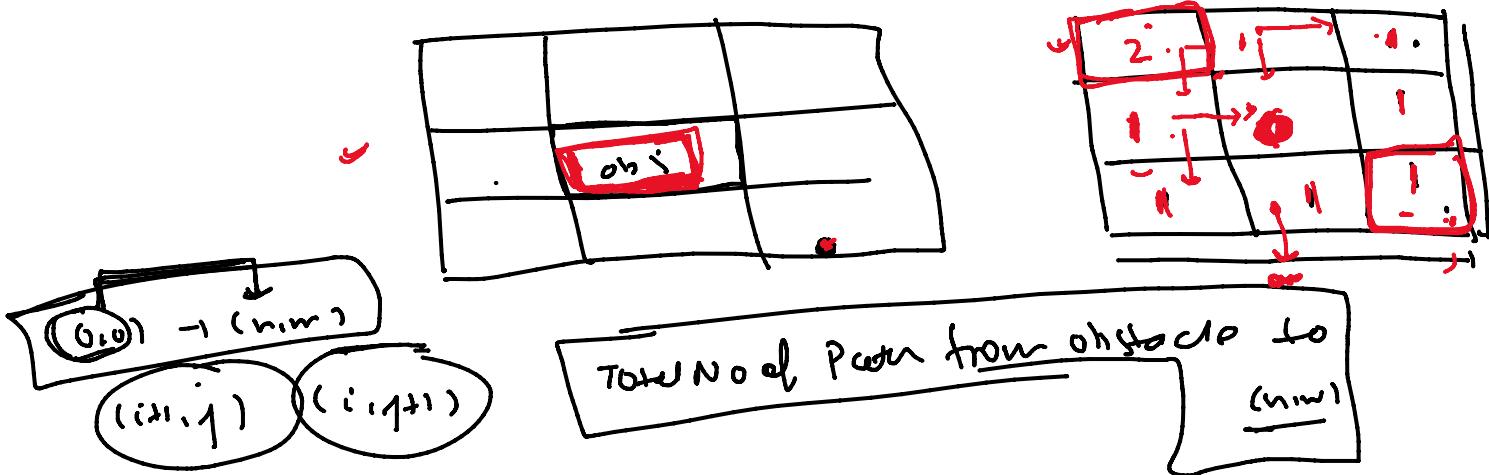


Class 87Dynamic ProgrammingUnique Path 2

vector<vector<int>> dp(n, vector<int>(m, 0));

$dp[n-1][m-1] = 1;$

$((0-0) \rightarrow (n,m))$

$(n,m) \rightarrow (n,m) = M$

for (i=n-1; i>=0; i--)

{ for (j=m-1; j>=0; j--)

 if (i == n-1 and j == m-1) continue;

 if (grid[i][j] == 1)

 if (i+1 < n)

 dp[i][j] += dp[i+1][j];

 }

 if (j+1 < m)

$i \leq j < m$

$dp[i][j] += dp[i][j+1];$

return $dp[0][0]$;

		0bs	0bs
	-	obs	-
(0,0)	(0,1)		
1	2	2	0
2	0	.	0
2	2	2	1
0	0	1	1

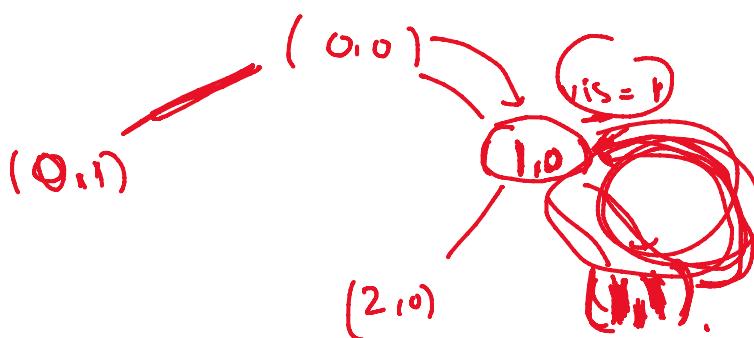
vis Array

DP

Backtracking

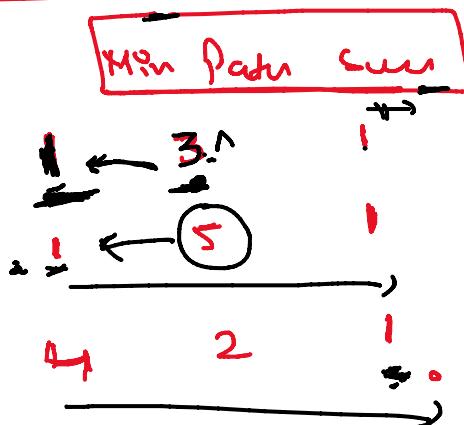
Backtracking DP?

