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## Exercises

Section 1.1: 1, 5(b,c,e), 8 (b), 12(c,d)

1. In high school, some students have been misled to believe that  $22/7$  is either the actual value of  $\pi$  or an acceptable approximation to  $\pi$ . Show that  $355/113$  is a better approximation in terms of both absolute and relative errors. Find some other simple rational fractions  $n/m$  that approximate  $\pi$ . For example, ones for which  $|\pi - n/m| < 10^{-9}$ . Hint: See Problem 1.1.4.

**Answer:**

$$\pi = 3.1415927$$

$$355/113 = 3.1415929$$

$$22/7 = 3.1428571$$

$$\varepsilon_1 = \left| \pi - \frac{22}{7} \right| \approx 1.2 \cdot 10^{-3}$$

$$\varepsilon_2 = \left| \pi - \frac{355}{113} \right| \approx 2.7 \cdot 10^{-7}$$

$355/113$  a better approximation of  $\pi$

$$\eta_1 = \left| \frac{\pi - \frac{22}{7}}{\pi} \right| \approx 4.0 \cdot 10^{-4}$$

$$\eta_2 = \left| \frac{\pi - \frac{355}{113}}{\pi} \right| \approx 8.5 \cdot 10^{-8}$$

$355/113$  has lower relative error in approximating  $\pi$

Example:

$\frac{n^x \pi}{n^x}$  we can use  $n$ ,  $x$  is any number like  $(9^n \pi) / 9^n = 3.141592654$ . Then  $|\pi - n/m| < 10^{-9}$

5. A given doubly subscripted array  $(a_{ij})_{n \times n}$  can be added in any order. Write the pseudocode segments for each of the following parts. Which is best?

**Answer:**

$$\text{b. } \sum_{j=1}^n \sum_{i=1}^n a_{ij}$$

**integer** sum, j, l, n

sum = 0

**for** j = 1 **to** n **do**

**for** i = 1 **to** n **do**

        sum = sum + a<sub>ji</sub>

**end for**

**end for**

$$c. \sum_{i=1}^n \left( \sum_{j=1}^i a_{ij} + \sum_{j=1}^{i-1} a_{ji} \right)$$

**integer** sum, j, l, n

sum = 0

**for** i = 1 **to** n **do**

**for** j = 1 **to** n **do**

        sum = sum + a<sub>ij</sub>

**end for**

**for** j = 1 **to** i-1 **do**

        sum = sum + a<sub>ji</sub>

**end for**

**end for**

$$e. \sum_{k=2}^{2n} \sum_{i+j=k}^n a_{ij}$$

**integer** sum, x, y, k, j, l, n

sum = 0

**for** i = 1 **to** n **do**

**for** j = 1 **to** n **do**

**for** k = i+j **to** n **do**

            sum = sum + a<sub>ij</sub>

**end for**

**end for**

**end for**

In each example it has there own algorithm. By the consider time worst, then sample b is les time than another one.

8. Show how these polynomials can be efficiently evaluated:

$$b. p(x) = 3(x - 1)^5 + 7(x - 1)^9$$

**Answer:**

$$\begin{aligned} b. p(x) &= 3(x - 1)^5 + 7(x - 1)^9 \\ &= 3(x - 1)^5 + 7(x - 1)^4(x - 1)^5 \\ &= (x-1)^5 + (3 + 7(x - 1)^4) \end{aligned}$$

Let  $z = (x-1)$

$$\text{Then } p(x) = z^5 + (3+7z^4)$$

We can compute  $z^5 = z^4 * z$  using one multiplication.

Then we compute  $3 + 7*z^4$  , and then multiply the result by  $z^5$ .

12. Using summation and product notation, write mathematical expressions for the following pseudocode segments:

```
c. integer  $i, n$ ; real  $v, x$ ; real array  $(a_i)_{0:n}$   
 $v \leftarrow a_0$   
for  $i = 1$  to  $n$  do  
   $v \leftarrow vx + a_i$   
end for
```

**Answer:**

$$v_0 = a_0.$$

$$v_1 = a_0 x + a_1.$$

$$v_2 = (a_0 x + a_1)x + a_2 = a_0 x^2 + a_1 x + a_2 .$$

...

$$V_n = a_0 x^n + a_1 x^{n-1} + \dots + a_n .$$

$$V = \sum_{i=0}^n a_{n-i} x^i$$

**d. integer  $i, n$ ; real  $V, x, z$ ; real array  $(a_i)_{0:n}$**

$V \leftarrow a_0$

$z \leftarrow x$

**for  $i = 1$  to  $n$  do**

$V \leftarrow V + z a_i$

$z \leftarrow xz$

**end for**

**Answer:**

$$V_0 = a_0.$$

$$V_1 = a_0 + a_1 x$$

$$V_2 = a_2 x^2 + a_1 x + a_0.$$

...

$$V_n = a_n x^n + a_{n-1} x^{n-1} + \dots + a_0.$$

$$V = \sum_{i=0}^n a_i x^i$$

### **Section 1.2: 1, 4(d,f), 14, 44**

**1. The Maclaurin series for  $(1+x)^n$  is also known as the binomial series. It states that**

$$(1+x)^n = 1 + nx + \frac{n(n-1)}{2!} x^2 + \frac{n(n-1)(n-2)}{3!} x^3 + \dots \quad (x^2 < 1)$$

**Derive this series. Then give its particular forms in summation notation by letting**

**$n = 2$ ,  $n = 3$ , and  $n = 1/2$ . Next use the last form to compute  $\sqrt{1.0001}$  correct to 15 decimal places (rounded).**

