

# On Chapter 7 of *Laws of Form*

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## 1 Introduction

In chapter 7 of *Laws of Form* Spencer-Brown extends the scope of his basic equations to expressions with any finite number of variables. Some of his arguments, when he provides them, are rigorous; others are mere sketches, and some possible generalizations are left unmentioned. This paper will present fully rigorous proofs of the propositions.

Below is a list of axioms and theorems referenced in subsequent proofs:

$$\overline{pr} \overline{qr} = \overline{p} \overline{q} r \quad (\text{J2})$$

$$\overline{pr} \overline{qr} = \overline{p} \overline{q} \overline{r} \quad (\text{J2.1})$$

$$\overline{a} = a \quad (\text{C1})$$

$$\overline{ab} b = \overline{a} b \quad (\text{C2})$$

$$\overline{a} \overline{b} c = \overline{ac} \overline{b} \quad (\text{C7})$$

$$\overline{a} \overline{br} \overline{cr} = \overline{a} \overline{b} \overline{c} \overline{a} \overline{r} \quad (\text{C8})$$

$$\overline{a} \overline{r} \overline{b} \overline{r} \overline{x} \overline{r} \overline{y} \overline{r} = \overline{r} \overline{ab} \overline{rxy} \quad (\text{C9})$$

$$\overline{a} \overline{r} \overline{x} \overline{r} = \overline{r} \overline{a} \overline{rx} \quad (\text{C9.1})$$

## 2 General theorems

Spencer-Brown begins the chapter by sketching an inductive generalization of J2. Here is the proof in full.

**Theorem (J2\*).**

$$\overline{a_1} \overline{a_2} \dots \overline{a_n} r = \overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r}$$

*Proof.* The proof proceeds by induction on  $n$ . The base case is J2, where  $n = 2$ . Let the induction hypothesis (J2h) be:

$$\overline{a_1} \overline{a_2} \dots \overline{a_n} r = \overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r}$$

The induction step:

$$\begin{aligned} & \overline{a_1} \overline{a_2} \dots \overline{a_n} \overline{a_{n+1}} \parallel r \\ &= \overline{\overline{a_1} \overline{a_2} \dots \overline{a_n}} \parallel \overline{a_{n+1}} \parallel r \end{aligned} \tag{C1}$$

$$= \overline{\overline{a_1} \overline{a_2} \dots \overline{a_n}} \parallel r \parallel \overline{a_{n+1} r} \parallel \tag{J2}$$

$$= \overline{\overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r}} \parallel \overline{a_{n+1} r} \parallel \tag{J2h}$$

$$= \overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r} \parallel \overline{a_{n+1} r} \parallel \tag{C1}$$

□

*Alternate proof.* A very similar and equally short proof, using the same induction hypothesis as above. The induction step:

$$\begin{aligned} & \overline{a_1} \overline{a_2} \dots \overline{a_n} \overline{a_{n+1}} \parallel r \\ &= \overline{a_1} \overline{a_2} \dots \overline{a_n} \parallel \overline{a_{n+1}} \parallel r \end{aligned} \tag{C1}$$

$$= \overline{a_1 r} \overline{a_2 r} \dots \overline{a_n} \parallel \overline{a_{n+1}} \parallel r \parallel \tag{J2h}$$

$$= \overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r} \parallel \overline{a_{n+1} r} \parallel \tag{J2}$$

$$= \overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r} \parallel \overline{a_{n+1} r} \parallel \tag{C1}$$

□

Before continuing, I prove a useful generalization of corollary J2.1.

**Theorem (J2.1\*).**

$$\overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r} = \overline{\overline{a_1} \overline{a_2} \dots \overline{a_n}} \parallel r \parallel$$

*Proof.*

$$\begin{aligned} & \overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r} \\ &= \overline{\overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r}} \parallel \end{aligned} \tag{C1}$$

$$= \overline{\overline{a_1} \overline{a_2} \dots \overline{a_n}} \parallel r \parallel \tag{J2*}$$

□

Spencer-Brown states the generalizations of C8 and C9 but omits the proofs, merely noting that they are similar to J2\*.

**Theorem (C8\*).**

$$\overline{a} \overline{b_1 r} \overline{b_2 r} \dots \overline{b_n r} = \overline{a} \overline{b_1} \overline{b_2} \dots \overline{b_n} \parallel \overline{a} \parallel r \parallel$$

*Proof.* The proof proceeds by induction on  $n$ . The base case is C8, where  $n = 2$ . Let the induction hypothesis (C8h) be:

$$\overline{a| \overline{b_1 r} | \overline{b_2 r} | \dots \overline{b_n r} |} = \overline{a| \overline{b_1} | \overline{b_2} | \dots \overline{b_n} |} \overline{a| \overline{r} |}$$