No cycles



$dp[0][1]$
 $dp[1][0]$
 $dp[1][2]$
 $dp[2][1]$

$(U, v) \rightarrow (C, v)$



$(U, v) \rightarrow (C, v)$ (min)

Recursion

vector<vector<int>> dp(n, vector<int>(n, INT_MAX));

For ($i = 0$; $i < n$; $i++$)

{ For ($j = 0$; $j < n$; $j++$)

{ if ($i == 0$ and $j == 0$)

 dp[0][0] = grid[0][0];

else if ($i == 0$)

{ dp[0][j] = dp[0][j-1] + grid[0][j]

 }

else if ($j == 0$)

 dp[i][0] = dp[i-1][0] + grid[i][0];

else

 dp[i][j] = min(dp[i-1][j], dp[i][j-1]) + grid[i][j];

y

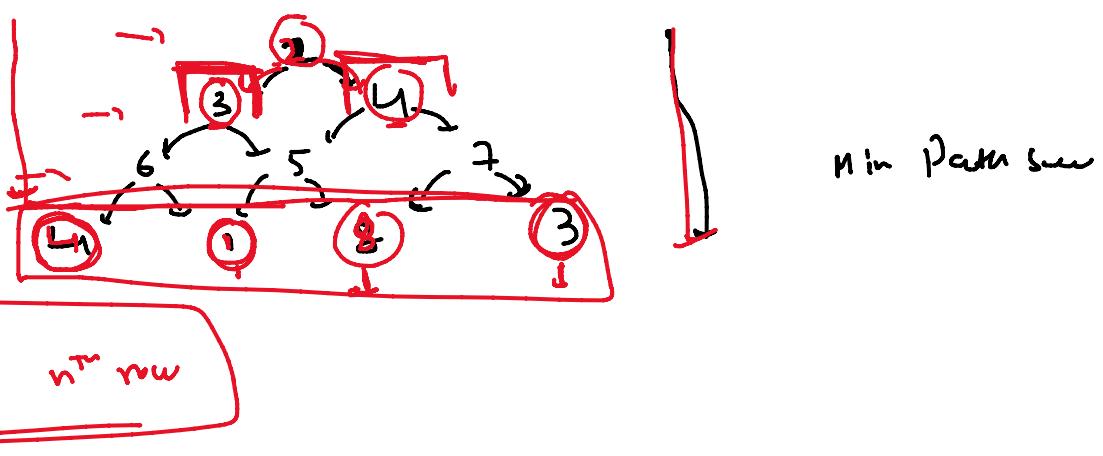
return dp[n-1][n-1];

1 2 3 4 5

1 3 1 -1
 1 5 1
 4 2 1

1	4	5
2	7	6
6	8	7

$dp[i][j] = \text{Min Path sum to travel from } (i, j) \text{ to } (0, 0)$



vector<vector<int>> dp(n+1);

for (i=0; i<n; i++)
 {
 & dp[i]. resize (n+1);
 }

for (i=0; i<n; i++)
 {
 & dp[n-1][i] = triangle[n-1][i];
 }

for (i=n-2; i>=0; i--)
 {
 & for (j=0; j<=i; j++)

2 For ($j=0$; $j < i$; $j+1$)

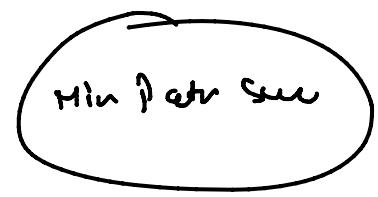
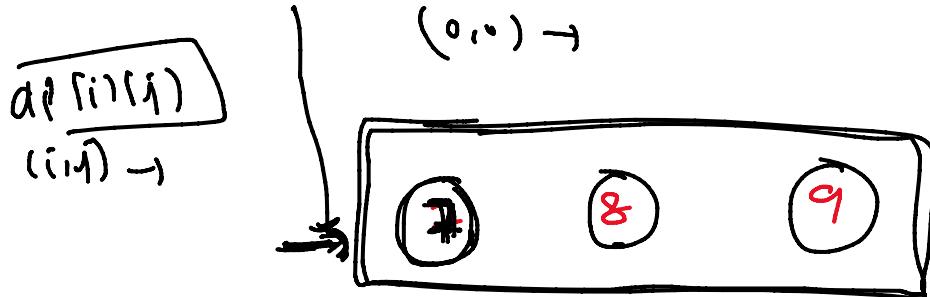
$$\leftarrow dp[i][j] = \min(dp[i+1][j], dp[i+1][j+1]) + \text{magn}[i]$$

;

;

return $dp[0][0]$;

Min Path Falling Sum



(7)

vector <vector<int>> dp (n, vector<int>(n, INT_MAX));

For ($i=0$; $i < n$; $i++$)

$$dp[n-1][i] = \text{and}[n-1][i]$$

for ($i=n-2$; $i >= 0$; $i--$)

 For ($j=0$; $j < n$; $j+1$)

$$\leftarrow dp[i][j] = \boxed{\text{and}[i][j] + dp[i+1][j]}$$

if ($j > 0$)

left diagonally

$$dp[i][j] = \min(\underbrace{dp[i][j-1]}, \boxed{dp[i+1][j-1]} + \underbrace{\text{and}[i][j-1]}).$$

$$dp[i][j] = \min(\underbrace{dp[i][j-1]}_{\text{and } f(i)[j]}, \overbrace{dp[i+1][j-1]) + f(i+1)[j]}^{\text{and } f(i)[j]});$$

$i \leq j < n-1$

$$dp[i][j] = \min(dp[i][j], dp[i+1][j] + \underbrace{f(i+1)[j]}_{\text{and } f(i)[j]}))$$



γ

```
for(i=0;i<n;i++)
    ans=min(ans,dp[0][i]);
return ans;
```