

# **Math 131AH: Homework #2**

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## Problem 1

*Proof.* We need to use induction to prove that  $\forall n \geq 1 : \sum_{k=1}^n k^3 = \left( \sum_{k=1}^n k \right)^2$

Let's call this predicate  $P(n)$ . First we show  $P(1)$ . For  $n = 1$ ,  $\text{LHS} = \text{RHS} = 1$ .

Assume  $P(n)$ , check  $P(n+1)$ .  $\text{LHS} = \sum_{k=1}^{n+1} k^3 = \sum_{k=1}^n k^3 + (n+1)^3$ . By  $P(n)$ , this equals  $\left( \sum_{k=1}^n k \right)^2 + (n+1)^3$ .

Now, we use the formula for the sum of the first  $n$  integers:  $\sum_{k=1}^n k = \frac{n(n+1)}{2}$ . This simplifies the LHS to  $\frac{n^2(n+1)^2}{4} + (n+1)^3 = \frac{(n+1)^2(n^2+4n+4)}{4} = \left( \frac{(n+1)(n+2)}{2} \right)^2 = \left( \sum_{k=1}^{n+1} k \right)^2 = \text{RHS}$ . Then, by the principle of mathematical induction  $P(n)$  is **TRUE**  $\forall n \geq 1$   $\square$

## Problem 2

Using Problem 1, we can express  $\sum_{k=1}^n k^3$  as a polynomial of degree 4. The hope here is to express  $\sum_{k=1}^n k^4$  as a polynomial of degree 5. Before solving this question, let's see if we can find an alternate way to directly prove Problem 1. Then, we can generalize this method.

Consider the polynomial  $(x+1)^4 = x^4 + 4x^3 + 6x^2 + 4x + 1$

$$\begin{aligned} \Rightarrow (x+1)^4 - x^4 &= 4x^3 + 6x^2 + 4x + 1 \\ \Rightarrow \sum_{x=1}^n (x+1)^4 - x^4 &= \sum_{x=1}^n (4x^3 + 6x^2 + 4x + 1) \\ \Rightarrow (n+1)^4 - 1 &= \sum_{x=1}^n (4x^3) + \sum_{x=1}^n (6x^2) + \sum_{x=1}^n (4x) + \sum_{x=1}^n (1) \end{aligned}$$

Here the last implication follows from the telescoping sum on the RHS and distribution of the sum on the RHS. Recall, we have closed formulae for the sum of consecutive squares and consecutive integers so we can rearrange the sum to get an expression for  $\sum_{x=1}^n x^3$  as required. We will use the same method to deduce a polynomial expression for  $\sum_{k=1}^n k^4$ .

Now consider the polynomial  $(x+1)^5 = x^5 + 5x^4 + 10x^3 + 10x^2 + 5x + 1$

$$\begin{aligned} \Rightarrow (x+1)^5 - x^5 &= 5x^4 + 10x^3 + 10x^2 + 5x + 1 \\ \Rightarrow \sum_{x=1}^n (x+1)^5 - x^5 &= \sum_{x=1}^n (5x^4 + 10x^3 + 10x^2 + 5x + 1) \\ \Rightarrow (n+1)^5 - 1 &= \sum_{x=1}^n (5x^4) + \sum_{x=1}^n (10x^3) + \sum_{x=1}^n (10x^2) + \sum_{x=1}^n (5x) + \sum_{x=1}^n (1) \\ \Rightarrow (n+1)^5 - 1 &= 5 \sum_{x=1}^n x^4 + 10 \left( \frac{n^2(n+1)^2}{4} \right) + 10 \left( \frac{n(n+1)(2n+1)}{6} \right) + 5 \left( \frac{n(n+1)}{2} \right) + n \\ \Rightarrow \sum_{x=1}^n x^4 &= \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30} \end{aligned}$$

Again, we can prove this using induction.

