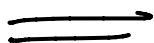


CLASS-31LPS Array

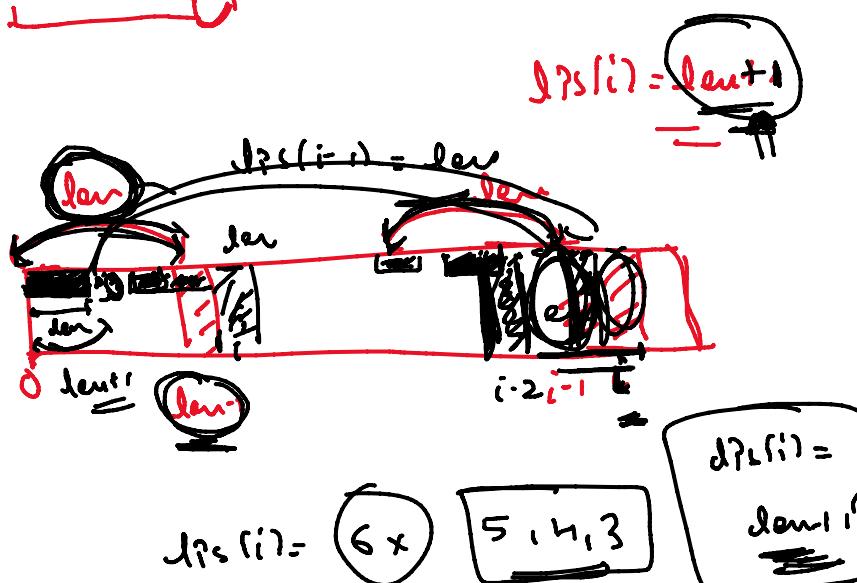
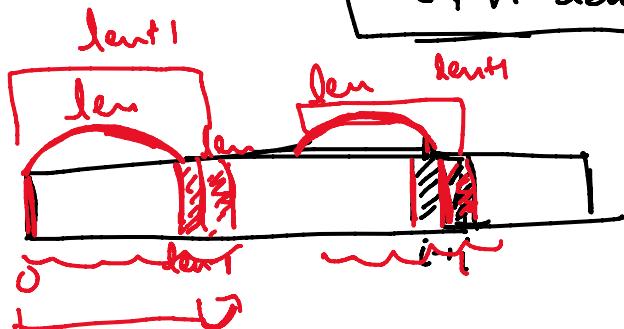
LPS Array

LPS[n]LPS[i] $O(n^3)$

(Brute Force)

len

i++

LPS[i]i, n - len - i $\text{LPS}[\text{len} - 1]$

$O(n)$

void fillLPS(int LPS[], String str, int len)

```
int n = str.size();
```

```
int len = 0;
```

```
int i = 1;
```

```
LPS[0] = 0;
```

```
while (i < n)
```

```
{ if (str[i] == str[len])
```

```
    { len++;
```

```
    LPS[i] = len;
```

```
    i++;
}
```

```
else
```

```
{ if (len == 0)
```

```
    { LPS[i] = 0;
```

```
    i++;
}
```

```
else
```

lps[i-1]

y

else

lps[i] = lps[i-1];

y

7q

KMP Algorithm

Knuth Morris Pattern Matching

Pattern Matching

str = a a a b a c a d (n)

Pat = a a b (m)

Brute Force = O(n * m)

