

Question 1:

(a) $f(n) = 6n + 3n^2 + n \log n + 3\sqrt{n}$

the order of $f(n)$ is $O(n^2)$

(b) prove: there is a constant c

such that $6n + 3n^2 + n \log n + 3\sqrt{n} \leq cn^2, \forall n \geq n_0$

$\because \log n \leq n \quad \therefore n \log n \leq n^2$

$3\sqrt{n} \leq 3n^2 \quad 6n \leq 6n^2$

$\therefore 6n + 3n^2 + n \log n + 3\sqrt{n} \leq 6n^2 + 3n^2 + n^2 + 3n^2 = 13n^2$

$\therefore f(n) \leq 13n^2, \forall n \geq n_0$

Question 2:

(a) Procedure $F(n)$

if $n == 0$ or $n == 1$ or $n == 2$ or $n == 3$ then

return 1

else

return $F(n-1) + F(n-2) + F(n-3)$

(b) The time complexity of this algorithm is exponential $O(3^n)$

\because for each value equal or greater than 4, the function make 3 recursive calls to $n-1, n-2, n-3$,
and each of these calls makes 3 more recursive calls and so on.

So the number of function calls grow exponentially with an exponential time complexity.

Question 3:

(a) the number of swapping operation is 3

(b) the number of key comparisons is 15

Question 4:

(a) a b c d e f

a 0 1 0 1 0 0

b 1 0 1 0 1 0

c 0 1 0 0 1 1

d 1 0 0 0 1 0

e 0 1 1 1 0 1

f 0 0 1 0 1 0

$a \rightarrow b \rightarrow d$

$b \rightarrow a \rightarrow c \rightarrow e$

$c \rightarrow b \rightarrow e \rightarrow f$

$d \rightarrow a \rightarrow e$

$e \rightarrow b \rightarrow d \rightarrow c \rightarrow f$

$f \rightarrow e \rightarrow c$

(b) a b c d e f

1 1 1 0 0 0

2 1 0 0 1 0

3 0 1 1 0 0

4 0 1 0 0 1

5 0 0 1 0 1

6 0 0 1 0 0

7 0 0 0 1 1

8 0 0 0 0 1

$1 \rightarrow a \rightarrow b$

$2 \rightarrow a \rightarrow d$

$3 \rightarrow b \rightarrow c$

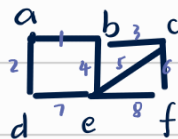
$4 \rightarrow b \rightarrow e$

$5 \rightarrow c \rightarrow e$

$6 \rightarrow c \rightarrow f$

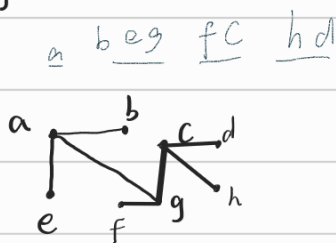
$7 \rightarrow d \rightarrow e$

$8 \rightarrow e \rightarrow f$

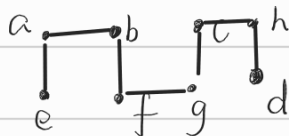


Question 5

(a)



(b)



Question 6 → 连通-个整体, 辐射-片

(a) order selected a(-, -) b(-, ∞) c(-, ∞) d(-, ∞) e(-, ∞) f(-, ∞) g(-, ∞) h(-, ∞)

a(0, -) b(a, 4) c(-, ∞) d(-, ∞) e(a, 1) f(-, ∞) g(-, ∞) h(-, ∞)

e(a, 1) b(a, 4) c(-, ∞) d(-, ∞) f(e, 2) g(-, ∞) h(-, ∞)

f(e, 2) b(f, 3) c(f, 10) d(-, ∞) g(f, 5) h(-, ∞)

b(f, 3) c(b, 6) d(-, ∞) g(f, 5) h(-, ∞)

g(f, 5) c(g, 12) d(g, 14) g(f, 5) h(-, ∞)

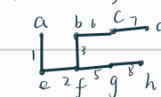
c(g, 12) d(c, 7) h(-, ∞)

d(c, 7) h(d, 8)

h(d, 8)



not the only MST



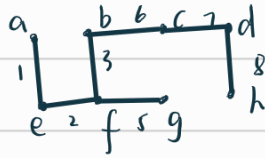
→ here is another

↑ height

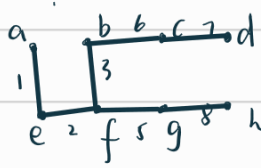
(b)

(a, e)	1
(e, f)	2
(f, b)	3
(a, b)	4
(f, g)	5
(b, c)	6
(c, d)	7
(d, h)	8
(g, h)	8
(f, c)	10
(g, c)	12
(g, d)	14

MST:

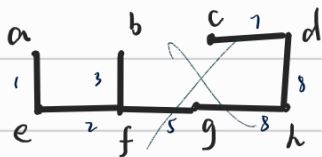


MST is not unique:



(c)

order selected	a (-, -)	b (-, ∞)	c (-, ∞)	d (-, ∞)	e (-, ∞)	f (-, ∞)	g (-, ∞)	h (-, ∞)
a (-, -)		b(a, 4)	c(-, ∞)	d(-, ∞)	e(a, 1)	f(-, ∞)	g(-, ∞)	h(-, ∞)
e(a, 1)		b(a, 4)	c(-, ∞)	d(-, ∞)		f(e, 3)	g(-, ∞)	h(-, ∞)
f(e, 3)		b(a, 4)	c ^(f, 3) (f, 3)	d(-, ∞)			g(f, 8)	h(-, ∞)
b(a, 4)			c(b, 10)	d(-, ∞)			g(f, 8)	h(-, ∞)
g(f, 8)			c ^(b, 10) (g, 20)	d ^(g, 22) (g, 22)				h(g, 16)
h(g, 16)			c(-, ∞)	d ^(c, 17) (h, 24)				h(g, 16)
h(g, 16)			c(-, ∞)	d(c, 17)				
d(h, 24)			c(d, 31)					
c(d, 31)								



$$\Rightarrow \frac{|E| \times \log |V|}{|E| \times \log |V|}$$

a → e → f → g → h
倒着

每行选又前一行最大值，重新计算各点权值

