Design analysis and algorithm Lab 10 Divide And Conquer Nearest Neighbour

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Course code: CSE3004

Nearest Neighbour:

```
import math
import copy
class Point():
    def __init__(self, x, y):
        self.x = x
        self.y = y
def dist(p1, p2):
    return math.sqrt((p1.x - p2.x) *
                      (p1.x - p2.x) +
                      (p1.y - p2.y) *
                      (p1.y - p2.y))
def bruteForce(P, n):
    min_val = float('inf')
    for i in range(n):
        for j in range(i + 1, n):
            if dist(P[i], P[j]) < min_val:</pre>
                min_val = dist(P[i], P[j])
    return min_val
def stripClosest(strip, size, d):
    min_val = d
    for i in range(size):
        j = i + 1
        while j < size and (strip[j].y -</pre>
                             strip[i].y) < min_val:</pre>
            min_val = dist(strip[i], strip[j])
            j += 1
    return min_val
def closestUtil(P, Q, n):
        return bruteForce(P, n)
    mid = n // 2
    midPoint = P[mid]
    Pl = P[:mid]
    Pr = P[mid:]
    dl = closestUtil(Pl, Q, mid)
    dr = closestUtil(Pr, Q, n - mid)
    d = min(dl, dr)
    stripP = []
    stripQ = []
    lr = Pl + Pr
    for i in range(n):
       if abs(lr[i].x - midPoint.x) < d:</pre>
```

```
stripP.append(lr[i])
        if abs(Q[i].x - midPoint.x) < d:</pre>
            stripQ.append(Q[i])
    stripP.sort(key = Lambda point: point.y)
    min_a = min(d, stripClosest(stripP, len(stripP), d))
    min_b = min(d, stripClosest(stripQ, len(stripQ), d))
    return min(min_a,min_b)
def closest(P, n):
    P.sort(key = Lambda point: point.x)
    Q = copy.deepcopy(P)
    Q.sort(key = lambda point: point.y)
    return closestUtil(P, Q, n)
P = [Point(2, 3), Point(12, 30),
    Point(40, 50), Point(5, 1),
     Point(12, 10), Point(3, 4)]
n = len(P)
print("The smallest distance is",closest(P, n))
```

Output:

The smallest distance is 2.23606797749979

Analysis:

Nearest Neighbour: To solve this we first split the given in points by an appropriately chosen virtical line into a natures S, & S, of since 1/2. By making & recursive calls for the wets & & & , we find the ruinimum distances d, & de in these subsets. Let d= frum {d, ,d} but this way takes a lot of three thus not providing us the goal. What we do is, part the strip into dx d squares as show that each such square contains at most four input points. · de, we stort in points by their x-coordinate & their split the resulting sorted list duto a habits S, & S. & size u/2. Calculate min of & d, & d, distances & check whether it is

Smaller than d. To perform such a check, we filter the initial point that & keep only those points whose x-distance to the middle time does not exceed d. Afterwards we sort the well points in the resulting this by their y wordinalis & sean the list for point we construct its discovery the season distance to the searn seebsequent points in this list of points. For the searn seebsequent points in this list of points. For the steam we excountered during this seam, Aftwords, we return nin { d, d'}

The running times of algorithm statisfies the recurrence

T(n) = 27 (n/2) + 0 (nlog n)

The O(n log n) term comes from sorting the points in the strip by their y-coordinates at every iteration.