**Answer:**

$$\text{Let } f(x) = (1+x)^n$$

$$\text{Then } f'(x) = n(1+x)^{n-1}.$$

$$\text{Then } f''(x) = n(n-1)(1+x)^{n-2}.$$

$$\text{Then } f'''(x) = n(n-1)(n-2)(1+x)^{n-3}$$

$$\text{Then } f^{(4)}(x) = n(n-1)(n-2)(n-3)(1+x)^{n-4}$$

We have Maclaurin:

$$f(x) = f(0) + f'(0)x + \frac{f''(0)}{2!}x^2 + \frac{f'''(0)}{3!}x^3 + \dots$$

So:  $f(0) = 1$ ;  $f'(0) = n$ ;  $f''(0) = n(n-1)$ ;  $f'''(0) = n(n-1)(n-2)$ ;  $f^{(4)}(x) = n(n-1)(n-2)(n-3)$

Therefore,

$$f(x) = 1 + nx + \frac{n(n-1)x^2}{2!} + \frac{n(n-1)(n-2)x^3}{3!} + \dots$$

$$\text{Hence, } (1+x)^n = 1 + nx + \frac{n(n-1)x^2}{2!} + \frac{n(n-1)(n-2)x^3}{3!} + \dots \quad (x^2 < 1)$$

$$\text{For } n = 2: \text{ then } (1+x)^2 = 1 + 2x + x^2$$

$$\text{For } n = 3: \text{ then } (1+x)^3 = 1 + 3x + 3x^2 + x^3$$

$$\text{For } n = 1/2: \text{ then } (1+x)^{1/2} = 1 + (1/2)x - (1/2^3)x^2 + (1/2^4)x^3 + \dots$$

Then with  $n = 1/2$ , choose  $x = 0.0001$

$$(1+0.0001)^{1/2} = 1 + (1/2)0.0001 - (1/2^3)0.0001^2 + (1/2^4)0.0001^3 + \dots = 1.0000499987500625$$

Round off value we have  $(1.0001)^{1/2} = \mathbf{1.000049998750062}$

#### 4. Why do the following functions not possess Taylor series expansions at $x = 0$ ?

**Answer:**

d.  $f(x) = \cot x$

$$\rightarrow f(x) = f(0) + f'(0)x + \frac{f''(0)x^2}{2!} + \frac{f'''(x)x^3}{3!} + \dots$$

We have:

$$f(x) = \cot x$$

$$f'(x) = -\csc^2 x \text{ then } f'(0) = -\infty$$

Hence, Taylor Series about  $x = 0$  does not exist as the first derivative  $f'(0) = -\infty$

f.  $f(x) = x^\pi$

$$\rightarrow f(x) = f(0) + f'(0)x + \frac{f''(0)x^2}{2!} + \frac{f'''(x)x^3}{3!} + \dots$$

We have:

$$f(x) = x^\pi$$

$$f'(x) = \pi (x)^{3.14 - 1}$$

$$f''(x) = \pi (2.14) (x)^{3.14 - 2}$$

$$f'''(x) = \pi (2.14)(1.14) (x)^{3.14 - 3}$$

$$f^{(4)}(x) = \pi (2.14)(1.14)(0.14) (x)^{3.14 - 4}$$

When  $x = 0$ , at  $f^{(4)}(0) = -\infty$

at  $x=0$ , the function become  $-\infty$

Hence, Taylor Series about  $x = 0$  does not exist as the function become  $-\infty$  when  $x = 0$

14. Write the Taylor series for the function  $f(x) = x^3 - 2x^2 + 4x - 1$ , using  $x = 2$  as the point of expansion; that is, write a formula for  $f(2 + h)$ .

**Answer:**

$$f(2+h) = f(2) + hf'(2) + (h^2/2!)f''(2) + (h^3/3!)f'''(2) + (h^4/4!)f^{(4)}(2)$$

$$f(x) = x^3 - 2x^2 + 4x - 1 \quad \text{then } f(2) = 7$$

$$f'(x) = 3x^2 - 4x + 4 \quad \text{then } f'(2) = 8$$

$$f''(x) = 6x - 4 \quad \text{then } f''(2) = 8$$

$$f'''(x) = 6 \quad \text{then } f'''(2) = 6$$

$$f^{(4)}(x) = 0 \quad \text{then } f^{(4)}(2) = 0$$

$$\begin{aligned} f(2+h) &= 7 + 8h + (h^2/2!)8 + (h^3/3!)6 + (h^4/4!)0 \\ &= 7 + 8h + 4h^2 + h^3 \end{aligned}$$

$$\text{Hence, } f(2+h) = 7 + 8h + 4h^2 + h^3$$

44. How many terms are needed in the series

$\operatorname{arccot} x = \frac{\pi}{2} - x + \frac{x^3}{3} - \frac{x^5}{5} + \frac{x^7}{7} - \dots$  to compute  $\operatorname{arccot} x$  for  $x^2 < 1$  accurate to 12 decimal places (rounded)?

**Answer:**

Consider:

$$\operatorname{arccot} x = \frac{\pi}{2} - x + \frac{x^3}{3} - \frac{x^5}{5} + \frac{x^7}{7} - \dots$$

That is:

$$\operatorname{arccot} x = \frac{\pi}{2} \sum_{n=1}^{\infty} (-1)^{-n} \frac{x^{2n-1}}{(2n-1)}$$

$$\text{Therefore, the } (n+1)^{\text{th}} \text{ term } (-1)^{-n} \frac{x^{2n-1}}{(2n-1)}$$

$$\text{We have for accurate to 12 decimal places : } \left| (-1)^{-n} \frac{x^{2n-1}}{(2n-1)} \right| < \frac{1}{2} \cdot 10^{-12}$$

$$\text{Now: } \frac{x^{2n-1}}{(2n-1)} < \frac{1}{2} \cdot 10^{-12}$$

$$\text{For } x^2 < 1: \frac{1}{(2n-1)} < \frac{1}{2} \cdot 10^{-12}$$

$$\text{Then: } n > \frac{1}{2} + 10^{12}$$

Hence, at least  $10^{12} + 1$  terms are need in the series to compute arc cot x, for accurate to 12 decimal places.

