

---

6. Show that the following problem is NP-complete:  
Input: An undirected graph  $G$  and an edge  $e$ .  
Question: Does  $G$  have a Hamiltonian cycle that passes through the edge  $e$ .

**Solution:**

To prove that this problem is NP-complete:  
let this problem be  $y$ .

1. I will first prove that this problem is in NP, so  $(y \in \text{NP})$

If we are given a certificate, so a Hamiltonian cycle that passes through the edge  $e$ , we can verify this certificate in polynomial time. As a verification algorithm, we can simply iterate through the cycle given and check if every node is visited once at the meanwhile we can also check if edge  $e$  is included in the cycle. This can be done in polynomial time  $O(n)$ , and the problem is in NP.

2. We know from the lectures that the Hamiltonian cycle problem is NP-complete (slide 34 of Intractability-II), so let the Hamiltonian cycle problem be  $x \in \text{NP-complete}$ .

3. I will reduce the Hamiltonian cycle problem to this problem, so  $(x \leq_p y)$ .  
For an input  $G$  for Hamiltonian cycle problem, we can solve the problem with the algorithm below:

For every edge  $e$  in  $G$

    Call problem  $y$  with graph  $G$  and edge  $e$ , if the output is YES, output YES and terminate.

output NO

Given that we can solve problem  $x$  with polynomial calls to problem  $y$  then we can say that  $(x \leq_p y)$ .

After completing the three steps above we can say that the problem is NP-complete.