Encryption

Introduction

Before electronic communications

Codes

A few fundamentals

Key-based encryption

Cracking the code

Brute force

Block or stream

Private-key methods

Encryption keys

Passing keys

Public-key encryption

One-way hash

Encrypting disks

PGP encryption









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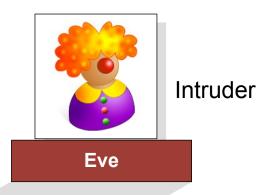






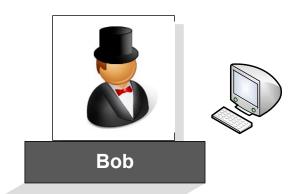


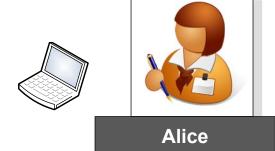
Introduction





Bob and Alice – good Eve – bad Trust – trusted?

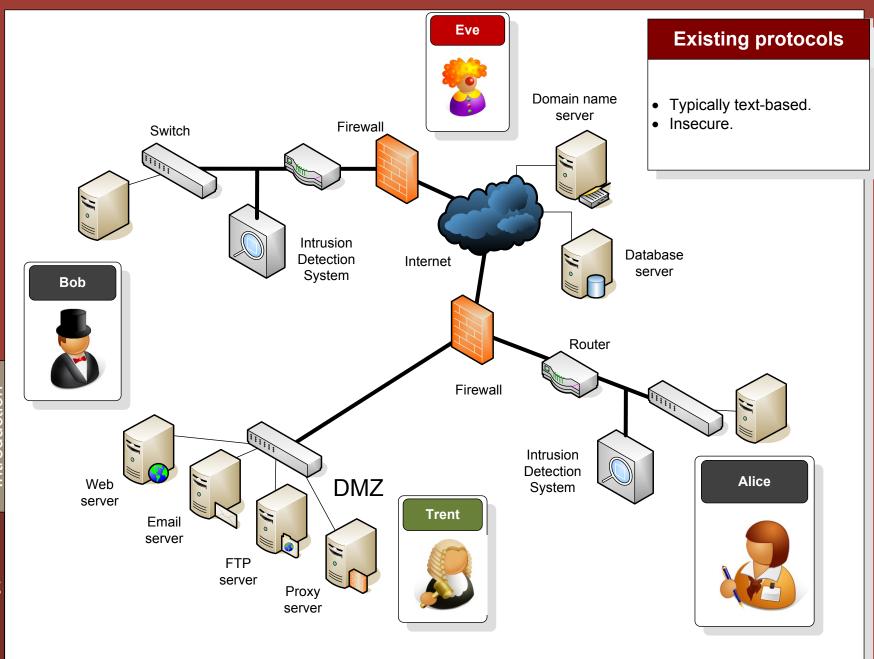


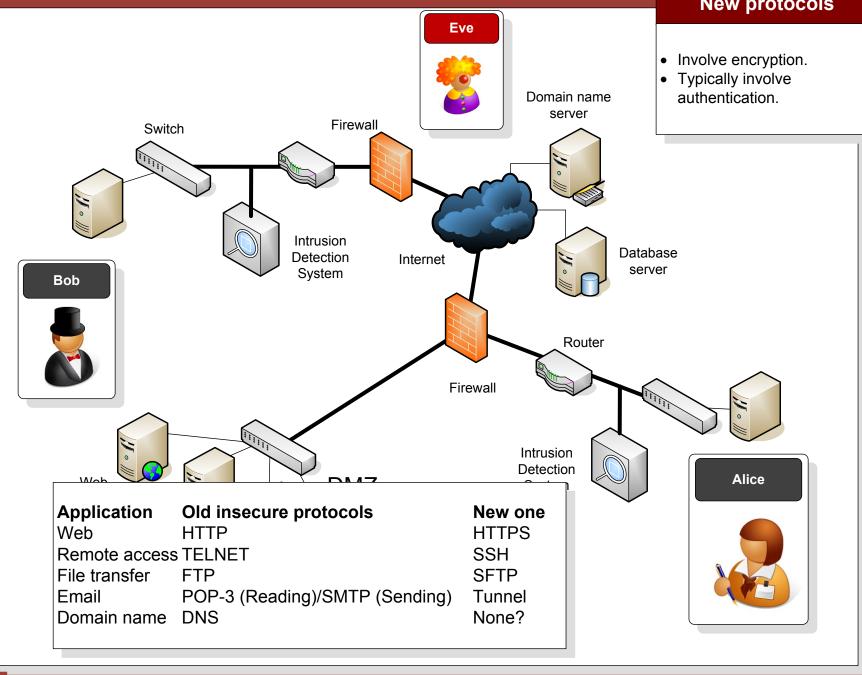




Trusted third party

Trent







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Before electronic communications





Secret Communications

- Quilts
- Carrier pigeon
- Smoke signals
- Etc...

Carrier pigeon

Quilt patterns (used by slaves to escape)



Microfiche



Smoke signals



Code talkers: Navajo words

Before electronic communications

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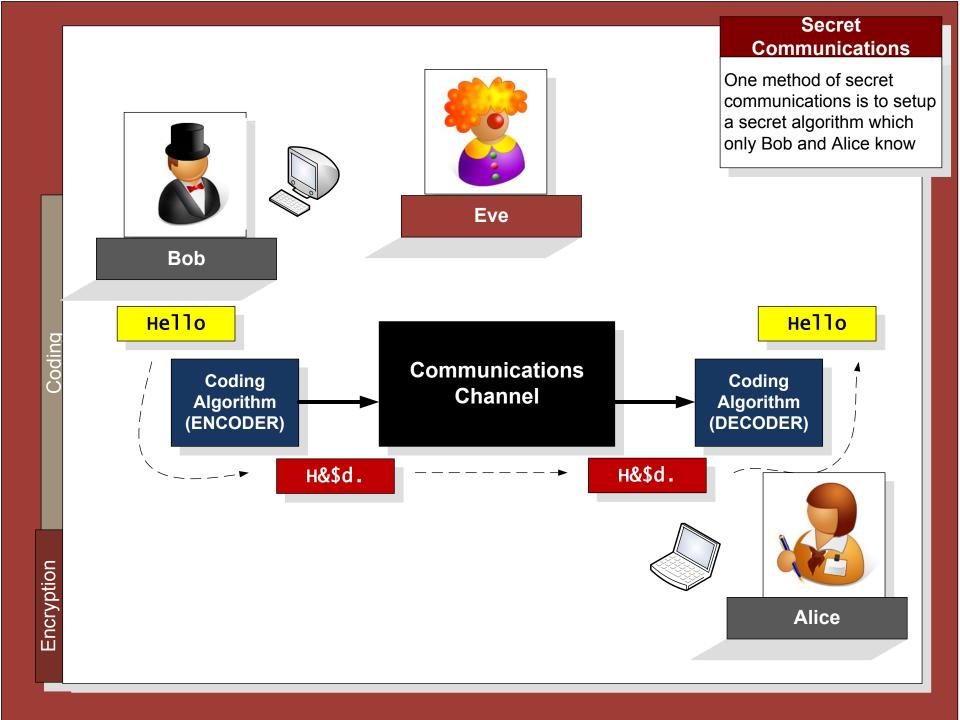








Codes



Simple alphabet shifting

abcdefghijklmnopqrstuvwxyz YZABCDEFGHIJKLMNOPQRSTUVWX

RFC ZMW QRMMB ML RFC ZSPLGLE BCAI



Caesar code

25 code mappings



Code Mapping

Code mapping scrambles the alphabet ..

403 million billion billion codes.

Code mapping

abcdefghijklmnopqrstuvwxyz MGPOAFZBCDIEHXJKLNTQRWSUVY

QBCT CT MX AUMHKEA KCAPA JF QAUQ



4.03×10²⁶ codes



Letters (%)	Digrams (%)	Trigrams (%)	Words (%)
E 13.05	тн 3.16	THE 4.72	THE 6.42
т 9.02	IN 1.54	ING 1.42	OF 4.02
0 8.21	ER 1.33	AND 1.13	AND 3.15
A 7.81	RE 1.30	ION 1.00	TO 2.36
N 7.28	AN 1.08	ENT 0.98	A 2.09
I 6.77	HE 1.08	FOR 0.76	IN 1.77
R 6.64	AR 1.02	TIO 0.75	THAT 1.25
S 6.46	EN 1.02	ERE 0.69	IS 1.03
н 5.85	TI 1.02	HER 0.68	I 0.94
D 4.11	TE 0.98	ATE 0.66	IT 0.93
L 3.60	AT 0.88	VER 0.63	FOR 0.77
C 2.93	ON 0.84	TER 0.62	AS 0.76
F 2.88	на 0.84	THA 0.62	WITH 0.76
U 2.77	OU 0.72	ATI 0.59	WAS 0.72
м 2.62	IT 0.71	HAT 0.55	HIS 0.71
P 2.15	ES 0.69	ERS 0.54	не 0.71
Y 1.51	ST 0.68	HIS 0.52	BE 0.63
W 1.49	OR 0.68	RES 0.50	NOT 0.61
G 1.39	NT 0.67	ILL 0.47	BY 0.57
в 1.28	ні 0.66	ARE 0.46	BUT 0.56
V 1.00	EA 0.64	CON 0.45	HAVE 0.55
К 0.42	VE 0.64	NCE 0.43	YOU 0.55
x 0.30	co 0.59	ALL 0.44	WHICH 0.53
J 0.23	DE 0.55	EVE 0.44	ARE 0.50
Q 0.14	RA 0.55	ITH 0.44	ON 0.47
Z 0.09	RO 0.55	TED 0.44	OR 0.45

Code Mapping

Code mapping can typically be easily cracked by analysing the probability of the mapped letters.



