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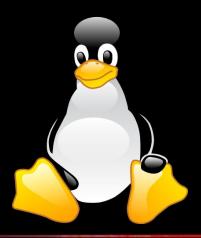
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OPERATING SYSTEM

Part 18: POSIX Threads





Thread



- Multiple strands of execution in a single program are called threads. A more precise definition is that a thread is a sequence of control within a process.
- When we create a new thread in a process, the new thread of execution gets its own stack (and hence local variables) but shares global variables, file descriptors, signal handlers, and its current directory state with the process that created it.



POSIX



- Linux first acquired thread support around 1996, with a library often referred to as "LinuxThreads."
 This was very close to the POSIX standard
- POSIX: "Portable Operating System Interface for Unix" is the name of a family of related standards specified by the IEEE to define the application programming interface (API), along with shell and utilities interfaces, for software compatible with variants of the Unix operating system, although the standard can apply to any operating system.



Advantages of threads



- Sometimes it is very useful to make a program appear to do two things at once. The classic example: MS-Word.
- The performance of an application that mixes input, calculation, and output may be improved by running these as three separate threads
- Now that multi-cored CPUs are common even in desktop and laptop machines, using multiple threads inside a process can, if the application is suitable, enable a single process to better utilize the hardware resources available.
- In general, switching between threads requires the operating system to do much less work than switching between processes.



Drawbacks of thread



- Writing multithreaded programs requires very careful design. The potential for introducing. Alan Cox (the well respected Linux guru) has commented that threads are also known as "how to shoot yourself in both feet at once."
- Debugging a multithreaded program is much, much harder than debugging a single-threaded one, because the interactions between the threads are very hard to control.
- A program that splits a large calculation into two and runs the two parts as different threads will not necessarily run more quickly on a single processor machine.





- Based on IEEE POSIX 1003.1c standard.
- Almost all vendors support POSIX threads in one way or another
- Shared memory based
- Characteristics
 - Thread creation, management and destruction
 - Synchronization
- Many people refer them as light weight processes
 - Independent flow of control
 - Maintain stack pointer, registers, scheduling, set of pending and block signals and thread specific data.
 - Share any other process resources such as (shared) memory.





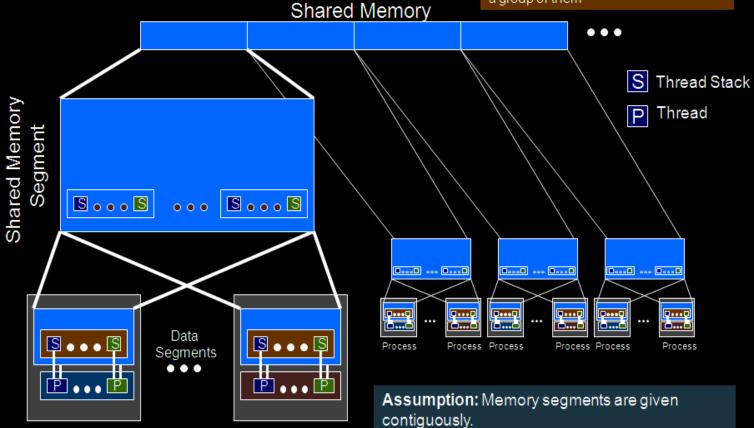
System V IPC

A main structure that contains all pages in the shared memory

A Logical View

POSIX IPC

Memory is defined as a file system and each page can be accessed as a file or a group of them







- Independent flow of control → Functions
- Data can be passed by value or by reference
 - In Pthreads all arguments are passed as a void pointer. Thus it is usually to be considered pass by reference, even though you can trick this!!!!





