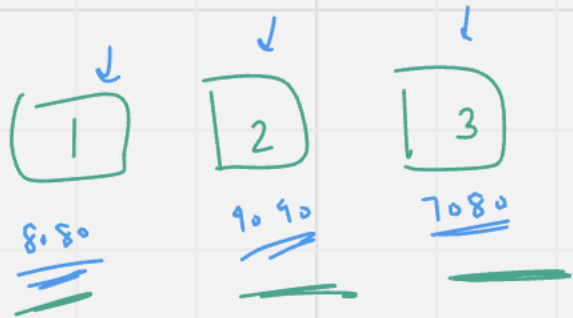
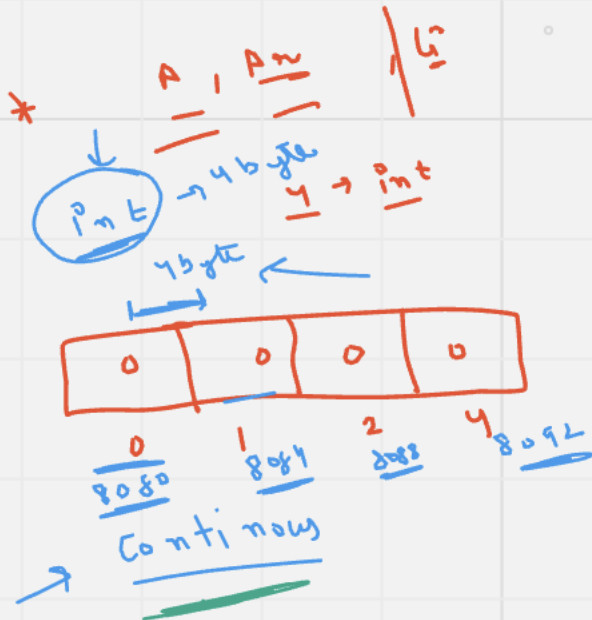


LinkedList

Size / no size

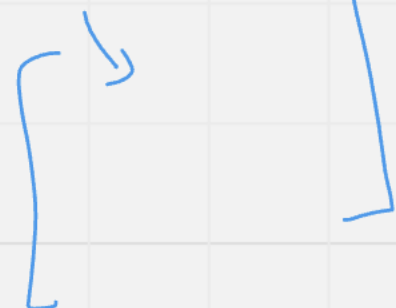


3

arr

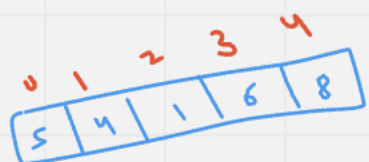
Size

[1]



arr, list

inc / dec

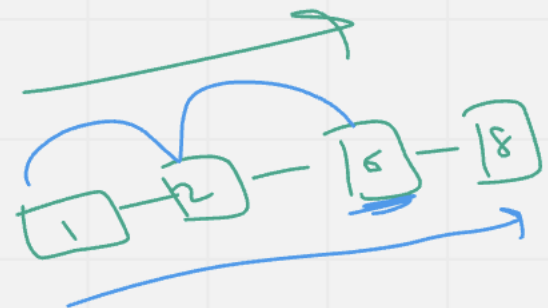


3

arr[3] = 6

o(1)

