

MATH 154: Midterm 2

Professor Kane

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Question 1

Let the number of vertices be v , edges be e and faces be f . Further, let number of triangle faces be x and square faces be y . Then, $x+y = f$. Then, for any vertex $a \in V$, since it is connected to exactly 5 faces, we have $\deg(a) = 5$. Then, by the handshake lemma, $\sum_{a \in V} \deg(v) = 2 \cdot e \implies e = \frac{5 \cdot v}{2}$. Further, since triangles has 3 sides and a square has 4 sides, using the dual-handshake lemma, $\sum_{a \in F} \text{sides}(a) = 2 \cdot e \implies 3 \cdot x + 4 \cdot y = 5 \cdot v$. Now, since polyhedra are planar graphs, using euler's formula $v - e + f = 2 \implies v - \frac{5 \cdot v}{2} + f = 2 \implies f = x + y = 2 + \frac{3 \cdot v}{2}$. Plugging $v = 72$, we have a system of two equations, we solve for the 2 variables: $x + y = 110 \implies 3x + 3y = 330$ and $3x + 4y = 72 * 5 = 360$. Subtracting, $y = 30$. Plugging back in, $x = 110 - 30 = 80$. Therefore, the polyhedra has 80 triangle faces and 30 square faces.

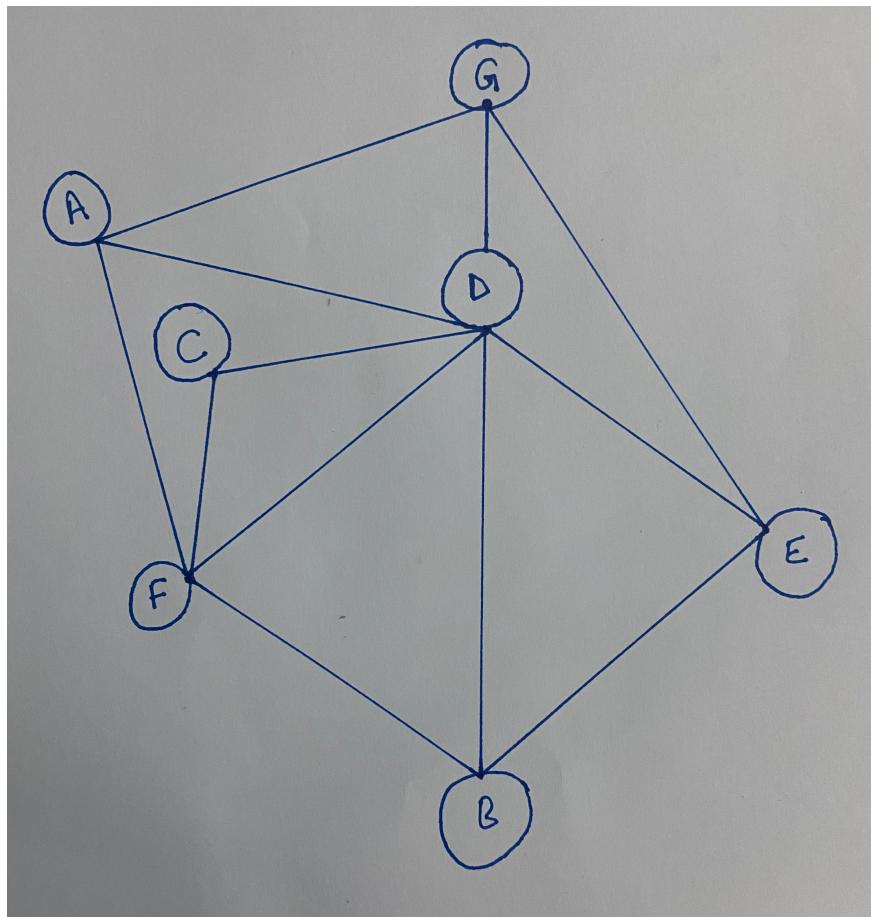
Question 2

The set of nodes with color 1 is $\{A, E, F, H\}$.

