

A specification of the algorithm described in *Paxos Made Simple*. This specification is a modification of: <https://lamport.azurewebsites.net/tla/PConProof.tla> Look there for comments.

EXTENDS *Integers, FiniteSets*

CONSTANT *Value, Acceptor, Quorum*

ASSUME  $QA \triangleq \wedge \forall Q \in Quorum : Q \subseteq Acceptor$   
 $\wedge \forall Q1, Q2 \in Quorum : Q1 \cap Q2 \neq \{\}$

*Ballot*  $\triangleq Nat$

ASSUME *BallotAssump*  $\triangleq (Ballot \cup \{-1\}) \cap Acceptor = \{\}$

*None*  $\triangleq \text{CHOOSE } v : v \notin Value$

*Message*  $\triangleq$   
 $\cup [type : \{"1a"\}, bal : Ballot]$   
 $\cup [type : \{"1b"\}, acc : Acceptor, bal : Ballot,$   
 $mbal : Ballot \cup \{-1\}, mval : Value \cup \{None\}]$   
 $\cup [type : \{"2a"\}, bal : Ballot, val : Value]$   
 $\cup [type : \{"2b"\}, acc : Acceptor, bal : Ballot, val : Value]$

```
--algorithm PCon{
variables maxBal = [a ∈ Acceptor ↦ -1],
          maxVbal = [a ∈ Acceptor ↦ -1],
          maxVVal = [a ∈ Acceptor ↦ None],
          msgs = {}

define {
  sentMsgs(t, b)  $\triangleq \{m \in msgs : (m.type = t) \wedge (m.bal = b)\}$ 

  ShowsSafeAt(Q, b, v)  $\triangleq$ 
    LET Q1b  $\triangleq \{m \in sentMsgs("1b", b) : m.acc \in Q\}$ 
    IN  $\wedge \forall a \in Q : \exists m \in Q1b : m.acc = a$ 
       $\wedge \forall m \in Q1b : m.mbal = -1$ 
         $\vee \exists m \in Q1b :$ 
           $\wedge m.mval = v$ 
           $\wedge \forall m1 \in Q1b : m1.mbal \leq m.mbal$ 
}

macro Phase1a() { msgs := msgs  $\cup \{[type \mapsto "1a", bal \mapsto self]\}$ ; }

macro Phase1b(b) {
  when b > maxBal[self]  $\wedge sentMsgs("1a", b) \neq \{\}$ ;
  maxBal[self] := b;
  msgs := msgs  $\cup \{[type \mapsto "1b", acc \mapsto self, bal \mapsto b,$ 
    mbal  $\mapsto maxVbal[self], mval \mapsto maxVVal[self]]\}$ ;
}
}
```

```

macro Phase2a(v){
  when  $\wedge \text{sentMsgs}(\text{"2a"}, \text{self}) = \{\}$ 
     $\wedge \exists Q \in \text{Quorum} : \text{ShowsSafeAt}(Q, \text{self}, v);$ 
     $\text{msgs} := \text{msgs} \cup \{[type \mapsto \text{"2a"}, bal \mapsto \text{self}, val \mapsto v]\};$ 
}

macro Phase2b(b){
  when  $b \geq \text{maxBal}[\text{self}];$ 
  with  $(m \in \text{sentMsgs}(\text{"2a"}, b))\{$ 
     $\text{maxBal}[\text{self}] := b;$ 
     $\text{maxVVal}[\text{self}] := b;$ 
     $\text{maxVVal}[\text{self}] := m.val;$ 
     $\text{msgs} := \text{msgs} \cup \{[type \mapsto \text{"2b"}, acc \mapsto \text{self}, bal \mapsto b, val \mapsto m.val]\}$ 
  }
}

process (acceptor  $\in \text{Acceptor}$ ){
  acc: while (TRUE){
    with  $(b \in \text{Ballot})\{\text{either Phase1b}(b) \text{ or Phase2b}(b) \}$ 
  }
}

process (leader  $\in \text{Ballot}$ ){
  ldr: while (TRUE){
    either Phase1a()
    or with  $(v \in \text{Value})\{\text{Phase2a}(v)\}$ 
  }
}
}

```

BEGIN TRANSLATION ( $\text{chksum}(\text{pcal}) = \text{"39408c33"} \wedge \text{chksum}(\text{tla}) = \text{"ed3338c5"}$ )

VARIABLES  $\text{maxBal}, \text{maxVVal}, \text{maxVVal}, \text{msgs}$

**define statement**

$\text{sentMsgs}(t, b) \triangleq \{m \in \text{msgs} : (m.type = t) \wedge (m.bal = b)\}$

$\text{ShowsSafeAt}(Q, b, v) \triangleq$

LET  $Q1b \triangleq \{m \in \text{sentMsgs}(\text{"1b"}, b) : m.acc \in Q\}$

IN  $\wedge \forall a \in Q : \exists m \in Q1b : m.acc = a$

$\wedge \forall m \in Q1b : m.mbal = -1$

$\vee \exists m \in Q1b :$

$\wedge m.mval = v$

$\wedge \forall m1 \in Q1b : m1.mbal \leq m.mbal$

$\text{vars} \triangleq \langle \text{maxBal}, \text{maxVVal}, \text{maxVVal}, \text{msgs} \rangle$

$\text{ProcSet} \triangleq (\text{Acceptor}) \cup (\text{Ballot})$



THEOREM  $Spec \Rightarrow \Box Correctness$

$Liveness \stackrel{\Delta}{=} \forall b \in Ballot :$

$\Box(\forall a \in Acceptor : maxBal[a] \leq b) \Rightarrow \Diamond(\exists v \in Value : Chosen(v))$

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\\* Modification History

\\* Last modified *Thu Jul 18 15:30:44 PDT 2024* by *nano*

\\* Created *Thu Sep 03 22:58:03 EDT 2015* by *nano*