																											٦
Plain	a	b	С	d	е	f	g	h	i	j	k	1	m	n	0	р	q	r	s	t	u	v	W	x	У	Z	
1	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	Ν	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	А	
2	С	D	Ε	F	G	Н	Ι	J	K	L	М	Ν	0	Ρ	Q	R	S	Τ	U	V	M	Χ	Y	Z	Α	В	
3	D	Ε	F	G	Н	Ι	J	K	L	М	Ν	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	
4	Ε	F	G	Н	Ι	J	K	L	М	Ν	0	Р	Q	R	S	Т	U	V	W	Χ	Υ	Z	Α	В	С	D	
5	F	G	Н	Ι	J	K	L	М	Ν	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	
6	G	Н	I	J	K	L	Μ	Ν	0	Р	Q	R	S	Т	U	V	M	Χ	Y	Z	Α	В	С	D	Ε	F	
7	Н	I	J	K	L	М	N	0	Р	Q	R	S	Т	U	V	M	Х	Y	Z	Α	В	С	D	Ε	F	G	T
8	I	J	K	L	Μ	Ν	0	P	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	E	F	G	Н	
9	J	K	L	Μ	Ν	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	Ι	
10	K	L	Μ	Ν	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J	
11	L	Μ	Ν	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	Ι	J	K	
12	Μ	N	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	А	В	С	D	Ε	F	G	Н	I	J	K	L	
13	N	0	P	Q	R	S	Т	U	V	W	Х	Y	Z	A	В	С	D	Ε	F	G	Н	I	J	K	L	М	T
14	0	P	Q	R	S	Т	U	V	W	Χ	Y	Z	A	В	С	D	Ε	F	G	Н	I	J	K	L	М	N	T
15	Р	Q	R	S	Т	U	V	W	Х	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	N	0	
16	Q	R	S	Т	U	V	W	Х	Y	Ζ	Α	В	С	D	E	F	G	Н	Ι	J	K	L	Μ	Ν	0	Ρ	
17	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	М	N	0	Р	Q	
18	S	Т	U	V	W	Х	Y	Z	А	В	С	D	Ε	F	G	Н	I	J	K	L	М	N	0	Р	Q	R	T
19	Т	U	V	W	Х	Y	Z	А	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	N	0	Р	0	R	S	
20	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	N	0	Р	Q	R	S	Т	
21	V	W	Х	Y	Z	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	Ν	0	Р	0	R	S	Т	U	
22	W	Х	Y	Z	А	В	С	D	E	F	G	Н	I	J	K	L	Μ	Ν	0	Р	0	R	S	Т	U	V	
23	Х	Y			В	С		E	F		Н						N			0	R		Т	U	V	W	
24	Y	Z	A	В	С	D	Ē													~			U	V	W	Χ	
25			В																~								
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Moves the mapping depending on a keyword (in this case "GREEN")

Hello GREEN



Plain	a	b	С	d	e	f	g	h	i	j	k	1	m	n	0	р	q	r	s	t	u	v	w	x	У	z
1	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	Ν	0	Ρ	Q	R	S	Т	U	V	M	Χ	Y	Z	A
2	С	D	Ε	F	G	Η	Ι	J	K	L	Μ	Ν	0	Ρ	Q	R	S	Τ	U	V	M	Χ	Y	Z	Α	В
3	D	Ε	F	G	Н	I	J	K	L	М	Ν	0	Р	Q	R	S	Τ	U	V	M	Χ	Υ	Z	Α	В	С
4	Ε	F	G	Н	I	J	K	L	М	Ν	0	Р	Q	R	S	Т	U	\forall	M	Χ	Y	Z	Α	В	С	D
5	F	G	Н	Ι	J	K	L	М	N	0	Р	Q	R	S	Τ	U	V	W	Χ	Y	Ζ	Α	В	С	D	Ε
6	G	Н	Ι	J	K	L	М	N	0	Р	Q	R	S	Τ	U	V	W	Χ	Y	Ζ	Α	В	С	D	Ε	F
7	Η	I	J	K	L	Μ	N	0	Р	Q	R	S	Τ	U	V	W	Χ	Y	Ζ	Α	В	С	D	Ε	F	G
8	I	J	K	L	Μ	N	0	Р	Q	R	S	Т	U	V	M	Χ	Y	Ζ	Α	В	С	D	Ε	F	G	Н
9	J	K	L	Μ	Ν	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I
10	K	L	Μ	Ν	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J
11	L	Μ	N	0	Р	Q	R	S	Т	U	V	M	Χ	Υ	Z	Α	В	С	D	Ε	F	G	Н	I	J	K
12	Μ	Ν	0	P	Q	R	S	Т	U	V	W	Х	Y	Ζ	Α	В	С	D	Ε	F	G	Н	I	J	K	L
13	N	0	P	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	L	Μ
14	0	Ρ	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	E	F	G	Н	I	J	K	L	Μ	N
15	Р	Q	R	S	Т	U	V	W	Х	Y	Ζ	Α	В	С	D	E	F	G	Н	Ι	J	K	L	Μ	Ν	0
16	Q	R	S	Т	U	V	M	Х	Y	Z	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	Ν	0	Р
17	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	Ν	0	Р	Q
18	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	N	0	Р	Q	R
19	Т	U	V	W	Χ	Y	Ζ	Α	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	Ν	0	Р	Q	R	S
20	U	V	W	Χ	Y	Ζ	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	N	0	Р	Q	R	S	Т
21	V	M	Χ	Υ	Ζ	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	N	0	Ρ	Q	R	S	Т	U
22	M	Χ	Y	Ζ	Α	В	С	D	E	F	G	Н	Ι	J	K	L	Μ	N	0	Р	Q	R	S	Т	U	V
23	Χ	Y	Ζ	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	N	0	Р	Q	R	S	Т	U	V	W
24	Y	Z	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	N	0	Р	Q	R	S	Т	U	V	W	Χ
25	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	Ν	0	Ρ	Q	R	S	Т	U	V	W	Χ	Y

Moves the mapping depending on a keyword (in this case "GREEN")

Hello GREEN N



Plain	a	b	С	d	е	f	g	h	i	j	k	1	m	n	0	р	q	r	s	t	u	v	w	x	У	z
1	В	С	D	Ε	F	G	Н	I	J	K	L	М	N	0	Р	Q	R	S	Τ	U	V	M	Χ	Υ	Ζ	Α
2	С	D	Ε	F	G	Н	I	J	K	L	Μ	N	0	Ρ	Q	R	S	Τ	U	V	M	Χ	Y	Z	Α	В
3	D	Ε	F	G	Н	I	J	K	L	M	N	0	Р	Q	R	S	Т	U	V	M	Χ	Υ	Z	Α	В	С
4	E	F	G	Н	I	J	K	L	Μ	N	0	Р	Q	R	S	Т	U	V	M	Χ	Y	Z	Α	В	С	D
5	F	G	Н	Ι	J	K	L	М	Ν	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	E
6	G	Н	I	J	K	L	Μ	Ν	0	Р	Q	R	S	Τ	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F
7	Н	I	J	K	L	М	N	0	Р	Q	R	S	Т	U	V	M	Χ	Y	Z	Α	В	С	D	Ε	F	G
8	I	J	K	L	М	N	0	Р	Q	R	S	Τ	U	V	\mathbb{W}	Χ	Y	Ζ	Α	В	С	D	Ε	F	G	Н
9	J	K	L	Μ	N	0	Ρ	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	E	F	G	Н	I
10	K	L	Μ	N	0	Р	Q	R	S	Τ	U	V	M	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J
11	L	Μ	N	0	Р	Q	R	S	Т	U	V	M	Χ	Υ	Z	Α	В	С	D	Ε	F	G	Н	Ι	J	K
12	Μ	N	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Ζ	Α	В	С	D	Ε	F	G	Н	I	J	K	L
13	N	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	L	Μ
14	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Ζ	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	N
15	P	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	N	0
16	Q	R	S	Т	U	V	W	Χ	Y	Ζ	Α	В	С	D	E	F	G	Н	I	J	K	L	Μ	Ν	0	P
17	R	S	Т	U	v	W	Х	Y	Z	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	М	N	0	Р	Q
18	S	Т	U	V	W	Х	Y	Z	А	В	С	D	Ε	F	G	Н	I	J	K	L	М	N	0	Р	Q	R
19	Т	U	V	W	Х	Y	Z	Α	В	С	D	E	F	G	Н	I	J	K	L	Μ	N	0	P	Q	R	S
20	U	V	M	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	N	0	Р	Q	R	S	Т
21	V	W	Χ	Υ	Z	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	N	0	Р	Q	R	S	Т	U
22	W	Х	Y	Z	А	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	N	0	Р	Q	R	S	Т	U	V
23	Χ	Y	Ζ	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	N	0	Р	Q	R	S	Τ	U	V	W
24	Y	Ζ	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	M	N	0	Р	Q	R	S	Т	U	V	W	Χ
25	Z	Α	В	С	D	Е	F	G	Н	Ι	J	K	L	Μ	N	0	Р	Q	R	S	Τ	U	V	W	Χ	Y
					I	1												_								

Moves the mapping depending on a keyword (in this case "GREEN")

Hello GREEN NV



Plain	a	b	С	d	e	f	g	h	i	j	k	1	m	n	0	p	q	r	s	t	u	v	w	x	У	z	
1	В	С	D	Ε	F	G	Н	Ι	J	K	L	М	Ν	0	Р	Q	R	S	Т	U	V	M	Χ	Υ	Z	A	
2	С	D	Ε	F	G	Н	Ι	J	K	L	М	Ν	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	
3	D	Ε	F	G	Н	I	J	K	L	Μ	Ν	0	P	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	
4	Ε	F	G	Н	Ι	J	K	L	М	Ν	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	А	В	С	D	Ī
5	F	G	Н	Ι	J	K	L	М	N	0	Р	Q	R	S	Τ	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	Γ
6	G	Н	Ι	J	K	L	Μ	N	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	\mathbf{E}	F	
7	Н	Ι	J	K	L	Μ	N	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Ζ	Α	В	С	D	Ε	F	G	
8	I	J	K	L	Μ	N	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	
9	J	K	L	Μ	Ν	0	Р	Q	R	S	Τ	U	V	M	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	
10	K	L	Μ	N	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	Ι	J	
11	L	Μ	Ν	0	Р	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	
12	Μ	Ν	0	Р	Q	R	S	Т	U	V	W	Х	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	L	
13	N	0	Ρ	Q	R	S	Т	U	V	M	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	L	M	
14	0	Ρ	Q	R	S	Т	U	V	M	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	N	
15	Р	Q	R	S	Т	U	V	W	Χ	Y	Ζ	Α	В	С	D	Ε	F	G	Н	I	J	K	L	М	N	0	
16	Q	R	S	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	Ν	0	P	
17	R	S	Т	U	v	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	L	М	Ν	0	Р	Q	
18	S	Т	U	V	W	Χ	Y	Ζ	Α	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	N	0	Р	Q	R	
19	Т	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	N	0	P	Q	R	S	
20	U	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	Ν	0	Ρ	Q	R	S	T	
21	V	W	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	Ν	0	Р	Q	R	S	Т	U	
22	M	Χ	Y	Z	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	Ν	0	Ρ	Q	R	S	Т	U	V	
23	Χ	Y	Ζ	Α	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	N	0	Р	Q	R	S	Τ	U	V	M	
24	Y	Ζ	Α	В	С	D	Ε	F	G	Н	Ι	J	K	L	Μ	N	0	Р	Q	R	S	Т	U	V	W	Χ	
25	Z	Α	В	С	D	Ε	F	G	Н	I	J	K	L	Μ	N	0	Р	Q	R	S	Т	U	V	M	Χ	Y	

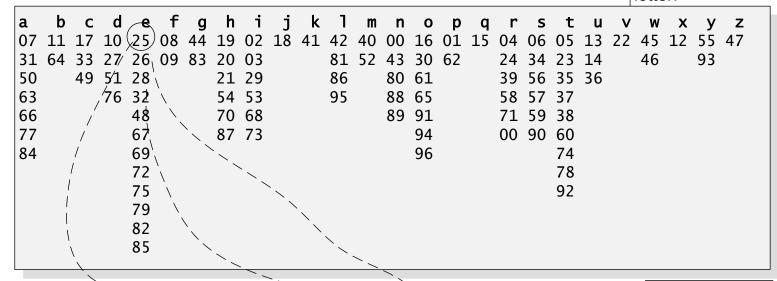
Moves the mapping depending on a keyword (in this case "GREEN")

Hello GREEN NVP



Homophonic substitution code

Number of codes varies with the probability of the letter.







