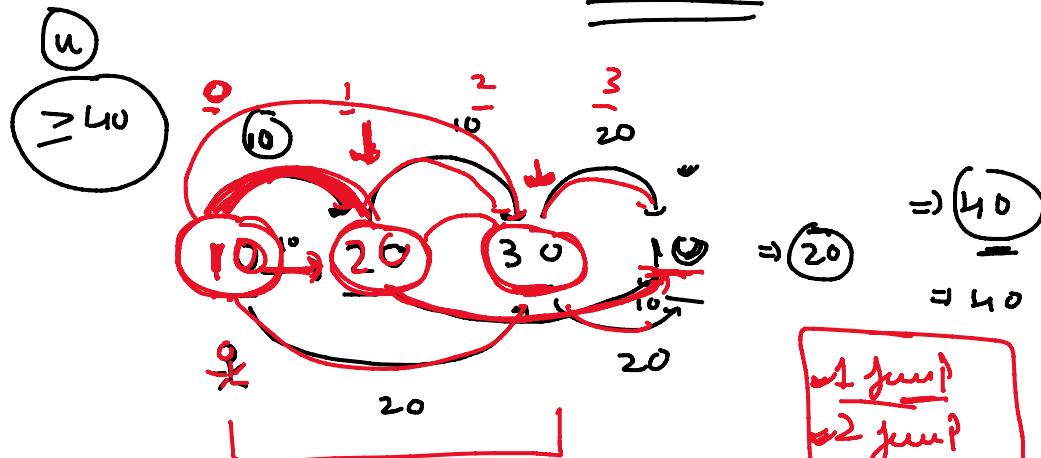


Class 80FrogJump PMin Energy Spent

1) Index

2) Apply all operations on given index

3) Calculate all using tree operations  
 $\min(o_1, o_2)$ 

Recursion

```
int frogJumpP(vector<int> &energy, int idn) {
```

Vector<int>  
Ldp  
MinValue

```
    if (idn == 0)
```

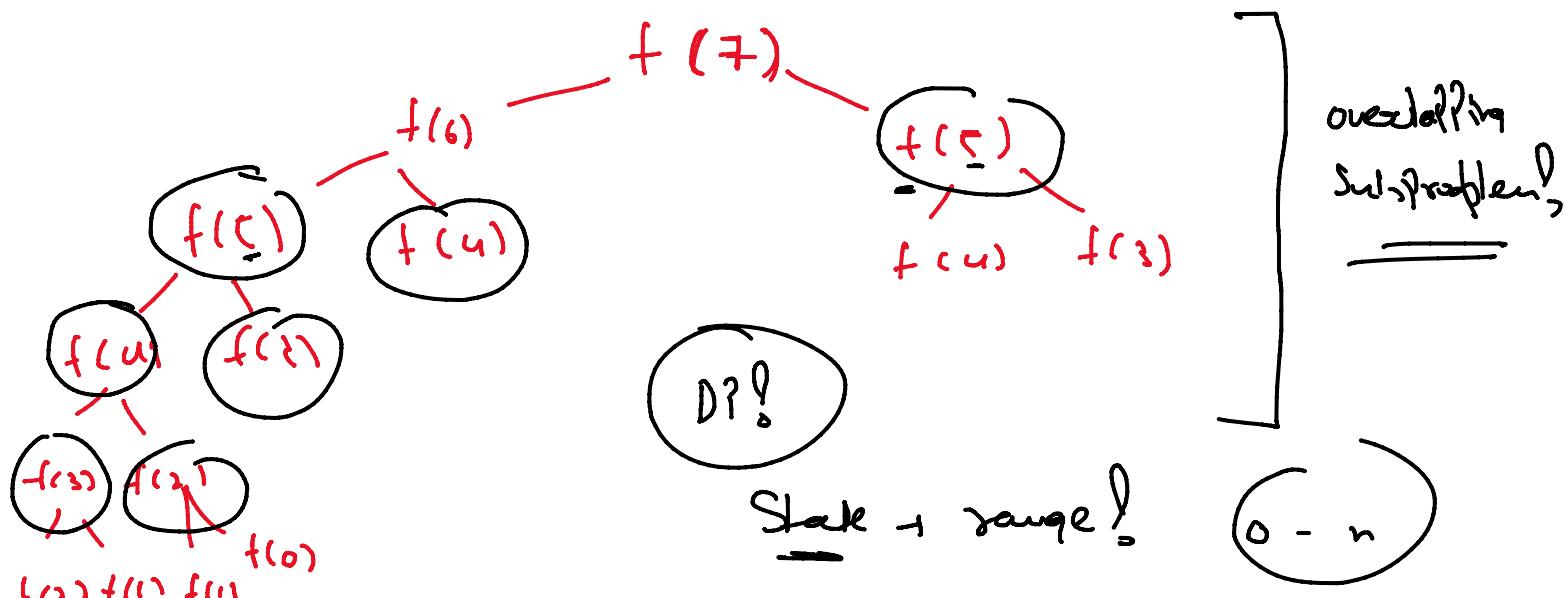
```
        return 0;
```

```
    if (idn == 1) return abs(energy[0] - energy[1]);
```

```
    if (dP[idn] != -1) return dP[idn];
```

```
    int op1 = frogJumpP(energy, idn - 1) + abs(energy[idn] - energy[idn - 1]);
```