arr[4] = 8

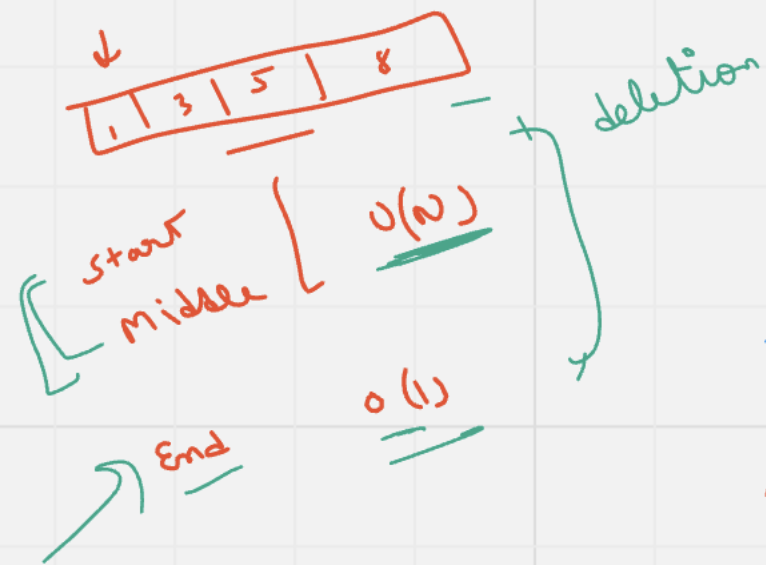


4 = 8

6

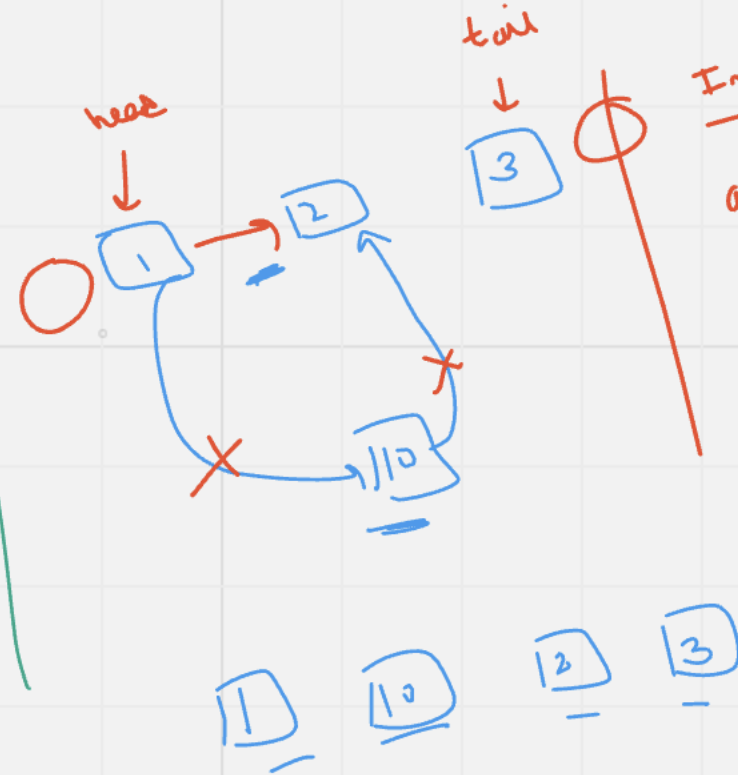
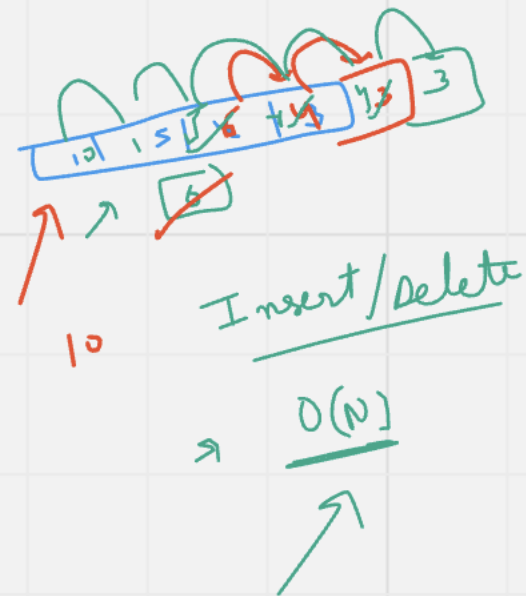
O(N)

