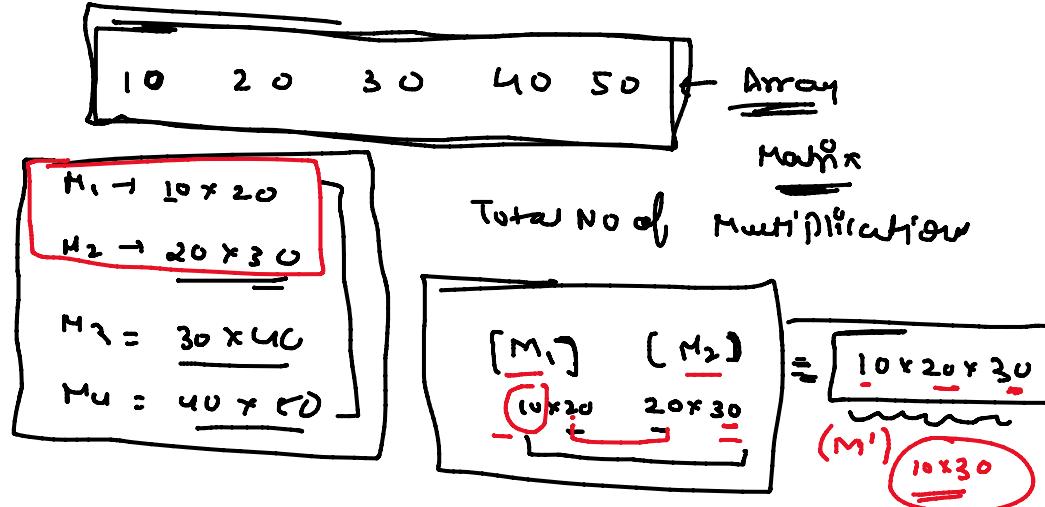
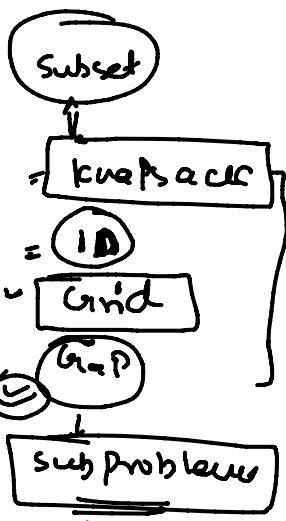


Class 88Dynamic ProgrammingMatrix chain multiplication / Partition DP

Order  
Matters

$$\begin{aligned}
 & M_1 \quad M_2 \quad M_3 \quad M_4 \\
 & [M_1] [M_2 \quad M_3 \quad M_4] \Rightarrow [M_1] [M' \quad [M_3 \quad M_4]] \\
 & 20 \times 30 \times 40 = 24,000 \\
 & (20 \times 30) \times (40 \times 10) = 20 \times 40 \times 10 = 40,000 \\
 & M'' = 20 \times 10 \quad M' = 10 \times 30 \\
 & = [M_1 \quad M_2] [M_3 \quad M_4] = 81,000 \\
 & 10 \times 20 \times 30 = 6,000 \\
 & 30 \times 40 \times 20 = 24,000
 \end{aligned}$$

$$\begin{aligned}
 & M'' = 20 \times 10 \quad M' = 10 \times 30 \\
 & = 10 \times 20 \times 10 = 10,000 \Rightarrow 6,000 + 10,000 = 16,000 \\
 & M'' = 20 \times 10 \quad M' = 10 \times 30
 \end{aligned}$$

MCM

$$\begin{aligned}
 & M_1 \quad M_2 \quad M_3 \quad M_4 \\
 & [M_1] [M_2 \quad M_3 \quad M_4] \quad [M_1 - M_4] \\
 & [M_1 \quad M_2] [M_3 \quad M_4] \quad [(M_1 \quad M_2) \quad M_3] \quad (M_4)
 \end{aligned}$$

Partition DP?

$$\begin{aligned}
 & [M_2] [M_3 \quad M_4] \\
 & (M_1 \quad M_2) [M_3 \quad M_4]
 \end{aligned}$$

$i > i$   $\vdots$   $\vdots$



`int solve( int i, int j, vector<int> arr )`

```
< if (i >= j)      // Base case
  return 0;
```

