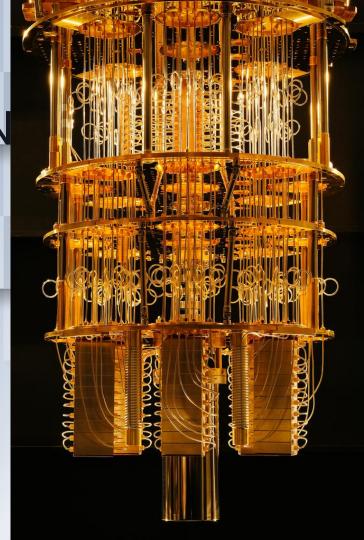
INTEGER FACTORIZATION **USING SHOR'S ALGORITHM AND ITS** IMPLEMENTATION ON A QUANTUM COMPUTER

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Integer Factorization

- □ An integer $I = p_1^{a_1} p_2^{a_2} ... p_n^{a_n}$ can be decomposed into unique primes: $p_i's$ and $a_i's$ are their respective power.
- Question: what are the prime factors?
- □ Classical Method (a): Divide I by all the values $2 \le x < I$ to find reminder.
- □ Classical method (b): Divide I by all the values $2 \le x < \sqrt{I}$ to find reminder

- Complexity: amount of resources(time or space needed)
- $\ \square$ Method(a) takes $\mathcal{O}(2^w)$ time complexity where $\ w = log_2 I$
- □ Method(b) takes $\mathcal{O}(2^{w/2})$ time complexity
- □ The best known algorithm for factorization is General Number Field Sieve(GNFS):

time complexity:
$$\mathcal{O}\left(\exp\left(\mathbf{c}\mathbf{w}^{\frac{1}{3}}(\log\mathbf{w})^{\frac{2}{3}}\right)\right)$$
 [2](Hamdi et al.,2014)

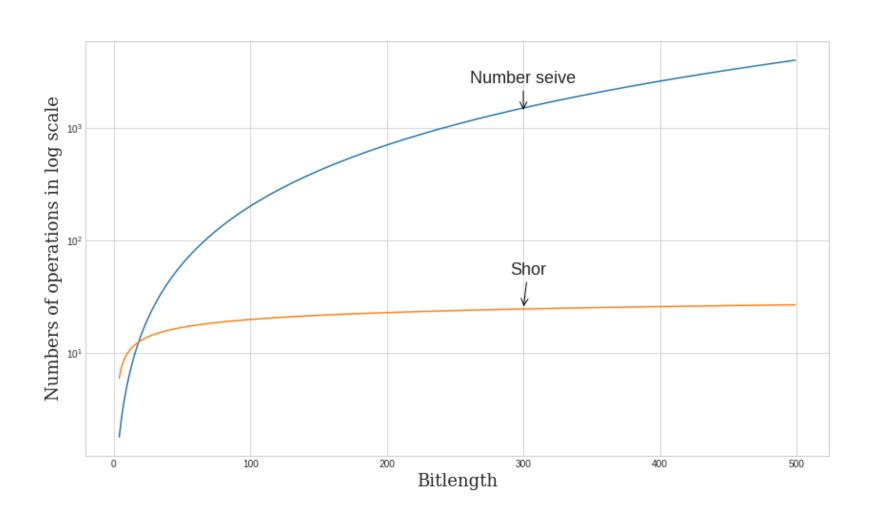
 RSA cryptography is based on the difficulty of factoring problem

Shor's Algorithm

It is a

- ☐ Quantum Algorithm
- □ Developed by Peter Shor in 1994
- ☐ For composite integer factorization
- ☐ In polynomial time with bounded error.