Plaintext h e l l o e v e r y o n e Ciphertext: 19 25 42 81 16 26 22 28 04 55 30 00 32

Before electronic communications

Codes

A few fundamentals

Key-based encryption

Cracking the code

Brute force

Block or stream

Private-key methods

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Encrypting disks

PGP encryption

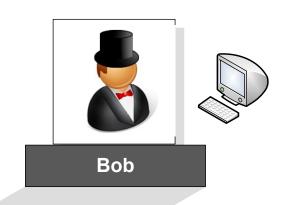








A few fundamentals



0101 1110 0010 0000 1110 0110 1010 1010

Bit stream

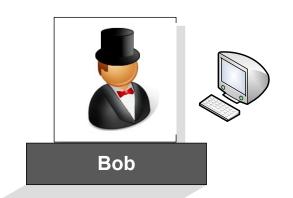
5 e 2 0 e 6 a a

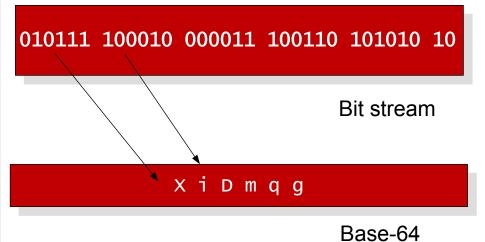
Hex

Viewing binary

With hexadecimal, the bit stream is split into groups of four, and converted into hex values (0-9,A-F)

	Decimal	Binary	Hex
	0	0000	0
	1	0001	1
	2	0010	2
	3	0011	3
L	4	0100	4
	5	0101	5
	6	0110	6
	7	0111	7
	8	1000	8
	9	1001	9
	10	1010	Α
	11	1011	В
	12	1100	С
	13	1101	D
	14	1110	Ε
	15	1111	F





Viewing binary

With Base-64, the bits are split into groups of six, and then converted. Base-64 is used extensively on the Internet (such as in email).

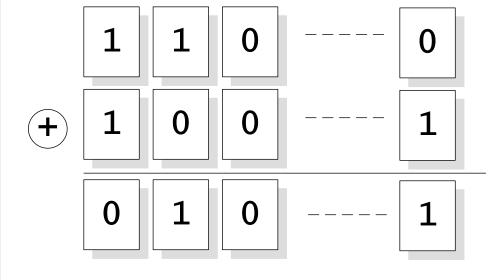
Val	Enc	Val	Enc	Val	Enc	Val	Enc
0	Α	16	Q	32	g	48	W
1	В	17	R	33	h	49	Χ
2	С	18	S	34	i	50	У
3	D	19	Τ	35	j	51	Z
4	Ε	20	U	36	k	52	0
5	F	21	V	37	1	53	1
6	G	22	W	38	m	54	2
7	Н	23	Χ	39	n	55	3
8	1	24	Υ	40	0	56	4
9	J	25	Z	41	p	57	5
10	K	26	а	42	q	58	6
11	L	27	b	43	r	59	7
12	M	28	С	44	S	60	8
13	Ν	29	d	45	t	61	9
14	0	30	е	46	u	62	+
15	Р	31	f	47	٧	63	/

Bob

The two main operators used in encryption are Ex-OR and ROR/ROL

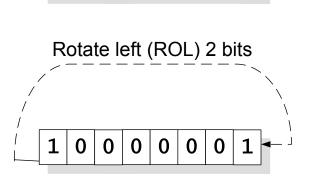
Encryption operators

The two main operators used in encryption are Ex-OR and ROL/ROR, as they are fast, and preserve info.



Exclusive-OR

operation



1

1

0 0

0

0

Rotate left (ROL)

Rotate right (ROR)

Before electronic communications

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Cracking the code

Brute force

Block or stream

Private-key methods

Encryption keys

Passing keys

Public-key encryption

One-way hash

Encrypting disks

PGP encryption

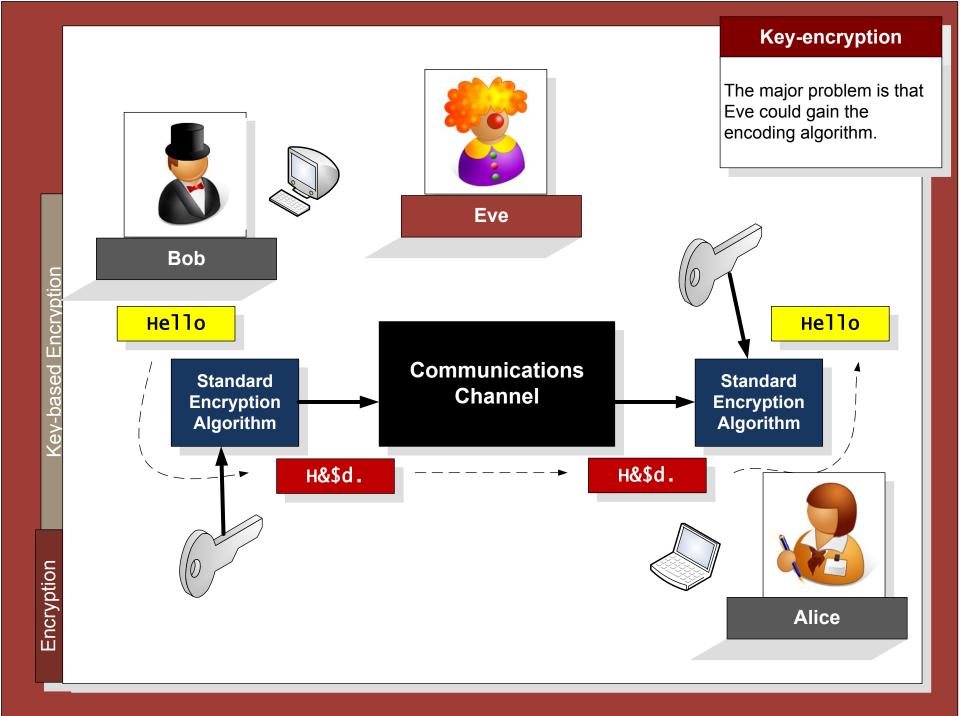


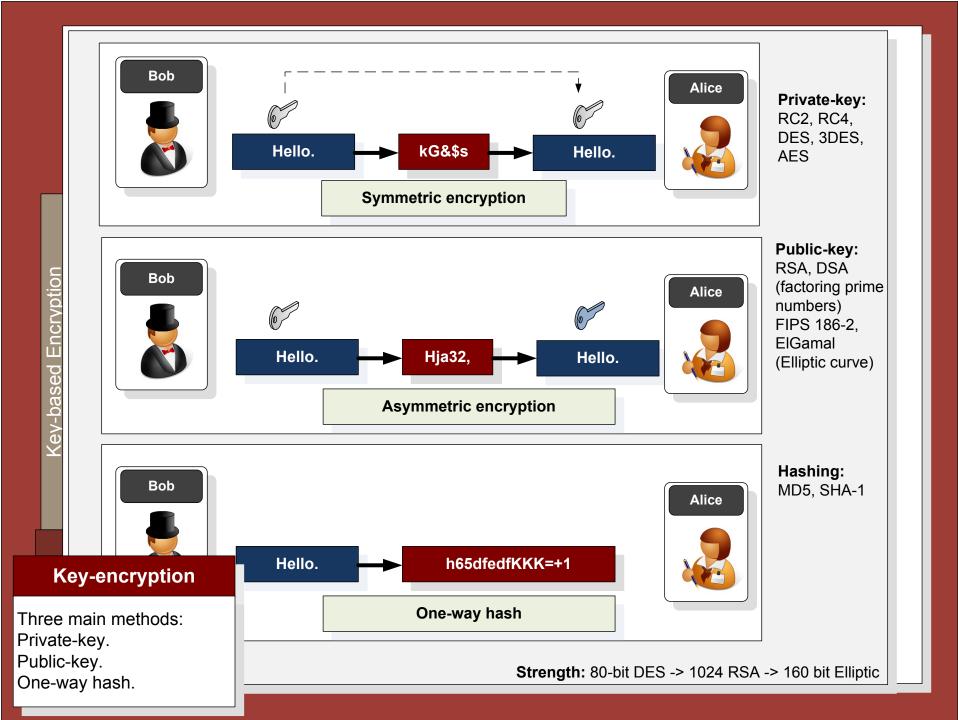






Key-based encryption







For example, if we have a key with four notches ... each which can exist or not ... how many keys can we have?

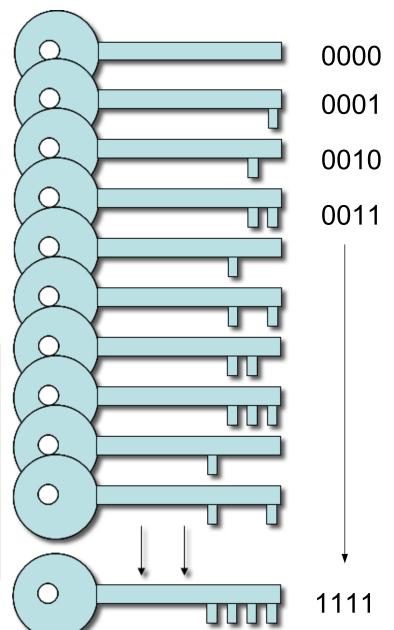
How safe is the key?

- the more keys ... the less likely it is to find the key.





16 key combinations
2 to the power of 4
(2⁴)



Width of Napier (100m)



Width of Edinburgh (6 miles)



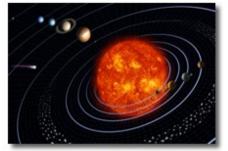
Earth to the Moon 93,000,000 miles

If each key was 1mm, and each key was laid end-on-end, what is the distance spanned for all the possible 64-bit electronic keys?





