

Journal Club

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Part I

Real Analysis

TEXTBOOK: [Analysis 1](#) by Terence Tao

Chapter 1

Natural Numbers

Numbers were built to count. A system for counting was made, and that system is the number system.

Definition 1.0.1

A natural number is an element of the set \mathbb{N} of the set

$$\mathbb{N} = \{0, 1, 2, 3 \dots\}$$

is obtained from 0 and counting forward indefinitely.

1.1 Peano Axioms

We start with axioms to help clarify this.

- Axiom 1 : $0 \in \mathbb{N}$
- Axiom 2: If $n \in \mathbb{N}$, then $n++ \in \mathbb{N}$
- Axiom 3: 0 is not an increment of any other natural number $n \in \mathbb{N}$
- Axiom 4: If $n \neq m$, $n++ \neq m++$
- Axiom 5: (Principle Of Mathematical Induction) Let $P(n)$ be any property pertaining to a natural number n . Suppose that $P(0)$ is true, and suppose that whenever $P(n)$ is true, $P(n++)$ is also true. Then $P(n)$ is true for every natural number.

We then make an assumption: That the set \mathbb{N} which satisfies these five axioms is called the set of natural numbers. With these 5 axioms, we can construct sequences

1.2 Recursive Definitions

Proposition 1.2.1 (Recursive Definitions). Suppose for each natural number n , we have some function $f_n : \mathbb{N} \rightarrow \mathbb{N}$ from the natural numbers to the natural numbers. Then we can assign a unique natural number a_n to each natural number n , such that $a_0 = c$ and $a_{n++} = f_n(a_n)$ for each natural number n .

1.3 Addition

Definition 1.3.1: Addition Of Natural Numbers

Let n be a natural number. ($n \in \mathbb{N}$). To add zero to m , we define $0 + m := m$. Now suppose inductively that we have defined how to add n to m . Then we can add $n++$ to m by defining $(n++) + m := (n+m)++$

Lemma 1.3.1. For any natural number $n + 0 = n$

Proof. We use induction,
The base case, $n = 0$,

$$\begin{aligned} n = 0, 0 + 0 &= 0 \\ n + 0 &= n \\ (n++) + 0 &= (n + 0)++ = (n++) \end{aligned}$$

Suppose inductively, that $n + 0 = n$,
For $n = n++$,

$$\begin{aligned} (n++) + 0 &= (n + 0)++ \\ \text{We know that } n + 0 &= n \\ (n++) + 0 &= (n++) \end{aligned}$$

□

Lemma 1.3.2. For any natural numbers n and m ,

$$n + (m++) = (n + m)++$$

Proof. Inducting on n while keeping m fixed,

$$\begin{aligned} n = 0, \\ 0 + (m++) &= (0 + m)++ \\ 0 + (m++) &= (m++) \end{aligned}$$

This we know is true from the definition of addition ($0 + m := m$)

Suppose inductively, that $n + (m++) = (n + m)++$ is true. For $n = (n++)$,

$$\begin{aligned} (n++) + (m++) &= ((n++) + m)++ && \text{From the definition of addition} \\ &= (n + (m++))++ \\ &= ((n + m)++)++ \end{aligned}$$

□

Putting $m = 0$, we get $n + 1 = n++$

Proposition 1.3.1 (Addition is commutative). For any natural numbers n and m , $n + m = m + n$

Proof. We induct over n , For the base case, $n = 0$,

We must show that $m + 0 = 0 + m$ From the definition of addition, we have

$$0 + m = m$$

As shown earlier, we have

$$m + 0 = m$$

This is clearly true for $n = 0$.

Now suppose inductively that $m + n = n + m$

For $n = n + +$, we must show that $m + (n + +) = (n + +) + m$

We know from the definition of addition that,

$$(n + +) + m := (m + n) + +$$

And we proved earlier that,

$$m + (n + +) = (m + n) + +$$

Therefore,

$$m + (n + +) = (n + +) + m$$

□

Proposition 1.3.2 (Addition is associative). For any natural numbers, a, b and c , we have $(a + b) + c = a + (b + c)$

Proof. We take $(a + b) + n = a + (b + n)$

Inducting over n ,

For $n = 0$,

We have in the LHS,

$$\begin{aligned} &= (a + b) + 0 \\ &= a + b \end{aligned}$$

$$\text{Since } n + 0 = n$$

On the RHS,

$$\begin{aligned} &= a + (b + 0) \\ &= a + b \end{aligned}$$

$$\text{Since } n + 0 = n$$

Suppose inductively that $(a + b) + n = a + (b + n)$,

For $n = n + +$, We have to show that $(a + b) + (n + +) = a + (b + (n + +))$

On the LHS we have,

$$\begin{aligned} &= (a + b) + (n + +) \\ &= (a + b + n) + + \end{aligned}$$

$$\text{(From the lemma } m + (n + +) = (m + n) + +)$$

On the RHS we have,

$$\begin{aligned}
 &= a + (b + (n + +)) \\
 &= a + (b + n) + + && \text{(From the lemma } m + (n + +) = (m + n) + + \text{)} \\
 &= (a + b + n) + +
 \end{aligned}$$

LHS = RHS

□

Proposition 1.3.3 (Cancellation Law). Let a, b, c be natural numbers such that $a + b = a + c$. Then we have $b = c$.