Try to optimize

KMP Algorithm

lps

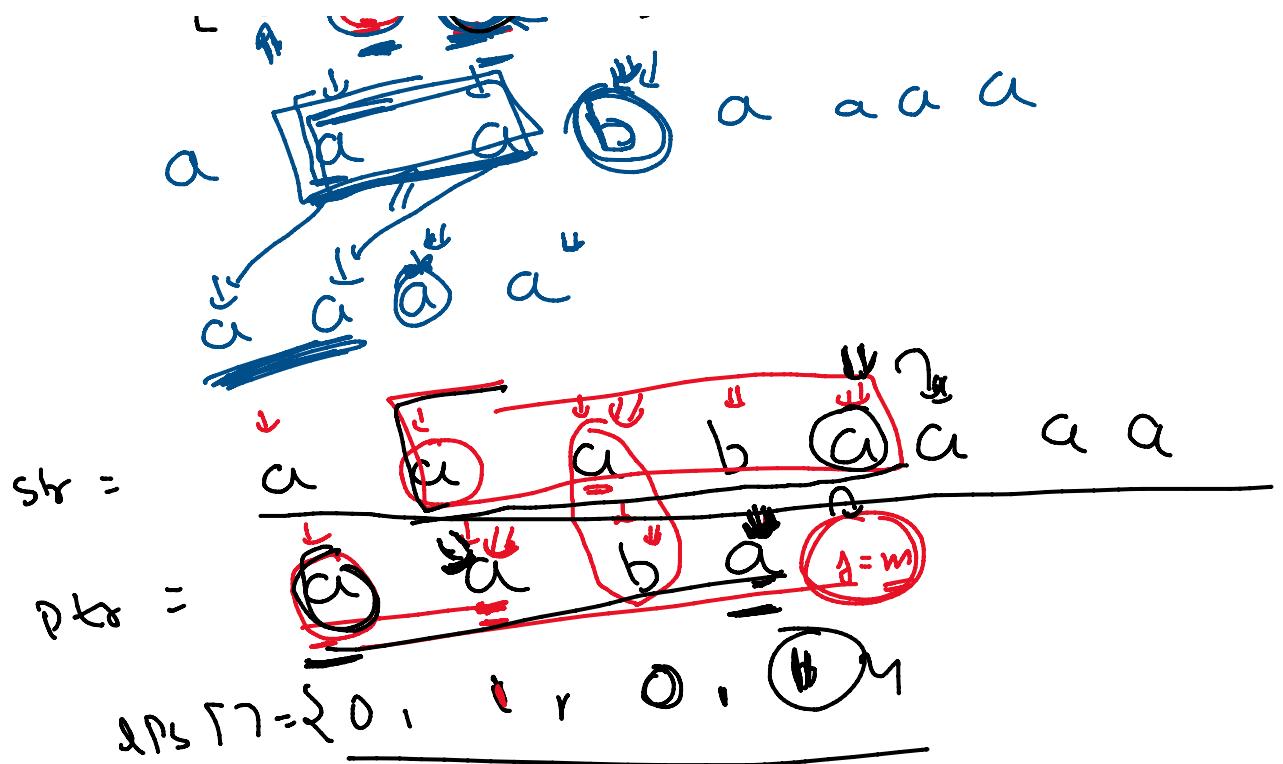
Rabin Karp Algorithm

Str = a a a b a
~~a a a~~ (1=0) ~~b~~ a
 Pat = a a a

lps[7] = [0, 1, 2, 3]

ptr = a a a b a a a a
~~a a a~~ a a a a
 [0 1 2 3] j = lps[4-1]

a a a a



void KMP (String str, String pat)

```

    int n = str.size();
    int m = pat.size();

```

```
int dp[m];
```

O(m)

~~fillLPS (dp, pat);~~

$T.C = O(n+m)$

```
i=0; j=0;
```

```
while (i < n)
```

```

    if (str[i] == pat[j])
        i++; j++;

```

O(N)

Brute Force $O(n^m)$

if (j == m)

cout << "we found the pattern" << i-m;

j = dp[i-j];

$$f = \text{EPS}f_{i-1} - j$$

γ

else if ($i < n$ and $\text{str}(i) \neq \text{Pat}(i)$)
 & if ($j = 0$) $i \leftarrow i + 1$
 else $f = \underline{\text{EPS}f_{i-1} - j}$

γ

γ

$\text{str} = a \quad \begin{array}{ccccccccc} & i & & i & & i & & i & \\ a & a & a & | & a & a & a & a & a \end{array}$

$\text{Pat} = \begin{array}{ccccccccc} & i & & i & & i & & i & \\ a & a & a & | & a & & & & \\ & & & & & m & & & \\ & & & & & & j = m & & \end{array}$