Basic Datatypes:

pthread_t var;

pthread_attr_t var;

pthread_mutex_t var;

pthread_mutexattr_t var;

pthread cond t var;

pthread_condattr_t var;

A variable that will be used as an ID for a thread

A variable that defines the thread behavior (i.e. scheduling, priority, stack size, scope, etc)

The definition of a Mutually Exclusive lock in Pthreads

Defines the mutex behavior. Either PTHREAD_PROCESS_PRIVATE (the default) and PTHREAD_PROCESS_SHARED

The definition of a Conditional variable in Pthreads

Defines the conditional variable behavior. Either PTHREAD_PROCESS_PRIVATE (the default) and PTHREAD_PROCESS_SHARED





A Side Note:

Threads keep an independent stack pointer

Stacks are not shared among threads

Auto variables are as if they were local to the thread that declare them

If the parameters are passed by reference to the thread, then they will be considered as shared





Thread creation:

Return a non zero value in success

pthread_t *thr

Will contain the newly created thread's id. Must be passed by reference

const pthread_attr_t *attr

Give the attributes that this thread will have. Use NULL for the default ones

void *(*start_routine)(void)

The name of the function that the thread will run. Must have a void pointer as its return and parameters values

void *arg

The argument for the function that will be the body of the Pthreads

Pointers of the type void can reference **ANY** type of data, but they **CANNOT** be used in any type of operations that reads or writes its data without a cast







Miscellaneous Useful functions:

pthread_t pthread_self(void)

Return the id of the calling thread. Returns a pthread_t type which is usually an integer type variable

OpenMP Counterpart
int omp_get_thread_num(void);

void pthread_exit(void *arg);

This function will indicate the end of a Pthread and the returning value will be put in *arg*







Example 1:

```
#include <pthread.h>
#define NUM THREADS 4
void *work(void *i){
 printf("Hello, world from %i\n", pthread self());
 pthread exit(NULL);
int main(int argc, char **argv){
 int i:
 pthread tid[NUM THREADS];
 for(i = 0; i < NUM THREADS; ++i){}
   if(pthread_create(&id[i], NULL, work, NULL)){
     printf("Error creating the thread\n"); exit(19);
 printf("After creating the thread. My id is: %i\n",
          pthread self());
return 0;}
```

Hello, world from 2
Hello, world from 3
After creating the thread.
My id is: 1
Hello, world from 4

What happened to thread 5???

Hello, world from 2
Hello, world from 3
Hello, world from 4
After creating the thread.
My id is: 1
Hello, world from 5





- Single Argument Passing
 - Cast its value as a void pointer (a tricky pass by value)
 - Cast its address as a void pointer (pass by reference).
 - The value that the address is pointing should NOT change between Pthreads creation

- Multiple Argument Passing
 - Heterogonous: Create an structure with all the desired arguments and pass an element of that structure as a void pointer.
 - Homogenous: Create an array and then cast it as a void pointer









Example 2a:

```
#include <pthread.h>
#define NUM THREADS 10
void *work(void *i){
 int f = *((int *)(i));
 printf("Hello, world from %i with value %i\n",
          pthread_self(), f);
 pthread exit(NULL);
int main(int argc, char **argv){
 int i:
 pthread tid[NUM THREADS];
 for(i = 0; i < NUM THREADS; ++i){
   if(pthread_create(&id[i], NULL, work, (void *)(&i))){
     printf("Error creating the thread\n"); exit(19);}
 return 0;}
```

Hello, world from 2 with value 1
Hello, world from 3 with value 2
Hello, world from 6 with value 5
Hello, world from 5 with value 5
Hello, world from 4 with value 4
Hello, world from 8 with value 9
Hello, world from 9 with value 9
Hello, world from 10 with value 9
Hello, world from 7 with value 6
Hello, world from 11 with value 10

Wrong Method!!!!







