__/__/____

Assignment -I

PES2UG20CS209
Naman Choudhary

1. Plaintext = FIREWALL
Key = OCCURENCES

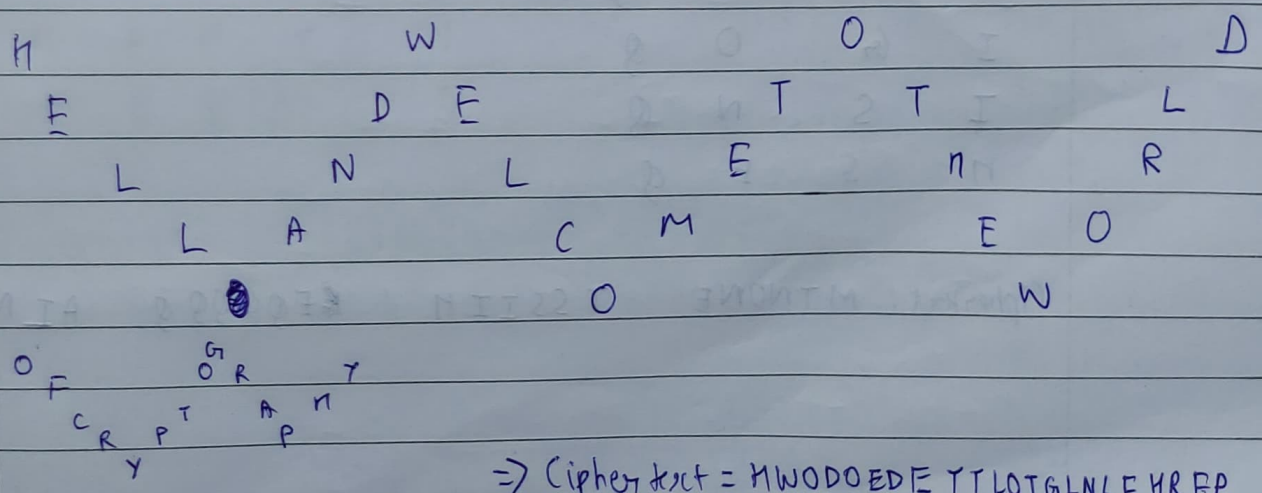
Playfair square:

Encypted Message = GKEOXSPVPV

O	C	U	R	E
N	S	A	B	D
F	G	H	I/J	K
L	M	P	Q	T
V	W	X	Y	Z

2. Plaintext: WELCOME TO PES UNIVERSITY
key: 7
Ciphertext: DLSJVTL AV WLZ BUPCLYZPAF

3. HELLO AND WELCOME TO THE WORLD OF CRYPTOGRAPHY



=> Ciphertext = HWODOEDE TTLOTGLNLE HR FP
RYLACMEOCYAHOOWRP

__/__/____

4. Key 1: BRIDGE

Key 2: OVER

Plaintext: THIS IS ASSIGNMENT ONE

Cipher:

①	⑥	⑤	②	④	③	
B	R	I	D	G	E	key 1
T	H	I	S	I	S	
A	S	S	I	G	N	
M	E	N	T	O	N	
E	Q	Q	Q	Q	Q	

Ciphertext: TAME [illegible]SITQ [illegible]SNNQ IGOQ ISNQ HSEQ

②	④	①	③	
O	V	E	R	key 2
T	A	M	E	
S	I	T	Q	
[illegible]	[illegible]	[illegible]		
S	N	N	Q	
I	G	O	Q	
I	S	N	Q	
H	S	E	Q	

Ciphertext: MTNONE TSSIIH [illegible]EQQQQQ AINGSS

__/__/____

Decipher:

Key: OVER

②	④	①	③	
O	V	E	R	Key 2
T	A	M	E	
S	I	T	Q	
S	N	N	Q	
I	G	O	Q	
I	S	N	Q	
M	S	E	Q	

⇒ Plaintext: TAME SITQ SNNQ IGOQ ISNQ MSEQ

Key: BRIDGE

①	⑥	⑤	②	④	③
B	R	I	D	G	E
T	H	I	S	I	S
A	S	S	I	G	N
M	E	N	T	O	N
E	Q	Q	Q	Q	Q

⇒ Plaintext: THISISASSIGNMENTONEQQQQQ

__/__/____

5. Known Plaintext	Chosen cipher text
• Plaintext cannot be selected but can observe plain-text — ciphertext pairs	• Plaintext can be selected, and encrypted to observe ciphertext & reverse the entire process
• Comparitively harder	• Comparitively easier

6. $Key = \begin{bmatrix} 6 & 24 & 1 \\ 13 & 16 & 10 \\ 2 & 12 & 15 \end{bmatrix}$, Plaintext = APT $= \begin{bmatrix} 0 \\ 15 \\ 19 \end{bmatrix}$