`int ans = INT_MAX;`

`for( k = i; k <= j - 1; k++ )` sb1

$\downarrow$  `int tempAns = colup( i, k, arr ) +`  
 $\quad\quad\quad$  `solve( k+1, j, arr )` sb2

`ans = max( ans, tempAns );`

merge

$\downarrow$  
$$\text{ans}[i-1] * \text{ans}[j] * \text{ans}[j]$$

$\nwarrow$   
 $\searrow$   
 rechungen:

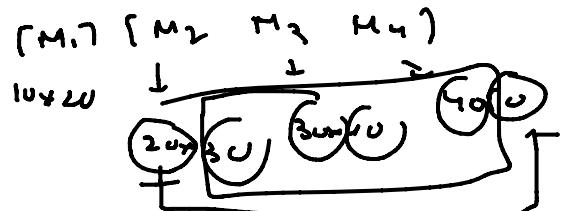
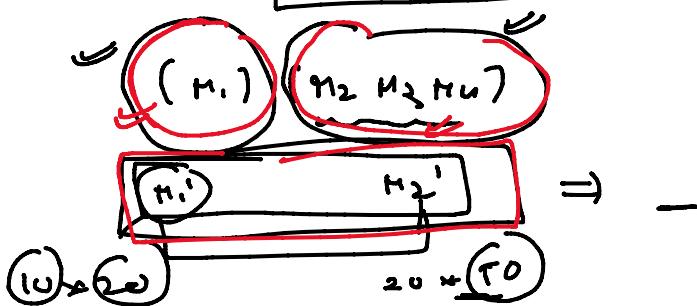


$(M_1, M_2)$   
 $(M_3, M_4)$   
 $20 \times 30$

$\text{ans}[i-1] * \text{ans}[i]$

$M_1'$   
 $10 \times 20$   
 $20 \times 50$

$10 \times 20 \times 50$





```
int solve ( int i , int j , vector<int> arr )
```

if ( i >= j )  
return 0;

if ( dp[i][j] != -1 ) return dp[i][j];

```
int ans = INT_MAX ;
```

```
for ( int k = i ; k <= j - 1 ; k++ )
```

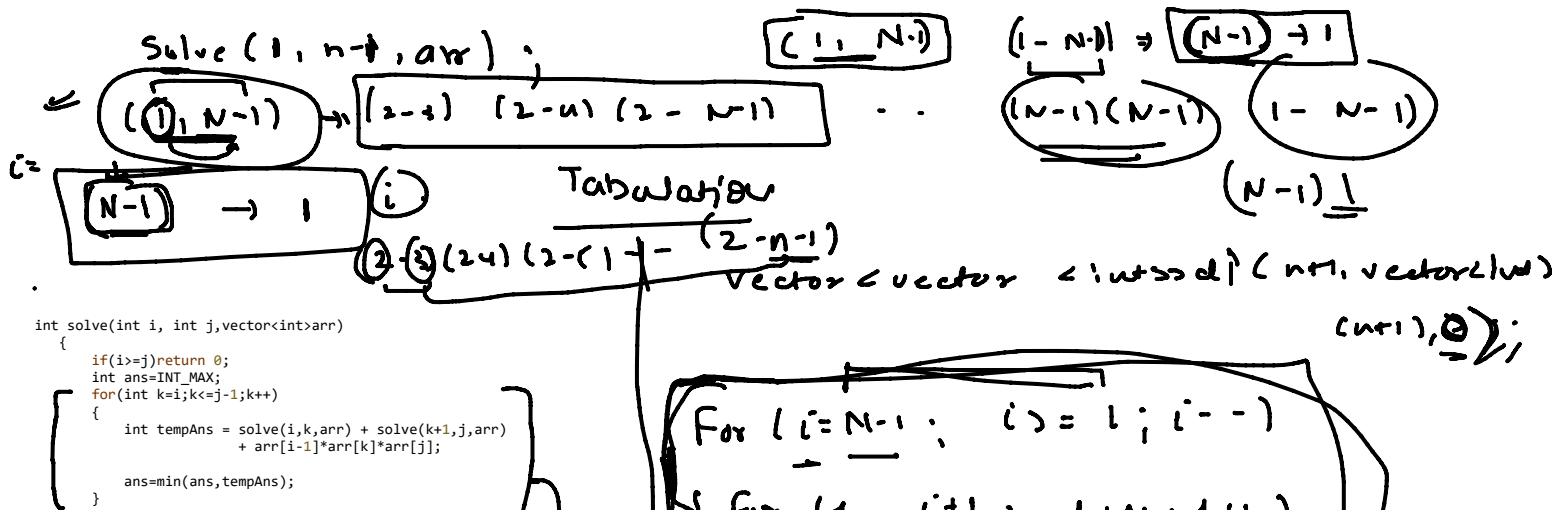
int tempAns = solve( i , k , arr ) + solve( k + 1 , j , arr ) +