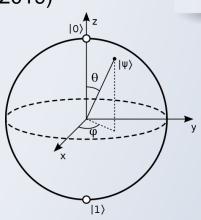
Introduction to Quantum Computing

- Act of leveraging quantum mechanical
 properties to perform computing [3](Hidary, 2019)
- Qubit

Basic unit of information

Superposition of $|0\rangle$ and $|1\rangle$ states

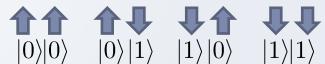
$$\begin{aligned} |\Psi\rangle &= \alpha |0\rangle + \beta |1\rangle where |\alpha|^2 + |\beta|^2 = 1 \\ &= \left(\cos \frac{\theta}{2} |0\rangle + e^{i\phi} \sin \frac{\theta}{2} |1\rangle\right) \end{aligned}$$



Measurement of $|\Psi\rangle=\alpha\,|0\rangle+\beta\,|1\rangle\,$ collapse superposition to single

$$P(|0\rangle) = |\alpha|^2 \qquad P(|1\rangle) = |\beta|^2$$

For 2 qubit system, system can have superposition of all four states



 \square For n qubit system, there are 2^n different possible states

$$|\Psi\rangle = \sum_{m=0}^{2^n-1} a_m |m\rangle$$

■ Entanglement between qubits is a special type of correlation such that change in one instantaneously triggers an effect on other.[4]

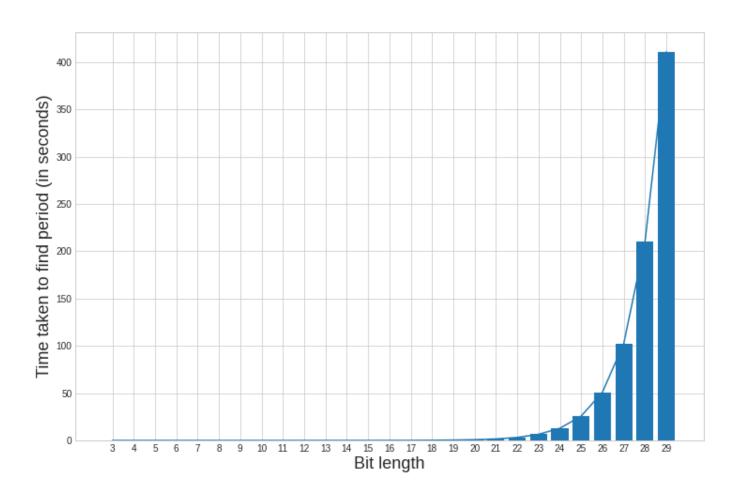
Theory of Shor's Algorithm

Say $M = p \cdot q$ be a composite odd integer and $M \neq p^k$, for prime p,

- 1. Choose a random number : 1 < x < M
- 2. If $GCD(x, M) \neq 1$, then factor = GCD(x, M) where GCD = greatest common divisor
- 3. If gcd(x, M) = 1, find period, a of MEF function $f(r) = x^r \mod M$, $r \in \mathbb{Z}(M)$
- 4. If period: a is odd or $x^{a/2} \equiv 1 \pmod{M}$, we restart the algorithm from step 1
- 5. If a is even and $x^{a/2} \not\equiv 1 \pmod{M}$ then, factors of M are: $p = GCD(x^{a/2} + 1, M)$ or/and $q = GCD(x^{a/2} 1, M)$

An example

- \Box Let M = 15 and x = 2 then GCD(2,15)=1, $\Box f(r) = 2^r \mod 15, r \in \mathbb{Z}(15)$ would give $f(0) = 1 \mod 15 = 1$, $f(1) = 2 \mod 15 = 2$, f(2) = 1 $4 \mod 15 = 4, \dots$ i.e. $f(r) = 1, 2, 4, 8, 1, 2, 4, 8, 1, 2, 4, 8 \dots$ So period, a = 4□ Factors are $GCD(2^{\frac{4}{2}} + 1,15) = GCD(5,15) = 5$, $GCD(2^{\frac{4}{2}} - 1.15) = GCD(3.15) = 3$
- \Box Hence factors are 5 and 3:5*3 =15