Example 2b:

```
#include <pthread.h>
#define NUM THREADS 10
void *work(void *i){
     int f = (int)(i);
     printf("Hello, world from %i with value %i\n",
          pthread_self(), f);
     pthread exit(NULL);
int main(int argc, char **argv){
     int i:
     pthread tid[NUM THREADS];
     for(i = 0; i < NUM_THREADS; ++i){</pre>
          if(pthread_create(&id[i], NULL, work, (void *)(i))){
               printf("Error creating the thread\n");
               exit(19);
     return 0;
```

Hello, world from 2 with value 0 Hello, world from 3 with value 1 Hello, world from 4 with value 2 Hello, world from 5 with value 3 Hello, world from 6 with value 4 Hello, world from 7 with value 5 Hello, world from 8 with value 6 Hello, world from 10 with value 8 Hello, world from 11 with value 9

Right Method 1









Example 2c:

```
#include <pthread.h>
#define NUM THREADS 10
void *work(void *i){
     int f = *((int *)(i));
     printf("Hello, world from %i with value %i\n",
          pthread self(), f);
     pthread exit(NULL);}
int main(int argc, char **argv){
     int i:
     int y[NUM THREADS]:
     pthread tid[NUM THREADS];
     for(i = 0; i < NUM_THREADS; ++i){</pre>
          y[i] = i
          if(pthread_create(&id[i], NULL, work, (void *)(&y[i]))){
               printf("Error creating the thread\n");
               exit(19);
     return 0;}
```

Hello, world from 2 with value 0
Hello, world from 4 with value 2
Hello, world from 5 with value 3
Hello, world from 6 with value 4
Hello, world from 7 with value 5
Hello, world from 8 with value 6
Hello, world from 9 with value 7
Hello, world from 3 with value 1
Hello, world from 10 with value 8
Hello, world from 11 with value 9

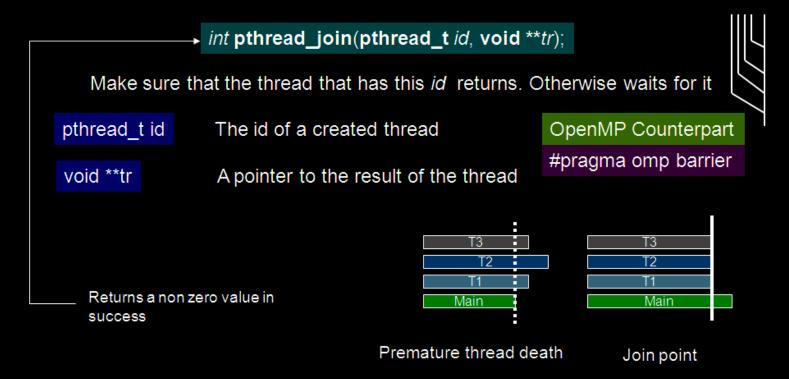
Right Method 2







The Joining of All Loose Ends: pthread_join



Why use it? If the main thread dies, then all other threads will die with it. Even if they have not completed their work







Example 3:

```
#include <pthread.h>
#define NUM THREADS4
void *work(void *i){
    printf("Hello, world from %i\n", pthread self());
    pthread exit(NULL);
int main(int argc, char **argv){
    int i:
    pthread tid[NUM THREADS]:
    for(i = 0; i < NUM THREADS; ++i){
         if(pthread_create(&id[i], NULL, work, NULL)){
               exit(19):
     printf("After creating the thread. My id is: %i\n", pthread self());
    for(i = 0; i < NUM_THREADS; ++i){
         if(pthread_join(id[i], NULL)){
               exit(19);
    printf("After joining\n");
    return 0:
```

Hello, world from 2
Hello, world from 3
Hello, world from 4
After creating the thread.
My id is: 1
Hello, world from 5
After joining







Attributes Revisited: Useful Functions (1)

```
int pthread_attr_init(pthread_attr_t *attr);
```

Initialize an attribute with the default values for the attribute object

- •Default Schedule: SCHED_OTHER (?)
- Default Scope: PTHREAD_SCOPE_SYSTEM (?)
- Default Join State: PTHREAD_CREATE_JOINABLE (?)