$\underbrace{arr[i-1] * arr[k] * arr[j]}_{\text{mergerow}}$  ;

ans = min( ans , tempAns );

q

return dp[i][j] = ans ;



```

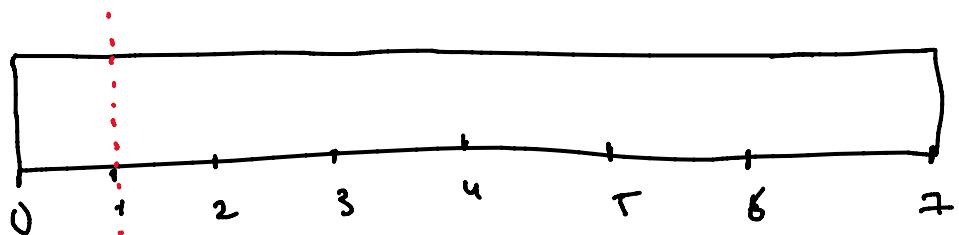
    int tempAns = solve(i, k, arr) + solve(k+1, j, arr)
        + arr[i-1]*arr[k]*arr[j];
    ans=min(ans,tempAns);
}
int matrixMultiplication(vector<int> &arr, int N)
{
    // Write your code here.
    return solve(1, N-1, arr);
}

```

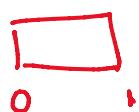
$$[2-2] = O(M_1) = \Theta$$

for ( $i = N-1$ ;  $i = l, l-1$ )  
 for ( $j = i+1$ ;  $j \leq N$ ;  $j+1$ )  
 "int ans = INT-MAX",  
 for ( $k = i$ ;  $k \leq j-1$ ;  $k+1$ )  
 ↳ int tempAns = dp[i][r] +  
 dp[k+1][f] +  
 arr[i-1] \* arr[k] \* arr[f];  
 ↳ ans = min(ans, tempAns);  
 dp[i][j] = ans;  
 ↳  
 ↳ return dp[1][N-1];