The induction step:

$$\begin{aligned} & \overline{a| \overline{b_1 r} | \overline{b_2 r} | \dots \overline{b_n r} | \overline{b_{n+1} r} |} \\ &= \overline{a| \overline{b_1 r} |} \overline{b_2 r} | \dots \overline{b_n r} | \overline{b_{n+1} r} | \end{aligned} \quad (\text{C1})$$

$$= \overline{a| \overline{b_1 r} |} \overline{b_2} | \dots \overline{b_n} | \overline{b_{n+1}} | \overline{a| \overline{b_1 r} |} \overline{r} | \quad (\text{C8h})$$

$$= \overline{b_2} | \dots \overline{b_n} | \overline{b_{n+1}} | \overline{r} | \overline{a| \overline{b_1 r} |} \quad (\text{J2.1})$$

$$= \overline{b_2} | \dots \overline{b_n} | \overline{b_{n+1}} | \overline{r} | \overline{a| \overline{b_1 r} |} \quad (\text{C1 twice})$$

$$= \overline{b_2} | \dots \overline{b_n} | \overline{b_{n+1}} | \overline{b_1} | \overline{r} | \overline{a} | \quad (\text{J2.1})$$

$$= \overline{b_1} | \overline{b_2} | \dots \overline{b_n} | \overline{b_{n+1}} | \overline{r} | \overline{a} | \quad (\text{C1})$$

$$= \overline{b_1} | \overline{b_2} | \dots \overline{b_n} | \overline{b_{n+1}} | \overline{r} | \overline{a} | \quad (\text{C1})$$

$$= \overline{a| \overline{b_1} | \overline{b_2} | \dots \overline{b_n} | \overline{b_{n+1}} |} \overline{a| \overline{r} |} \quad (\text{J2})$$

$$= \overline{a| \overline{b_1} | \overline{b_2} | \dots \overline{b_n} | \overline{b_{n+1}} |} \overline{a| \overline{r} |} \quad (\text{C1})$$

□

J2.1\* allows for a quicker direct proof.

*Alternate proof.*

$$\begin{aligned} & \overline{a| \overline{b_1 r} | \overline{b_2 r} | \dots \overline{b_n r} |} \\ &= \overline{a| \overline{b_1} | \overline{b_2} | \dots \overline{b_n} |} \overline{r} | \end{aligned} \quad (\text{J2.1*})$$

$$= \overline{a| \overline{b_1} | \overline{b_2} | \dots \overline{b_n} |} \overline{r} | \quad (\text{C1})$$

$$= \overline{a| \overline{b_1} | \overline{b_2} | \dots \overline{b_n} |} \overline{a| \overline{r} |} \quad (\text{J2})$$

$$= \overline{a| \overline{b_1} | \overline{b_2} | \dots \overline{b_n} |} \overline{a| \overline{r} |} \quad (\text{C1})$$

□

**Theorem (C9\*).**

$$\begin{aligned} & \overline{\overline{a_1 | r | a_2 | r | \dots a_n | r | x_1 | r | x_2 | r | \dots x_m | r |}} \\ &= \overline{r | a_1 a_2 \dots a_n | r x_1 x_2 \dots x_m |} \end{aligned}$$

*Proof.*

$$\begin{aligned} & \overline{\overline{a_1 | r | a_2 | r | \dots a_n | r | x_1 | r | x_2 | r | \dots x_m | r |}} \\ &= \overline{\overline{\overline{a_1 | a_2 | \dots a_n | r | x_1 | x_2 | \dots x_m | r |}} \quad (\text{J2.1* twice}) \\ &= \overline{\overline{a_1 a_2 \dots a_n | r | x_1 x_2 \dots x_m | r |}} \quad (\text{C1 } n+m \text{ times}) \\ &= \overline{r | a_1 a_2 \dots a_n | r x_1 x_2 \dots x_m |} \quad (\text{C9.1}) \end{aligned}$$

□

Next we prove a generalization of C2.

**Theorem (C2\*).**

$$\overline{\overline{\overline{a_n b | \dots | a_2 | a_1 |} b} = \overline{\overline{\overline{a_n | \dots | a_2 | a_1 |} b}}$$

*Proof.* The proof proceeds by induction on  $n$ . The base case is C2, where  $n = 1$ . Let the induction hypothesis be:

$$\overline{\overline{\overline{a_n b | \dots | a_2 | a_1 |} b} = \overline{\overline{\overline{a_n | \dots | a_2 | a_1 |} b}}$$

Substitute  $\overline{a_{n+1} b | a_n}$  for  $a_n$ . The induction step then follows immediately:

$$\overline{\overline{\overline{a_{n+1} b | a_n b | \dots | a_2 | a_1 |} b} = \overline{\overline{\overline{a_{n+1} b | a_n | \dots | a_2 | a_1 |} b}}$$

□

Spencer-Brown does not mention a generalized C7. Here is one possible version.