int pthread_attr_destroy(pthread_attr_t *attr);

De-allocate any memory and state that the attribute object occupied. It is safe to delete the attribute object after the thread has been created

```
int pthread_attr_setdetachstate(pthread_attr_t *attr, int JOIN_STATE);
```

Set the attached parameter on the attribute object with the JOIN_STATE variable

- •PTHREAD_CREATE_JOINABLE: It can be joined at a join point. State must be saved after function ends
- PTHREAD_CREATE_DETACHED: It cannot be joined at a join point. State and resources are de-allocated immediately





Attributes Revisited: Useful Functions (2)

int pthread_attr_setschedpolicy(pthread_attr_t *attr, int policy)

Set the scheduling policy of the thread:

•SCHED_OTHER → Regular scheduling •SCHED_RR → Round-robin (SU) •SCHED_FIFO → First-in First-out (SU)

int pthread_attr_setschedparam(pthread_attr_t *attr, const struct sched_param *pr)

Contains the schedule priority of the thread **Default: 0**

int pthread_attr_setinheritsched(pthread_attr_t *attr, int inherit)

Tell if the scheduling parameters will be inherit from the parent or the ones in the attribute object will be used

PTHREAD_EXPLICIT_SCHED → Scheduling parameters from the attribute object will be used.

PTHREAD_INHERIT_SCHED → inherit the attributes from its parent.

int pthread_attr_setscope(pthread_attr_t *attr, int scope)

Contention parameter

- PTHREAD SCOPE SYSTEM
- PTHREAD_SCOPE_PROCESS





```
#include <pthread.h>
#define NUM_THREADS 4
struct args{int a; float b; char c;};
void *work(void *i){
     struct args *a = (struct args *)(i);
     printf("(%3i, %.3f, %3c) --> %i\n", a->a, a->b, a->c, pthread_self());
     pthread exit(NULL);
int main(int argc, char **argv){
     int i:
     struct args a[NUM THREADS];
     pthread tid[NUM THREADS];
     pthread attr t ma;
     pthread attr init(&ma);
     pthread attr setdetachstate(&ma, PTHREAD CREATE JOINABLE);
     for(i = 0; i < NUM THREADS; ++i){
         a[i].a = i; a[i].b = 1.0 /(i+1); a[i].c = 'a' + (char)(i);
         pthread create(&id[i], &ma, work, (void *)(&a[i]));
     pthread attr destroy(&ma);
     for(i = 0; i < NUM_THREADS; ++i){pthread_join(id[i], NULL);}
     return 0;}
```

Example 4:

```
( 0, 1.000, a) --> 2
( 3, 0.250, d) --> 5
( 2, 0.333, c) --> 4
( 1, 0.500, b) --> 3
```





Synchronization Types





Mutual Exclusion Lock. Only the thread that has the lock can access the protected region

Semaphores

^{yকুর} Monitors

Act as a guard of some resource. Consists of a mutex with some kind of notification scheme

ুদ্রা^{ঠুঠি} Conditional Variables

Synchronization occurs when a condition is met. Always used in conjunction with a mutex. Interthread (process) communication.

हार^{हरी} Reader / Writer Locks

Permits only reads on a data or writes on data in a group. In other words, lock out the writers when the readers are on the shared data and vice versa.





Synchronization (1): Mutex

Pthread

OpenMP

 $int \, pthread_mutex_init(pthread_mutex_t \, *m, \\ const \, pthread_mutexattr_t \, *ma);$

int omp_lock_init(omp_lock_t */kc);

int pthread_mutex_lock(pthread_mutex_t *m);

int omp_lock_set(omp_lock_t */kc);

int pthread_mutex_unlock(pthread_mutex_t *m);

int omp_lock_unset(omp_lock_t */kc);

int pthread_mutex_destroy(pthread_mutex_t
*m);

int omp_lock_destroy(omp_lock_t
*/kc);