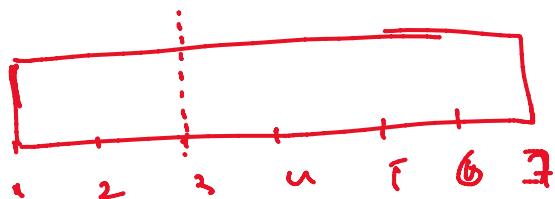
Cut the Stick



- 7



[ 0 3 4 5 ]  
=

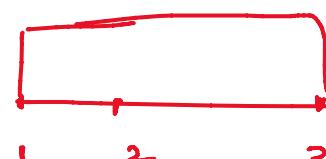


$$7 + 6 + 4 + 3$$

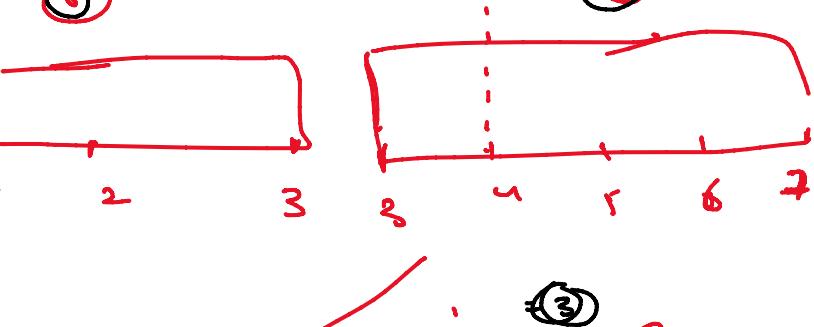
$$= 20$$

$$4 \quad 7 \quad \overline{13}$$

6

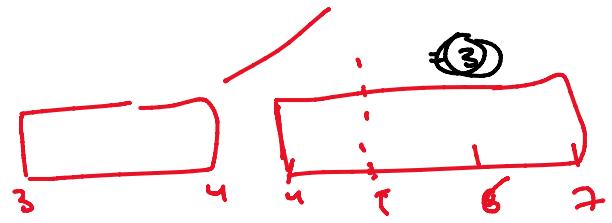


4

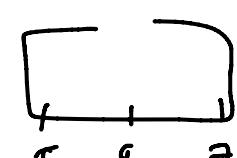
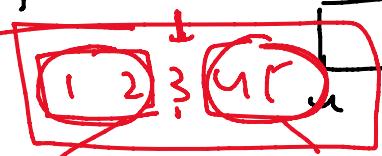


3

11. 15

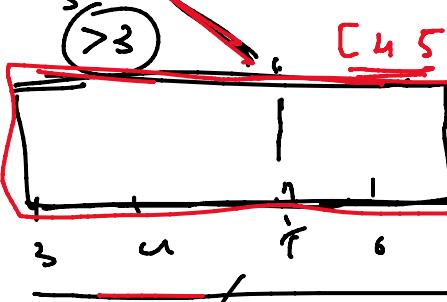
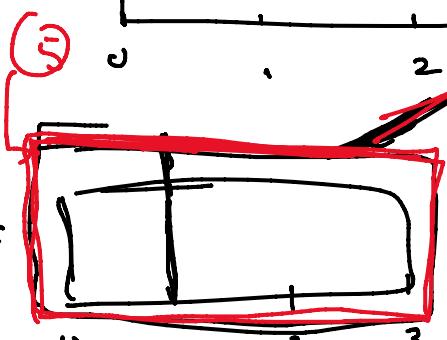


[5, 3, 1, 4, 2]  
[5, 3, 1, 4, 2]



(1, 2)

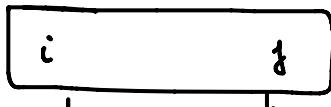
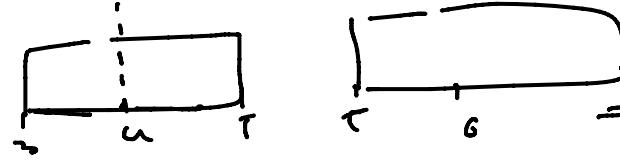
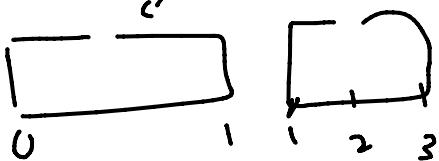
[0 - 3)



[3, 4, 5, 6, 7]

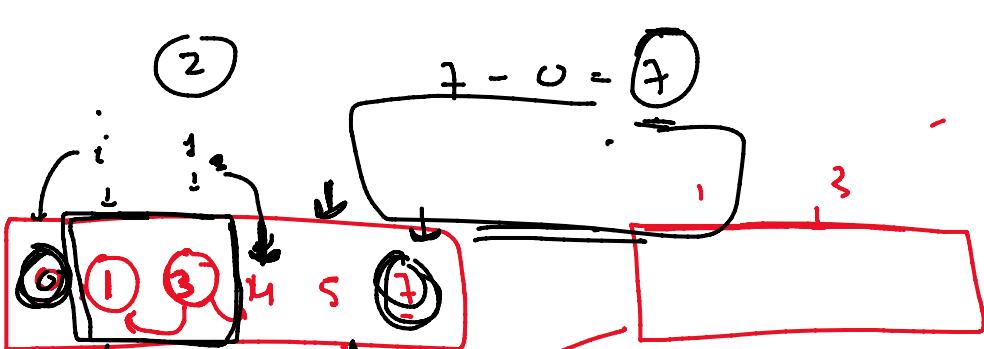
(2)

$$\begin{array}{r} + 2 \\ + 2 \\ \hline = 16 \end{array}$$

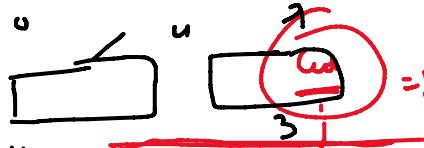
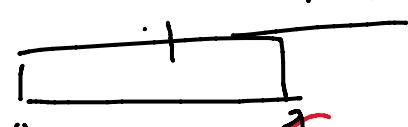