Shor's algorithm on quantum computer

- 1. Initialize two quantum registers $|0\rangle^n |1\rangle^r$ Where $M^2 < 2^n < 2M^2$, r > log(M) + 1
- 2. Prepare uniform superposition by applying Hadamard gates
 - ($H^{\otimes n}$)in first register
- 3. Apply the periodic function $f(x) = z^x \pmod{M}$ on second register
- 4. Apply QFT on first register to obtain the frequency of the superimposed state in 1st register and
- 5. Measure the input register

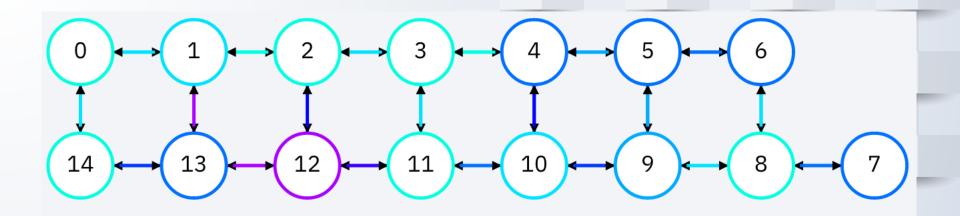


Fig: Circuit diagram for Shor's Algorithm Implementation

- 5. Classical analysis: Use Continued Fraction Algorithm and Euclidean Algorithm to find period a and then find the factors
- Shor's algorithm has a bounded error
- Complexity
- Hadamard gates $(H^{\otimes n})$: $O(\log M)$ Oracle Function : $O(\log^3 M)$
- QFT: $O(log^2 M)$ Euclidian Algorithm: $O(log^3 M)$
- □ Complexity of algorithm= $O(log^3 M)$
- Hence Shor's algorithm is a Bounded error Quantum polynomial (BQP) algorithm.

Methodology: Implementation

- Programming language: Python; package: Qiskit
- Qiskit: an open source SDK to work with quantum circuits and algorithms, and acts as interface for running quantum circuits on quantum computers at IBM [8]
- Hardware: ibmq_qasm_simulator and ibmq_16_Melbourne
- ibmq_qasm_simulator: quantum simulator at IBM
- ibmq_16_Melbourne: a 15 qubit quantum computer at IBM
- Qubits used: 0,1,2,3,13,14



Single-qubit U2 error rate

CNOT error rate

4.451e-4

4.292e-3

1.152e-2

8.551e-2

Topography diagram and coupling map of *ibmq_16_Melbourne* [8]

Circuit design

- ❖ Factor an integer I=15 with base z = 2 for MEF
- Created a simplified compiled version of Shor's algorithm (Geller and Zhau,2013)