*Proof.* Let  $P(n)$  denote the predicate we wish to prove.  $P(1) : \text{LHS} = 1, \text{RHS} = \frac{1 \times 2 \times 3 \times 5}{30} = 1$ . Assume  $P(n)$ , check  $P(n+1)$ . As before, the LHS reduces to  $\frac{n(n+1)(2n+1)(3n^2+3n-1)}{30} + (n+1)^4$   
 $= (n+1) \left( \frac{n(2n+1)(3n^2+3n-1) + 30(n+1)^3}{30} \right) = \frac{(n+1)(n+2)(2n+3)(3n^2+9n+5)}{30} = \text{RHS} \quad \square$

Then, by the principle of mathematical induction  $P(n)$  is TRUE  $\forall n \geq 1$

### Problem 3

We are given  $P(n) : \sum_{k=1}^n k = \frac{1}{2}(n+\frac{1}{2})^2$ . Let's assume  $P(n)$  and try to prove  $P(n+1)$ .  $\text{LHS} = \sum_{k=1}^n k + (n+1) = \frac{1}{2}(n+\frac{1}{2})^2 + (n+1) = \frac{1}{2}(n+\frac{3}{2})^2 = \text{RHS}$ . Therefore,  $P(n) \Rightarrow P(n+1)$ . However, this does not imply that  $P(n)$  is TRUE for all  $n \in \mathbb{N}$  as the base case  $P_0$  is clearly false. In fact, one can see that  $P(n)$  is FALSE for all  $n \in \mathbb{N}$  as the LHS is always an integer while the RHS always has a fractional component.

### Problem 4

We define  $\mathbb{J} := \left\{ n + \frac{m}{m+1} : n, m \in \mathbb{N} \right\}$  and  $S : \mathbb{J} \rightarrow \mathbb{J}$  as  $S\left(n + \frac{m}{m+1}\right) = n + \frac{m+1}{m+2}$ .  
 Let us first verify that  $(\mathbb{J}, 0, S)$  satisfy (P1)-(P4).

1. Let  $n = 0, m = 0$ . Therefore,  $0 \in \mathbb{J}$ .
2.  $S$  is clearly defined over all of  $\mathbb{J}$ .
3.  $n + \frac{m+1}{m+2} \geq \frac{m+1}{m+2} \geq \frac{1}{2}$  so  $0 \notin \text{Ran}(S)$ .
4. Assume  $a + \frac{b+1}{b+2} = c + \frac{d+1}{d+2}$ . Denote LHS as  $\alpha$  and RHS as  $\beta$ . We have  $a < \alpha < a+1, c < \beta < c+1$  and  $\alpha = \beta$  with  $a, c \in \mathbb{N}$ . Every rational has a unique representation as a natural number and a proper fraction. Therefore,  $a = c$  and  $a+1 = c+1$ . As a result  $a + \frac{b}{b+1} = c + \frac{d}{d+1}$  so  $S$  is injective.

We now show that (P5) fails by giving a subset  $A \subset \mathbb{J} : 0 \in A \wedge S(A) \in A \wedge (A \neq \mathbb{J})$ . Let  $A = \text{Ran}(S) \cup \{0\}$ . Then, by construction,  $0 \in A$ . Since  $S(\mathbb{J}) \subset A$ , it is clear that  $S(A) \subset A$ . However,  $1 \in \mathbb{J}$  but  $1 \notin A$  so  $A \neq \mathbb{J}$ .

Now, consider  $\mathbb{K} \subset \mathbb{J}, \mathbb{K} = \left\{ \frac{m}{m+1} : m \in \mathbb{N} \right\}$ . In other words,  $\mathbb{K}$  is the subset obtained when  $n$  is fixed at 0.  $S$  is defined like above  $S\left(\frac{m}{m+1}\right) = \frac{m+1}{m+2}$ . Again, (P1) holds as  $0 \in \mathbb{K}$  by setting  $m = 0$ . (P2) to (P4) hold as they hold in  $\mathbb{J}$  and the properties are preserved under taking subsets. Finally, (P5) holds here as it holds for  $\mathbb{N}$ .

### Problem 5

Recall that we have defined  $m+n$  for  $m, n \in \mathbb{N}$  recursively as  $m+0 = m$  and  $\forall n \in \mathbb{N} : m+S(n) = S(m+n)$ . We are asked to show that  $\forall k, m, n \in \mathbb{N} : n + (m+k) = (n+m) + k$ . We will prove this via induction on  $k$ . Fix  $n, m \in \mathbb{N}$ . Then  $P(0) : n + (m+0) = n+m = (n+m) + 0$  by the definition of addition with 0. Assume  $P(k)$ . We need to show  $P(S(k)) : n + (m+S(k)) = (n+m) + S(k)$ . We start from the LHS:  $n + (m+S(k)) \stackrel{\text{def}}{=} n + S(m+k) \stackrel{\text{def}}{=} S(n+(m+k)) \stackrel{P(k)}{=} S((n+m)+k) \stackrel{\text{def}}{=} (n+m) + S(k)$ . Therefore  $P(k) \Rightarrow P(S(k))$ . By (P5) this is true for all  $k \in \mathbb{N}$ .