### Section 2.1: 1(a,b), 3, 5(e-h), 12, 20, 36

1. Determine the machine representation in single precision on a 32-bit word-length computer for the following decimal numbers.

a.  $2^{-30}$

b. 64.01562 5

Answer:

$$\text{a. } 2^{-30} = (-1)^0 * 2^{127-30} * (1.0)$$

$$97/2 = 1$$

$$48/2 = 0$$

$$24/2 = 0$$

$$12/2 = 0$$

$$6/2 = 0$$

$$3/2 = 1$$

$$\frac{1}{2} = 1$$

Read from bottom to top  $(1100001)_2$ . We have substitute binary form  $2^{97}$  is

0 01100001 000000000000000000000000

Hence, the single precision machine representation in a 32 bit word length computer  $2^{-30}$  is

0 01100001 000000000000000000000000

b. 64.015625

$$(64)_{10} = (1000000)_2$$

$$64/2 = 0$$

$$32/2 = 0$$

$$16/2 = 0$$

$$8/2 = 0$$

$$4/2 = 0$$

$$2/2 = 0$$

$$\frac{1}{2} = 1$$

$$(.015625)_{10} = (000001)_2$$

$$.015625 * 2 = 0$$

$$.03125 * 2 = 0$$

$$.0625 * 2 = 0$$

$$0.125 * 2 = 0$$

$$0.25 * 2 = 0$$

$$0.5 * 2 = 1$$

$$(64.015625)_{10} = (1000000.000001)_2 = (1.000000000001 * 2^6)_2 = (-1)^0 * 2^{127+6} * 2^{-127} (1.000000000001) \\ = (-1)^0 * 2^{133} * 2^{-127} (1.000000000001)$$

$$(133)_{10} = (10000101)_2$$

$$133/2 = 1$$

$$66/2 = 0$$

$$33/2 = 1$$

$$16/2 = 0$$

$$8/2 = 0$$

$$4/2 = 0$$

$$2/2 = 0$$

$$\frac{1}{2} = 1$$

We have substitute binary form  $2^{133}$  is

0 10000101 000000 000001000000000000

Hence, the single precision machine representation in a 32 bit word length computer 64.015625 is

0 10000101 000000 000001000000000000



a.  $10^{403}$

b.  $1 + 2^{-32}$

c.  $1/5$

d.  $1/10$

e.  $1/256$

a.  $10^{403}$  Since  $10^3 \approx 2^{10}$ ,  $10^{403} \approx 2^{1340}$ . This is beyond range of the exponent even in double -precision.  
So  $10^{403}$  is not a machine number

c.  $1/5$  is not a finite of power of 2 because must have denominator that is power of 2 in reduce form. Hence  $1/5$  is not a finite of power of 2. We consider that  $1/5$  is not a machine learning.

e.  $1/256$  is  $1 \cdot 2^{-8}$  and can be represented exactly within 23 bit mantissa. This is a machine number

e. 0 00000001 000000000000000000000000

h. 0 01111011 10011001100110011001100

e. 0 00000001 000000000000000000000000

h. 0 01111011 10011001100110011001100

Sign is positive

01111011 = 243 the actual biased exponent is  $123 - 127 = -4$ .

1.10011001100110011001100

Since 1.100 re present for  $3/5 = 1.5$

The geometric series evaluate to  $3/2(1+1/16 + 1/16^2 + 1/16^2 + 1/16^3 + 1/16^4 + 1/16^5) = 1.599999905$

Hence, we have  $1.599999905 * 2^{-4} = 0.099999995$

**12. What are the machine numbers immediately to the right and left of  $2^m$ ? How far is each from  $2^m$ ?**

**Answer:**

$1 + \text{epsilon} \neq 1$

On a floating point number on a 32-bit word-length machine,  $\text{epsilon} = 2^{-23}$ .

$2^m$  represented on the machine by an effective exponent of  $m$ , and mantissa of 1.0 (empty mantissa).

. The smallest machine number to its right given by machine by incrementing the rightmost bit in mantissa:

$$x = 2^m + (1 + \text{epsilon}).$$

$$x - 2^m = 2^m * \text{epsilon}$$

On the a floating point number on 32 bit word length machine, this distance is  $2^{m-23}$ .

. The largest machine number to the left is given by decrementing the exponent, but using a full mantissa (1.111.....1):

$$y = 2^{m-1}(2 - \text{epsilon}) \text{ the distance from } 2^m$$

$$2^m - y = 2^{m-1} * \text{epsilon}$$

On the a floating point number on 32 bit word length machine, this distance is  $2^{m-24}$ .

**20. What is the roundoff error when we represent  $2^{-1} + 2^{-25}$  by a machine number? Note:**

**This refers to absolute error, not relative error.**

**Answer:**

$x = 2^{-1} + 2^{-25} = 2^{-1} (1 + 2^{-24})$  is not a machine number since the mantissa require at least 24 bits, whereas we only have 23 bits allocated in a single precision floating point number representation.

The two closet machine number are the one to its right  $x^+ = 2^{-1} (1 + 2^{-23})$

and the one to its left.  $x^- = 2^{-1} (1 + 0)$

Turn out that they are equally distance from  $x$ , and the distance is

$$2^{-1}(1+2^{-24}) - 2^{-1} (1 + 0) = 2^{-25}$$

36. Show by an example that in computer arithmetic  $a + (b+c)$  may differ from  $(a + b) + c$ .

**Answer:**

$$1 + \text{epsilon} \neq 1$$

Let  $a = 1$ ;  $b = \text{epsilon}$ ;  $c = -\text{epsilon}$  (something can cause overflow)

$$\text{Consider } a + (b + c) = 1$$

$$\text{Consider } (a + b) = 1 + \text{epsilon} \text{ (cause overflow)}$$

$$(a + b) + c = 1 - \text{epsilon} \text{ (error here already) } \neq 1$$

Hence,  $a + (b + c) = 1$  while  $(a + b) + c \neq 1$ , then  $a + (b + c) \neq (a + b) + c$

## Section 2.2: 2, 3, 9, 14, 17

2. Calculate  $f(10^{-2})$  for the function  $f(x) = e^x - x - 1$ . The answer should have five significant figures and can easily be obtained with pencil and paper. Contrast it with the straightforward evaluation of  $f(10^{-2})$  using  $e^{0.01} \approx 1.0101$ .

