

# Midterm topics

- Computer arithmetic
- Root finding
- Interpolation
- Integration

## Not covered topics

- Splines
- Adaptive Simpson
- Quadrature

## Midterm exam

- 5 major problems with subproblems
- Some conceptual questions (T/F and multiple choice)
- No pseudo code

# Computer Arithmetic

- Absolute error v.s. Relative error
- Taylor series and Taylor Theorem
- Deflating  $p(x) = (x - r)q(x) + p(r)$  (synthetic division)
- Floating point representation
- Machine number
- Loss of significance

# Float-point rep.

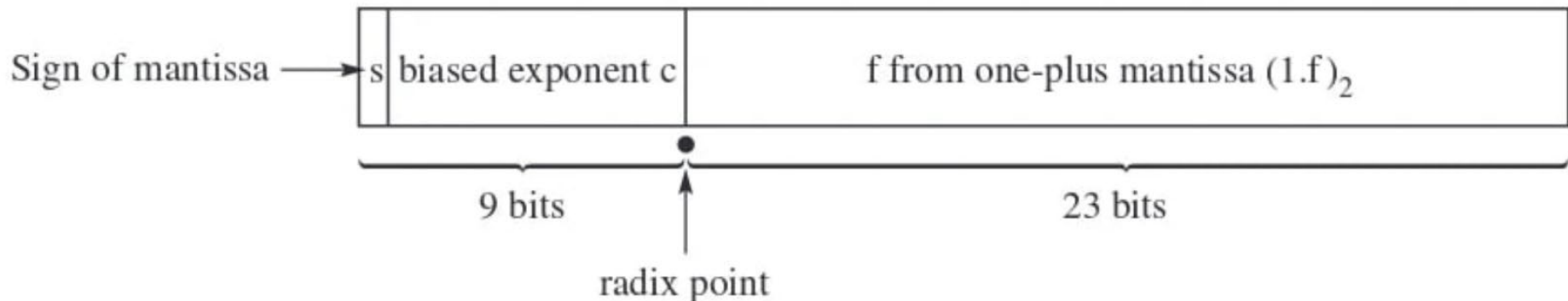
- If  $x \neq 0$ , it can be written as

$$x = \pm q \times 2^m \left( \frac{1}{2} \leq q < 1 \right)$$

- The **mantissa** would be expressed a sequence of binary values (0 or 1)

$$q = (0.b_1b_2b_3 \cdots)_2$$

- $b_1 \neq 0 \rightarrow b_1 = 1 \rightarrow q \geq \frac{1}{2}$ .



# Loss of significance

- It occurs in the subtraction of two nearly equal numbers, which produces a result much smaller than either one.
- Theorem

Let  $x$  and  $y$  be normalized floating-point machine numbers, where  $x > y > 0$ . If  $2^{-p} \leq 1 - (y/x) \leq 2^{-q}$  for some positive integers  $p$  and  $q$ , then at most  $p$  and at least  $q$  significant binary bits are lost in the subtraction  $x - y$ .

- How to avoid?
  - Double precision
  - Taylor series
  - Rationalization
  - Trigonometric identities
  - Logarithmic properties
  - Range reduction

# Root finding

Bisection

Newton's

Secant

# Algorithms

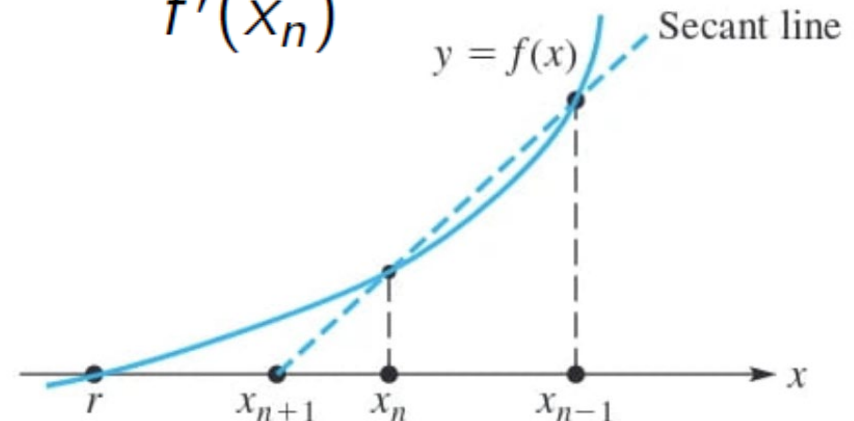
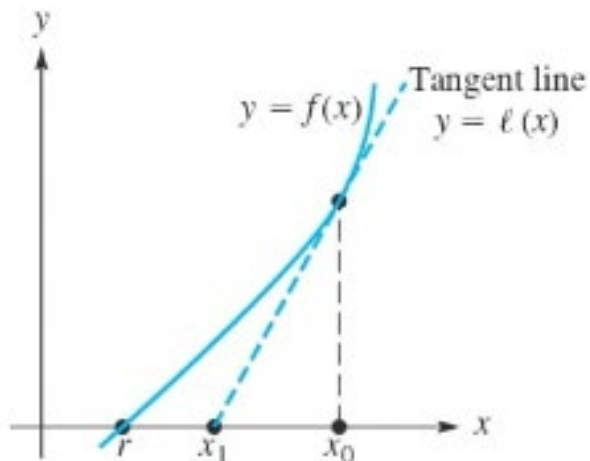
- Bisection method