Width of the Milky Way 90,000 light years across



Width of the Solar System 3,666,000,000 miles

Width of Napier (100m)



Width of Edinburgh (6 miles)



somewhere between the Milky Way and the Universe

Size would be

Secret
Communications

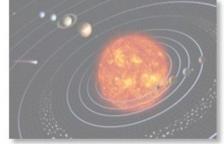
Earth to the Moon 93,000,000 miles



If each key was 1mm, and each key was laid end-on-end, what is the distance spanned for all the possible 64-bit electronic keys? (1,300,000,000,000,000 miles)



Width of the Milky Way 90,000 light years across



Width of the Solar System 3,666,000,000 miles

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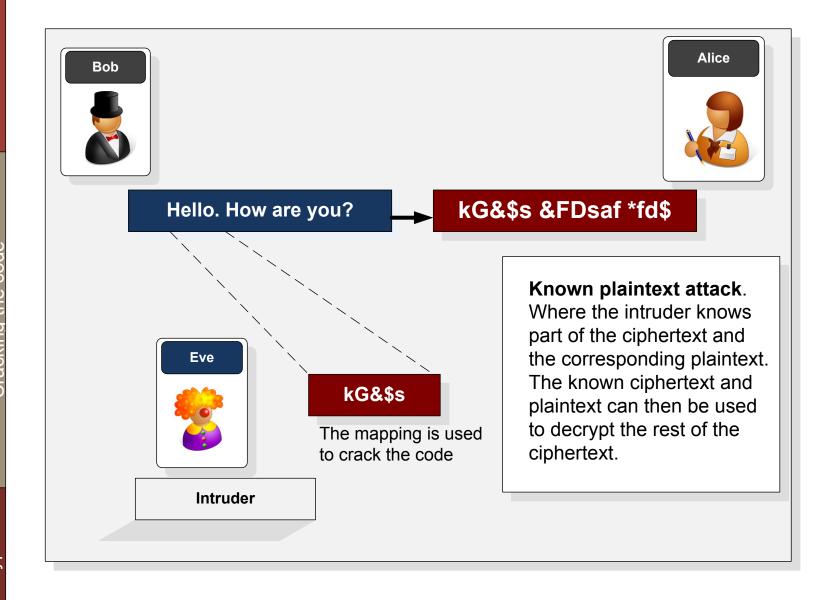


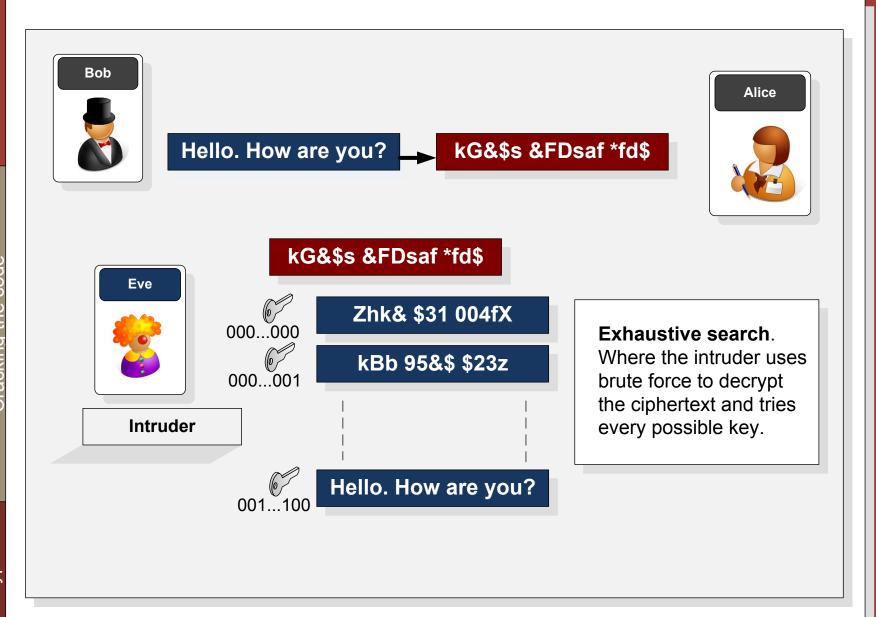


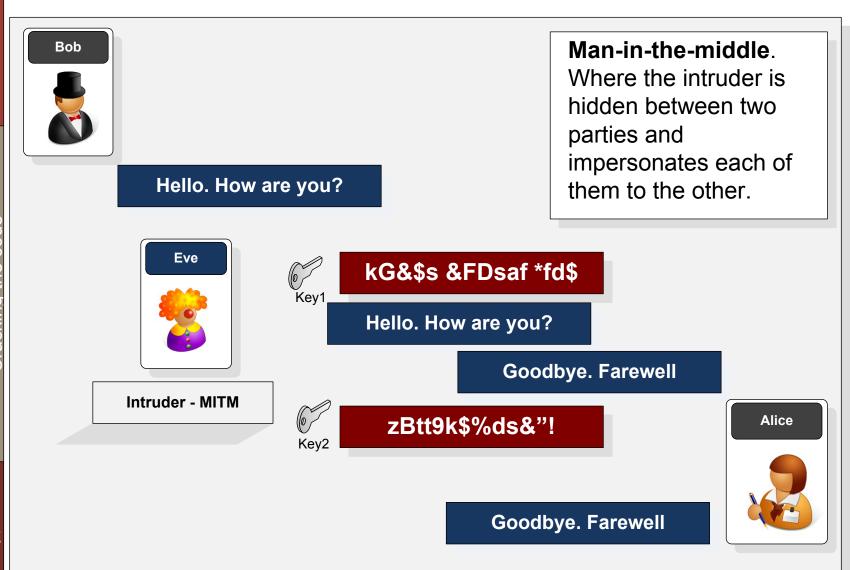


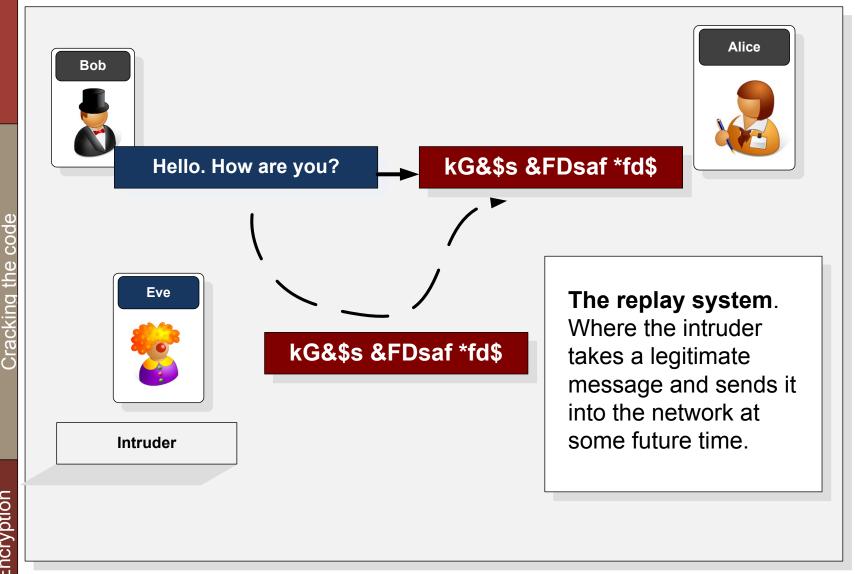


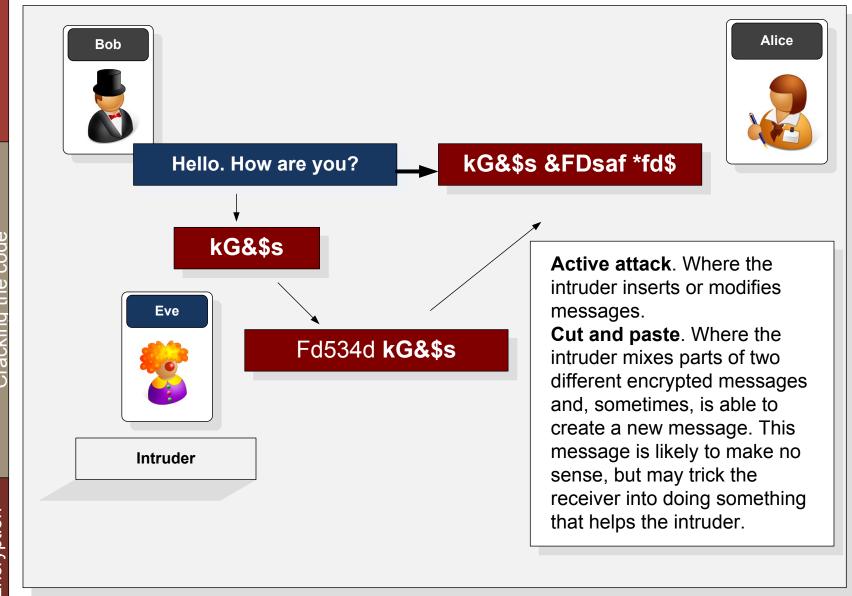
Cracking the code

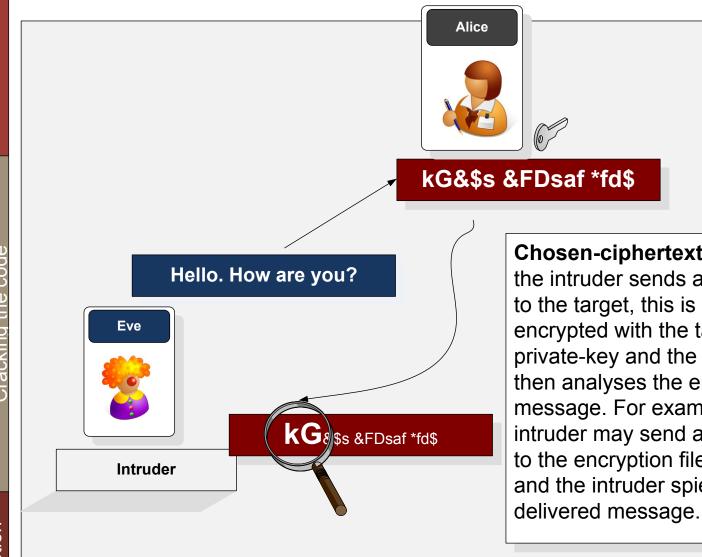












Chosen-ciphertext. Where the intruder sends a message to the target, this is then encrypted with the target's private-key and the intruder then analyses the encrypted message. For example, an intruder may send an e-mail to the encryption file server and the intruder spies on the

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Brute force





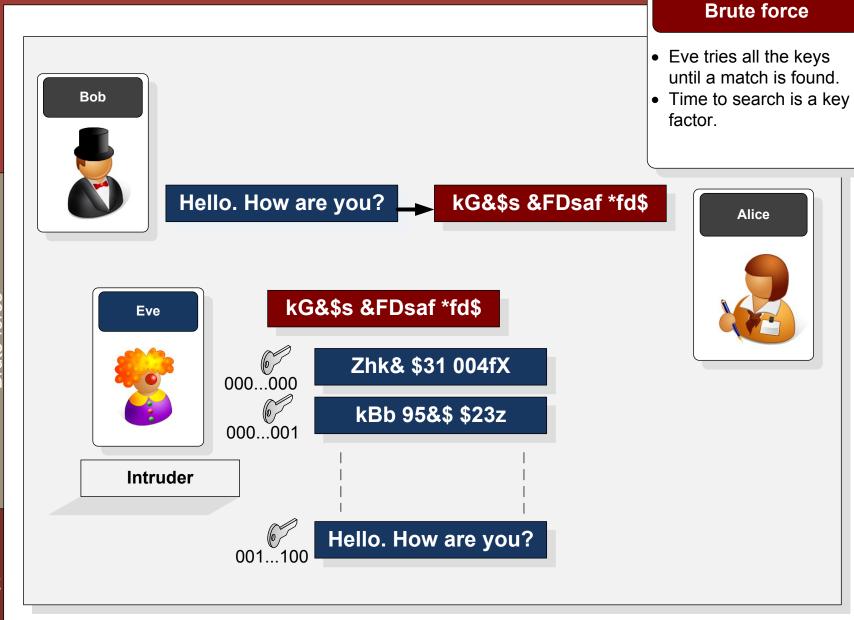




Number of keys

The larger the key, the greater the key space.

