Heaps with STL	<pre>- int u = minDistance(dist, sptSet);</pre>	>getRightChildPtr()))	For example, in the following graph,
<pre>Complexity: 0(log n) vector<int> v;</int></pre>	- sptSet[u] = true;	<pre>- colorFlip(subTreePtr); - return subTreePtr;</pre>	we start traversal from vertex 2. When
<pre>vector(int) v; v.push_back(randomNum); cout << v[i]</pre>	<pre>- for (int v = 0; v < V; v++) - if (!sptSet[v] && graph[u][v] &&</pre>	- return subtreeptr;	we come to vertex 0, we look for all adjacent vertices of it. 2 is also
";</td <td>dist[u] != INT_MAX && dist[u] +</td> <td>MST/K-th Smallest Element</td> <td>an adjacent vertex of 0. If we don't</td>	dist[u] != INT_MAX && dist[u] +	MST/K-th Smallest Element	an adjacent vertex of 0. If we don't
<pre>make_heap(v.begin(), v.end());</pre>	graph[u][v] < dist[v])	<pre>int ksmallestElementSumRec(Node *root,</pre>	,
<pre>v.push_back(avgTotal); push_heap(v.</pre>	<pre>- dist[v] = dist[u] + graph[u][v];</pre>	int k, int &count)	processed again and it will become a
<pre>begin(), v.end()); //add vector for (size_t i=0; i<v.size(); i++)<="" pre=""></v.size();></pre>	Priority Queue	<pre>- if (root == NULL) - return 0;</pre>	<pre>non-terminating process. void Graph::addEdge(int v, int w)</pre>
cout << "Max" << v.front();	q.top(), q.pop(), q.push(n), q.size()	- if (count > k)	- adj[v].push_back(w);
<pre>pop_heap(v.begin(), v.end());</pre>	while(!q.empty())	- return 0;	void Graph::DFSUtil(int v, bool vis-
v.pop_back();	print g.top, g.pop	<pre>- int res = ksmallestElementSumRec(-</pre>	ited[])
<pre>sort_heap(v.begin(), v.end());</pre>	<pre>prime for loop (i=2; n/2; i++)</pre>	root->left, k, count);	- visited[v] = true; - cout << v << " ";
<pre>find-max/min, delete/extract-max/min, replace, insert, heapify, merge(k)/</pre>	Merge Sort	<pre>- if (count >= k) - return res;</pre>	- list <int>::iterator i;</int>
meld(r), isHeap, reverse	<pre>void merge(int arr[], int 1, int m,</pre>	- res += root->data;	- for (i = adj[v].begin(); i != ad-
	int h)	- count++;	j[v].end(); ++i)
Hashing	- int i, j, k;	- if (count >= k)	- if (!visited[*i])
<pre>bool search(int X) { - if (X >= 0)</pre>	- int n1 = m - l + 1; - int n2 = h - m;	<pre>- return res; - return res + ksmallestElementSum-</pre>	<pre>- DFSUtil(*i, visited); void Graph::DFS(int v)</pre>
- if (has[X][0] == 1)	- int L[n1], R[n2];	Rec(root->right, k, count);	<pre>- bool *visited = new bool[V];</pre>
- return true;	- for (i = 0; i < n1; i++)		- for (int i = 0; i < V; i++)
-X = abs(X);	- L[i] = arr[l + i];	RBT Insert	<pre>- visited[i] = false;</pre>
<pre>- if (has[X][1] == 1) - return true; }</pre>	- for (j = 0; j < n2; j++) - R[j] = arr[m + 1+ j];	<pre>RedBlackNode<itemtype>* LeftLean- ingRedBlackTree<itemtype>::inser-</itemtype></itemtype></pre>	<pre>- DFSUtil(v, visited);</pre>
void insert(int a[], int n)	-i = 0; j = 0; k = 1;	tRec(RedBlackNode <itemtype> *sub-</itemtype>	Breadth-First Search
- for (int i = 0; i < n; i++)	- while (i < n1 && j < n2)	TreePtr, RedBlackNode <itemtype></itemtype>	<pre>void Graph::addEdge(int v, int w)</pre>
- if (a[i] >= 0)	<pre>- if (L[i] <= R[j])</pre>	*newNodePtr)	- adj[v].push_back(w);
- has[a[i]][0] = 1;	- arr[k] = L[i]; i++;	- RedBlackNode <itemtype> *tempPtr =</itemtype>	void Graph::BFS(int s)
<pre>- else - has[abs(a[i])][1] = 1; }</pre>	<pre>- else - arr[k] = R[j]; j++;</pre>	<pre>nullptr; - if (subTreePtr == nullptr)</pre>	<pre>- bool *visited = new bool[V]; - for(int i = 0; i < V; i++)</pre>
int main()	- k++;	- return newNodePtr;	- visited[i] = false;
- int a[] = { -1, 9, -5, -8 }	- while (i < n1)	<pre>- if (subTreePtr->getItem() > newN-</pre>	- list <int> queue;</int>
<pre>- int n=sizeof(a)/sizeof(a[0]);</pre>	- arr[k] = L[i]; i++; k++;	odePtr->getItem())	<pre>- visited[s] = true;</pre>
- insert(a, n); - int X = -5;	- while (j < n2) - arr[k] = R[j]; j++; k++;	 subTreePtr->setLeftChildPtr(inser- tRec(subTreePtr->getLeftChildPtr(), 	<pre>- queue.push_back(s); - list<int>::iterator i;</int></pre>
- if (search(X) == true)	void mergeSort(int arr[], int 1, int	newNodePtr));	- while(!queue.empty())
- cout << "Present";	h)	- else	<pre>- s = queue.front();</pre>
	- if (1 < h)	- subTreePtr->setRightChildPtr(inser-	- cout << s << " ";
TSP #include <bits stdc++.h=""></bits>	<pre>- int m = 1+(h-1)/2; - mergeSort(arr, 1, m);</pre>	<pre>tRec(subTreePtr->getRightChildPtr(), newNodePtr));</pre>	<pre>- queue.pop_front(); - for (i = adj[s].begin(); i != ad-</pre>