The set of nodes with color 2 is $\{B, C, D, G\}$.

The set of nodes with color 3 is $\{I, J\}$.

Question 3



Question 4

Let V_G represent the set of vertices in a graph G and E_G represent the set of edges in G , i.e. $G = (V_G, E_G)$. Also, for $u \in V_G$, let $\text{color}(u, G) =$ the color of node u in a coloring corresponding to chromatic number of graph G .

Lemma 1: For a complete graph H and any graph G' , in any valid colouring of $H \cup G'$, for any $u, v \in V_H$, $\text{color}(u, H \cup G') \neq \text{color}(v, H \cup G')$.

This is because for any $u, v \in V_H$, since H is a complete subgraph, there is an edge between vertex u and v . Therefore, any valid coloring cannot use the same color for both the vertices.

WLOG, let $\chi(H \cup E) \leq \chi(H \cup F)$.

Consider the coloring of $H \cup E$ corresponding to $\chi(H \cup E)$. Using lemma 1, we know that for any $u, v \in V_H$, $\text{color}(u, H \cup E) \neq \text{color}(v, H \cup E)$. Therefore, each node $u \in V_H$ has a different color, i.e. $|\{\text{color}(u, H \cup E) \forall u \in V_H\}| = |V_H|$. WLOG, let the set of colors $\{\text{color}(u, H \cup E) \forall u \in V_H\}$ be the set $\{1, 2, \dots, |V_H|\}$. This is always possible because we can just permute the set of colors so the labels/colors $\forall u \in V_H$ corresponds to $\{1, 2, \dots, |V_H|\}$.

Similarly, using the same reasoning as $H \cup E$, in the coloring of $H \cup F$, let the set of colors $\{\text{color}(u, H \cup F) \forall u \in V_H\}$ be the set $\{1, 2, \dots, |V_H|\}$. Further, we can permute the set $\{\text{color}(u, H \cup F) \forall u \in V_H\}$ such that $\text{color}(u, H \cup E) = \text{color}(u, H \cup F) \quad \forall u \in V_H$. This is always possible because we can set the color of u such that $\text{color}(u, H \cup E) = \text{color}(u, H \cup F)$ and permute the color of all vertices to obtain an equivalent valid coloring.

Now, using these two independent colorings, we can obtain a coloring of G based on three possibilities $\forall u \in V_G$: If $u \in V_H$, then $\text{color}(u, G) = \text{color}(u, H \cup E) = \text{color}(u, H \cup F)$. Else if $u \in E$, $\text{color}(u, G) = \text{color}(u, H \cup E)$. Else, $u \in F$, and $\text{color}(u, G) = \text{color}(u, H \cup F)$. Since E and F are separate components in $G \setminus H$, in this coloring, there are no edges (u, v) such that $u \in E$ and $v \in F$. So the set of neighbours for any node $u \in E \cup F$ is the same as the constituent colorings, where it had a valid color. For all nodes $u \in V_H$, since we obtained a coloring in both subgraphs $H \cup E$ and $H \cup F$ such that u had the same color, it must be a valid color in $G = H \cup E \cup F$. So, we have a coloring of G using $\chi(H \cup F)$ colors.

We have thus shown $\chi(G) \leq \chi(H \cup F)$.

Assume, for the sake of contradiction, that $\chi(G) < \chi(H \cup F)$. Then, consider the coloring of nodes in G corresponding to $\chi(G)$. Since $H \cup F$ is an induced subgraph of G , any coloring valid in G will also be valid in $H \cup F$. This is because the set of constraints defined by edges $E_{H \cup F} \subseteq E_G$. Therefore, we have obtained a coloring of $H \cup F$ such that we use fewer than $\chi(H \cup F)$ colors, which is a contradiction. Therefore, $\chi(G) \geq \chi(H \cup F)$.

Combining results, $\chi(G) = \chi(H \cup F)$.

Question 5

I have permission to take the exam at a non-standard time.

Start time: 7:30 AM Pacific Time. End time: 8:30 AM Pacific Time. Date: Saturday, May 23.