Code size	Number of keys	Code size	Number of keys	Code size	Number of keys
1	2	12	4,096	52	4.5×10 ¹⁵
2	4	16	65,536	56	7.21×10 ¹⁶
3	8	20	1,048,576	60	1.15×10 ¹⁸
4	16	24	16,777,216	64	1.84×10 ¹⁹
5	32	28	2.68×10 ⁸	68	2.95×10 ²⁰
6	64	32	4.29×10 ⁹	72	4.72×10 ²¹
7	128	36	6.87×10 ¹⁰	76	7.56×10 ²²
8	256	40	1.1×10 ¹²	80	1.21×10 ²⁴
9	512	44	1.76×10 ¹³	84	1.93×10 ²⁵
10	1024	48	2.81×10 ¹⁴	88	3.09×10^{26}





Okay... we select a **64-bit key** ... which has 1.84x10¹⁹ combinations

Time to crack

 It is important to understand the length of time that a message takes to crack as it may need to be secret for a certain time period.

How long will it take to cracked It by brute-force (on average)?



A 64-bit key has 1.84x10¹⁹ combinations and it could be cracked by brute-force in 0.9x10¹⁹ goes.

Time to crack

 It is important to understand the length of time that a message takes to crack as it may need to be secret for a certain time period.

If we use a fast computer such as 1GHz clock (1ns), and say it takes one clock cycle to test a code, the time to crack the code will be:

9,000,000,000 seconds (150 million minutes)

... **2.5 million hours** (285 years)



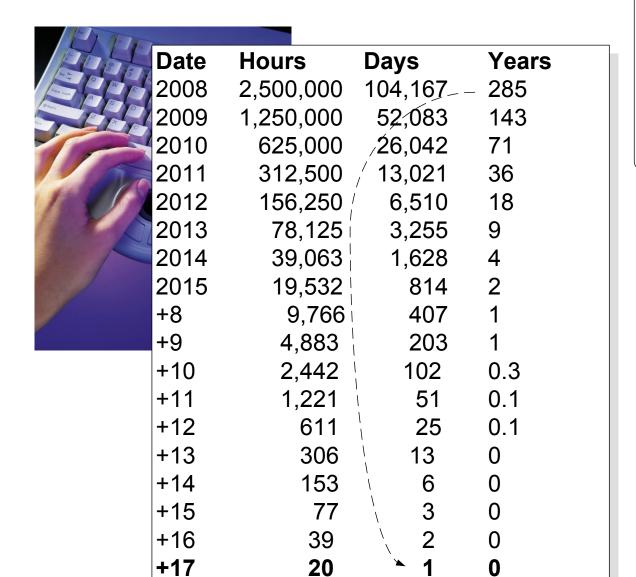
If it takes 2.5 million hours (285 years) to crack a code. How many years will it take to crack it within a day?

Time to crack

 It is important to understand the length of time that a message takes to crack as it may need to be secret for a certain time period.

Computers typically improve their performance every year ... so assume a **doubling** of performance each year.

Date	Hours	Days	Years
2008	2,500,000	104,167	285
2009	1,250,000	52,083	143



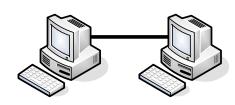
Time to crack

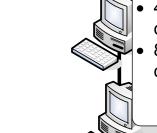
 From 285 years to 1 day, just by computers increasing their computing power.

56-bit DES:
Developed
1975
30 years ago!
... now easily
crackable

Parallel processing

2x1 = 2 element array

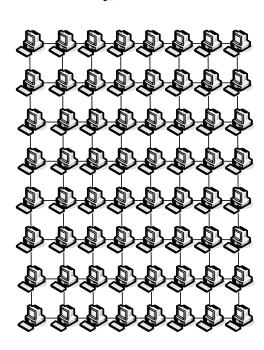




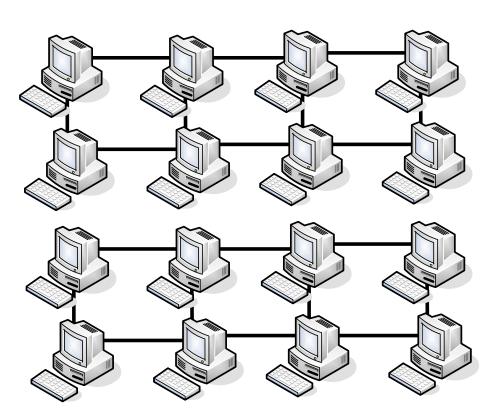
• 2x2 array = 4 computers.

 4x4 array = 16 computers.

8x8 array = 64 computers.



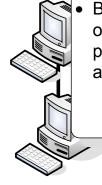
16x16 = 256 element array

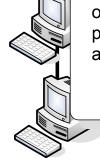


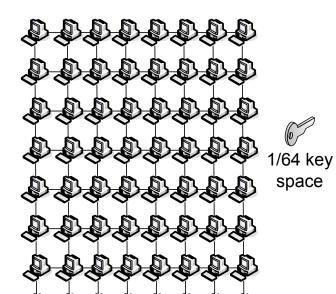
4x4 = 16 element array

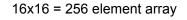
Parallel processing

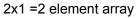
Brute-force cracking is one of the most scalable parallel processing applications.

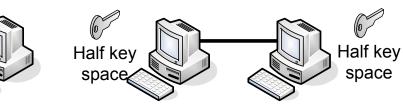


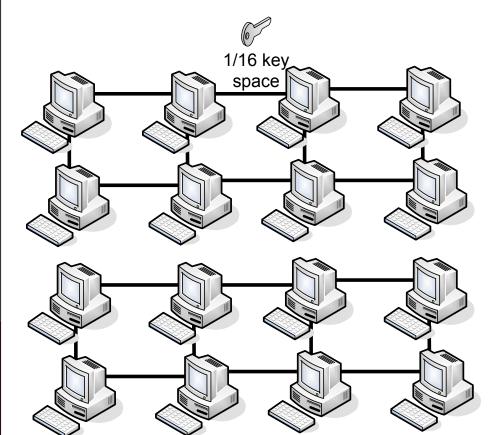






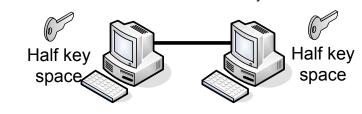






4x4 = 16 element array

2x1 = 2 element array



 64-bit key --- from 104,000 days (284 years) to one hour or less.



4x4 = 16 element array

 56-bit DES is seen as insecure as it can be cracked by enhanced

processors.





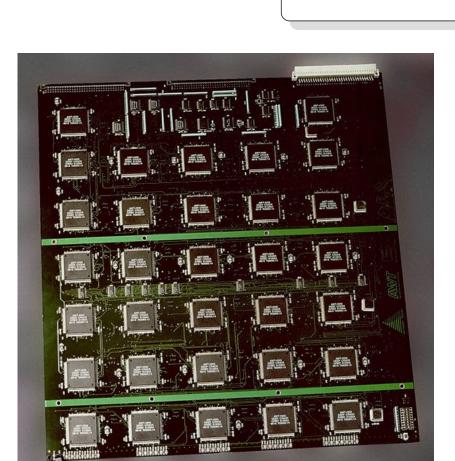
Year: 1998

Electronic Frontier
Foundation Cyberspace Civil
Rights Group
90,000,000 keys per
seconds

Array: 29 circuits of 64 chips

= 1856 elements

2.5 days



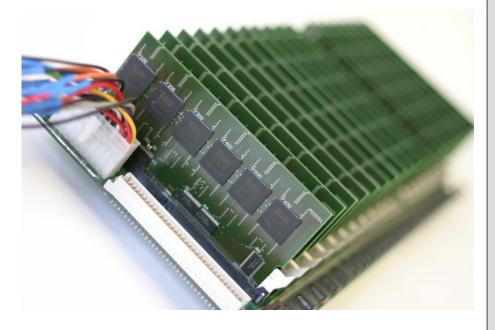


 Cracks 64-bit DES in less than nine days for less than \$10,000

Now

System: **COPACOBANA** (Cost-Optimized Parallel COde Breaker)

Time to crack: Less than 9 days for DES (64-bit code). Cost: Less than \$10,000



1997. RSA Lab's 56-bit RC5 Encryption Challenge - 250 days and 47% of the key space tested) – **distributed.net**



1998. RSA Lab's 56-bit DES II-1 Encryption Challenge - 39 days.

1998. RSA Lab's 56-bit DES II-2 Encryption Challenge - 2.5 days.

1999. RSA Lab's 56-bit DES-III Encryption Challenge - after 22.5 hours using EFF's Deep Crack custom DES cracker.



2002. RSA Lab's 64-bit RC5 Encryption Challenge — Completed 14 July 2002 – 1,757 days and 83% of the key space tested.

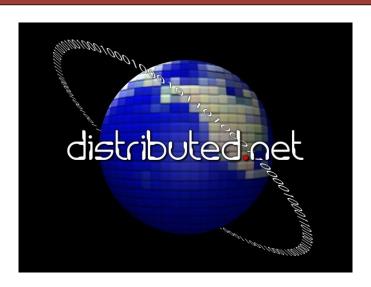
RSA Lab's 72-bit RC5 Encryption Challenge - In progress.



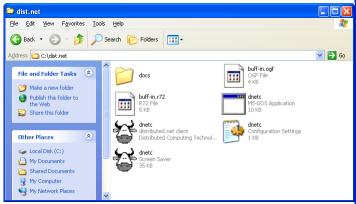
 RSA Labs have a number of challenges, each of which have been solved. The present challenge is 72-bit RC5.





