

A specification of the algorithm described in *Paxos Made Simple*. This specification is a modification of: <https://lampart.azurewebsites.net/tla/PConProof.tla> Look there for comments.

EXTENDS *Integers, FiniteSets, TLC*

CONSTANT *Value, Acceptor, Quorum*

ASSUME $QA \triangleq \bigwedge \forall Q \in Quorum : Q \subseteq Acceptor$
 $\bigwedge \forall Q1, Q2 \in Quorum : Q1 \cap Q2 \neq \{\}$

Ballot $\triangleq Nat$

ASSUME *BallotAssump* $\triangleq (Ballot \cup \{-1\}) \cap Acceptor = \{\}$

None $\triangleq \text{CHOOSE } v : v \notin Value$

Message \triangleq
 $\cup [type : \{ "1a" \}, bal : Ballot]$
 $\cup [type : \{ "1b" \}, acc : Acceptor, bal : Ballot,$
 $mbal : Ballot \cup \{-1\}, mval : Value \cup \{None\}]$
 $\cup [type : \{ "2a" \}, bal : Ballot, val : Value]$
 $\cup [type : \{ "2b" \}, acc : Acceptor, bal : Ballot, val : Value]$

--algorithm *PCon*{

variables *maxBal* = $[a \in Acceptor \mapsto -1]$,
maxVbal = $[a \in Acceptor \mapsto -1]$,
maxVval = $[a \in Acceptor \mapsto None]$,
msgs = $\{\}$

define {

sentMsgs(*t*, *b*) $\triangleq \{m \in msgs : (m.type = t) \wedge (m.bal = b)\}$

Max(*M*) \triangleq A message with the highest ballot number among the set of messages *ms*
 $\text{CHOOSE } maxM \in M : \forall m \in M : m.mbal \leq maxM.mbal$

HighestAcceptedValue(*Q1bMessages*) $\triangleq Max(Q1bMessages).mval$

ShowsSafeAt(*Q*, *b*, *v*) \triangleq

LET *Q1b* $\triangleq \{m \in sentMsgs("1b", b) : m.acc \in Q\}$
IN $\bigwedge \forall a \in Q : \exists m \in Q1b : m.acc = a$
 $\bigwedge \forall m \in Q1b : m.mbal = -1$
 $\vee v = HighestAcceptedValue(Q1b)$

}

macro *Phase1a*() {*msgs* := *msgs* $\cup \{[type \mapsto "1a", bal \mapsto self]\}$; }

macro *Phase1b*(*b*) {

when (*b* > *maxBal*[*self*]) $\wedge (sentMsgs("1a", b) \neq \{\})$;
maxBal[*self*] := *b*;

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    msgs := msgs ∪ {[type ↦ "1b", acc ↦ self, bal ↦ b,
                      mbal ↦ maxVVal[self], mval ↦ maxVVal[self]]];
  }

macro Phase2a(v){
  when ∧ sentMsgs("2a", self) = {}
    ∧ ∃ Q ∈ Quorum : ShowsSafeAt(Q, self, v);
  msgs := msgs ∪ {[type ↦ "2a", bal ↦ self, val ↦ v]};
}

macro Phase2b(b){
  when b ≥ maxBal[self];
  with (m ∈ sentMsgs("2a", b)){
    maxBal[self] := b;
    maxVVal[self] := b;
    maxVVal[self] := m.val;
    msgs := msgs ∪ {[type ↦ "2b", acc ↦ self, bal ↦ b, val ↦ m.val]}
  }
}

process (acceptor ∈ Acceptor){
  acc: while (TRUE){
    with (b ∈ Ballot){either Phase1b(b)or Phase2b(b) }
  }
}

process (leader ∈ Ballot){
  ldr: while (TRUE){
    either Phase1a()
    or with (v ∈ Value){Phase2a(v)}
  }
}
}

TypeOK ≜ ∧ maxBal ∈ [Acceptor → Ballot ∪ {-1}]
          ∧ maxVVal ∈ [Acceptor → Ballot ∪ {-1}]
          ∧ maxVVal ∈ [Acceptor → Value ∪ {None}]
          ∧ msgs ⊆ Message

ChosenIn(b, v) ≜
  ∃ Q ∈ Quorum : ∀ a ∈ Q :
    ∃ m ∈ sentMsgs("2b", b) :
      ∧ m.acc = a
      ∧ m.val = v

Chosen(v) ≜ ∃ b ∈ Ballot : ChosenIn(b, v)

Correctness ≜
  ∀ v1, v2 ∈ Value : Chosen(v1) ∧ Chosen(v2) ⇒ v1 = v2

```

THEOREM $Spec \Rightarrow \Box Correctness$

* Modification History
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