(2)

$$7 - 0 = 7$$

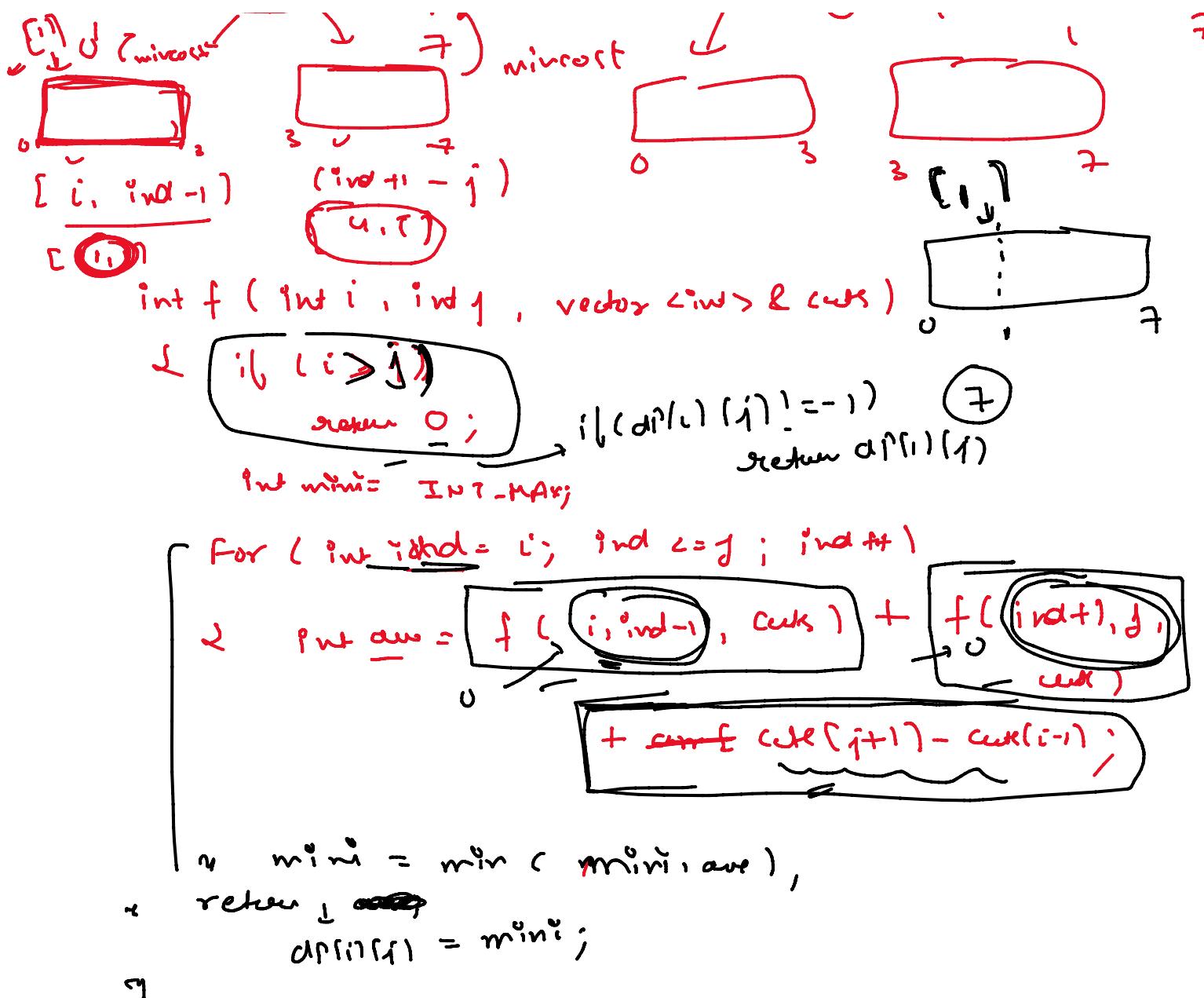


1, 3, 4, 5



$$4 - 0 = 4$$

mincost



```

    cuts . push_back (rod.length);
    cuts . insert (cuts . begin () , 0 );
    sort (cut . begin () , cuts . end ());
    f ( 1 , n - 2 , cuts );
  