- Secant method

$$x_{n+1} = x_n - \left( \frac{x_n - x_{n-1}}{f(x_n) - f(x_{n-1})} \right) f(x_n)$$

- Newton's method

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$$



# Convergence analysis

If the bisection algorithm is applied to a continuous function  $f$  on an interval  $[a, b]$ , where  $f(a)f(b) < 0$ , then, after  $n$  steps, an approximate root will have been computed with error at most  $(b - a)/2^{n+1}$ .

## Newton's Method Theorem

If  $f$ ,  $f'$ , and  $f''$  are continuous in a neighborhood of a root  $r$  of  $f$  and if  $f'(r) \neq 0$ , then there is a positive  $\delta$  with the following property: If the initial point in Newton's method satisfies  $|r - x_0| \leq \delta$ , then all subsequent points  $x_n$  satisfy the same inequality, converge to  $r$ , and do so quadratically; that is,

$$|r - x_{n+1}| \leq c(\delta)|r - x_n|^2$$

Fixed-point iteration

$$x_{n+1} = g(x_n)$$

Locally convergent if  $x^* = g(x^*)$ ,  $|g'| < 1$ .



# Convergence rate

- **Linear** convergence:  $C \in [0,1)$

$$|x_{n+1} - x^*| \leq C|x_n - x^*|$$

Bisection is  
not linear

- **Superlinear** convergence:  $\alpha \in (1,2)$

$$|x_{n+1} - x^*| \leq C|x_n - x^*|^\alpha$$

Secant

- **Quadratic** convergence

$$|x_{n+1} - x^*| \leq C|x_n - x^*|^2$$

Newton's

# Interpolation

Vandermonde matrix

Lagrange form

Divided difference

# Vandermonde Matrix

$x$	$x_0$	$x_1$	$\cdots$	$x_n$
$y$	$y_0$	$y_1$	$\cdots$	$y_n$

$$p_n(x) = c_0 + c_1x + c_2x^2 + \cdots + c_nx^n$$

$$\begin{bmatrix} 1 & x_0 & x_0^2 & \cdots & x_0^n \\ 1 & x_1 & x_1^2 & \cdots & x_1^n \\ 1 & x_2 & x_2^2 & \cdots & x_2^n \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & x_n & x_n^2 & \cdots & x_n^n \end{bmatrix} \begin{bmatrix} c_0 \\ c_1 \\ c_2 \\ \vdots \\ c_n \end{bmatrix} = \begin{bmatrix} y_0 \\ y_1 \\ y_2 \\ \vdots \\ y_n \end{bmatrix}$$

Can be used  
to prove the  
existence and  
uniqueness!

Vandermond matrix

# Lagrange Form

The Lagrange form of the interpolation polynomial is given by

$$p_n(x) = \sum_{i=0}^n \ell_i(x) f(x_i)$$

$$\ell_i(x) = \left( \frac{x - x_0}{x_i - x_0} \right) \left( \frac{x - x_1}{x_i - x_1} \right) \cdots \left( \frac{x - x_{i-1}}{x_i - x_{i-1}} \right) \left( \frac{x - x_{i+1}}{x_i - x_{i+1}} \right) \cdots \left( \frac{x - x_n}{x_i - x_n} \right)$$

# Example

$x$	$\frac{1}{3}$	$\frac{1}{4}$	$1$
$f(x)$	$2$	$-1$	$7$

$$\ell_0(x) = \frac{(x - \frac{1}{4})(x - 1)}{(\frac{1}{3} - \frac{1}{4})(\frac{1}{3} - 1)} = -18\left(x - \frac{1}{4}\right)(x - 1)$$

$$\ell_1(x) = \frac{(x - \frac{1}{3})(x - 1)}{(\frac{1}{4} - \frac{1}{3})(\frac{1}{4} - 1)} = 16\left(x - \frac{1}{3}\right)(x - 1)$$

$$\ell_2(x) = \frac{(x - \frac{1}{3})(x - \frac{1}{4})}{(1 - \frac{1}{3})(1 - \frac{1}{4})} = 2\left(x - \frac{1}{3}\right)\left(x - \frac{1}{4}\right)$$

$$p_2(x) = -36\left(x - \frac{1}{4}\right)(x - 1) - 16\left(x - \frac{1}{3}\right)(x - 1) + 14\left(x - \frac{1}{3}\right)\left(x - \frac{1}{4}\right)$$

# Divided Difference

- Concise notation

$$p_n(x) = \sum_{i=0}^n a_i \prod_{j=0}^{i-1} (x - x_j)$$

with  $\prod_{j=0}^{-1} (x - x_j) = 1$

- Nested form

$$p(x) = a_0 + (x - x_0)(a_1 + (x - x_1)(a_2 + \cdots + (x - x_{n-1})a_n)) \cdots)$$

$x$	$f[ \ ]$	$f[ \ , \ ]$	$f[ \ , \ , \ ]$	$f[ \ , \ , \ , \ ]$
$x_0$	$f[x_0]$			
$x_1$	$f[x_1]$	$f[x_0, x_1]$	$f[x_0, x_1, x_2]$	
$x_2$	$f[x_2]$	$f[x_1, x_2]$	$f[x_1, x_2, x_3]$	$f[x_0, x_1, x_2, x_3]$
$x_3$	$f[x_3]$	$f[x_2, x_3]$		

# Example

Let  $f(x) = x^3 + 2x^2 + x + 1$ .