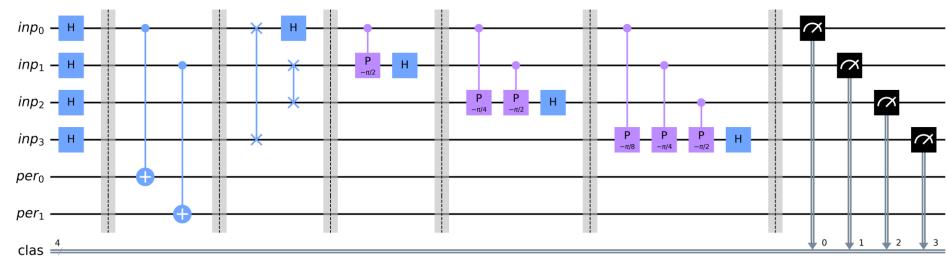
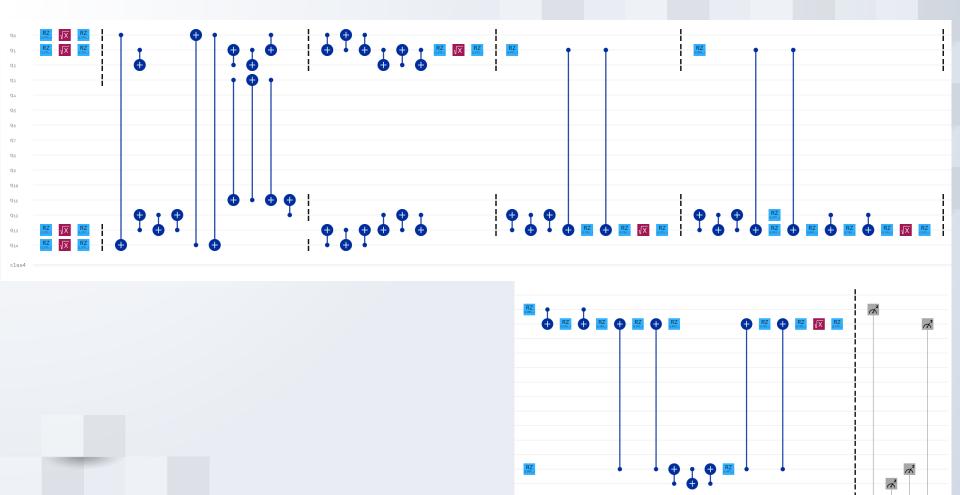
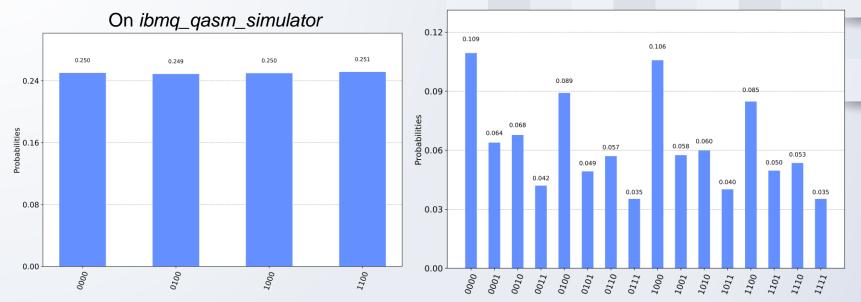


Fig: Quantum Circuit diagram for Shor's Algorithm to factor integer 15



Result

On ibmq_16_MelbourneQC



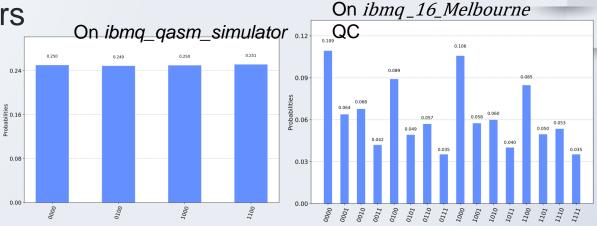
- Peaks at decimal equivalents: 0,4,8,12
- By classical processing, order a = 4
- And factors = $GCD(2^2 + 1,15) = 5$ and $GCD(2^2 1,15) = 3$
- Validation: 3*5=15.

Discussion

Dissimilarities in results

Quantum Errors

Scalability



Significance: Break RSA cryptography

- □ A widely used cryptographic system for online data transmission including emails and online payments
- Based upon the difficulty of factoring a large number
- □ 300 trillion years for classical computer to break 2048 bit key while it takes
 10 secs for 4099 qubit QC(Zhang, 2020)
- □ Largest experimentally factored number by quantum computer using Shor's algorithm: '21' (Martin-Lopez,2012)

RSA-(2048) = 2519590847565789349402718324004839857142928212620403202777713783 60436620207075955562640185258807844069182906412495150821892985591 49176184502808489120072844992687392807287776735971418347270261896 37501497182469116507761337985909570009733045974880842840179742910 06424586918171951187461215151726546322822168699875491824224336372 59085141865462043576798423387184774447920739934236584823824281198 16381501067481045166037730605620161967625613384414360383390441495 26344321901146575444541784240209246165157233507787077498171257724 67962926386356373289912154831438167899885040445364023527381951378 636564391212010397122822120720357

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