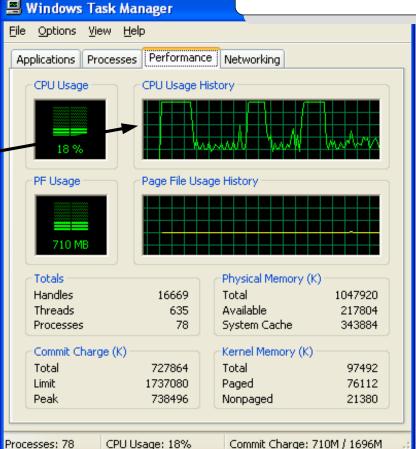


Distributed.net is starting and stopping (Max CPU when searching for possible keys)



distributed.net

- Tries to crack RSA Lab challenge by processing a range of possible keys while the screen save is on.
- Massive parallel processing system.





BlueGene/L – eServer Blue Gene Solution

DOE/NNSA/LLNL, IBM
Department of Energy's (DOE)
National Nuclear Security
Administration's (NNSA).
131,072 processors
367,000 Gigaflop= 367,000,000
Mflops

Super Computers

 BlueGene is 1.8million times more powerful than a standard PC.

Red Storm - Sandia/ Cray Red Storm

NNSA/Sandia National Laboratory United States, Opteron 2.4 GHz dual core Cray Inc.

26,544 processors 127,000 Gflops



Typical PC: 200 Mflop ... BlueGene is **1,835,000** times more powerful than a desktop.





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PGP encryption









Block or stream

Stream cipher

• Stream cipher (RC4)

110101 ...

1000 ...

Plaintext

Pseudoinfinite
key generate

Secret
key

Random

seed

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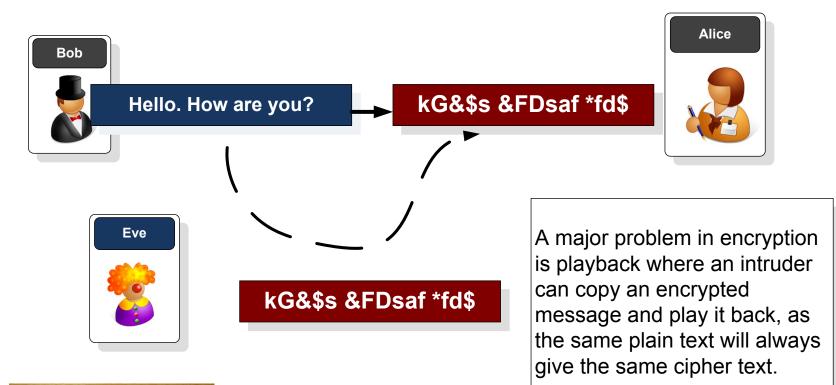








Private-key Methods





The solution is to add **salt** to the encryption key, as that it changes its operation from block-to-block (for block encryption) or data frame-todata frame (for stream encryption)



Block 1

- DES/3DES 64 bits
- RC2 64 bits
- AES/Rijndael 128 bits)

Block 2

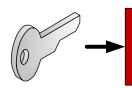
- DES/3DES 64 bits
- RC2 64 bits
- AES/Rijndael 128 bits)

Electronic Code Book (ECB) method. This is weak, as the

same cipher text appears for the same blocks.

Hello → 5ghd%43f=

Hello → 5ghd%43f=



Encrypted Block

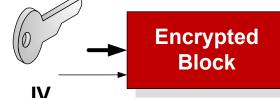
Encrypted Block

Block 1

- DES/3DES 64 bits
- RC2 64 bits
- AES/Rijndael 128 bits)

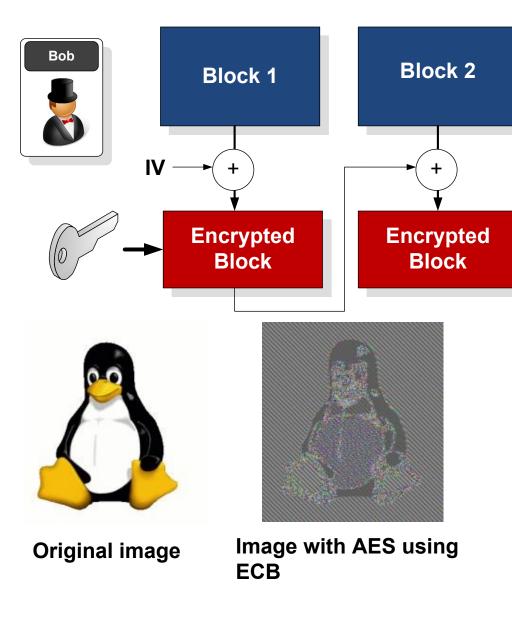
Block 2

- DES/3DES 64 bits
- RC2 64 bits
- AES/Rijndael 128 bits)



Encrypted Block

Adding salt. This is typically done with an IV (Initialisation Vector) which must be the same on both sides. In WEP, the IV is incremented for each data frame, so that the cipher text changes.



Cipher Block Chaining (CBC). This method uses the IV for the first block, and then the results from the previous block to encrypt the current block.

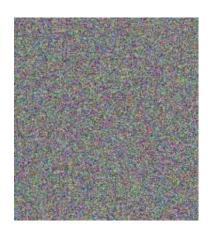


Image with AES using CBC

Image ref: http://en.wikipedia.org/wiki/Block_cipher_modes_of_operation

3-DES. The DES encryption algorithm uses a **64-bit block** and a 64-bit encryption key (of which only **56 bits** are actively used in the encryption process). Unfortunately DES has been around for a long time, and the 56-bit version is now easily crackable (in less than a day, on fairly modest equipment). An enhancement, and one which is still fairly compatible with DES, is the 3-DES algorithm. It has three phases, and splits the key into two. Overall the key size is typically **112 bits** (2x54 bits - with a combination of the three keys - of which two of the keys are typically the same). The algorithm is:

 $Encrypt_{K3}(Decrypt_{K2}(Encrypt_{K1}(message)))$

http://buchananweb.co.uk/security07.aspx

where K1 and K3 are typically the same (to keep compatibility).

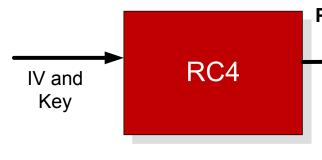
RC-2. RC2 ("Rivest Cipher") is seen as a replacement for DES. It was created by Ron Rivest in 1987, and is a **64-bit block code** and can have a key size from 40 bits to 128-bits (in increments of 8 bits). The 40-bit key version is seen as weak, as the encryption key is so small, but is favoured by governments for export purposes, as it can be easily cracked. In this case the key is created from a Key and an IV (Initialisation Vector). The key has 12 characters (96 bits), and the IV has 8 characters (64 bits), which go to make the overall key.

http://buchananweb.co.uk/security06.aspx

AES/Rijndael. AES (or Rijndael) is the new replacement for DES, and uses **128-bit blocks** with 128, 192 and 256 bit encryption keys. It was selected by NIST in 2001 (after a five year standardisation process). The name Rijndael comes from its Belgium creators: Joan Daemen and Vincent Rijmen.

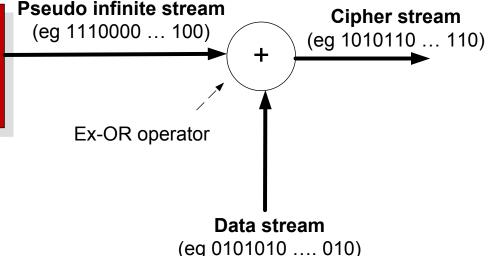
http://buchananweb.co.uk/security15.aspx

RC4. This is a **stream** encryption algorithm, and is used in wireless communications (such as in WEP) and SSL (Secure Sockets).



The IV (Initiation Vector) gives variation in the output for the same key





Data stream
Pseudo infinite stream
Cipher stream

0101010 ... 010 1110000 ... 100 + 1010110 ... 110 Introduction

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Encryption keys

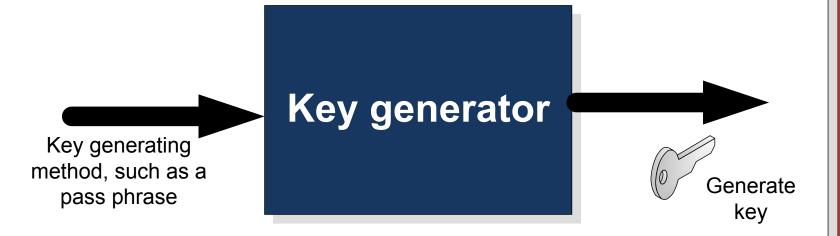
Key entropy: Relates to the equivalent number of bits given the range of phases used.

For example: if there were eight pass phrases – this would be equivalent to a 3-bit key.

Standard English gives 1.3 bits per character. Thus an **8 character word** gives **10.4 bits** for the key entropy.

Key enthropy

- 256 phrases -> 8 bit equivalent key.
- 1024 phases -> 10 bit equivalent key.
- 1,048,576 phrases -> 20 equivalent key.



Pass phrases might be: Napier, napier, napier1, napier11, napier123, and so on (the range of key will obviously be limited if the number of phrases are limited)

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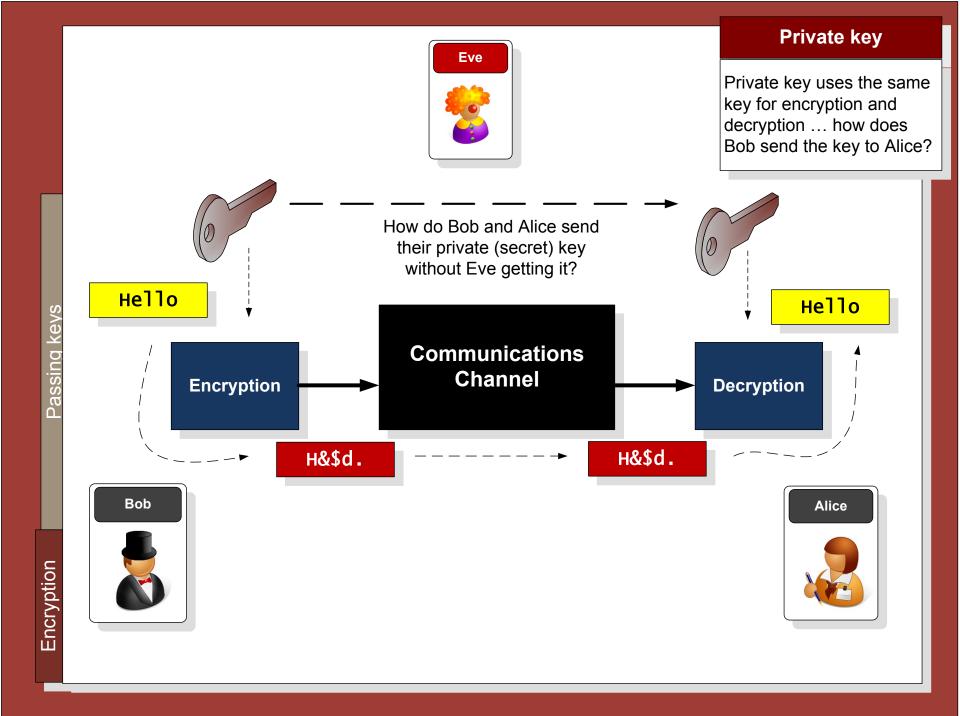