**Answer:**

$$f(x) = e^x - x - 1$$

Taylor series expansion for  $e^x$  is

$$e^x = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \frac{x^4}{4!} + \frac{x^5}{5!} + \dots$$

$$\text{then } f(x) = e^x - x - 1 = \frac{x^2}{2!} + \frac{x^3}{3!} + \frac{x^4}{4!} + \frac{x^5}{5!} + \dots$$

$$\text{The answer should have five significant figures } f(x) \approx \frac{x^2}{2!} + \frac{x^3}{3!} + \frac{x^4}{4!}$$

$$f(0.01) \approx \frac{0.01^2}{2!} + \frac{0.01^3}{3!} + \frac{0.01^4}{4!} = 5.0167 \cdot 10^{-5}$$

Instead, we use the approximation

$$e^{0.01} \approx 1.0101.$$

$$f(0.01) \approx 1.0101 - 0.01 - 1 = 0.0001.$$

3. What is a good way to compute values of the function  $f(x) = e^x - e$  if full machine precision is needed?

Note: There is some difficulty when  $x = 1$ .

**Answer:**

When  $x = 1$  then  $e^x = e$

$f(x) = e^x - e$  then function cannot compute reliability.

When  $x > 1$  then  $e^x > e$ . We use the loss of precision theorem to determine the range of  $x$  in which at miss one bit lost in the subtraction  $f(x) = e^x - e$ .

### THEOREM 1

### LOSS OF PRECISION THEOREM

Let  $x$  and  $y$  be normalized floating-point machine numbers, where  $x > y > 0$ . If  $2^{-p} \leq 1 - (y/x) \leq 2^{-q}$  for some positive integers  $p$  and  $q$ , then at most  $p$  and at least  $q$  significant binary bits are lost in the subtraction  $x - y$ .

$\frac{1}{2} < 1 - (e/e^x) \rightarrow 1 + \ln 2 \leq x$ . Hence, when  $x$  is not close to 1 in particular when  $x \leq 1 - \ln 2$  or  $1 + \ln 2 \leq x$  we can use original expression to compute  $f(x)$ .

$$e^x = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \frac{x^4}{4!} + \frac{x^5}{5!} + \dots$$

$$f(x) = e \left[ (x-1) + \frac{(x-1)^2}{2!} + \frac{(x-1)^3}{3!} + \frac{(x-1)^4}{4!} + \frac{(x-1)^5}{5!} + \dots \right]$$

Let  $z = x-1$

$$f(z) = e \left[ (z) + \frac{(z)^2}{2!} + \frac{(z)^3}{3!} + \frac{(z)^4}{4!} + \frac{(z)^5}{5!} + \dots \right]$$

We can take for example 100 terms to get the full machine precision

There will be some difficulty when  $x=1$  because the series may not be summable

9. For some values of  $x$ , the assignment statement  $y \leftarrow 1 - \cos x$  involves a difficulty. What is it, what values of  $x$  are involved, and what remedy do you propose?

**Answer:**

When  $x \approx 2\pi n \rightarrow \cos(x) \approx 1$ , so  $y = 1 - \cos(x)$  cannot be computed reliability

When  $x > 1$  then  $1 - \cos(x)$ . We use the loss of precision theorem to determine the range of  $x$  in which at miss one bit lost in the subtraction  $f(x) = 1 - \cos(x)$ .

Loss of precision theorem:  $\frac{1}{2} \leq 1 - \cos(x) \rightarrow \cos(x) \leq \frac{1}{2}$ . then  $x$  outside the range  $[-\pi/3 + 2n\pi, \pi/3 + 2n\pi]$

We have  $f(x) = (1 - \cos(x)) * 1$ .

$$= (1 - \cos(x)) * (1 + \cos(x)) / (1 + \cos(x))$$

$$= \sin^2 x / (1 + \cos(x))$$

When  $x$  in range  $[-\pi/3 + 2n\pi, \pi/3 + 2n\pi]$  with  $f(x) = \sin^2 x / (1 + \cos(x))$ . Otherwise, we can use the original expression as-is to compute  $f(x) = y$ .

14. How can values of the function  $f(x) = \sqrt{x+2} - \sqrt{x}$  be computed accurately when  $x$  is large?

**Answer:**

when  $x$  is large  $\sqrt{x+2} \approx \sqrt{x}$ . Since  $\sqrt{x+2} - \sqrt{x}$  cannot be computed reliably

To avoid subtraction of two close numbers,

$$\begin{aligned} f(x) &= \sqrt{x+2} - \sqrt{x} \cdot 1 \\ &= \sqrt{x+2} - \sqrt{x} \cdot \frac{\sqrt{x+2} + \sqrt{x}}{\sqrt{x+2} + \sqrt{x}} \\ &= \frac{2}{\sqrt{x+2} + \sqrt{x}} \end{aligned}$$

When  $x$  is large, we can further write

$$f(x) = \frac{2}{\sqrt{x+2} + \sqrt{x}} \approx \frac{2}{\sqrt{x} + \sqrt{x}} \approx \frac{1}{\sqrt{x}} \text{ to avoid numerical loss}$$

17. Without using series, how could the function

$$f(x) = \frac{\sin x}{x - \sqrt{x^2 - 1}}$$

be computed to avoid loss of significance?