```

( Tabulation )

$i = 1$	$n-2$	$i = n-1$
---------	-------	-----------

$i = n-2$

$\text{dp}(n-2) \dots$

$i = j$

For ( $i = n-2$ ;  $i \geq 1$ ;  $i--$ )

  ↓ For ( $j = i$ ;  $j <= n-2$ ;  $j+1$ )

    { ~~if~~ int  $\min_i = \text{INT-MIN}$ ,

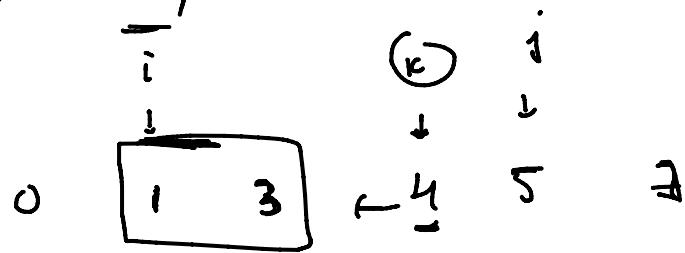
      For ( $\text{int } \text{ind} = i$ ;  $\text{prod } c = 1$ ;  $\text{ind } + 1$ )

        ↓  $\text{ave} = \text{ave}[j+1] - \text{cav}[i-1] + \text{dp}[i][\text{ind}-1]$   
                  +  $\text{dp}[\text{ind}+1][j]$ ;

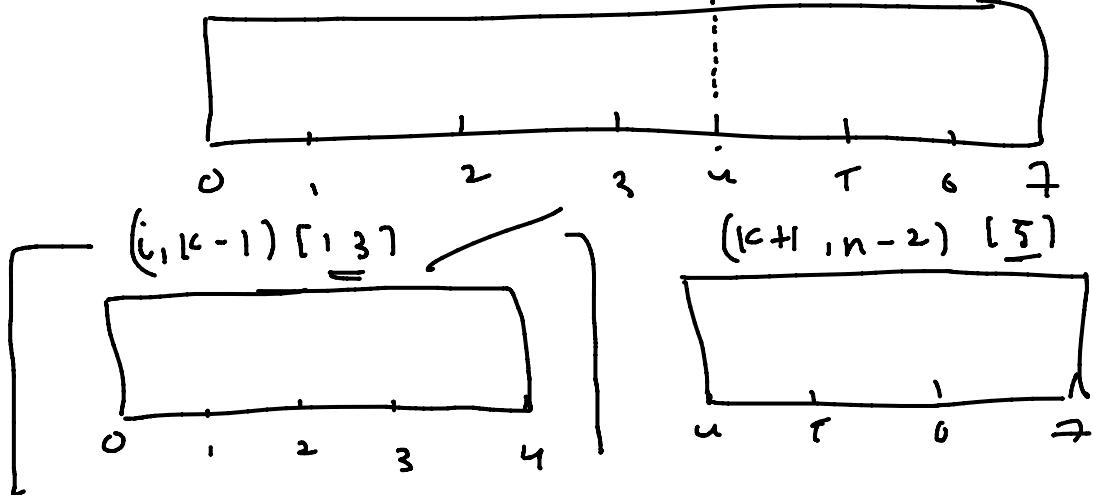
$\min_i = \min(\min_i, \text{ave})$ ;

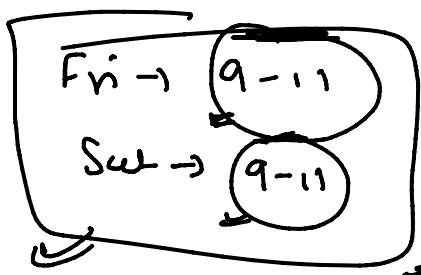
~~set~~  $\text{dp}[i][j] = \min_i$ ;

return  $\text{dp}[1][n-2]$ ;



$\leq k$





11-1

IPL Match

3:30