$\text{int } \alpha_2 = \text{frequp} / \text{energy}_{i\text{du}-2} + \text{abs}(\text{energy}_{i\text{du}} - \text{energy}_{i\text{du}-2})$   
 $\text{dp}[i\text{du}] = \min(\alpha_1, \alpha_2)$   
 return  $\min \text{ and } \max$   
State!: 1 State, 1 single state



Memoization and Tabulation

vector<int> dp(n, -1);

$\text{frequp}(\dots, \text{dp});$

Tabulation

vector<int> dp(n, -1);

$$dp[0] = 0;$$

$$dp[i] = \min(\text{energy}(i) - \text{energy}(i-1),$$

int prev1, int prev2

For ( $i=2$ ;  $i < n$ ;  $i++$ )

, Space optimized version -

OC11

$$\{ \quad \text{int } oP_1 = \min(dp[i-1], \text{abs}(\text{energy}(i-1) - \text{energy}(i))),$$

$$\text{int } oP_2 = \min(dp[i-2], \text{abs}(\text{energy}(i-2) - \text{energy}(i))),$$

$$dp[i] = \min(oP_1, oP_2);$$

?

i

$(i-1)(i-2)$

return  $dp[n-1]$ ;



Frog jump P with IC available jump

IC<sup>ith</sup> pos

1 2 3 4 --- IC

int frogjump (vector<int> &energy, int k, int idu)

if ( $idu == 0$ ) return 0;

if ( $dp[idu] != -1$ ) return dp[idu];

int mn = INT\_MAX;

For ( $i=1$ ;  $i < k$ ;  $i++$ )

if ( $idu - i \geq 0$ )

{  $mn = \min(mn, \text{frogjump}(\text{energy}, k, idu - i) +$   
     $\text{abs}(\text{energy}[idu - i] - \text{energy}[i]))$

?

meten dp [idk] = mn

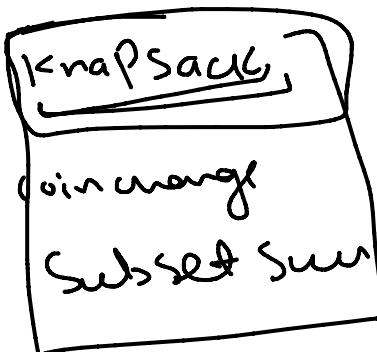
$T.C = O(n \times k)$

$S.C = \underline{O(n)}$

Tabulation

1D DP

Count all ways



→ 10 - 20

Subset

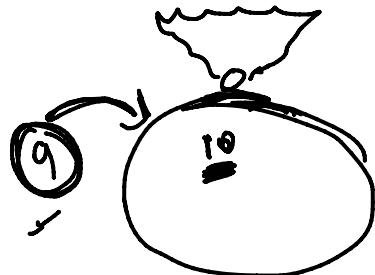
fractional knapsack  
(Greedy)

KnapSack

0 - 1 knapsack

unhandled knapsack

Recursion



KnapSack

DP  
Subset

Combination

D D N D

subset maximizes total weight

knapsack weights

100 50 90 150 10

Recursion

→ Index

→ Index

→ All objects at an Index ✓  
✗

→ Max value which object can → max( $oP_1$ ,  $oP_2$ )

subset,

0 1 2 ... n

int knap (vector<int>&weights, vector<int>&values,

int maxweight, vector<vector<int>>&dp)

index

0 1 2 3 4 ↓

if ( $i_{du} == 0$ )

return 0

if (maxweight  $\leq 0$ )

return 0;

if ( $i_{du} < maxweight$ )  
return  $dp[i_{du}] \leftarrow \max \{ weight[i_{du}-1] + \text{value}[i_{du}-1], \text{value}[i_{du}-1] \}$

int  $oP_1 = 0$ ,  $int oP_2 = 0$ ;

if (weight[i\_{du}-1]  $\leq maxweight$ )

$oP_1 = knap (weights, values, i_{du}-1, maxweight - weight[i_{du}-1]) +$   
 $value[i_{du}-1];$

$oP_2 = knap (weights, values, i_{du}-1, maxweight);$

return  $\max(oP_1, oP_2)$

return  $\max(oP_1, oP_2)$

State ? = 2

2D vector

knap (10, n)

knap (8, n-1)

knap (6, n-2)

knap (8, n-2)

knap (10, n-1)

knap (8, n-2)

$\text{knapsack}(n-2)$

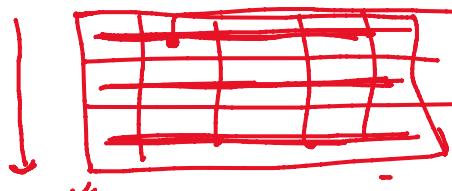
2D DP

$$5138 \Rightarrow (n+1) \times (w+1)$$

$10^{12}$  knap

`vector<vector<int>> dp(n+1, vector<int>(w+1, -1));`

I level



$$O = (n+1) \times (w+1)$$

w Tabulation

	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	-	-	-	-	-	-
2	0	-	-	-	-	-	-
3	0	-	-	-	-	-	-
4	0	-	-	-	-	-	-
5	0	-	-	-	-	-	-
6	0	-	-	-	-	-	-

frog for I, II

knapsack -> 0-1, fraction

unbounded

filtration, climbing stairs