

Ans 3. Given: $T(n) = n^2 + 3 \cdot T(n-1)$
 $T(0) = 1$

Proof:- $T(n) = n^2 + 3 \cdot T(n-1)$

$$T(n) = n^2 + 3 \cdot [(n-1)^2 + 3 \cdot T(n-2)]$$

$$T(n) = n^2 + 3^1 \cdot (n-2)^2 + 3^2 \cdot [(n-2)^2 + 3 \cdot T(n-3)]$$

⋮

$$T(n) = \sum_{i=0}^n 3^i (n-i)^2 + 3^n \cdot T(0)$$

$$T(n) = 3^n + \sum_{i=0}^n 3^i (n-i)^2$$

$$= 3^n + \sum_{i=0}^n 3^i [n^2 + i^2 - 2ni]$$

$$3^n + n^2 \left[\frac{3^{n+1} - 1}{2} \right] - 2n \left\{ \sum_{i=0}^n 3^i i + \sum_{i=0}^n i^2 3^i \right\}$$

~~Now obtaining the upper bound by considering~~

~~$S(n) = T(n) / 3^n$ which satisfies recurrence relation~~
 ~~$S(n) = S(n-1) + \frac{n^2}{3^n}$~~

~~and it implies $T(n) = O(3^n)$ as~~

$$T(n) = 3^n + \sum_{i=0}^n 3^i (n-i)^2$$

$$\leq 3^n + \sum_{i=0}^n 3^i n^2$$

$$T(n) \leq 3^n + n^2 \sum_{i=0}^n 3^i$$

$$T(n) \leq 3^n + n^2 \cdot \frac{(3n-1)}{2}$$

$$\Rightarrow \boxed{T(n) = O(3^n)} \quad \text{a}$$

Ans 2. (a). Let $n = 2^k$, $\sqrt{n} = 2^{k/2}$

$$k = \log n.$$

Substituting in above eqⁿ:

$$T(n) = \sqrt{n} \cdot T(\sqrt{n}) + 5$$

$$\Rightarrow T(2^k) = 2^{k/2} T(2^{k/2}) + 5 \quad \text{--- (1)}$$

~~Q. 2. (a) Let $n = 2^k$, $\sqrt{n} = 2^{k/2}$~~

~~Substituting in above eqⁿ:~~

~~$T(n) = \sqrt{n} \cdot T(\sqrt{n}) + 5$~~

(Master's Theorem: $T(n) = a \cdot T(\frac{n}{b}) + n^d$). --- (3)

dividing eqⁿ (1) by 2^k .

$$\frac{T(2^k)}{2^k} = \frac{T(2^{k/2})}{2^{k/2}} + \frac{5}{2^{k/2}} \quad \text{--- (2)}$$

~~Let $\frac{T(2^k)}{2^k} = y(k)$~~

$$\text{Let } \frac{T(2^k)}{2^k} = y(k)$$

$$\therefore y(k) = y(k/2) + \frac{5}{2^{k/2}}$$

Ans 2. (b). $T(n) = 3 T\left(\frac{n}{4}\right) + T\left(\frac{n}{6}\right) + \underline{n}$.

$T(1) = T(2) = \underline{5}$.

Using recurrence relⁿ we can see that the depth is $\log_6 n$.

$$\log_6 n = O(\log n).$$

We can also see that total work at each level is $11/2$ of upper level.

\therefore The recursion method tree implies the upper bound

$$T(n) \leq n \sum_{i=0}^{\log_6 n} \left(\frac{11}{2}\right)^i$$

$$T(n) \leq 12n$$

$$\boxed{T(n) = O(n)} \quad \text{as } \underline{\underline{12}}$$

Ans 3.

By Induction:-

Base case: $T(0) = O(1) \Rightarrow 1 \leq c \cdot 1$ (True)

Consider: $T(n) \leq c \cdot 3^n$. (Inductive step)

To prove: $T(n+1) \leq c \cdot 3^{n+1}$

Proof:- $T(n) \leq c \cdot 3^n$

$$3 \cdot T(n) \leq c \cdot 3^{n+1}$$

$$T(n+1) - (n+1)^2 \leq c \cdot 3^{n+1}$$

$$\Rightarrow T(n+1) \leq c \cdot 3^{n+1} + (n+1)^2$$

$\because (n+1)^2 > 0$, so it is clearly shown that

$$c \cdot 3^{n+1} + (n+1)^2 \geq c' \cdot \underline{3^{n+1}} \quad \text{for } c, c'$$

$$\therefore T(n+1) \leq c' \cdot 3^{n+1} \quad \therefore \exists c' > 0.$$

$$\text{such that } T(n+1) \leq \underline{c' \cdot 3^{n+1}}$$

and hence proved that,

~~of~~ $T(n)$ is of order 3^n

i.e. $T(n) = \underline{O(3^n)}$ as