## Problem 6

We can solve this by constantly dividing by powers of 6 and noting the quotients.

$$\begin{aligned}
 10000_{10} &= 7776 + 1296 + 4 \cdot 216 + 36 + 4 \cdot 6 + 4 \\
 &= 6^5 + 6^4 + 4 \cdot 6^3 + 6^2 + 4 \cdot 6 + 4 \\
 &= 114144_6
 \end{aligned}$$

Therefore,  $10000_{10} = 114144_6$ .

## Problem 7

We define the ordering relation  $\leq$  as  $m \leq n \iff \exists k \in \mathbb{N} : n = m + k$ .

1. Let  $n \in \mathbb{N}$ . We can always write  $n = 0 + n$  so  $0 \leq n$ .
2. Let  $n \in \mathbb{N}$ . Consider  $S(n) = S(n+0) \stackrel{\text{def}}{=} n + S(0)$ . Therefore,  $n \leq S(n)$ .
3. Let  $m, n \in \mathbb{N}$  and assume  $m \leq n$ . Then,  $\exists k \in \mathbb{N}$  such that  $n = m + k$ . Then,  $S(n) = S(m + k) \stackrel{\text{comm}}{=} S(k + m) \stackrel{\text{def}}{=} k + S(m) \stackrel{\text{comm}}{=} S(m) + k$ . Therefore  $S(m) \leq S(n)$ .

Now, we recursively define  $n \cdot m$  as  $0 \cdot m = 0$  and  $\forall n \in \mathbb{N} : S(n) \cdot m = n \cdot m + m$ .

We are required to prove that  $\forall m, n, r \in \mathbb{N} : m \leq n \Rightarrow r \cdot m \leq r \cdot n$ . Assume that  $m \leq n$  so that  $\exists k \in \mathbb{N} : n = m + k$ . Let  $r \in \mathbb{N}$ . Then, by distributivity,  $r \cdot n = r \cdot (m + k) = r \cdot m + r \cdot k$ . Therefore,  $r \cdot m \leq r \cdot n$ .

## Problem 8

Denote  $[0, n) := \{k \in \mathbb{N} : k < n\}$ . Let  $m, n \in \mathbb{N}$ ,  $h : [0, n) \rightarrow [0, m)$ .

1. Fix  $n \in \mathbb{N}$ . We will prove the result by induction on  $m$ .  $P(0) : h : [0, n) \rightarrow [0, 0)$  is injective  $\Rightarrow n \leq m$ . Since we cannot define such a function, the statement is vacuously true. Assume  $P(m)$ . Let  $h : [0, n) \rightarrow [0, m + 1)$  be an injection. If  $h([0, n)) \subset [0, m)$  then we can redefine  $h$  as an injection  $[0, n) \rightarrow [0, m)$ . Then,  $n \leq m \leq m + 1$  and we are done. Otherwise,  $\exists$  exactly one  $j \in [0, n) : h(j) = m$ . We can then reconstruct the map and create a new function  $\tilde{h} : [0, n - 1) \rightarrow [0, m)$  which is still an injection. Here we define  $n - 1$  as  $S(n - 1) = n$ . Then,  $n - 1 \leq m$  so  $n \leq m + 1$ . Therefore, by (P5), the statement is true for all  $m \in \mathbb{N}$ .
2. Fix  $n \in \mathbb{N}$ . Once again, we induct on  $m$ .  $P(0)$  is vacuously true like before. Assume  $P(m)$ . Consider  $h : [0, n) \rightarrow [0, m + 1)$  a surjection. Then,  $\exists k \in [0, n)$  such that  $h(k) = m$ . We can remove  $k$  from a domain and reconstruct the function  $\tilde{h}$  in a way such that  $\tilde{h} : [0, n - 1) \rightarrow [0, m)$  is surjective. Then, by the inductive hypothesis,  $m \leq n - 1$  so  $S(m) \leq n$ . By (P5), this holds for all  $m \in \mathbb{N}$ .
3. This follows for the previous parts. If  $h$  is bijective, then it is injective (so  $n \leq m$ ) and surjective (so  $m \leq n$ ). Then  $\exists k : m = n + k, \exists l : n = m + l$ . By associativity,  $n = n + (k + l)$ . Therefore,  $k + l = 0 \Rightarrow k = 0, l = 0$ . So,  $m = n$ .