✓



$O(1)$

$arr(4) = 8$



Insert

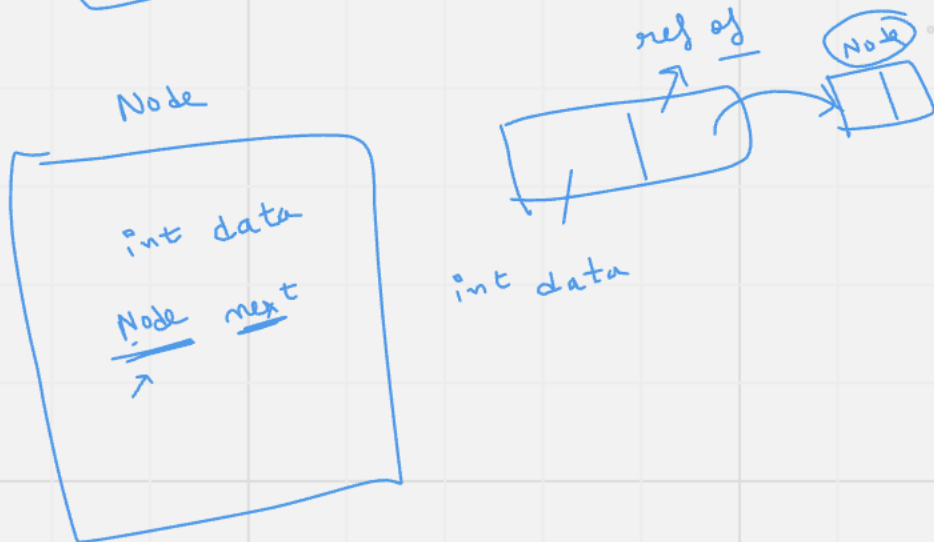
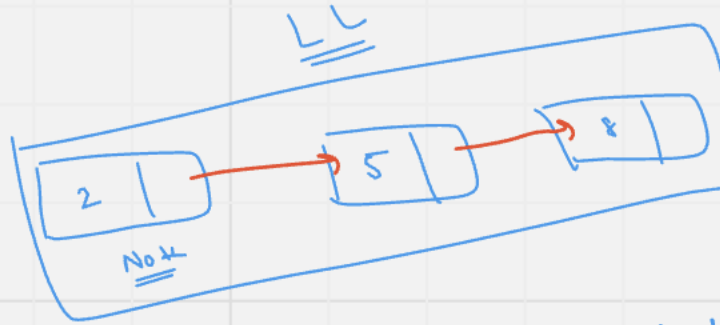
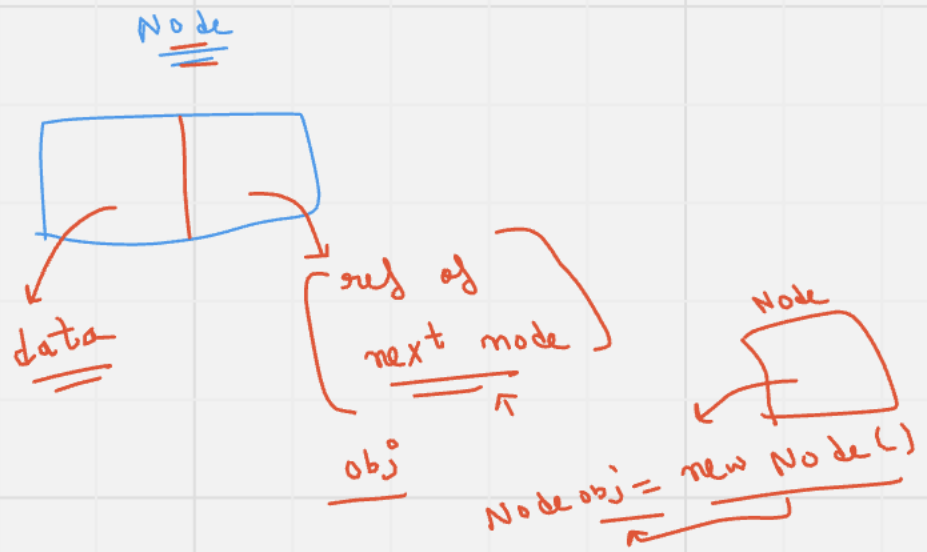
at start

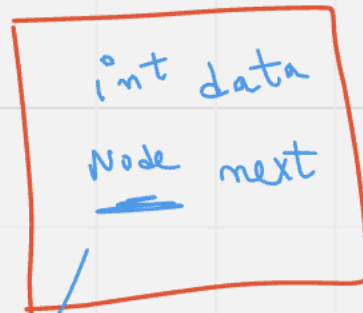
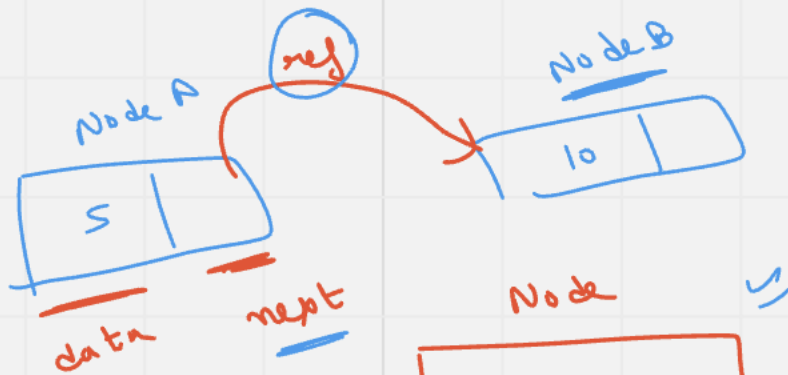
at end

