Comprehensive Report & Analysis (c-f)

c. How the system handles conflicts when jobs are put into a waiting queue, and there are jobs still entering the system. Which job goes first?

When jobs enter the waiting queue:

- Jobs that cannot be immediately allocated to memory are placed in a waiting queue (using Python's Queue data structure.
- · Each job records its queue entry time
- The gueue maintains FIFO (First In, First Out) ordering

How do we resolve conflict?

- 1. Jobs already in the waiting queue have priority over newly arriving jobs
- 2. When a memory block becomes free, the system:
 - First attempts to allocate jobs from the waiting queue before considering new arrivals
 - Processes gueued jobs in the order they entered the gueue
 - Only if a queued job cannot fit does it remain in the queue and the next queued job is checked

Which job goes first:

- · Jobs in the waiting queue are served before new incoming jobs
- Among waiting jobs: the job that entered the queue earliest gets priority (FCFS)
- If multiple jobs in the queue can fit in a freed block, the first one in the queue is selected
- New jobs entering the system only get allocated if no waiting jobs can fit the available memory

This prevents **starvation** and ensures fairness - a job that has been waiting will not be bypassed by newer arrivals.

```
FIXED PARTITION MEMORY ALLOCATION SIMULATION

RUNNING FIRST-FIT

[1-0.0] Job 1 allocated to Block 1 (wastes/37.0, wait-0.0) [1-0.0] Job 2 allocated to Block 2 (wastes/28.0, wait-0.0) [1-0.0] Job 2 allocated to Block 2 (wastes/28.0, wait-0.0) [1-0.0] Job 2 allocated to Block 2 (wastes/28.0, wait-0.0) [1-0.0] Job 3 allocated to Block 3 (wastes/28.0, wait-0.0) [1-0.0] Job 5 allocated to Block 5 (wastes/28.0, wait-0.0) [1-0.0] Job 5 allocated to Block 6 (wastes/28.0, wait-0.0) [1-0.0] Job 7 queued (queue size: 1) [1-0.0] Job 9 allocated to Block 6 (wastes/28.0, wait-0.0) [1-0.0] Job 9 allocated to Block 8 (wastes/28.0, wait-0.0) [1-0.0] Job 9 allocated to Block 8 (wastes/28.0, wait-0.0) [1-0.0] Job 9 allocated to Block 8 (wastes/28.0, wait-0.0) [1-0.0] Job 13 queued (queue size: 3) [1-0.0] Job 13 queued (queue size: 3) [1-0.0] Job 13 allocated to Block 9 (wastes/28.0, wait-0.0) [1-0.0] Job 15 allocated to Block 9 (wastes/28.0, wait-0.0) [1-0.0] Job 15 allocated to Block 9 (wastes/28.0, wait-0.0) [1-0.0] Job 15 allocated to Block 9 (wastes/28.0, wait-0.0) [1-0.0] Job 15 queued (queue size: 10) [1-0.0] Job 15 queued (queue size: 11) [1-0.0] Job 15 queued (queue size: 12) [1-0.0] Job 15 queued (queue size: 13) [1-0.0] Job 15 queued (queue size: 14) [1-0.0] Job 15 queued (queue size: 15) [1-0.0] Job 15 queued (queue size: 14) [1-0.0] Job 15 queued (queue size: 14) [1-0.0] Job 15 queued (queue size: 14) [1-0.0] Job 15 queued (queue size: 15) [1-0.0] Job 15 queued (queue size: 14) [1-0.0] Job 15 queued (queue size: 14) [1-0.0] Job 15 queued (queue size: 15) [1-0.0] Job 15 queued (queue size: 14) [1-0.0] Job 15 queued (queue size: 15) [1-0.0] Job 15 queued (queue size: 15) [1-0.0] Job 15 queued (queue size: 15) [1-0
```

d. How the system handles the "job clocks" and "wait clocks"

Our system uses SimPy's discrete-event simulation framework to handle timing: Job Clocks (Execution Time Tracking):

- Each job has a 'time' attribute representing CPU execution duration
- When a job is allocated:
 - SimPy automatically advances the simulation clock by this amount
 - The job remains in memory for its full execution time before deallocation
 - Current simulation time is tracked

Wait Clocks (Queue Wait Time Tracking):

- When a job enters the waiting queue, its exact simulation time is recorded
- When the job is finally allocated from the queue:

```
wait_time = self.env.now - job['queue_entry_time']job['wait_time'] = wait_time
```