int TSP(int graph[][V], int s)	<pre>- mergesort(arr, 1, m); - mergeSort(arr, m+1, h);</pre>	- if (isRed(subTreePtr->get-	<pre>i[s].end(); ++i)</pre>
<pre>- vector<int> vertex;</int></pre>	- merge(arr, 1, m, h);	RightChildPtr()) && !isRed(sub-	- if (!visited[*i])
- for (int i = 0; i < V; i++)		<pre>TreePtr->getLeftChildPtr()))</pre>	<pre>- visited[*i] = true;</pre>
- if (i != s)	Heap Sort	- subTreePtr = rotateLeft(sub-	<pre>- queue.push_back(*i);</pre>
<pre>- vertex.push_back(i); - int min_path = INT_MAX;</pre>	<pre>void heapify(int arr[], int n, int i) - int largest = i;</pre>	<pre>TreePtr); - if (isRed(subTreePtr->ge-</pre>	Prim's Algorithim
- do {	- int $1 = 2*i + 1$;	tLeftChildPtr()) && isRed(sub-	Start at any vertex, look for short-
<pre>- int current_pathweight = 0;</pre>	- int $r = 2*i + 2;$	TreePtr->getLeftChildPtr()->ge-	est on any visited vertex, without
- int k = s;	- if (1 < n && arr[1] > arr[largest])	tLeftChildPtr()))	cycles.
<pre>- for (int i = 0; i < vertex.size(); i++)</pre>	<pre>- largest = 1; - if (r < n && arr[r] > arr[largest])</pre>	<pre>- subTreePtr = rotateRight(sub- TreePtr);</pre>	Kruskal's MST
<pre>- current_pathweight += graph[k]</pre>	- largest = r;	- if (isRed(subTreePtr->ge-	Start at the smallest cost, keep as-
[vertex[i]];	- if (largest != i)	tLeftChildPtr()) && isRed(sub-	cending up without cycle.
<pre>- k = vertex[i];</pre>	<pre>- swap(arr[i], arr[largest]);</pre>	<pre>TreePtr->getRightChildPtr()))</pre>	void KruskalMST(Graph* graph)
<pre>- current_pathweight += graph[k][s]; min_pathmin(min_pathcurrent)</pre>	- heapify(arr, n, largest);	<pre>- colorFlip(subTreePtr); - return subTreePtr;</pre>	- int V = graph->V;
<pre>- min_path = min(min_path, current_ pathweight); }</pre>	<pre>void heapSort(int arr[],int n) - for (int i=n/2-1; i>=0; i)</pre>	- recurr subfreercr,	<pre>- Edge result[V]; - int e = 0;</pre>
- while (next_permutation(vertex.	heapify(arr, n, i);	RBT Rotate Left	- int i = 0;
<pre>begin(), vertex.end()));</pre>	- for (int i=n-1; i>=0; i)	RedBlackNode <itemtype>* LeftLean-</itemtype>	- qsort(graph->edge, graph->E, sizeo-
Didleston	- swap(arr[0], arr[i]);	ingRedBlackTree <itemtype>::ro-</itemtype>	f(graph->edge[0]), myComp);
Dijkstra Complexity: O(n^2)	<pre>- heapify(arr, i, 0);</pre>	<pre>tateLeft(RedBlackNode<itemtype> *subTreePtr)</itemtype></pre>	<pre>- subset *subsets = new subset[(V * sizeof(subset))];</pre>
int minDistance(int dist[], bool	RBT Fix-Up	- RedBlackNode <itemtype> *tempPtr =</itemtype>	- for (int v = 0; v < V; ++v)
sptSet[])	RedBlackNode <itemtype>* LeftLeanin-</itemtype>	<pre>subTreePtr->getRightChildPtr();</pre>	<pre>- subsets[v].parent = v;</pre>
- for (int v = 0; v < V; v++)	gRedBlackTree <itemtype>::fixUp(Red-</itemtype>	- subTreePtr->setRightChildPtr(-	- subsets[v].rank = 0;
<pre>- if (sptSet[v] == false && dist[v] <= min)</pre>	<pre>BlackNode<itemtype> *subTreePtr) - if (isRed(subTreePtr-</itemtype></pre>	<pre>tempPtr->getLeftChildPtr()); - tempPtr->setLeftChildPtr(sub-</pre>	<pre>- while (e < V - 1 && i < graph->E) - Edge next_edge = graph->edge[i++];</pre>
<pre>- min = dist[v], min_index=v</pre>	>getRightChildPtr()))	TreePtr);	int x = find(subsets, next_edge.
<pre>- return min_index;</pre>	- subTreePtr =	- tempPtr->setRed(sub-	src);
<pre>void dijkstra(int graph[V][V], int</pre>	<pre>rotateLeft(subTreePtr); if (i=Pat(subTreePtr);</pre>	<pre>TreePtr->getRed());</pre>	<pre>- int y = find(subsets, next_edge.</pre>
src)	<pre>- if (isRed(subTreePtr- >getLeftChildPtr()) &&</pre>	- subTreePtr->setRed(1);	dest); - if (x != y)
<pre>- int dist[V]; - bool sptSet[V];</pre>	isRed(subTreePtr->getLeftChildPtr()-	- return tempPtr;	- ir (x != y) - result[e++] = next_edge;
- for (int i = 0; i < V; i++)	>getLeftChildPtr()))	Depth-First Search	- Union(subsets, x, y);
<pre>- dist[i] = INT_MAX, sptSet[i] =</pre>	- subTreePtr =	The only catch here is, unlike trees,	- cout<<"Following are the edges in
false;	rotateRight(subTreePtr);	graphs may contain cycles, so we may	the constructed MST\n";
<pre>- dist[src] = 0; - for (int count = 0; count < V - 1;</pre>	<pre>- if (isRed(subTreePtr- >getLeftChildPtr())</pre>	come to the same node again. To avoid processing a node more than once, we	- for (i = 0; i < e; ++i) - cout< <result[i].src<<" "<<re-<="" td=""></result[i].src<<">
count++)	&& isRed(subTreePtr-	use a boolean visited array.	sult[i].dest<<" == "< <result[i].< td=""></result[i].<>
•	•	,	