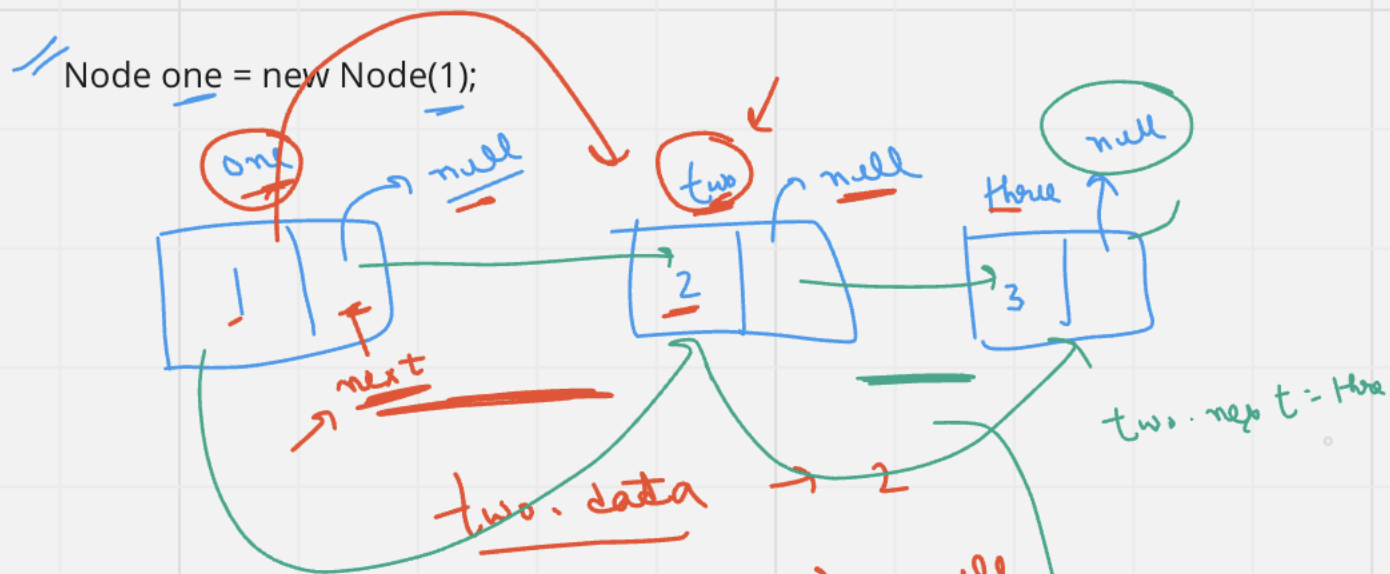
$O(1)$

in middle $O(n)$

delet





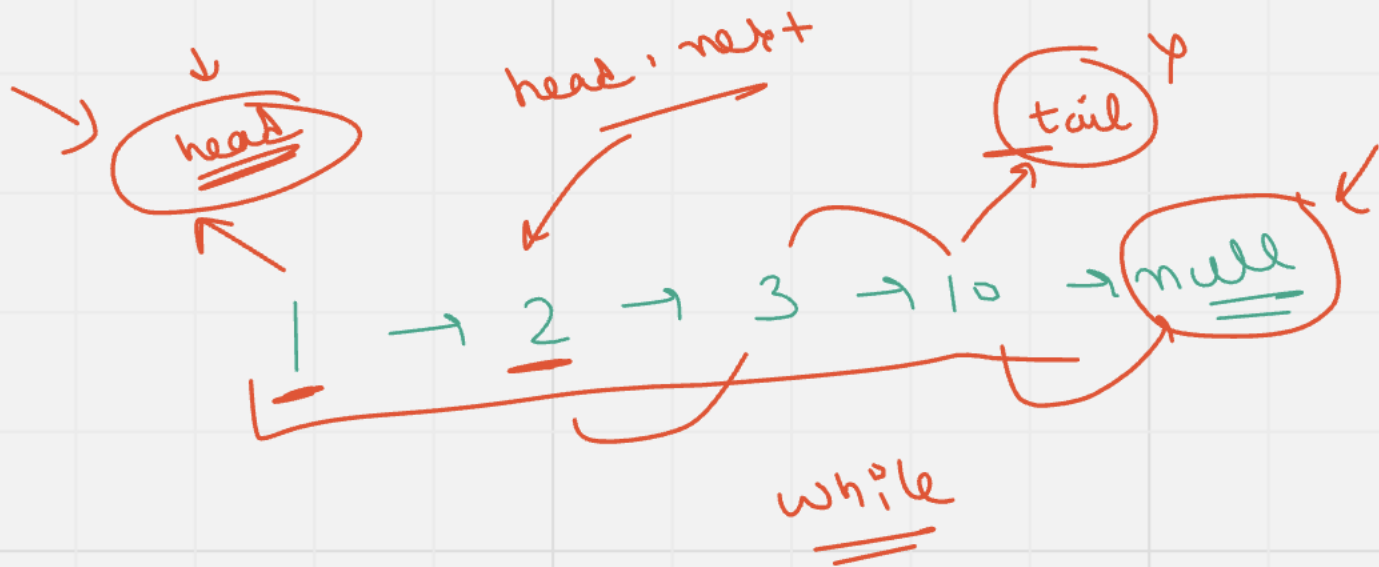


one.next.data

one.next.next.data = 3

one.next.next.data = null.data

null pointer
exception



head next

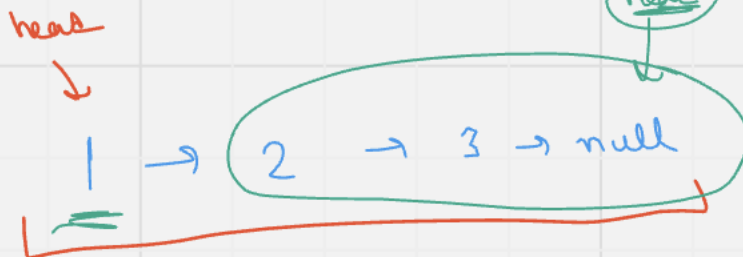
```
static void printLL(Node head) {  
  
}
```



```
print(head.data) ←
```

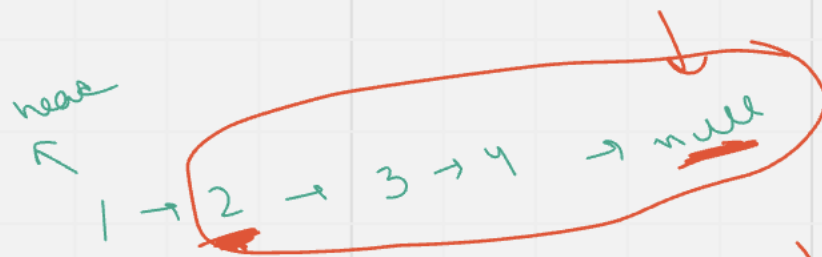
A diagram illustrating the state of a linked list. A variable named 'head' is shown on