Proof. We have,

$$n + b = n + c$$

Inducting over n , For the base case, $n = 0$

$$\begin{aligned}
 0 + b &= 0 + c \\
 b &= c
 \end{aligned}$$

Suppose inductively that $n + b = n + c$ For $n = n + +$,

$$(n + +) + b = (n + +) + c$$

On the LHS

$$\begin{aligned}
 &= (n + +) + b \\
 &= (n + b) + +
 \end{aligned}$$

On the RHS

$$\begin{aligned}
 &= (n + +) + c \\
 &= (n + c) + +
 \end{aligned}$$

We know from the inductive hypothesis that,

$$\text{If } n + b = n + c, \text{ then } b = c$$

Thus we have,

$$b + + = c + +$$

□

Definition 1.3.2: Positive natural number

All numbers where,

$$n \neq 0, n \in \mathbb{N}$$

Proposition 1.3.4. If a is a positive natural number and b is a natural number, then $a + b$ is positive.

Proof. Inducting over b ,

For $b = 0$,

$$a + 0 = a$$

This proves the base case, since we know a is positive.

Now, suppose inductively, that $(a + b)$ is positive.

For $(a + (n + +))$,

$$a + (n + +) = (a + n) + +$$

We know from Axiom 3 that $n + + \neq 0$. Thus we close the inductive loop. \square

Lemma 1.3.3. For every a , there exists a unique b such that $b + + = a$

Proof. Proof by contradiction, Suppose that there are two different increments, $m + +$, $n + +$ that equal to a ,

We have,

$$m + + = a$$

$$n + + = a$$

Then we can say,

$$m + + = n + +$$

$$m + 1 = n + 1$$

$$m = n$$

(By Cancellation Law)

But we said that m and n are different numbers which increment to a .

Therefore, we can conclude that there is only one number b which increments to a \square

1.4 Order

Definition 1.4.1: Order

Let n and m be natural numbers we say that n is greater than or equal to m , and write $n \geq m$ iff we have $n = m + a$ for some natural number a . We say that $n > m$ when $n \geq m$ and $n \neq m$

Proposition 1.4.1 (Basic properties of order for natural numbers). Let a, b, c be natural numbers then

1. (Order is reflexive) $a \geq a$
2. (Order is transitive) If $a \geq b$ and $b \geq c$, then $a \geq c$
3. (Order is antisymmetric) If $a \geq b$ and $b \geq a$ then $a = b$
4. (Addition preserves order) $a \geq b$ if and only if $a + c \geq b + c$
5. $a < b$ if and only if $a + + \leq b$

6. $a < b$ if and only if $b = a + d$ for some positive number d .

Proof. 1. Proving order is reflexive, $a \geq a$

We know that,

$$a = a + 0$$

From the definition of order, We can write that $a \geq b$ when $a = b + d$ where $d \in \mathbb{N}$

Thus $a \geq a$.

2. Proving order is transitive, $a \geq b$ and $b \geq c$ then $a \geq c$

We write,

$$a = b + d$$

$$b = c + e$$

$$a = c + e + d$$

We can say that since $(e + d) \in \mathbb{N}$

We define $f := (e + d)$ Where $f \in \mathbb{N}$

$$a = c + (f)$$

Thus we can say,

$$\text{If } a \geq b, b \geq c \text{ then } a \geq c$$

3. Proving order is antisymmetric, If $a \geq b$ and $b \geq a$ then $a = b$ We can say,

$$a = b + d$$

$$b = a + e$$

Where $d, e \in \mathbb{N}$

$$a = (a + e) + d$$

$$b = (b + d) + e$$

Then we can write,

$$a = a + (e + d)$$

$$b = b + (d + e)$$

Then we can say that $(e + d)$ and $(d + e)$ are 0.

We know that if $a + b = 0$ then $a, b = 0$

Thus d and e are 0.

$$\begin{aligned} a &= b + d \\ a &= b \end{aligned}$$

4. Proving $a < b$ if and only if $b = a + d$ for some positive number d If $b = a + d$ where d is a positive natural number, $d \neq 0$

Which means that $b \neq a + 0$ or $b \neq a$

This means that b is strictly greater than a

If $a < b$ then $a \geq b$ and $a \neq b$

So if $a \geq b$ Then,

$$a = b + d$$

But,

$$\begin{aligned} a &\neq b \\ a &\neq b + 0 \\ d &\neq 0 \end{aligned}$$

Thus d cannot be 0. d can only be a positive natural number.