**Theorem (C7\*).** *Let  $n$  be a positive even number. Then for all such  $n$  the following pair of equations holds:*

$$\begin{aligned} (i) \quad & \overline{\overline{\overline{a_n | \dots | a_2 | a_1 |} = \overline{a_n | a_{n-1} \dots a_3 a_1 |} \dots \overline{a_4 | a_3 a_1 |} \overline{a_2 | a_1 |}} \\ (ii) \quad & \overline{\overline{\overline{a_{n+1} | a_n | \dots | a_2 | a_1 |} = \overline{a_{n+1} a_{n-1} \dots a_3 a_1 |} \overline{a_n | a_{n-1} \dots a_3 a_1 |} \dots \overline{a_4 | a_3 a_1 |} \overline{a_2 | a_1 |}} \end{aligned}$$

*Proof.* Let equation (i) be the induction hypothesis. The base case is the identity  $\overline{a_2} \overline{a_1} = \overline{a_2} \overline{a_1}$ , where  $n = 2$ . Now substitute  $\overline{a_{n+1}} \overline{a_n}$  for  $a_n$ . Then,

$$\begin{aligned} & \overline{\overline{\overline{\overline{\overline{a_{n+1}} \overline{a_n}} \dots \overline{a_2} \overline{a_1}}}} \\ &= \overline{\overline{a_{n+1}} \overline{a_n}} \overline{a_{n-1} \dots a_3 a_1} \dots \overline{a_4} \overline{a_3 a_1} \overline{a_2} \overline{a_1} \end{aligned} \quad (i)$$

$$= \overline{a_{n+1} a_{n-1} \dots a_3 a_1} \overline{a_n} \overline{a_{n-1} \dots a_3 a_1} \dots \overline{a_4} \overline{a_3 a_1} \overline{a_2} \overline{a_1} \quad (C7)$$

proving the implication from (i) to (ii). In equation (ii) substitute  $\overline{a_{n+2}} \overline{a_{n+1}}$  for  $a_{n+1}$ . Then,

$$\overline{\overline{\overline{\overline{\overline{\overline{a_{n+2}} \overline{a_{n+1}}}} \dots \overline{a_2} \overline{a_1}}}} = \overline{a_{n+2}} \overline{a_{n+1} \dots a_3 a_1} \dots \overline{a_4} \overline{a_3 a_1} \overline{a_2} \overline{a_1} \quad (ii)$$

proving (i) for the succeeding even number. This proves the proposition for all  $n \geq 2$ , and hence for all echelons of depth greater than or equal to 2.  $\square$

**Theorem (T14).** *Any expression can be reduced to an equivalent expression not more than two crosses deep. Specifically, any expression  $E$  is equivalent to  $\overline{a_1} \overline{b_1} \overline{a_2} \overline{b_2} \dots \overline{a_n} \overline{b_n} \overline{c_1} \overline{c_2} \dots \overline{c_m} d$  where  $a_i, b_i, c_i, d$  are composed (at most) of juxtapositions of variables and the two constants,  $\overline{\quad}$  and  $\quad$ .*

*Proof.* Repeated applications of **C7\*** to any expression demonstrates the theorem. Spencer-Brown uses **C7** (not having proven a generalization), but it comes to the same thing.  $\square$

The final theorem follows Spencer-Brown closely.

**Theorem (T15).** *Given any expression  $E$  and any variable  $v$ ,  $E$  can be reduced to an equivalent expression containing not more than two appearances of  $v$ .*

*Proof.* In the case where  $v$  is not in  $E$ , the theorem is trivially true, since  $E = \overline{v} \overline{v} E$  by **J1**. So let us suppose that  $v$  appears in  $E$ . Using **C7\*** as many times as necessary, we rewrite  $E$ :

$$E = \overline{va_1} \overline{b_1} \overline{va_2} \overline{b_2} \dots \overline{va_n} \overline{b_n} \overline{vc_1} \overline{vc_2} \dots \overline{vc_m} d$$

where  $a_i, b_i, c_i$ , and  $d$  are expressions free of  $v$ . Then, by  $n$  applications of **C8.1**,

$$\begin{aligned} E &= \overline{v} \overline{b_1} \overline{a_1} \overline{b_1} \overline{v} \overline{b_2} \overline{a_2} \overline{b_2} \dots \overline{v} \overline{b_n} \overline{a_n} \overline{b_n} \overline{vc_1} \overline{vc_2} \dots \overline{vc_m} d \\ &= \overline{v} \overline{b_1} \overline{v} \overline{b_2} \dots \overline{v} \overline{b_n} \overline{vc_1} \overline{vc_2} \dots \overline{vc_m} f \\ &\quad (\text{where } f = \overline{a_1} \overline{b_1} \overline{a_2} \overline{b_2} \dots \overline{a_n} \overline{b_n} d \text{ is free of } v.) \\ &= \overline{\overline{b_1} \overline{b_2} \dots \overline{b_n} \overline{v}} \overline{\overline{c_1} \overline{c_2} \dots \overline{c_m} \overline{v}} f \end{aligned} \quad (J2.1^* \text{ twice})$$

$\square$