- a) Find the polynomial of degree 4 that interpolates the values of  $f$  at  $\pm 2, \pm 1, 0$ .
- b) Find the polynomial of degree 2 that interpolates the values of  $f$  at  $\pm 1, 0$ .

$$p_4(x) = -1 + 2(x + 2) - (x + 2)(x + 1) + (x + 2)(x + 1)x$$

$$p_2(x) = 1 + 2(x + 1)x$$

# Interpolation errors

## First Interpolation Error Theorem

If  $p$  is the polynomial of degree at most  $n$  that interpolates  $f$  at the  $n + 1$  distinct nodes  $x_0, x_1, \dots, x_n$  belonging to an interval  $[a, b]$  and if  $f^{(n+1)}$  is continuous, then for each  $x$  in  $[a, b]$ , there is a  $\xi$  in  $(a, b)$  for which

$$f(x) - p(x) = \frac{1}{(n+1)!} f^{(n+1)}(\xi) \prod_{i=0}^n (x - x_i) \quad (2)$$

## Second Interpolation Error Theorem

Let  $f$  be a function such that  $f^{(n+1)}$  is continuous on  $[a, b]$  and satisfies  $|f^{(n+1)}(x)| \leq M$ . Let  $p$  be the polynomial of degree  $\leq n$  that interpolates  $f$  at  $n + 1$  equally spaced nodes in  $[a, b]$ , including the endpoints. Then on  $[a, b]$ ,

$$|f(x) - p(x)| \leq \frac{1}{4(n+1)} M h^{n+1} \quad (6)$$

where  $h = (b - a)/n$  is the spacing between nodes.



# Runge's phenomenon

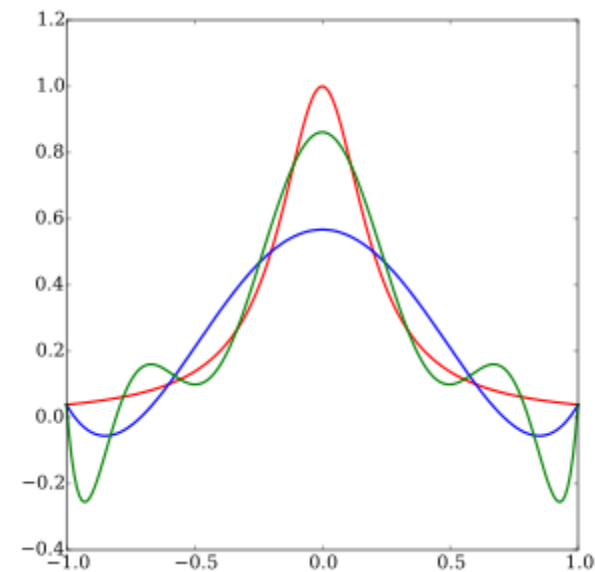
- A polynomial of degree  $n$  has  $n$  roots.

If all roots are real, the curve crosses the  $x$ -axis  $n$  times.

- These many turns result in **wide oscillations**.
- A specific example is provided by **Runge function**

$$f(x) = (1 + 25x^2)^{-1}$$

- Replace the **equispaced** nodes by **Chebyshev** nodes



$$\lim_{n \rightarrow \infty} \left( \max_{-1 \leq x \leq 1} |f(x) - P_n(x)| \right) = \infty.$$

# Integration

Trapezoid

Simpson's

- Trapezoid

$$\int_a^b f(x) dx \approx \frac{1}{2}(b-a)[f(a) + f(b)]$$

- Simpson's 1/3 Rule

$$\int_a^b f(x) dx \approx \frac{1}{6}(b-a) \left[ f(a) + 4f\left(\frac{a+b}{2}\right) + f(b) \right]$$

- No need to memorize 3/8 rule

## Theorem on Precision of Composite Trapezoid Rule

If  $f''$  exists and is continuous on the interval  $[a, b]$  and if the composite trapezoid rule  $T$  with uniform spacing  $h$  is used to estimate the integral  $I = \int_a^b f(x) dx$ , then for some  $\zeta$  in  $(a, b)$ ,

$$I - T = -\frac{1}{12}(b-a)h^2 f''(\zeta) = \mathcal{O}(h^2)$$



Simpson's 1/3 Rule:  $-\frac{1}{180}(b-a)h^4 f^{(4)}(\xi)$

Simpson's 3/8 rule has the same order but lower in accuracy.