

# Mastering the game of Go with deep neural networks and tree search

Understanding Monte Carlo Tree Search (MCTS)

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- The algorithm introduced is easily generalizable to other games.
- It makes use of the existing Monte Carlo Tree Search (MCTS) algorithm and enhances it using heuristics learnt through deep learning.

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- Derive and explore the optimal gameplay algorithm (minimax).
- Derive and explore vanilla MCTS
- Understand the Exploration-Exploitation tradeoff using Upper Confidence Trees (UCT)
- Describe MCTS as used in the paper

# Greedy-Coins Game

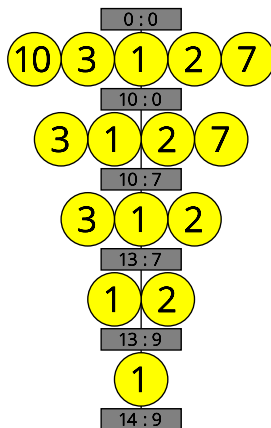
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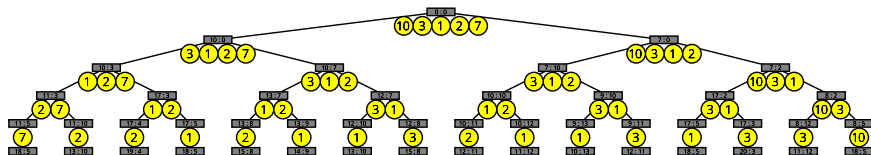
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- The aim is to collect coins worth more than that collected by the opponent
- Only one coin is picked on each turn from either the left or the right



# The Game Tree



# Minimax

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See **examples/minimax**

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- Even our game becomes infeasible if we start with a large number of coins.
- **However**, minimax requires no knowledge beyond the rules of the game!

# What if we run minimax on a subset of games?

Good moves are much rarer than bad ones and are easy to miss when sampled randomly.

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- We have now defined the RHS in terms of a smaller LHS. We can solve for  $P^*(a|s)$ .



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- Now, the next time we play a few more games don't choose the first move randomly but choose it according to  $P_e(\text{win}|a, s)$ .
- Monte Carlo tells us that this process (repeated) guarantees that  $P(a|s) \rightarrow P^*(a|s)$ .

# Monte Carlo Tree Search: Example

See **examples/vanilla**

# Monte Carlo Tree Search: Limitations

It can take a long time for  $P(a|s) \rightarrow P^*(a|s)$



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$$\frac{1}{N_i} \sum_{j=1}^{N_i} x_{ij} + \sqrt{\frac{2 \ln \sum_i N_i}{N_i}}$$

$N_i$  = number of times machine  $i$  is played

$x_{ij}$  = win or loss when playing machine  $i$  the  $j$ th time

# Speeding up MCTS: Example

See **examples/uct**

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- Second, when a move is selected, UCT uses no prior knowledge of gameplay to make a better decision.
- AlphaGo uses a prior distribution of moves learnt by a deep learning model using reinforcement learning and human-expert moves.

# Speeding up MCTS: AlphaGo Example

See **`examples/alphago`**

# Thank You

[https://github.com/nanonaren/PWL\\_AlphaGo](https://github.com/nanonaren/PWL_AlphaGo)