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Topic \rightarrow

DAA Assignment - 2

Sec/Roll no \rightarrow M/38

Q1) \rightarrow bool linearsearch (int arr, int n, int key)

```
{
    for (int i = 0; i < n; i++)
    {
        if (arr[i] == key)
        {
            return true;
        }
    }
    return false;
}
```

Q2 \rightarrow Iterative insertion sort

```
for (int i = 1; i < n; i++)
{
    int t = a[i];
    int j = i - 1;
    while (j > 0 && a[j] > t)
    {
        a[j + 1] = a[j];
        j--;
    }
    a[j + 1] = t;
}
```

Recursive insertion sort

```
void sort(arr, n)
```

```
{
```

```
    if ( $n \leq 1$ )
```

```
        return;
```

```
    sort(arr,  $n-1$ );
```

```
    int last = arr[n-1];
```

```
    int j =  $n-2$ ;
```

```
    while ( $j \geq 0$  &&  $arr[j] > last$ )
```

```
    {
```

```
        arr[j+1] = arr[j];
```

```
        j = j-1;
```

```
    }
```

```
    arr[j+1] = last;
```

```
}
```

* It can sort elements while receiving new ones that's why it is called online sorting

* other sorting techniques like merge, quick, selection can't do this

| Q3 → | Sorting technique | Best | Complexity avg | Worst |
|------|-------------------|---------------|----------------|---------------|
| | Bubble | $O(n)$ | $O(n^2)$ | $O(n^2)$ |
| | Selection | $O(n^2)$ | $O(n^2)$ | $O(n^2)$ |
| | Insertion | $O(n)$ | $O(n^2)$ | $O(n^2)$ |
| | Count | $O(n+k)$ | $O(n+k)$ | $O(n+k)$ |
| | Quick | $O(n \log n)$ | $O(n \log n)$ | $O(n^2)$ |
| | Merge | $O(n \log n)$ | $O(n \log n)$ | $O(n \log n)$ |
| | Heap | $O(n \log n)$ | $O(n \log n)$ | $O(n \log n)$ |
| | Radix | $O(d*(n+k))$ | $O(d*(n+k))$ | $O(d*(n+k))$ |

| Q4 → | Inplace | Stable | Outline sorting |
|-----------|---------|--------|-----------------|
| Bubble | ✓ | | |
| Selection | ✓ | | |
| Insertion | ✓ | ✓ | ✓ ★ |
| Count | | ✓ | |
| Quick | ✓ | | |
| Merge | | ✓ | |
| Heap | ✓ | | |
| Radix | | ✓ | |

Q5 → Recursive binary search

$T(n) \rightarrow$ int Search (arr, target, low, high)

{
if (low > high)
return -1;

mid = (low + high) / 2;

if (arr[mid] == target)

{ return mid; }

else if (arr[mid] < target)

{ return Search (arr, target, mid + 1, high); }

$T(n/2)$

else
 $T(n/2) \leftarrow$ { return Search (arr, target, low, mid - 1); }

Iterative binary search

int Search (arr, target)

{

low = 0;

high = length(arr) - 1;

while (low <= high)

{

mid = (low + high) / 2;

if (arr[mid] == target)

return mid;

else if (arr[mid] < target)

low = mid + 1;

else

high = mid - 1;

} return -1;

| | Recursive | | Iterative | |
|--------|-------------|-------------|-------------|--------|
| Linear | $O(n)$ | $O(n)$ | $O(n)$ | $O(1)$ |
| Binary | $O(\log n)$ | $O(\log n)$ | $O(\log n)$ | $O(1)$ |

Q6 → Refer Q5

$$T(n) = T(n/2) + C$$

Q7 → int findPairWithSumK(arr, K)

```

{
    sort(arr);
    int left = 0;
    int right = length(arr) - 1;
    while (left < right)
    {
        if (arr[left] + arr[right] == K)
        {
            return 1("found");
        }
        else if (arr[left] + arr[right] < K)
        {
            left = left + 1;
        }
        else
            right = right - 1;
    }
    return -1;
}
    ↪ (not found)