**Answer:**

When  $x$  is large  $\sqrt{x^2 - 1} \approx x$ . Since  $x - \sqrt{x^2 - 1}$  cannot be computed reliably

Loss of precision theorem:

$\frac{1}{2} \leq 1 - \frac{\sqrt{x^2 - 1}}{x} \rightarrow x \leq \sqrt{\frac{4}{3}}$  in the range in which the original expression will not suffer too much loss of precision.

$$\text{We have: } x - \sqrt{x^2 - 1} = x - \sqrt{x^2 - 1} \cdot 1 = x - \sqrt{x^2 - 1} \cdot \frac{x + \sqrt{x^2 - 1}}{x + \sqrt{x^2 - 1}} = \frac{1}{x + \sqrt{x^2 - 1}}$$

When  $x \leq \sqrt{\frac{4}{3}}$  we compute  $f(x) = \frac{\sin x}{x - \sqrt{x^2 - 1}}$  but use instead

$$f(x) = \frac{\sin x}{\frac{1}{x + \sqrt{x^2 - 1}}} = \sin x (x + \sqrt{x^2 - 1}) \text{ otherwise to avoid numerical loss.}$$

## Computing Exercises

### Section 1.1: 1,2, 21(a,b)

1. Write and run a computer program that corresponds to the pseudocode program First described in the text (p. 10) and interpret the results.

Answer:

Code (Mathlab):

```
n=30;

h=1;

emax=0;

x=0.5;

for i = 1:n
    h=0.25 .* h;
    y=(sin(x + h) - sin(x))/h;
    error = abs(cos(x) - y);
    fprintf('%d %f %f %f\n',i,h,y,err);
    if error > emax
        emax=error;
        imax=i;
    end
end

disp('imax and emax values are');
fprintf('%d %f\n',imax,emax);
```

Output:

```

1 0.250000 0.808853 0.068730
2 0.062500 0.862034 0.015548
3 0.015625 0.873801 0.003781
4 0.003906 0.876644 0.000939
5 0.000977 0.877348 0.000234
6 0.000244 0.877524 0.000059
7 0.000061 0.877568 0.000015
8 0.000015 0.877579 0.000004
9 0.000004 0.877582 0.000001
10 0.000001 0.877582 0.000000
11 0.000000 0.877583 0.000000
12 0.000000 0.877583 0.000000
13 0.000000 0.877583 0.000000
14 0.000000 0.877583 0.000000
15 0.000000 0.877583 0.000000
16 0.000000 0.877583 0.000000
17 0.000000 0.877583 0.000000
18 0.000000 0.877583 0.000000
19 0.000000 0.877579 0.000004
20 0.000000 0.877563 0.000019
21 0.000000 0.877441 0.000141
22 0.000000 0.877930 0.000347
23 0.000000 0.878906 0.001324
24 0.000000 0.875000 0.002583
25 0.000000 0.875000 0.002583
26 0.000000 0.750000 0.127583
27 0.000000 0.000000 0.877583
28 0.000000 0.000000 0.877583
29 0.000000 0.000000 0.877583
30 0.000000 0.000000 0.877583
27 0.877583

```

2. (Continuation) Select a function  $f$  and a point  $x$  and carry out a computer experiment like the one given in the text. Interpret the results. Do not select too simple a function. For example, you might consider  $1/x$ ,  $\log x$ ,  $e^x$ ,  $\tan x$ ,  $\cosh x$ , or  $x^3 - 23x$ .

**Answer:**

Code Matlab:

```
n=30;
```

```
h=1;
```

```
emax=0;
```

```
x=0.5;
```

```
for i = 1:n
```

```

h=0.25 .* h;

y=(exp(x + h) - exp(x))/h;

error = abs(exp(x) - y);

fprintf('%d %f %f %f\n',i,h,y,err);

if error > emax

emax=error;

imax=i;

end

end

disp('imax and emax values are');

fprintf('%d %f\n',imax,emax);

```

Output:

```

1 0.250000 1.873115 0.877583
2 0.062500 1.701334 0.877583
3 0.015625 1.661669 0.877583
4 0.003906 1.651946 0.877583
5 0.000977 1.649527 0.877583
6 0.000244 1.648923 0.877583
7 0.000061 1.648772 0.877583
8 0.000015 1.648734 0.877583
9 0.000004 1.648724 0.877583
10 0.000001 1.648722 0.877583
11 0.000000 1.648721 0.877583
12 0.000000 1.648721 0.877583
13 0.000000 1.648721 0.877583
14 0.000000 1.648721 0.877583
15 0.000000 1.648721 0.877583
16 0.000000 1.648721 0.877583
17 0.000000 1.648720 0.877583
18 0.000000 1.648712 0.877583
19 0.000000 1.648682 0.877583
20 0.000000 1.648682 0.877583
21 0.000000 1.648438 0.877583
22 0.000000 1.648438 0.877583
23 0.000000 1.640625 0.877583
24 0.000000 1.625000 0.877583
25 0.000000 1.500000 0.877583
26 0.000000 2.000000 0.877583
27 0.000000 0.000000 0.877583
28 0.000000 0.000000 0.877583

```



```
29 0.000000 0.000000 0.877583
30 0.000000 0.000000 0.877583
imax and emax values are
27 1.648721
```

21. Write a computer code that contains the following assignment statements exactly as shown. Analyze the results.

a. Print these values first using the default format and then with an extremely large format field:

**real** p, q, u, v, w, x, y, z

$x \leftarrow 0.1$

$y \leftarrow 0.01$

$z \leftarrow x - y$

$p \leftarrow 1.0/3.0$

$q \leftarrow 3.0p$

$u \leftarrow 7.6$

$v \leftarrow 2.9$

$w \leftarrow u - v$

**output** x, y, z, p, q, u, v, w

Answer:

Code:

Mathlab

```
format
x=0.1;
y=0.01;
z=x-y;
p=1.0/3.0;
q=3.0*p;
u=7.6;
v=2.9;
w=u-v;
disp('The given data in default format');
%Enable default format
format
fprintf('x =');disp(x)
fprintf('y =');disp(y)
fprintf('z =');disp(z)
fprintf('p =');disp(p)
fprintf('q =');disp(q)
```

```

fprintf('u =');disp(u)
fprintf('v =');disp(v)
fprintf('w =');disp(w)
disp('The given data in extreamly large format');
%Enable long format
format long
fprintf('x =');disp(x)
fprintf('y =');disp(y)
fprintf('z =');disp(z)
fprintf('p =');disp(p)
fprintf('q =');disp(q)
fprintf('u =');disp(u)
fprintf('v =');disp(v)
fprintf('w =');disp(w)