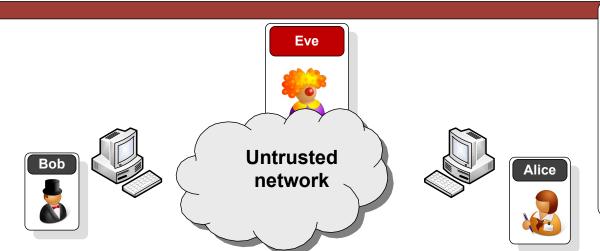




Passing keys



Diffie-Hellman



Diffie-Hellman

Eve can listen to the values of A and B, but should not be able to determine the secret key

- 1. Both nodes agree on two values (G and n)
- 2. Generate a random value (x)
- 2. Generate a random value (y)

3. $A = G^x \mod n$

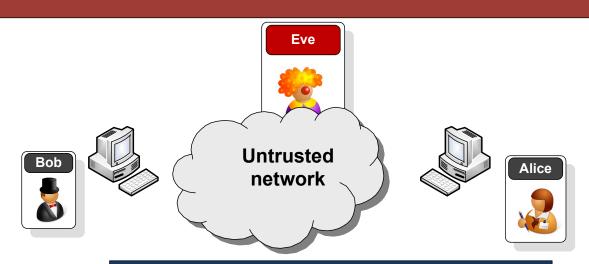
5. $K1 = B^x \mod n$

3. $B = G^y \mod n$



5. $K2 = A^y \mod n$

K1 and K2 should be the same and are the secret key



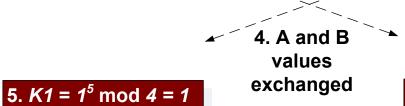
Diffie-Hellman

Eve can listen to the values of A and B, but should not be able to determine the secret key

- 1. Both nodes agree on two values (5 and 4)
- 2. Generate a random value (3)
- 2. Generate a random value (4)

3.
$$A = 5^3 \mod 4 = 5$$

3.
$$B = 5^4 \mod 4 = 1$$



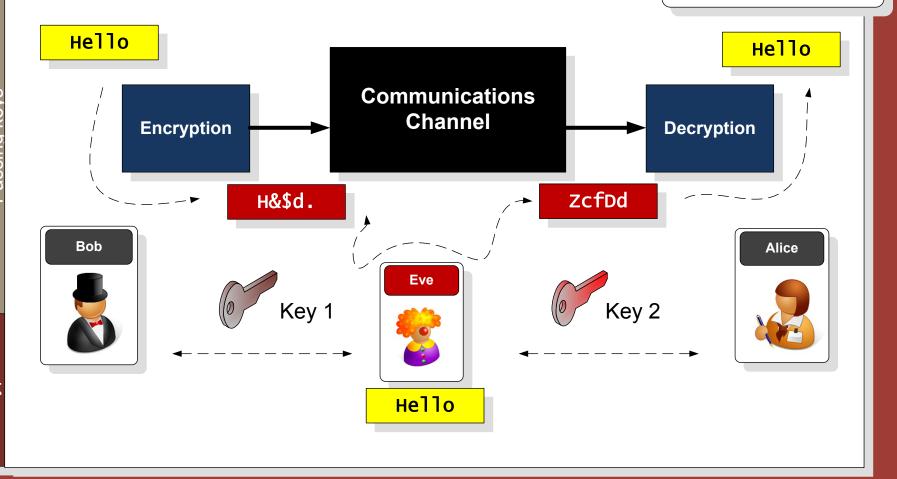
5.
$$K2 = 5^4 \mod 4 = 1$$

K1 and *K2* should be the **same** and are the secret key

Diffie-Hellman suffers from a man-in-themiddle attack, where Eve negotiates for each side, and creates two encryption channels

Man-in-the-middle

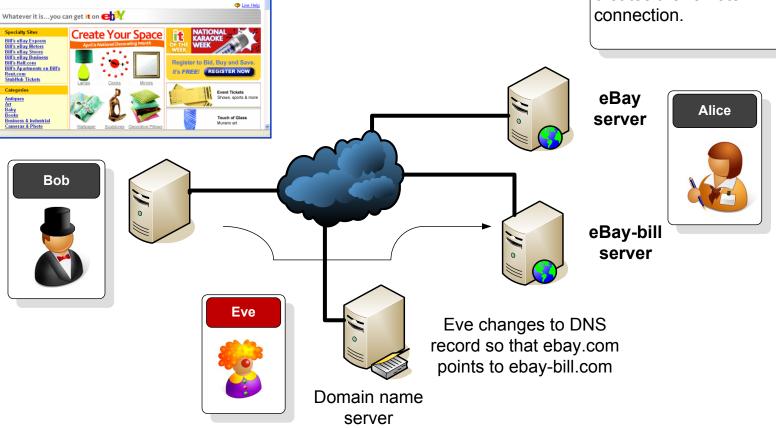
Diffie-Hellman suffers from Eve intercepting the key interchange, so that Bob thinks he's talking to Alice for the key exchange.



DNS poisoning A man-in-the-middle is

where Eve modifies the DNS, so that Bob things he is communicating with the remote server, but Eve which sign infort.

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Encrypting disks

PGP encryption









Public-key encryption

Public-key



Encryption/

Decryption

With Diffie-Hellman we need the other side to be active before we send data. Can we generate a special one-way function which allows is to distribute an encryption key, while we have the decryption key? RSA is still one of the most widely used encryption algorithms, and still stands up for secure communication, but is relatively slow in encrypting and decrypting.

Alice



Communications Channel

Encryption/ Decryption



Solved in 1977, By Ron Rivest, Adi Shamir, and Len Aldeman created the RSA algorithm for public-key encryption.



Select two prime numbers: a and b

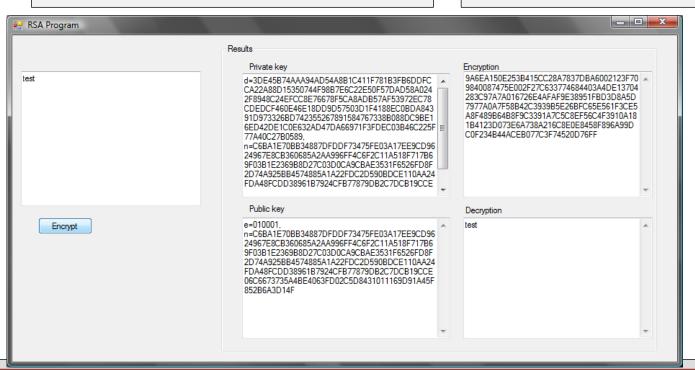
 $n = a \times b$

e is chosen so that e and (a-1)x(b-1) are relatively prime (no common factor greater than 1)

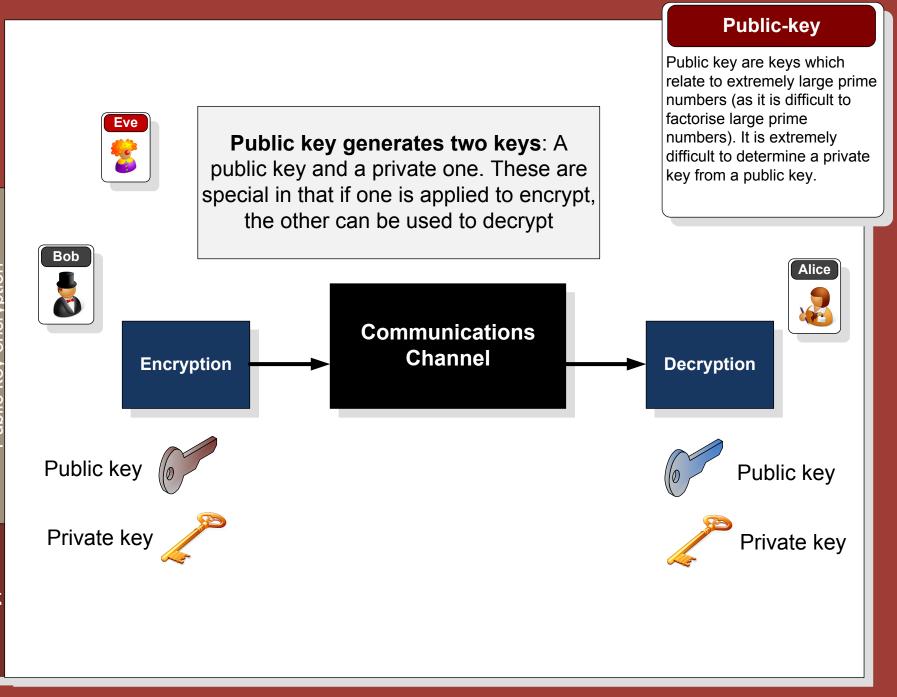
Public key is now: <e,n>

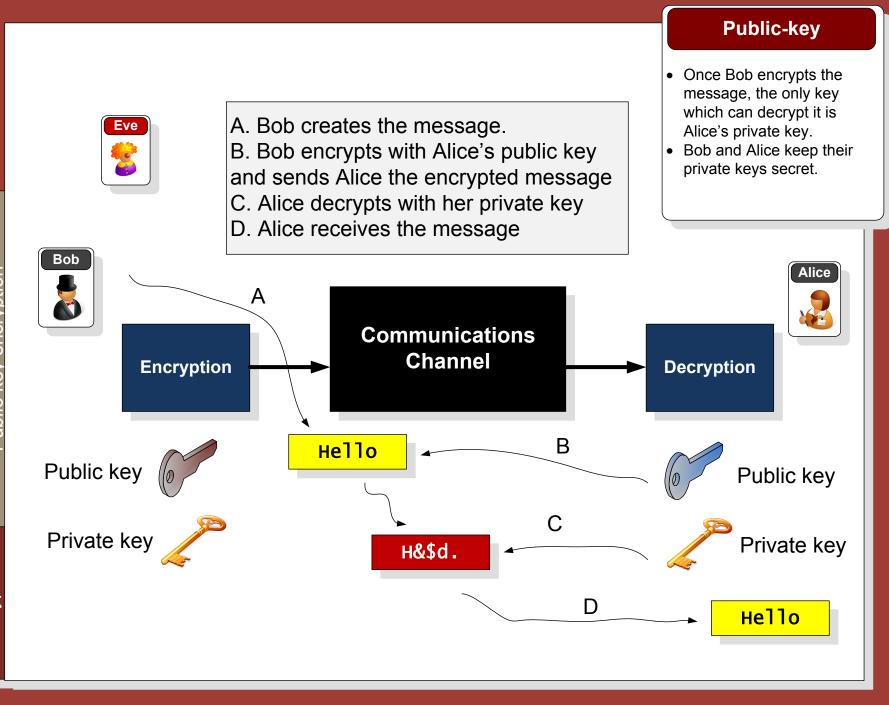
 $d = e-1 \mod [(a-1)x(b-1)]$

Private key is now: <d,n>



Author: Prof Bill Buchanan





Before electronic communications

Codes

A few fundamentals

Key-based encryption

Cracking the code

Brute force

Block or stream

Private-key methods

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Passing keys

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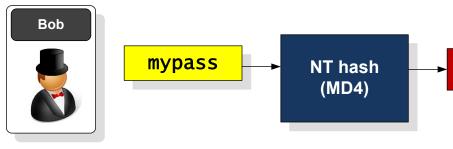






