

EXTENDS *FiniteSets*, *Integers*, *TLC*

CONSTANTS

P the set of parties
 , $Faulty$ the set of faulty parties
 , V the set of value that may be broadcast

$N \triangleq Cardinality(P)$
 $F \triangleq Cardinality(Faulty)$

ASSUME $Faulty \subseteq P \wedge N > 3 * F$

$CeilDiv(a, b) \triangleq \text{IF } a \% b = 0 \text{ THEN } a \div b \text{ ELSE } (a \div b) + 1$

$Message \triangleq$ the set of possible messages in the network
 $[src : P, dst : P, type : \{\text{"propose"}, \text{"echo"}, \text{"vote"}, \text{"ready"}\}, val : V]$

--algorithm Broadcast{

variables

$broadcaster \in \{\text{CHOOSE } p \in Faulty : \text{TRUE}, \text{CHOOSE } p \in P \setminus Faulty : \text{TRUE}\}$, the distinguished broadcaster
 $bcastValue = \text{CHOOSE } v \in V : \text{TRUE}$, the value to broadcast (faulty nodes will ignore)
 $msgs = \{\}$; the set of sent messages

define {

$Msgs(self, v, type) \triangleq \{m \in msgs : m.type = type \wedge m.val = v \wedge m.dst = self\}$
 $Echos(self, v) \triangleq Msgs(self, v, \text{"echo"})$
 $Votes(self, v) \triangleq Msgs(self, v, \text{"vote"})$
 $Readys(self, v) \triangleq Msgs(self, v, \text{"ready"})$

}

macro SendAll(type, value) {

$msgs := msgs \cup \{[src \mapsto self, dst \mapsto d, type \mapsto type, val \mapsto value] : d \in P\}$

}

fair process (correctParty $\in P \setminus Faulty$)

variable $delivered = \langle \rangle$; the delivered value

{

l0: **while (TRUE) with ($v \in V$) {**

either { send proposal

when $self = broadcaster$;

when $\forall m \in msgs : \neg(m.src = self \wedge m.type = \text{"propose"})$;

$SendAll(\text{"propose"}, bcastValue)$

}

or { send echo

when $\forall m \in msgs : \neg(m.src = self \wedge m.type = \text{"echo"})$;

await $[src \mapsto broadcaster, dst \mapsto self, type \mapsto \text{"propose"}, val \mapsto v] \in msgs$;

$SendAll(\text{"echo"}, v)$

}

```

    or { fast delivery
      await Cardinality(Echos(self, v) \ {broadcaster}) ≥ CeilDiv(N + 2 * F - 2, 2);
      delivered := v
    }
    or { send vote
      when ∀ m ∈ msgs : ¬(m.src = self ∧ m.type = "vote");
      await Cardinality(Echos(self, v) \ {broadcaster}) ≥ CeilDiv(N, 2);
      SendAll("vote", v)
    }
    or { send ready
      when ∀ m ∈ msgs : ¬(m.src = self ∧ m.type = "ready");
      await
        ∨ Cardinality(Echos(self, v) \ {broadcaster}) ≥ CeilDiv(N + F - 1, 2)
        ∨ Cardinality(Votes(self, v) \ {broadcaster}) ≥ CeilDiv(N + F - 1, 2)
        ∨ Cardinality(Readys(self, v)) ≥ F + 1;
      SendAll("ready", v)
    }
    or { slow delivery
      await Cardinality(Readys(self, v)) ≥ 2 * F + 1;
      delivered := v
    }
  }
}
process ( faultyParty ∈ Faulty ) {
l1:  with ( v ∈ V, t ∈ {"propose", "echo", "vote", "ready"}, d ∈ P \ Faulty )
      msgs := msgs ∪ {[src ↦ self, dst ↦ d, type ↦ t, val ↦ v]}
}
}

```

Correctness properties:

$Agreement \triangleq \forall p1, p2 \in P \setminus Faulty :$
 $delivered[p1] \neq \langle \rangle \wedge delivered[p2] \neq \langle \rangle \Rightarrow delivered[p1] = delivered[p2]$

$Liveness \triangleq$
 $\wedge (broadcaster \notin Faulty \Rightarrow \forall p \in P \setminus Faulty : \Diamond (delivered[p] = bcastValue))$
 $\wedge \Box ((\exists p \in P \setminus Faulty : delivered[p] \neq \langle \rangle) \Rightarrow \forall p \in P \setminus Faulty : \Diamond (delivered[p] \neq \langle \rangle))$

$Symm \triangleq Permutations(P \setminus (Faulty \cup \{CHOOSE p \in P \setminus Faulty : TRUE\}))$