5. Proving addition preserves order, $a \geq b$ if and only if $a + c \geq b + c$ Proving $a \geq b$ if $a + c \geq b + c$

Where $d \in \mathbb{N}$

$$\begin{aligned} a + c &= b + c + d && \text{By definition} \\ a + c &= (b + d) + c \\ a &= (b + d) && \text{By cancellation law} \\ a &\geq b \end{aligned}$$

Proving $a + c \geq b + c$ if $a \geq b$

We know,

$$a = b + d$$

Where $d \in \mathbb{N}$

We write $a + c$ using what we know from above,

$$\begin{aligned} a + c &= b + d + c \\ a + c &= b + c + d \\ (a + c) &= (b + c) + d \\ a + c &\geq b + c \end{aligned}$$

6. Proving $a < b$ if and only if $a++ \leq b$ Proving $a < b$ if $a++ \leq b$

We can write,

$$\begin{aligned} a++ &= b+d & \text{Where } d \in \mathbb{N} \\ a+++d &= b \\ a+(d++) &= b \end{aligned}$$

Since from Axiom 3, we know that 0 is not an increment of any natural number, $(d++ \neq 0)$
Therefore,

$$a < b$$

□

Proposition 1.4.2 (Trichotomy of order for natural numbers). Let a and b be natural numbers. Then exactly one of the following statements is true: $a < b, a = b$ or $a > b$

Proof. First we show that no more than one of the statements is true. If $a < b$ then $a \neq b$ by definition. If $a > b$ then $a \neq b$ by definition. If $a > b$ and $a < b$ then $a = b$, which we proved earlier.

Now to show that exactly one of these statements are true. We induct on a ,

When $a = 0$, We know that,

$$\begin{aligned} b &= 0 + b (\forall b \in \mathbb{N}) \\ b &\geq 0 \end{aligned}$$

Suppose inductively that exactly one of the above statements are true for a and b . For $a++$, We take each statement. First for $a > b$

$$\begin{aligned} a &> b \\ a &= b + d \\ (a++) &= (b + d)++ \\ (a++) &= b + d++ \\ (a++) &> b & \text{If } d \in \mathbb{N} \text{ then } d++ \in \mathbb{N} \end{aligned}$$

For $a = b$

$$\begin{aligned} a &= b \\ (a++) &= (b)++ \\ (a++) &= b + 1 \\ a &> b \end{aligned}$$

For $a < b$

$$\begin{aligned}a &< b \\a + d &= b \\(a + d) + + &= b + + \\(a + +) + d &= b + + \\(a + +) + d &= b + 1\end{aligned}$$

We have two cases, If $d = 1$, Then by cancellation law

$$a + + = b$$

If $d \neq 1$ Then

$$a + + < b$$

But never both, which concludes the inductive loop. \square

1.5 Special Forms Of Induction

Theorem 1.5.1 (Strong Principle Of Induction). Let m_0 be a natural number, and let $P(m)$ be a property pertaining to an arbitrary natural number m . Suppose that for each $m \geq m_0$, we have the following implication: if $P(m')$ is true for all natural numbers $m_0 \leq m' < m$, then $P(m)$ is also true. (In particular this means that $P(m_0)$ is true, since in this case the hypothesis is vacuous.) Then we can conclude that $P(m)$ is true for all natural numbers $m \geq m_0$.

Proof. For a property $Q(n)$, which is the property that $P(m')$ is true for $m_0 \leq m' < n$, then $P(n)$ is true... Then it is true $\forall m \geq m_0$

For $Q(0)$, we can say that the statement is vacuous since the conditions are not satisfied for both when $m_0 = 0$ and when $m_0 < 0$

Suppose inductively that $Q(n)$ is true. Then \square

1. Backward Induction

2. Induction Starting From The Base Case n Let n be a natural number, and let $P(m)$ be a property pertaining to the natural numbers such that whenever $P(m)$ is true, $P(m++)$ is true. Show that if $P(n)$ is true, then $P(m)$ is true for all $m \geq n$. (This principle is sometimes referred to as the principle of induction starting from the base case n .)

Proof. Take a property $P(n)$, $m \geq n$

Inducting over n , \square

1.6 Multiplication

Definition 1.6.1: Multiplication

Let m be a natural number. To multiply zero to m , we define $0 \times m := 0$. Now suppose inductively that we have defined how to multiply n to m . Then we can multiply $n++$ to m by defining $(n++) \times m := (n \times m) + m$

We can say $0 \times m = 0$, $1 \times m = 0 + m$, $2 \times m = 0 + m + m$ and so on.

Lemma 1.6.1. Prove that multiplication is commutative

For simplicity we now write $a \times b$ as ab

1.7 Exercise

1. Prove the identity $(a + b)^2 = a^2 + 2ab + b^2$
2. (Euclid's division lemma)
3. Backward Induction $m \in \mathbb{N}, P(m), P(m++) \Rightarrow P(m)$, Suppose $P(n)$ is true, then $P(m) \forall m \leq n$
For the base case, $n = 0$, $P(0) \Rightarrow P(0)$, *soit' true*.

For the inductive step, supposing $Q(n)$ is true,

4. Strong induction
5. Distributive Law
6. Multiplication
 - (a) Cancellation Law
 - (b) Associativity
 - (c) If $a < b$, and c is positive then $ac < bc$