the left, with an arrow pointing to a box labeled 'head.next'. This box is part of a larger structure representing a node, which also contains 'data' and 'next' fields. The 'head.next' field is highlighted with a blue oval, indicating it is the target of the next operation.

rec



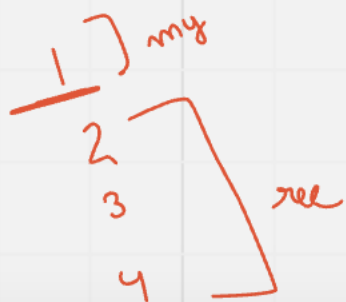
head.next

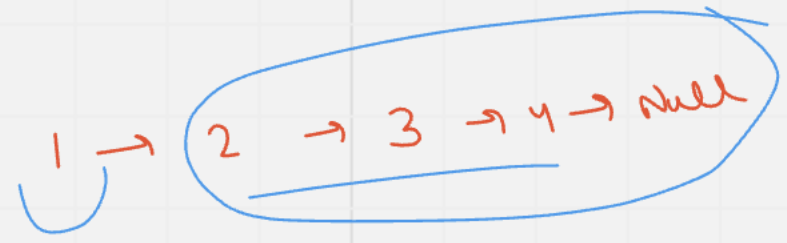
say rec



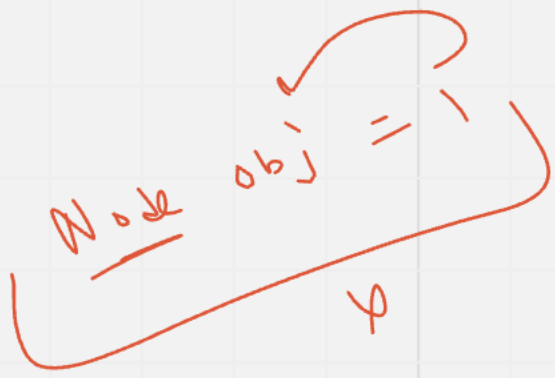
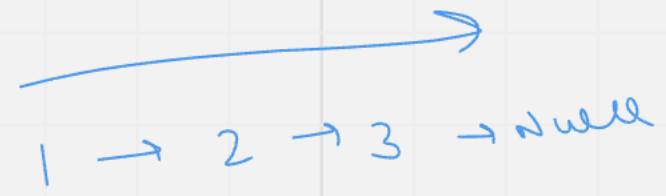
our task

