Definitions

## Example of non-deterministic algorithm (For searching):

**Algorithm** NDSearch (int a, int n, int key)

```
{  
    j= choice(a, n);  
    if(A[j]==x) then  
    {  
        write(j);  
        success();  
    }  
    write(0);  
    failure();  
}
```

# Key Terminologies and Definitions

**3. Polynomial time algorithms:** An algorithm solves a problem in **polynomial time** if its **worst-case time efficiency** belongs to  **$O(p(n))$**  where  **$p(n)$**  is a polynomial of the problem's **input size  $n$** .

**Examples:**

- Linear/Binary Search
- All sorting algorithms
- Normal/ Strassen's Matrix Multiplication

Problems that can be solved in polynomial are called **tractable** (relatively easy)

Problems that cannot be solved in polynomial time are called **intractable** (relatively hard).

**4. Conjunctive Normal Form (CNF):** In Propositional logic, a statement is in CNF if it is a **conjunction (AND) of clauses**, where a clause is a **disjunction (OR)** of literals, and a literal is a variable or its negation.

$$(A \vee B) \wedge (!A \vee C)$$
$$A \vee B$$

# Key Terminologies and Definitions

**6. CNF Satisfiability problem:** Given a Boolean formula of  $n$  variables  $f(x_1, x_2, \dots, x_n)$ , represented in CNF, the problem is to find the values of these variables, on which the formula takes on the value *true*.

**5. Problem Reduction:** A decision problem **D1** is said to be polynomially reducible to a decision problem **D2** (represented as  $D1 \leq D2$ ) if there exists a function **T** that transforms instances of **D1** to instances of **D2** such that:

- **T** maps all yes instances of D1 to yes instances of D2 and all no instances of D1 to no instances of D2
- **T** is computable by a polynomial time algorithm

# Class P and Class NP

**Class P** is a class of decision problems that can be solved in polynomial time by (deterministic) algorithms.

This class of problems is called **deterministic polynomial time problems**.

**Examples:** Searching, Sorting, Shortest path in weighted graphs, Computing Minimum Spanning Tree, Matrix Multiplication etc.

**Class NP** is the class of decision problems that can be solved by nondeterministic polynomial algorithms.

This class of problems is called **non-deterministic polynomial time problems**.

**Examples:** Hamiltonian circuit problem, the partition problem, decision versions of the traveling salesman, the knapsack, graph coloring, and many hundreds of other difficult combinatorial optimization problems cataloged.

# Class NP-Hard and NP-Complete

**Definition:** A decision problem **X** is said to be **NP-complete** if:

1. **X** belongs to class **NP**
2. **X** is **NP-Hard**

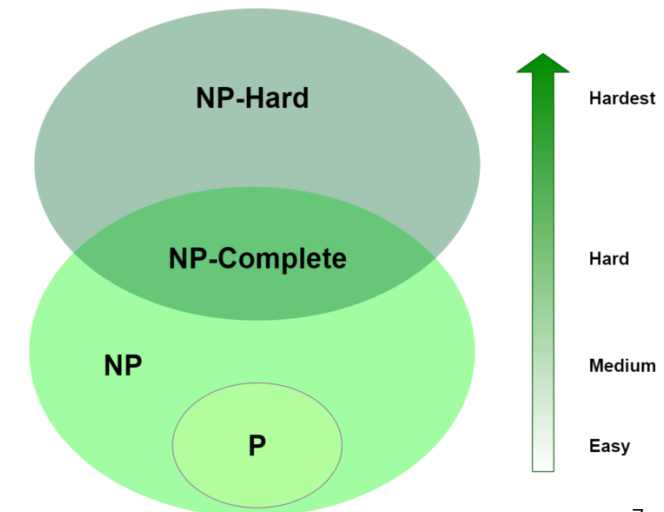
**Definition:** The decision problem **X** is **NP-Hard** if every problem **Y** belonging to **NP** reduces to **X**.

To prove the above point (2), it is enough to prove that:

- **X is in NP** and a **known problem** in **NP-Complete** can **be reduced to P**.

Examples of NP-Hard problems:

CNF satisfiability, Knapsack, Graph coloring, Travelling salesman Problem, Hamiltonian Circuit problem, Sudoku puzzle, Sub-set sum problem etc.



# Thank you!

## Any queries?