1→ Initialization

2 → Setting

3 → Unsetting

4 → Destroying



pthread_mutexattr_t * ma can be left NULL for the default values





Example 5

100010. Balance: Balance: Final Initial

```
#include <pthread.h>
#define NUM THREADS 2
#define CYCLE 100
double b:
void *deposit(void *i){ double m = *(double *)(i); int j;
     for(j = 0; j < CYCLE; ++j){ b += m; sleep(1); }}
void *withdraw(void *i){double m = *(double *)(i); int j;
     for(j = 0; j < CYCLE; ++j){ b -= m; sleep(1); }}
int main(int argc, char **argv){
     int i; double bi; double q = 10.0; pthread t id[NUM THREADS];
     b = 100000; bi = b;
     pthread create(&id[0], NULL, deposit, (void *)(&q));
     pthread_create(&id[1], NULL, withdraw, (void *)(&q));
     for(i = 0; i < NUM_THREADS; ++i){pthread_join(id[i], NULL);}
     printf("The Initial Balance: %.2lf\nThe Final Balance: %.2lf\n",
bi, b);
     return 0; }
```







Example 5

100000.00 8 100000. Balance: Balance:

```
pthread_mutex_tmt;
void *deposit(void *i){
    double m = *(double *)(i);
    int j;
    for(j = 0; j < CYCLE; ++j){
           pthread_mutex_lock(&mt); b += m; pthread_mutex_unlock(&mt);
sleep(1); }
void *withdraw(void *i){
    double m = *(double *)(i);
    int j;
    for(j = 0; j < CYCLE; ++j){
           pthread_mutex_lock(&mt); b -= m; pthread_mutex_unlock(&mt);
sleep(1); }
int main(int argc, char **argv){
    pthread_mutex_init(&mt, NULL);
    pthread_mutex_destroy(&mt);
    return 0;
```







Synchronization (2): Conditional Variables

int pthread_cond_init(pthread_cond_t *cond, pthread_condattr_t *attr)

Initialize a conditional variable *cond* with the attributes given by *attr*. The *attr* can be left NULL so it will use the default variable

int pthread_cond_destroy(pthread_cond_t *cond)

De-allocated any resources associated with the conditional variable cond

int pthread_condattr_init (pthread_condattr_t *attr)

Initialize the conditional attribute variable attr with the default value.

int pthread_condattr_destroy (pthread_condattr_t *attr)

De-allocate the conditional attribute variable *attr*. It is safe to deallocate the variable just after the conditional variable has been initialized





Synchronization (2): Conditional Variables

Tricky Conditional Variables

int pthread_cond_wait(pthread_cond_t *condition, pthread_mutex_t *m)

Make the calling thread wait for a signal in the conditional variable. Must be called when the associated mutex is locked and it will unlock it while the thread blocks. It will also unlock the mutex if the signal has been received.

int pthread_cond_signal(pthread_cond_t *condition)

Signal a thread that has been blocked waiting for condition to become true. It must be called **after** its associated mutex has been **locked** and the mutex must be **unlocked** after the signal has been **issued**

int pthread_cond_broadcast (pthread_cond *condition)

Similar to pthread_cond_signal but it signals all the waiting threads for this conditional variable





Example 6

```
#include <pthread.h>
#define COUNT_DONE 10
#define COUNT_HALT1 3
#define COUNT_HALT2 6
pthread_mutex_t count_mutex = PTHREAD_MUTEX_INITIALIZER;
pthread_mutex_t condition_mutex = PTHREAD_MUTEX_INITIALIZER;
pthread_cond_t condition_cond = PTHREAD_COND_INITIALIZER;
int count = 0;
```