This line of code above calculates the total time spent waiting

 Additionally, we have a class that accumulates and calculates total waiting time for all jobs: MemorySimulatorMetrics

e. How you define "event" and what happens when the event occurs

In our **event-driven simulation**, an event is any occurrence that changes the system state. Our implementation defines these key events:

Event Types:

- 1. Job Arrival Event
- 2. Job Allocation Event
- 3. Job Completion Event
- 4. Job enters Waiting Queue

We calculate and track the following:

- Memory block states (free ↔ occupied)
- Job states (waiting → queued → running → completed)
- Metrics (throughput, wait times, queue sizes)

f. Results comparison: First-Fit vs Best-Fit - Analysis and Recommendations

Based on our code implementation, we observed the following:

```
COMPREHENSIVE COMPARISON REPORT
_____
                                                      _____
1. THROUGHPUT (Jobs per time unit)
   First-Fit: 1.0435 jobs/time unit
                 1.0000 jobs/time unit
   Winner: First-Fit (by 4.35%)
2. STORAGE UTILIZATION
   Blocks Never Used:
      First-Fit: 0 (0.0%)
Best-Fit: 0 (0.0%)
   Average Block Usage:
First-Fit: 2.40 times
       Best-Fit: 2.40 times
       First-Fit: 4 times
Best-Fit: 4 times
3. WAITING QUEUE METRICS
   Maximum Queue Length:
       First-Fit: 15 jobs
       Best-Fit: 15 jobs
   Average Queue Length:
       First-Fit: 8.00 jobs
Best-Fit: 8.00 jobs
   Jobs That Required Queuing:
       First-Fit: 15 jobs
Best-Fit: 15 jobs
4. WAITING TIME IN QUEUE
   Average Wait Time:
       First-Fit: 7.21 time units
Best-Fit: 7.43 time units
   Maximum Wait Time:
       First-Fit: 15.00 time units
Best-Fit: 16.00 time units
   Winner: First-Fit (lower by 0.21 time units)
5. INTERNAL FRAGMENTATION
   Average Internal Fragmentation per Job:
       First-Fit: 1621.67 memory units
       Best-Fit: 1517.50 memory units
   Total Internal Fragmentation:
       First-Fit: 38920 memory units
   Best-Fit: 36420 memory units Winner: Best-Fit (reduced fragmentation by 6.42%)
6. OVERALL PERFORMANCE SUMMARY
   Average Turnaround Time:
       First-Fit: 10.29 time units
Best-Fit: 10.42 time units
   Simulation Duration:
       First-Fit: 23.00 time units
Best-Fit: 24.00 time units
   Job Completion Rate:
       Best-Fit: 96.0%
```

First-Fit Advantages

First-Fit proved faster at allocation time because it stopped at the first available block.

Even though it still has an **O(n)** time complexity, it filled memory quickly, which led to **higher internal fragmentation**.

Some blocks got used repeatedly, while others stayed idle, leaving many jobs waiting in the queue for earlier ones to finish. Overall, First-Fit is **better suited for systems that prioritize speed over optimization**, especially when quick allocations are more important than perfect memory utilization.

Best-Fit Advantages

Best-Fit performed better in terms of **memory space efficiency and utilization**.

It also shares the same **O(n)** time complexity since it has to search the entire list, but it resulted in **much less internal fragmentation** compared to First-Fit.

However, the need to find the "most optimal" block increased its overhead, slightly reducing throughput.

In short, Best-Fit is **slower but more space-efficient**, making it a better choice when conserving memory is the main priority.

Recommendations

Use First-Fit when:

- Real-time or near-instant response is critical
- Memory is relatively abundant compared to job sizes
- Job sizes are mostly uniform
- The system needs to handle a high rate of job arrivals efficiently

Use Best-Fit when:

- Memory is limited and must be managed carefully
- Job sizes vary widely
- Long-term efficiency and stability are more important than immediate response
- Minimizing internal fragmentation is a key performance goal

Final Thoughts

This analysis doesn't hold universally for all systems.

While **Best-Fit** tends to produce **less fragmentation**, **First-Fit** is generally **faster and simpler** — though that speed can come at the cost of wasted memory over time.

Both algorithms have trade-offs depending on system goals and workload patterns, but for the given simulation, **First-Fit performed slightly better overall (3/5 vs 2/5)** in throughput and waiting time.