```

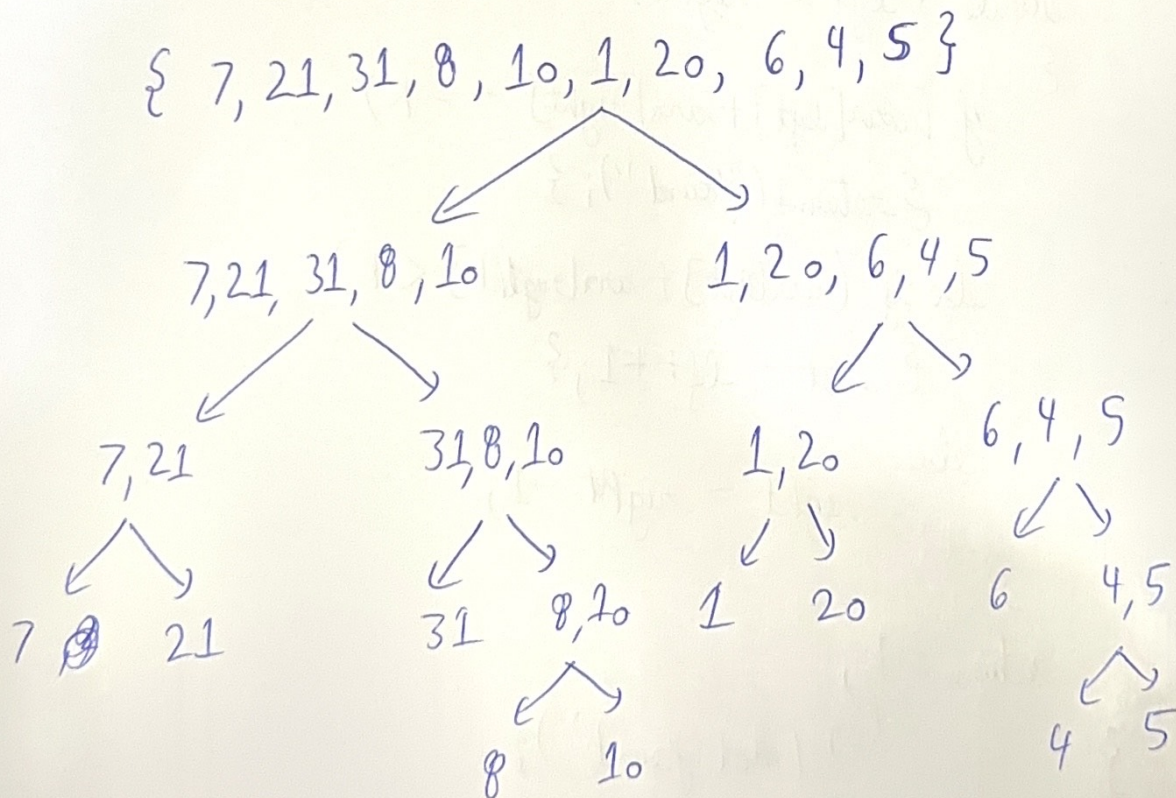

Q8 → Quick sort → Fastest sorting algorithm especially for large dataset

Merge sort → TC remains $O(n \log n)$ in all cases
Divide conquer nature

Heap sort → TC remains $O(n \log n)$
doesn't require extra space

Insertion sort → Inplace, stable, online sorting

Q9 → Inversion → When smaller element is after larger element or vice versa



Q10 → Best case → When first divides array in equal halves

Worst case → When first divides array in very unbalanced way

Q11 → Merge Sort

★ → Best case : $T(n) = 2T\left(\frac{n}{2}\right) + O(n)$

★ → Worst case : $T(n) = 2T\left(\frac{n}{2}\right) + O(n)$

Quick Sort

★ → Best case : $T(n) = T\left(\frac{n}{2}\right) + T\left(\frac{n}{2}\right) + O(n)$

★ → Worst case : $T(n) = T(n-1) + O(n)$

Similarities

★ → Both have divide & conquer

★ → $TC = O(n \log n)$

Difference

* → In quick sort TC varies

Q12 → void SelectionSort (arr, n)

```
{  
  for (i=0; i<n; i++)
```

```
{  
  int minIndex = i;
```

```
  for (j=i+1; j<n; j++)
```

```
  {  
    if (arr[j] < arr[minIndex])
```

```
    {  
      minIndex = j;
```

```
    }
```

```
  }  
  minVal = arr[minIndex];
```

```
  while (minIndex > i)
```

```
  {  
    arr[minIndex] = arr[minIndex - 1];
```

```
    minIndex = minIndex - 1;
```

```
  }
```

```
  arr[i] = minVal;
```

```
}
```

```
}
```


Q13 → void sort (arr, n)

```
{
    bool swapped;
    for (int i = 0; i < n-1; ++i)
    {
        swapped = false;
        for (int j = 0; j < n-1-i; ++j)
        {
            if (arr[j] > arr[j+1])
            {
                swap(arr[j], arr[j+1]);
                swapped = true;
            }
        }
        if (!swapped)
            break;
    }
}
```

Q14 → Merge Sort → Optimised for external sorting

* → Divide and conquer approach

* → Requires small portion of memory data to fit in memory



External sorting

Internal sorting

* → On virtual memory

* → ON RAM

* → ex → Quick and Merge
sort

* → Bubble, selection and
insertion sort