$\text{EPS}f = 2, 0, 1, 2, 3, 4$

$\underline{j = 3}$

$f = \underline{\text{EPS}f_{i-1} - j = 3}$

0
1
2
3
4

$\underline{\underline{f = \text{EPS}f_{i-1} - j}}$

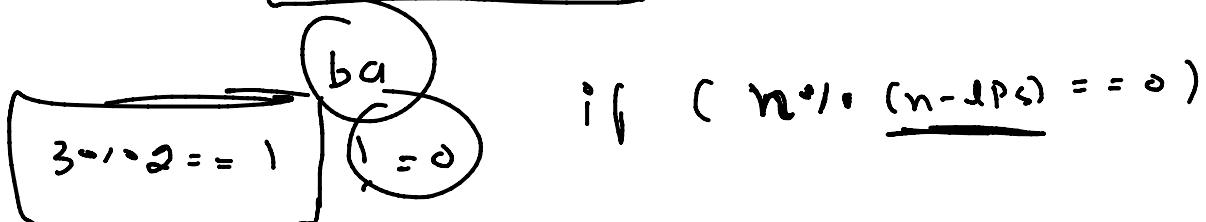
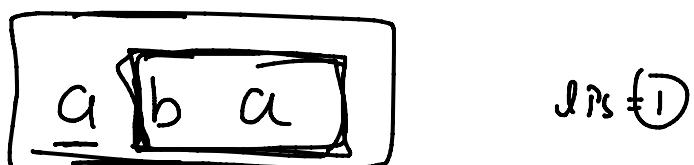
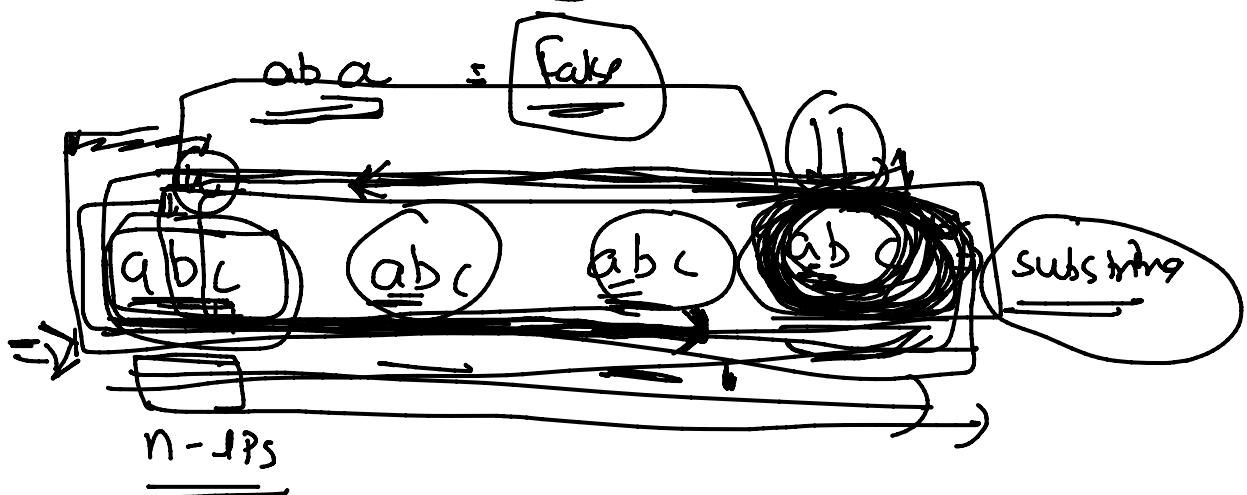
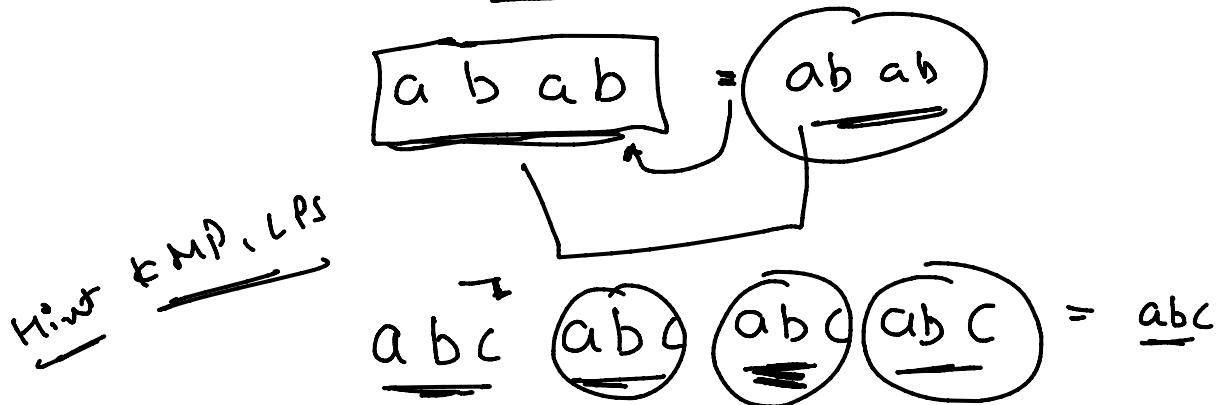
$\text{str} = \begin{array}{ccccccccc} & i & & i & & i & & i & \\ A & A & B & | & A & C & A & B & A \end{array}$