weight<<endl;
- return:</pre>

BST Destroy Tree / Traversals

Remember: if (rootPtr != nullptr) { }
void BinarySearchTree<!temType>::de stroyTree(BinaryNode<!temType>
 *subTreePtr)

- if (subTreePtr != nullptr)
- destroyTree(subTreePtr->getLeftChildPtr());
- destroyTree(subTreePtr->get-RightChildPtr());
- RightChildPtr());
 subTreePtr.reset();

Preorder: C;P;P;
Inorder: I;C;I;
Postorder: P;P;C;

BST Remove Value

auto BinarySearchTree<ItemType>::removeValue(BinaryNode<ItemType> *subTreePtr, const ItemType target, bool
&isSuccessful)

- BinaryNode<ItemType> *tempPtr =
 nullptr;
- if (subTreePtr == nullptr)
- isSuccessful = 0;
- else if (subTreePtr->getItem() ==
 target)
- subTreePtr = removeNode(sub-TreePtr);
- isSuccessful = 1;
- else if (subTreePtr->getItem() >
 target)
- tempPtr = removeValue(sub-TreePtr->getLeftChildPtr(), target, isSuccessful);
- subTreePtr->setLeftChildPtr(tempPtr);
- else
- tempPtr = removeValue(sub-TreePtr->getRightChildPtr(), target, isSuccessful);
- subTreePtr->setRightChildPtr(tempPtr);
- return subTreePtr;