```

Output:

The given data in default format

x = 0.1000

y = 0.0100

z = 0.0900

p = 0.3333

q = 1

u = 7.6000

v = 2.9000

w = 4.7000

The given data in extreamly large format

x = 0.1000000000000000

y = 0.0100000000000000

z = 0.0900000000000000

p = 0.3333333333333333

q = 1

u = 7.6000000000000000

v = 2.9000000000000000

w = 4.6999999999999999

b. What values would be computed for x, y, and z if this code is used?

```
integer n; real x, y, z
```

```
for n = 1 to 10 do
```

```
  x ← (n - 1)/2
```

```
  y ← n2/3.0
```

```
  z ← 1.0 + 1/n
```

```
output x, y, z
```

```
end for
```

**Answer:**

Code: MathLab

```
for n = 1:10  
x = (n - 1)/2;  
y = n^2 / 3.0;  
z = 1.0 + 1/n;  
end
```

```
disp("x = " + num2str(x))
```

```
disp("y = " + num2str(y))  
disp("z = " + num2str(z))
```

Output:

```
x = 4.5  
y = 33.3333  
z = 1.1
```

### Section 2.2: 6

Write a procedure to compute  $f(x) = \sin x - 1 + \cos x$ . The routine should produce nearly full machine precision for all x in the interval  $[0, \pi/4]$ . Hint: The trigonometric identity  $\sin^2 \theta = 1/2 (1 - \cos 2\theta)$  may be useful.

**Answer:**

```
clc;  
% create vector for value of x in range 0 to pi/4 with each step of 0.001  
x = [ 0 : 0.001 : pi / 4 ];
```

```

% create vector f of same length as x
% store the value of f(x) = sin(x)-1+cos(x)
f = [1 : length(x)];

fprintf('%5s %10s\n\n', 'x', 'f(x)');

% find the value of f(x) for each x

for i = 1 : length(x)

f(i) = sin( x(i) )-1+cos( x(i) );

fprintf('%5.3f %10.3f\n', x(i), f(i));

end

% Display diagram
xlabel('x');
ylabel('f(x) = sin(x)-1+cos(x)');
hold
plot(x, f);

```

Output:

x	f(x)
0.000	0.000
0.001	0.001
0.002	0.002
0.003	0.003
0.004	0.004
0.005	0.005
0.006	0.006
0.007	0.007
0.008	0.008
0.009	0.009
0.010	0.010
0.011	0.011
0.012	0.012
0.013	0.013
0.014	0.014
0.015	0.015
0.016	0.016
0.017	0.017
0.018	0.018
0.019	0.019
0.020	0.020
0.021	0.021
0.022	0.022
0.023	0.023
0.024	0.024
0.025	0.025

0.026	0.026
0.027	0.027
0.028	0.028
0.029	0.029
0.030	0.030
0.031	0.031
0.032	0.031
0.033	0.032
0.034	0.033
0.035	0.034
0.036	0.035
0.037	0.036
0.038	0.037
0.039	0.038
0.040	0.039
0.041	0.040
0.042	0.041
0.043	0.042
0.044	0.043
0.045	0.044
0.046	0.045
0.047	0.046
0.048	0.047
0.049	0.048
0.050	0.049
0.051	0.050
0.052	0.051
0.053	0.052
0.054	0.053
0.055	0.053
0.056	0.054
0.057	0.055
0.058	0.056
0.059	0.057
0.060	0.058
0.061	0.059
0.062	0.060
0.063	0.061
0.064	0.062
0.065	0.063
0.066	0.064
0.067	0.065
0.068	0.066
0.069	0.067
0.070	0.067
0.071	0.068
0.072	0.069
0.073	0.070
0.074	0.071
0.075	0.072
0.076	0.073

0.077	0.074
0.078	0.075
0.079	0.076
0.080	0.077
0.081	0.078
0.082	0.079
0.083	0.079
0.084	0.080
0.085	0.081
0.086	0.082
0.087	0.083
0.088	0.084
0.089	0.085
0.090	0.086
0.091	0.087
0.092	0.088
0.093	0.089
0.094	0.089
0.095	0.090
0.096	0.091
0.097	0.092
0.098	0.093
0.099	0.094
0.100	0.095
0.101	0.096
0.102	0.097
0.103	0.098
0.104	0.098
0.105	0.099
0.106	0.100
0.107	0.101
0.108	0.102
0.109	0.103
0.110	0.104
0.111	0.105
0.112	0.106
0.113	0.106
0.114	0.107
0.115	0.108
0.116	0.109
0.117	0.110
0.118	0.111
0.119	0.112
0.120	0.113
0.121	0.113
0.122	0.114
0.123	0.115
0.124	0.116
0.125	0.117
0.126	0.118
0.127	0.119

0.128	0.119
0.129	0.120
0.130	0.121
0.131	0.122
0.132	0.123
0.133	0.124
0.134	0.125
0.135	0.125
0.136	0.126
0.137	0.127
0.138	0.128
0.139	0.129
0.140	0.130
0.141	0.131
0.142	0.131
0.143	0.132
0.144	0.133
0.145	0.134
0.146	0.135
0.147	0.136
0.148	0.137
0.149	0.137
0.150	0.138
0.151	0.139
0.152	0.140
0.153	0.141
0.154	0.142
0.155	0.142
0.156	0.143
0.157	0.144
0.158	0.145
0.159	0.146
0.160	0.147
0.161	0.147
0.162	0.148
0.163	0.149
0.164	0.150
0.165	0.151
0.166	0.151
0.167	0.152
0.168	0.153
0.169	0.154
0.170	0.155
0.171	0.156
0.172	0.156
0.173	0.157
0.174	0.158
0.175	0.159
0.176	0.160
0.177	0.160
0.178	0.161