$E = PK \bmod 26$

$$(0 \quad 15 \quad 19) \begin{bmatrix} 6 & 24 & 1 \\ 13 & 16 & 10 \\ 2 & 12 & 15 \end{bmatrix} = \begin{bmatrix} 103 & 308 & 335 \end{bmatrix}$$

%26

$= [25 \quad 22 \quad 23] = [Z \quad W \quad X]$

Decryption: $C = [Z \quad W \quad X]$

$PK^{-1} \cdot K^{-1} \; C \bmod 26$

$|k| = 6(240 - 120) - 24(195 - 20) + (156 - 32) = (3356)$

$AB = 1 \bmod 26$

$(-3356)A = 1 \bmod 26$

$\Rightarrow$ A cannot be found

7. a) gcd (6150, 704)

Euclidean algorithm. [gcd(a,b) = gcd(b,r)]

a	b	r
6150	704	518
704	518	186
518	186	146
186	146	40
146	40	26
40	26	14
26	14	12
14	12	2
12	(2)	0

$\Rightarrow \gcd(6150, 704) = 2$

Stein's algorithm.

$A_1 = 6150,\ B_1 = 704,\ C_1 = 1$

(both even) $\Rightarrow A_2 = 3075,\ B_2 = 352,\ C_2 = 2$

(Odd/Even) $\Rightarrow A_3 = 3075,\ B_3 = 176,\ C_3 = 2$

$A_4 = 3075,\ B_4 = 88,\ C_4 = 2$

$A_5 = 3075,\ B_5 = 44,\ C_5 = 2$

$A_6 = 3075,\ B_6 = 22,\ C_6 = 2$

$A_7 = 3075,\ B_7 = 11,\ C_7 = 2$

(Odd/Odd) $\Rightarrow A_8 = 3064,\ B_8 = 11,\ C_8 = 2$

(Even/Odd) $\Rightarrow A_9 = 1532,\ B_9 = 11,\ C_9 = 2$

$A_{10} = 766,\ B_{10} = 11,\ C_{10} = 2$

$A_{11} = 383,\ B_{11} = 11,\ C_{11} = 2$

(Odd/Odd) $\Rightarrow A_{12} = 372,\ B_{12} = 11,\ C_{12} = 2$

(Even/Odd) $\Rightarrow A_{13} = 186,\ B_{13} = 11,\ C_{13} = 2$

$A_{14} = 93$, $B_{14} = 11$, $C_{14} = 2$

$A_{15} = 82$, $B_{15} = 11$, $C_{15} = 2$

$A_{16} = 41$, $B_{16} = 11$, $C_{16} = 2$

$A_{17} = 30$, $B_{17} = 11$, $C_{17} = 2$

$A_{18} = 15$, $B_{18} = 11$, $C_{18} = 2$

$A_{19} = 4$, $B_{19} = 11$, $C_{19} = 2$

$A_{20} = 2$, $B_{20} = 11$, $C_{20} = 2$

$A_{21} = 1$, $B_{21} = 11$, $C_{21} = 2$

$A_{22} = 10$, $B_{22} = 1$, $C_{22} = 2$

$A_{23} = 5$, $B_{23} = 1$, $C_{23} = 2$

$A_{24} = 4$, $B_{24} = 1$, $C_{24} = 2$

$A_{25} = 2$, $B_{25} = 1$, $C_{25} = 2$

$A_{26} = 1$, $B_{26} = 1$, $C_{26} = 2$

$A_{26} = B_{26}$

$\therefore, \gcd(A,B) = A_{26} \cdot C_{26} = 1 \times 2$

$\therefore, \gcd(6150, 704) = 2$

b) Euclidean Algorithm uses repeated modulo operator, while in Stein's algorithm, repeated bitwise shifts are used, which is faster, implies they have slightly better efficiency

__/__/____

8.

a) $4x \equiv 2 \pmod 3$

$\Rightarrow 4(2) \bmod 3 = 2$

$\Rightarrow 8 \bmod 3 = 2$

$\therefore$, $\boxed{x = 2}$

b) $7x \equiv 4 \pmod 9$

$\Rightarrow 7(4) \bmod 9$

$\Rightarrow 7(x) \bmod 9 = 4$

$\Rightarrow 7(7) \bmod 9 = 4 \quad \Rightarrow 49 \bmod 9 = 4$

$\Rightarrow \boxed{x = 7}$

c) $5x \equiv 3 \pmod{11}$

$\Rightarrow 5(5) \bmod 11 = 3 \Rightarrow 25 \bmod 11 = 3$

$\Rightarrow \boxed{x = 5}$