One-way hash

Windows login/ authentication



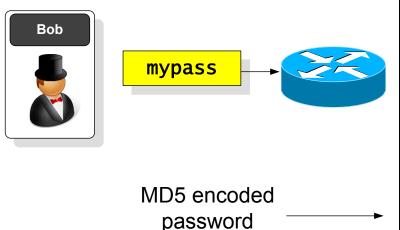
One-way hash

- Hashes are used for digital fingerprints (see the next unit) and for secure password storage.
- Typical methods are NT hash, MD4, MD5, and SHA-1.

fa1bfa14fa13fa12fa10fa1ffa14fa12

NT-password hash for Windows NT, XP and Vista

Cisco password storage (MD5)



```
# config t
(config)# enable secret test

Current configuration : 542 bytes
!
version 12.1
no service single-slot-reload-enable
service timestamps debug uptime
service timestamps log uptime
no service password-encryption
!
hostname Router
!
enable secret 5 $1$/Nwk$knsEQYxZVenGjWOGj/TGkO
```

A major factor with hash signatures is:

- **Collision.** This is where another match is found, no matter the similarity of the original message. This can be defined as a **Collision attack**.
- **Similar context**. This is where part of the message has some significance to the original, and generates the same hash signature. The can be defined as a Pre-image attack.
- **Full context**. This is where an alternative message is created with the same hash signature, and has a direct relation to the original message. This is an extension to a Pre-image attack.

In 2006 it was shown that MD5 can produce collision within less than a minute.

A 50% probability of a collision is:

$$\sqrt{N(signatures)} = \sqrt{2^n} = 2^{\frac{n}{2}}$$



where n is the number of bits in the signature. For example, for MD5 (128-bit) the number of operations that would be required for a better-than-50% chance of a collision is:

 2^{64}

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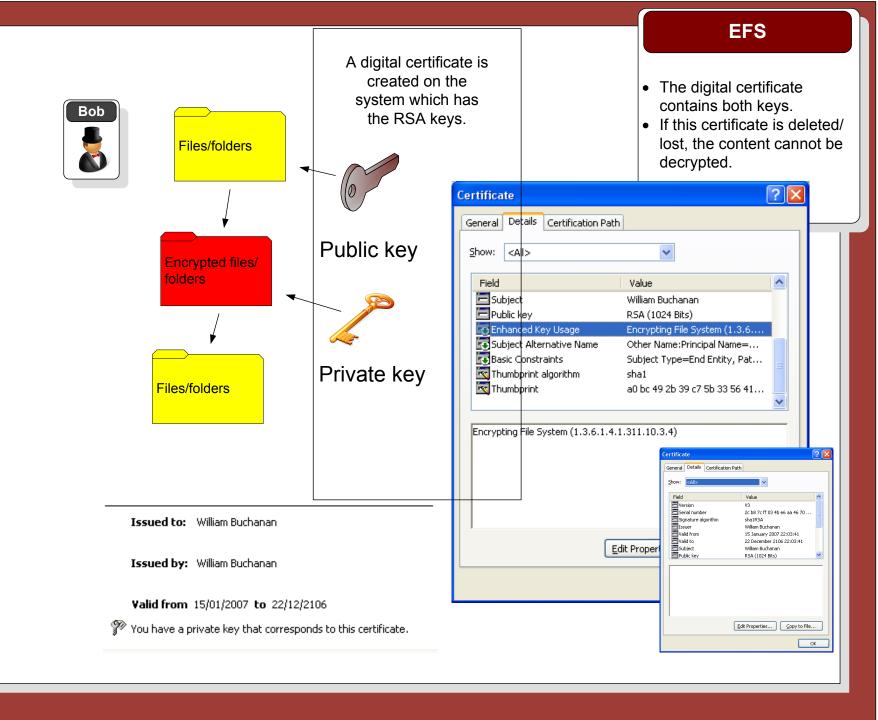


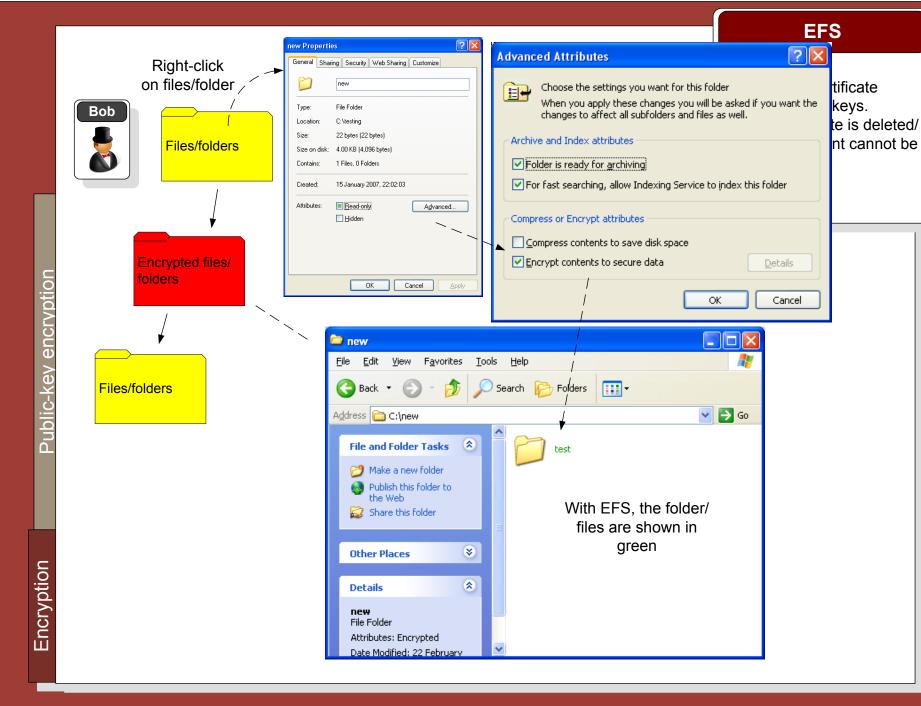


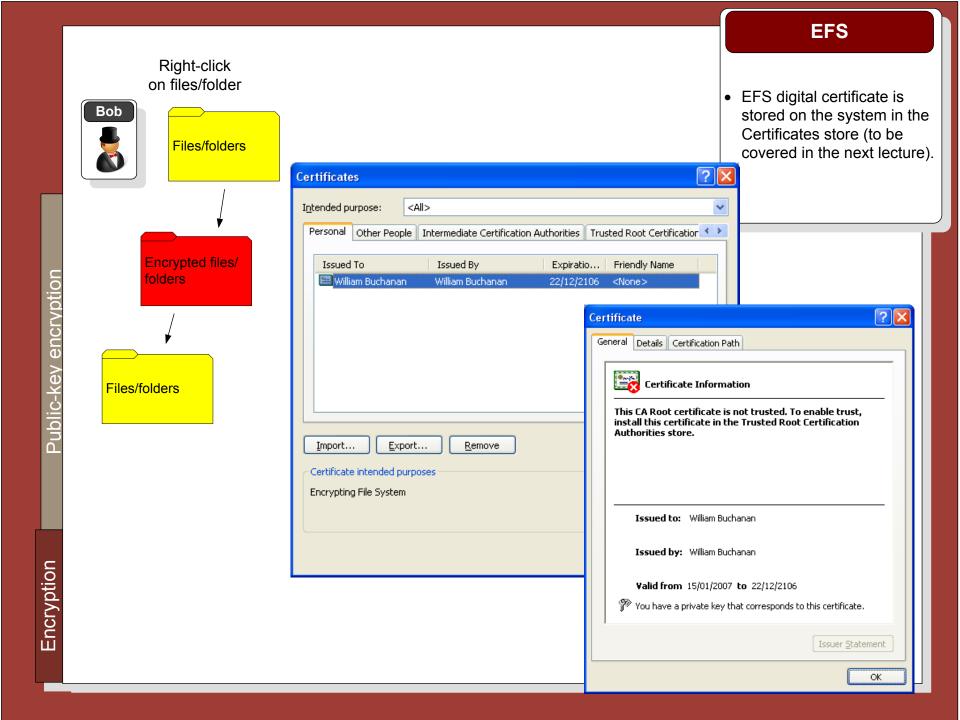




Encrypting disks







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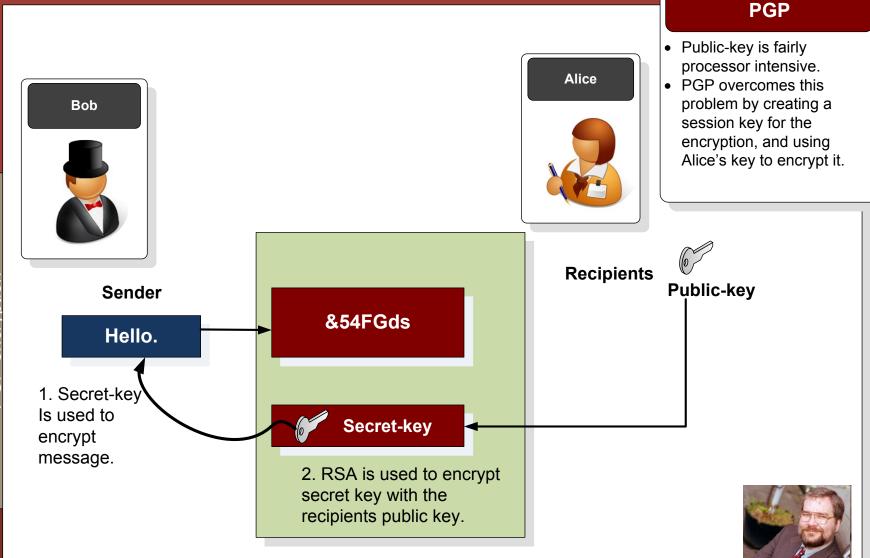








PGP encrypting



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Conclusions