0.179	0.162
0.180	0.163
0.181	0.164
0.182	0.164
0.183	0.165
0.184	0.166
0.185	0.167
0.186	0.168
0.187	0.168
0.188	0.169
0.189	0.170
0.190	0.171
0.191	0.172
0.192	0.172
0.193	0.173
0.194	0.174
0.195	0.175
0.196	0.176
0.197	0.176
0.198	0.177
0.199	0.178
0.200	0.179
0.201	0.180
0.202	0.180
0.203	0.181
0.204	0.182
0.205	0.183
0.206	0.183
0.207	0.184
0.208	0.185
0.209	0.186
0.210	0.186
0.211	0.187
0.212	0.188
0.213	0.189
0.214	0.190
0.215	0.190
0.216	0.191
0.217	0.192
0.218	0.193
0.219	0.193
0.220	0.194
0.221	0.195
0.222	0.196
0.223	0.196
0.224	0.197
0.225	0.198
0.226	0.199
0.227	0.199
0.228	0.200
0.229	0.201



0.230	0.202
0.231	0.202
0.232	0.203
0.233	0.204
0.234	0.205
0.235	0.205
0.236	0.206
0.237	0.207
0.238	0.208
0.239	0.208
0.240	0.209
0.241	0.210
0.242	0.211
0.243	0.211
0.244	0.212
0.245	0.213
0.246	0.213
0.247	0.214
0.248	0.215
0.249	0.216
0.250	0.216
0.251	0.217
0.252	0.218
0.253	0.218
0.254	0.219
0.255	0.220
0.256	0.221
0.257	0.221
0.258	0.222
0.259	0.223
0.260	0.223
0.261	0.224
0.262	0.225
0.263	0.226
0.264	0.226
0.265	0.227
0.266	0.228
0.267	0.228
0.268	0.229
0.269	0.230
0.270	0.231
0.271	0.231
0.272	0.232
0.273	0.233
0.274	0.233
0.275	0.234
0.276	0.235
0.277	0.235
0.278	0.236
0.279	0.237
0.280	0.237

0.281	0.238
0.282	0.239
0.283	0.239
0.284	0.240
0.285	0.241
0.286	0.241
0.287	0.242
0.288	0.243
0.289	0.244
0.290	0.244
0.291	0.245
0.292	0.246
0.293	0.246
0.294	0.247
0.295	0.248
0.296	0.248
0.297	0.249
0.298	0.250
0.299	0.250
0.300	0.251
0.301	0.252
0.302	0.252
0.303	0.253
0.304	0.253
0.305	0.254
0.306	0.255
0.307	0.255
0.308	0.256
0.309	0.257
0.310	0.257
0.311	0.258
0.312	0.259
0.313	0.259
0.314	0.260
0.315	0.261
0.316	0.261
0.317	0.262
0.318	0.263
0.319	0.263
0.320	0.264
0.321	0.264
0.322	0.265
0.323	0.266
0.324	0.266
0.325	0.267
0.326	0.268
0.327	0.268
0.328	0.269
0.329	0.269
0.330	0.270
0.331	0.271

0.332	0.271
0.333	0.272
0.334	0.273
0.335	0.273
0.336	0.274
0.337	0.274
0.338	0.275
0.339	0.276
0.340	0.276
0.341	0.277
0.342	0.277
0.343	0.278
0.344	0.279
0.345	0.279
0.346	0.280
0.347	0.280
0.348	0.281
0.349	0.282
0.350	0.282
0.351	0.283
0.352	0.283
0.353	0.284
0.354	0.285
0.355	0.285
0.356	0.286
0.357	0.286
0.358	0.287
0.359	0.288
0.360	0.288
0.361	0.289
0.362	0.289
0.363	0.290
0.364	0.290
0.365	0.291
0.366	0.292
0.367	0.292
0.368	0.293
0.369	0.293
0.370	0.294
0.371	0.295
0.372	0.295
0.373	0.296
0.374	0.296
0.375	0.297
0.376	0.297
0.377	0.298
0.378	0.298
0.379	0.299
0.380	0.300
0.381	0.300
0.382	0.301

0.383	0.301
0.384	0.302
0.385	0.302
0.386	0.303
0.387	0.303
0.388	0.304
0.389	0.305
0.390	0.305
0.391	0.306
0.392	0.306
0.393	0.307
0.394	0.307
0.395	0.308
0.396	0.308
0.397	0.309
0.398	0.309
0.399	0.310
0.400	0.310
0.401	0.311
0.402	0.312
0.403	0.312
0.404	0.313
0.405	0.313
0.406	0.314
0.407	0.314
0.408	0.315
0.409	0.315
0.410	0.316
0.411	0.316
0.412	0.317
0.413	0.317
0.414	0.318
0.415	0.318
0.416	0.319
0.417	0.319
0.418	0.320
0.419	0.320
0.420	0.321
0.421	0.321
0.422	0.322
0.423	0.322
0.424	0.323
0.425	0.323
0.426	0.324
0.427	0.324
0.428	0.325
0.429	0.325
0.430	0.326
0.431	0.326
0.432	0.327
0.433	0.327

0.434	0.328
0.435	0.328
0.436	0.329
0.437	0.329
0.438	0.330
0.439	0.330
0.440	0.331
0.441	0.331
0.442	0.332
0.443	0.332
0.444	0.333
0.445	0.333
0.446	0.334
0.447	0.334
0.448	0.334
0.449	0.335
0.450	0.335
0.451	0.336
0.452	0.336
0.453	0.337
0.454	0.337
0.455	0.338
0.456	0.338
0.457	0.339
0.458	0.339
0.459	0.340
0.460	0.340
0.461	0.340
0.462	0.341
0.463	0.341
0.464	0.342
0.465	0.342
0.466	0.343
0.467	0.343
0.468	0.344
0.469	0.344
0.470	0.344
0.471	0.345
0.472	0.345
0.473	0.346
0.474	0.346
0.475	0.347
0.476	0.347
0.477	0.347
0.478	0.348
0.479	0.348
0.480	0.349
0.481	0.349
0.482	0.350
0.483	0.350
0.484	0.350