$\text{Pat} = \begin{array}{ccccccccc} & i & & i & & i & & i & \\ A & A & B & | & A & C & A & B & A \\ & & & & & m & & & \\ & & & & & & j = m & & \end{array}$

0
6
9

$f = \underline{\text{EPS}f_{i-1}}$
 $\underline{j = 0}$

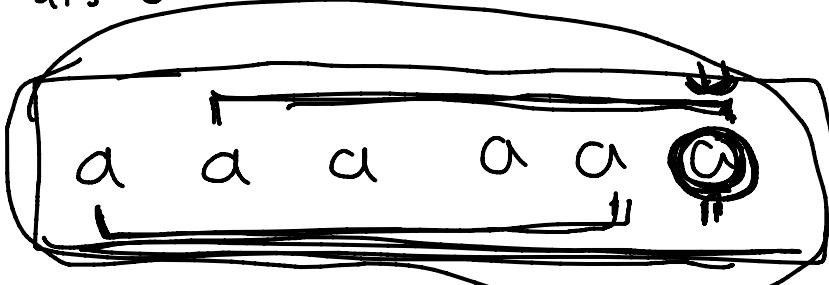
Repeated Substring Patterns



$\overbrace{a b c d a b c e}$

a b c d a b c e

$\lambda_{PS} = 0$ if ($n - \lambda_{PS} = 0$) return false



$$6 - 1 \circ 1 = = 0$$

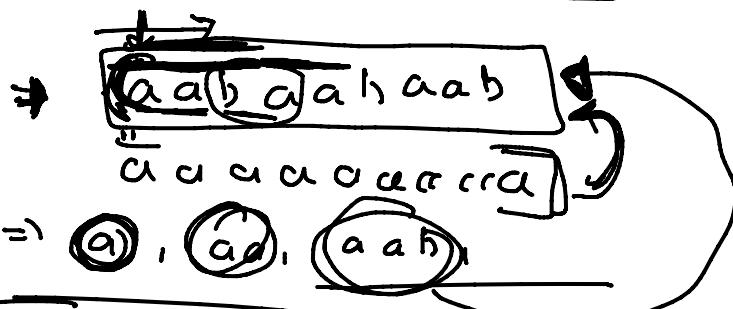
$n - \lambda_{PS}$

$\lambda_{PS}(n)$

$n - \lambda_{PS}(n-1)$

Brute Force !

String \rightarrow SubString



ab ab ab

as ab

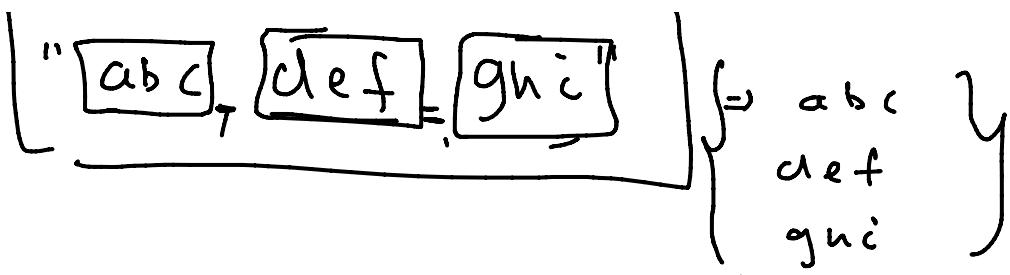
SubString

λ_{PS}

Rabin Karp

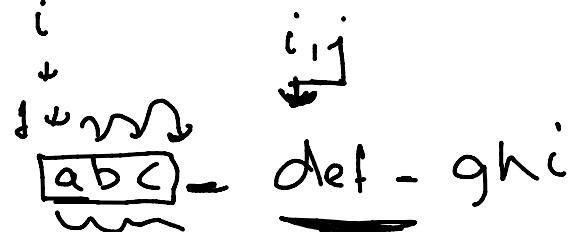
StringStream Object

"abc def ghi" \hookrightarrow abc 1



```

for (i=0; i<n; i++)
  {
    i = 0; j = 0;
    while (i < n)
  }
  
```



while (str[i] != "\n")

{
j++;
}

cout << str.substr(i, j-i);

i = j + 1;
j++;