Statically initialize the mutexes and the conditional variables





```
void *f1(){
 for(;;){
   pthread mutex lock( &condition mutex );
   while( count >= COUNT_HALT1 && count <= COUNT_HALT2 ){
     pthread cond wait( &condition cond, &condition mutex );
   pthread mutex unlock( &condition mutex );
   pthread mutex lock( &count mutex );
   count++;
   printf("Counter value %i (f1): %d\n",pthread self(), count);
   pthread mutex unlock( &count mutex );
   if(count >= COUNT_DONE) pthread_exit(NULL);
```

First function: If the count variable is between COUNT_HALT1 and COUNT_HALT2, wait for a signal from f2. Otherwise increment count







```
void *f2(){
  for(;;){
    pthread mutex lock( &condition mutex );
    if( count < COUNT_HALT1 || count > COUNT_HALT2 ){
      pthread cond signal( &condition cond );
    pthread_mutex_unlock( &condition_mutex );
    pthread mutex lock( &count mutex );
    count++;
    printf("Counter value %i (f2): %d\n",pthread self(), count);
    pthread mutex unlock( &count mutex );
    if(count >= COUNT_DONE) pthread_exit(NULL);
```

Second function: Send a signal to f1 if count is less than COUNT_HALT1 or greater than COUNT_HALT2. Later, increment the count





Counter value 2 (f1): 1
Counter value 2 (f1): 2
Counter value 2 (f1): 3
Counter value 3 (f2): 4
Counter value 3 (f2): 5
Counter value 3 (f2): 6
Counter value 3 (f2): 7
Counter value 3 (f2): 8
Counter value 3 (f2): 9
Counter value 3 (f2): 10
Counter value 2 (f1): 11

```
int main(int argc, char **argv)
{
   pthread_t thread1, thread2;
   pthread_create( &thread1, NULL, &f1, NULL);
   pthread_create( &thread2, NULL, &f2, NULL);
   pthread_join( thread1, NULL);
   pthread_join( thread2, NULL);
   return(0);
}
```

F1 is effectively halted between COUNT_HALT1 and COUNT_HALT2. Otherwise random

FBUNDATION



Example taken from YoLinux.com. Copyright of Greg Ippolito



Miscellaneous Useful Functions

int pthread_attr_getstackaddr (const pthread_attr_t *attr, void **stackaddr)

Return the stack address that this P-thread will be using

int pthread_attr_getstacksize (const pthread_attr_t *attr, size_t *stacksize)

Return the stack size that this P-thread will be using

int pthread_detach (pthread_t thr, void **value_ptr)

Make the thread that is identified by thr not joinable

int pthread_once (pthread_once_t *once_control, void (*init_routine)(void));

Make sure that the init_routine is executed by a single thread and only once. The once_control is a synchronization mechanism that can be defined as:

pthread_once_tonce_control=PTHREAD_ONCE_INIT;

void pthread_yield ()

Relinquish the use of the processor



OpenMP Counterpart

#pragma omp single



Resources



- LLNL Pthread Tutorial: http://www.llnl.gov/computing/tutorials/pthreads/
- The YoLinux Pthread tutorial: http://www.yolinux.com/TUTORIALS/LinuxTutorialPosixT hreads.html#SYNCHRONIZATION
- Linux Parallel Processing HOWTO by Hank Dietz: http://yara.ecn.purdue.edu/~pplinux/PPHOWTO/pphowto
 .html#toc2

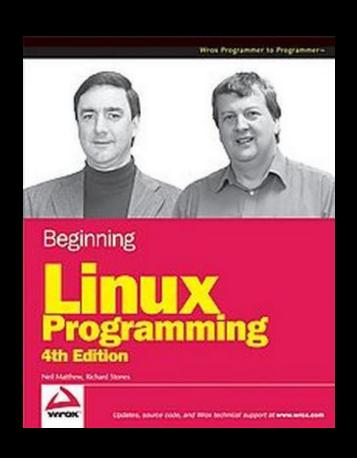






Reference Books





"Beginning Linux Programming", 4th Edition, Neil Mathew, Richard Stones, Wrox Publication.

Rating:











