Homework 3 NICK Alvarez

QL

- a) This algorithm determines if the elements of an array are unique. If not, it returns "false".
- b) for i = 0 to n-2 do

 All elements |

 for $j \in i+1$ to n-1 do

 All elements |

 if A[i] == A[j]return false

Worst case: Entire array traversed

Average case: Some of the array is pearched

So we can conclude the running time is $O(n^2)_r$ due to the nested "for" loops.

Q2

- a) See problem 2. CPP
- b) The answer will still be 9. As the left side

 1s evaluated first, the first "9" will be the

 max the computer reports, even though there

 1s a "9" on the right and middle.

 The > operator ensures this as left max

 1s compared to right Max, instead of using >=.

c)
$$(omparisons(h) = T(n)$$

 $T(n) = 2T(n/2) + 2$ when $n \ge 2$
 $T(1) = 0$

$$T(n) = 2T(n/2) + 1$$

$$L_{3} 2[2T(n/4) + 1] + 1$$

$$L_{4} 2[2[2T(n/8) + 1] + 1] + 1$$

$$\vdots$$

$$L_{5} 2^{i}T(n/8) + 4 + 2 + 1$$

$$\vdots$$

$$L_{5} 2^{i}T(n/2^{i}) + 2^{i-1} + 2^{i-2} + \dots + 2 + 1$$

$$L_{5} 2^{i-1} + 2^{i-2} + \dots + 2 + 1$$

$$L_{5} 2^{i-1} + 2^{i-2} + \dots + 2 + 1$$

$$L_{5} 2^{i-1} + 2^{i-2} + \dots + 2 + 1$$

$$L_{5} 2^{i-1} + 2^{i-2} + \dots + 2 + 1$$

We can equate this to Case 1 of Master's method and find that $T(n) = O(\log n)$