BST Remove

bool BinarySearchTree<ItemType>::remove(const ItemType &target)

- bool isSuccessful = 0;
- this->rootPtr = removeValue(this->rootPtr, target, isSuccessful);
- return isSuccessful;

BST Copy Tree

auto BinarySearchTree<ItemType>:: copyTree(const BinaryNode<ItemType>
 *oldTreeRootPtr) const

- BinaryNode<ItemType> *newTreePtr =
 nullptr;
- if (oldTreeRootPtr != nullptr)
- newTreePtr = new BinaryNode(Item-Type>(oldTreeRootPtr->getItem(), nullptr, nullptr);
- newTreePtr->setLeftChildPtr(copyTree(oldTreeRootPtr->getLeftChildPtr()));
- newTreePtr->setRightChildPtr(copyTree(oldTreeRootPtr->getRightChildPtr()));
- return newTreePtr;

BST Add

bool BinarySearchTree<ItemType>::add(const ItemType &newData)

- BinaryNode<ItemType> *newNodePtr = new BinaryNode<ItemType>(newData);
- this->rootPtr = placeNode(rootPtr,

newNodePtr);
- return 1;

BST Place Node

auto BinarySearchTree<ItemType>::placeNode(BinaryNode<ItemType>
 *subTreePtr, BinaryNode<ItemType>
 *newNodePtr)

- BinaryNode<ItemType> *tempPtr =
 nullptr;
- if (subTreePtr == nullptr)
- return newNodePtr;
- else if (subTreePtr->getItem() >
 newNodePtr->getItem())
- tempPtr = placeNode(subTreePtr->getLeftChildPtr(), newNodePtr);
- subTreePtr->setLeftChildPtr(tempPtr);
- else
- tempPtr = placeNode(sub-
- TreePtr->getRightChildPtr(), newNodePtr);
- subTreePtr->setRightChildPtr(tempPtr);
- return subTreePtr;

Binary Search Tree

The depth of a node is the number of edges from the root to the node. The height of a node is the number of edges from the node to the deepest leaf.

The height of a tree is a height of the root.

A full binary tree is a binary tree in which each node has exactly zero or two children.

A complete binary tree is a binary tree, which is completely filled, with the possible exception of the bottom level, which is filled from left to right.

 $h = O(\log n)$

Search/Insert/Delete = avg: O(log n)
worst: O(n)

AVL/Balanced Search Tree

Height difference between left and
right node cannot exceed 1
Search/Insert/Delete = O(logn)

Red Black Tree

Every node has a color either red or black.

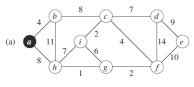
Root of tree is always black.
There are no two adjacent red nodes (A red node cannot have a red parent or red child).

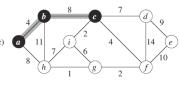
 $h = O(\log n)$

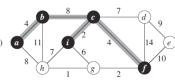
Search/Insert/Delete = O(logn)
Color Change = O(1) * height

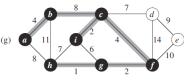
2-3-4 Trees

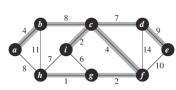
Insertion into a 2-3-4 Tree begins with a single node where values are inserted until it becomes full (ie. until it becomes a 4-node). The next value that is inserted will cause a split into two nodes: one containing values less than the median value and the other containing values greater than the median value. The median value is then stored in the parent node. It's possible that insertion could cause splitting up to the root node

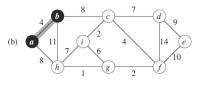


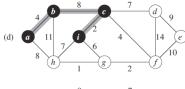


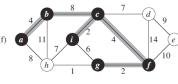


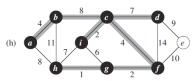


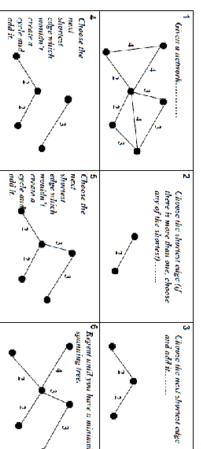


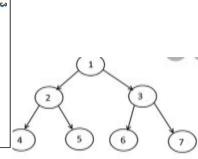












Inorder Traversal: 4251637 Preorder Traversal: 1245367 Postorder Traversal: 7635421 Breadth-First Search: 1234567 Depth-First Search: 1245367