0.485	0.351
0.486	0.351
0.487	0.352
0.488	0.352
0.489	0.353
0.490	0.353
0.491	0.353
0.492	0.354
0.493	0.354
0.494	0.355
0.495	0.355
0.496	0.355
0.497	0.356
0.498	0.356
0.499	0.357
0.500	0.357
0.501	0.357
0.502	0.358
0.503	0.358
0.504	0.359
0.505	0.359
0.506	0.359
0.507	0.360
0.508	0.360
0.509	0.361
0.510	0.361
0.511	0.361
0.512	0.362
0.513	0.362
0.514	0.362
0.515	0.363
0.516	0.363
0.517	0.364
0.518	0.364
0.519	0.364
0.520	0.365
0.521	0.365
0.522	0.365
0.523	0.366
0.524	0.366
0.525	0.367
0.526	0.367
0.527	0.367
0.528	0.368
0.529	0.368
0.530	0.368
0.531	0.369
0.532	0.369
0.533	0.369
0.534	0.370
0.535	0.370

0.536	0.370
0.537	0.371
0.538	0.371
0.539	0.372
0.540	0.372
0.541	0.372
0.542	0.373
0.543	0.373
0.544	0.373
0.545	0.374
0.546	0.374
0.547	0.374
0.548	0.375
0.549	0.375
0.550	0.375
0.551	0.376
0.552	0.376
0.553	0.376
0.554	0.377
0.555	0.377
0.556	0.377
0.557	0.377
0.558	0.378
0.559	0.378
0.560	0.378
0.561	0.379
0.562	0.379
0.563	0.379
0.564	0.380
0.565	0.380
0.566	0.380
0.567	0.381
0.568	0.381
0.569	0.381
0.570	0.382
0.571	0.382
0.572	0.382
0.573	0.382
0.574	0.383
0.575	0.383
0.576	0.383
0.577	0.384
0.578	0.384
0.579	0.384
0.580	0.384
0.581	0.385
0.582	0.385
0.583	0.385
0.584	0.386
0.585	0.386
0.586	0.386

0.587	0.386
0.588	0.387
0.589	0.387
0.590	0.387
0.591	0.388
0.592	0.388
0.593	0.388
0.594	0.388
0.595	0.389
0.596	0.389
0.597	0.389
0.598	0.389
0.599	0.390
0.600	0.390
0.601	0.390
0.602	0.390
0.603	0.391
0.604	0.391
0.605	0.391
0.606	0.392
0.607	0.392
0.608	0.392
0.609	0.392
0.610	0.393
0.611	0.393
0.612	0.393
0.613	0.393
0.614	0.393
0.615	0.394
0.616	0.394
0.617	0.394
0.618	0.394
0.619	0.395
0.620	0.395
0.621	0.395
0.622	0.395
0.623	0.396
0.624	0.396
0.625	0.396
0.626	0.396
0.627	0.397
0.628	0.397
0.629	0.397
0.630	0.397
0.631	0.397
0.632	0.398
0.633	0.398
0.634	0.398
0.635	0.398
0.636	0.398
0.637	0.399



0.638	0.399
0.639	0.399
0.640	0.399
0.641	0.399
0.642	0.400
0.643	0.400
0.644	0.400
0.645	0.400
0.646	0.400
0.647	0.401
0.648	0.401
0.649	0.401
0.650	0.401
0.651	0.401
0.652	0.402
0.653	0.402
0.654	0.402
0.655	0.402
0.656	0.402
0.657	0.403
0.658	0.403
0.659	0.403
0.660	0.403
0.661	0.403
0.662	0.403
0.663	0.404
0.664	0.404
0.665	0.404
0.666	0.404
0.667	0.404
0.668	0.404
0.669	0.405
0.670	0.405
0.671	0.405
0.672	0.405
0.673	0.405
0.674	0.405
0.675	0.406
0.676	0.406
0.677	0.406
0.678	0.406
0.679	0.406
0.680	0.406
0.681	0.407
0.682	0.407
0.683	0.407
0.684	0.407
0.685	0.407
0.686	0.407
0.687	0.407
0.688	0.408

0.689	0.408
0.690	0.408
0.691	0.408
0.692	0.408
0.693	0.408
0.694	0.408
0.695	0.408
0.696	0.409
0.697	0.409
0.698	0.409
0.699	0.409
0.700	0.409
0.701	0.409
0.702	0.409
0.703	0.409
0.704	0.410
0.705	0.410
0.706	0.410
0.707	0.410
0.708	0.410
0.709	0.410
0.710	0.410
0.711	0.410
0.712	0.410
0.713	0.411
0.714	0.411
0.715	0.411
0.716	0.411
0.717	0.411
0.718	0.411
0.719	0.411
0.720	0.411
0.721	0.411
0.722	0.411
0.723	0.411
0.724	0.412
0.725	0.412
0.726	0.412
0.727	0.412
0.728	0.412
0.729	0.412
0.730	0.412
0.731	0.412
0.732	0.412
0.733	0.412
0.734	0.412
0.735	0.412
0.736	0.412
0.737	0.413
0.738	0.413
0.739	0.413

0.740	0.413
0.741	0.413
0.742	0.413
0.743	0.413
0.744	0.413
0.745	0.413
0.746	0.413
0.747	0.413
0.748	0.413
0.749	0.413
0.750	0.413
0.751	0.413
0.752	0.413
0.753	0.413
0.754	0.414
0.755	0.414
0.756	0.414
0.757	0.414
0.758	0.414
0.759	0.414
0.760	0.414
0.761	0.414
0.762	0.414
0.763	0.414
0.764	0.414
0.765	0.414
0.766	0.414
0.767	0.414
0.768	0.414
0.769	0.414
0.770	0.414
0.771	0.414
0.772	0.414
0.773	0.414
0.774	0.414
0.775	0.414
0.776	0.414
0.777	0.414
0.778	0.414
0.779	0.414
0.780	0.414
0.781	0.414
0.782	0.414
0.783	0.414
0.784	0.414
0.785	0.414

Current plot released

Figure 1 × +

