

Common definitions for DAG-based consensus protocols

EXTENDS *FiniteSets*, *Integers*

CONSTANTS

N The set of nodes
 F Byzantine nodes
 R set of rounds
 $Quorum$ The set of quorums (*e.g.* cardinality $\geq 2f+1$)
 $Blocking$ The set of blocking set (*e.g.* cardinality $\geq f+1$)
 $Leader(-)$ operator mapping each round to its leader

ASSUME $\exists n \in R : R = 0 \dots n$

DAG notions

DAG vertices are just pairs consisting of a node and a round:
 $V \triangleq N \times R$
 $Node(v) \triangleq v[1]$
 $Round(v) \triangleq v[2]$

A *digraph* is just a set of edges:
 $IsDigraph(digraph) \triangleq \forall e \in digraph :$
 $\quad \wedge e = \langle e[1], e[2] \rangle$
 $\quad \wedge \{e[1], e[2]\} \subseteq V$

$Vertices(digraph) \triangleq \text{UNION } \{\{e[1], e[2]\} : e \in digraph\}$

$Children(v, digraph) \triangleq$
 $\{c \in V : \langle v, c \rangle \in digraph\}$

RECURSIVE $Reachable(-, -, -)$
 $Reachable(v1, v2, dag) \triangleq$
 $\quad \vee v1 = v2$
 $\quad \vee \exists c \in Children(v1, dag) : Reachable(c, v2, dag)$

$Parents(v, digraph) \triangleq$
 $\{e[1] : e \in \{e \in digraph : e[2] = v\}\}$

$SubDAG(vs, es) \triangleq$ vertices *vs* form a sub-DAG (no missing children) of *DAG es*
 $\forall v \in vs : Children(v, es) \subseteq vs$

Other stuff

An arbitrary ordering of the nodes:
 $NodeSeq \triangleq \text{CHOOSE } s \in [1 \dots Cardinality(N) \rightarrow N] :$
 $\forall i, j \in 1 \dots Cardinality(N) : i \neq j \Rightarrow s[i] \neq s[j]$

$$NodeIndex(n) \triangleq \text{CHOOSE } i \in 1 \dots Cardinality(N) : NodeSeq[i] = n$$

An arbitrary ordering of the nodes with the leader last:

$$NodeSeqLeaderLast(r) \triangleq \text{CHOOSE } s \in [1 \dots Cardinality(N) \rightarrow N] :$$

$$\wedge s[Cardinality(N)] = Leader(r)$$

$$\wedge \forall i, j \in 1 \dots Cardinality(N) : i \neq j \Rightarrow s[i] \neq s[j]$$

$$NodeIndexLeaderLast(n, r) \triangleq \text{CHOOSE } i \in 1 \dots Cardinality(N) : NodeSeqLeaderLast(r)[i] = n$$