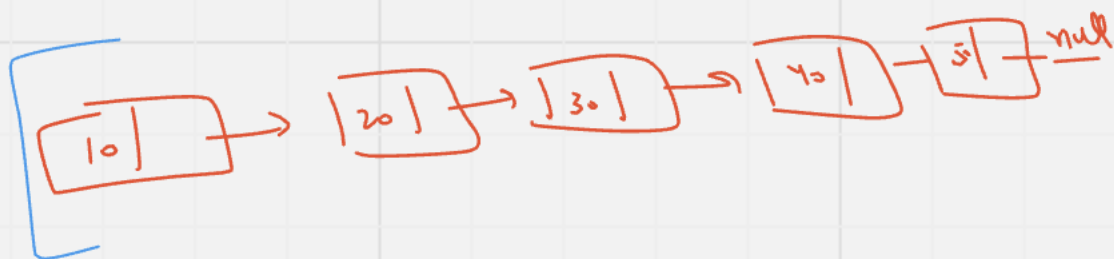
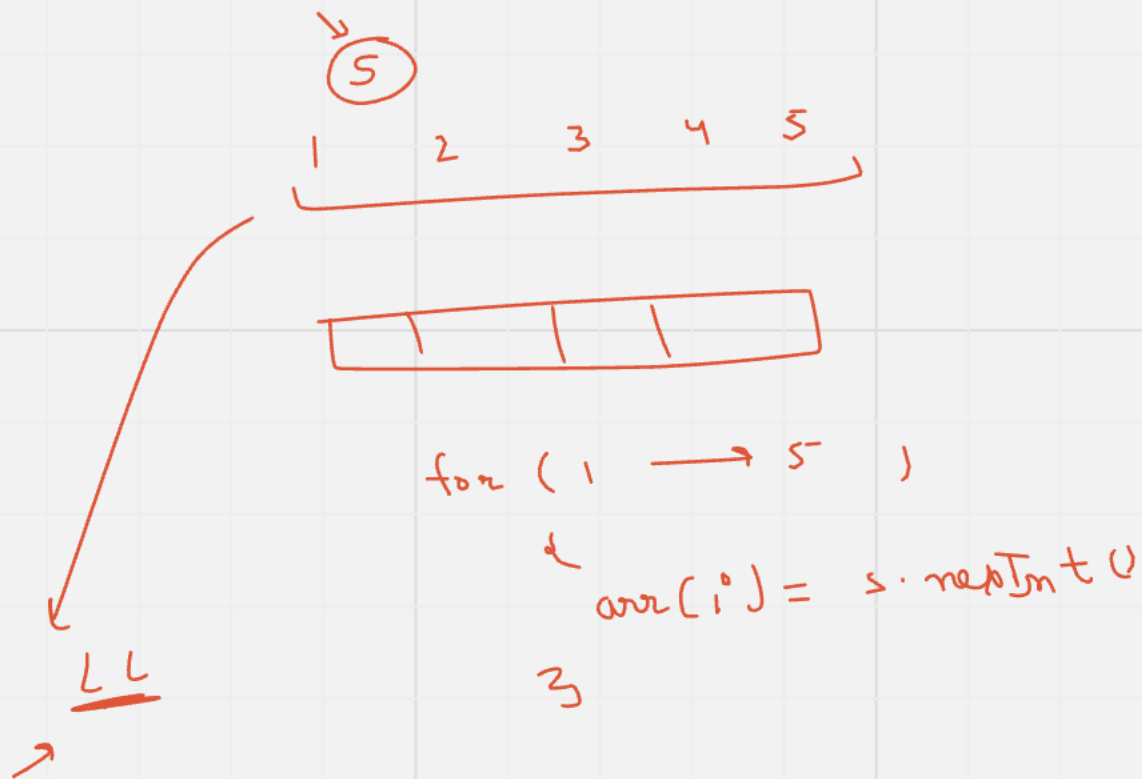
(head.next)

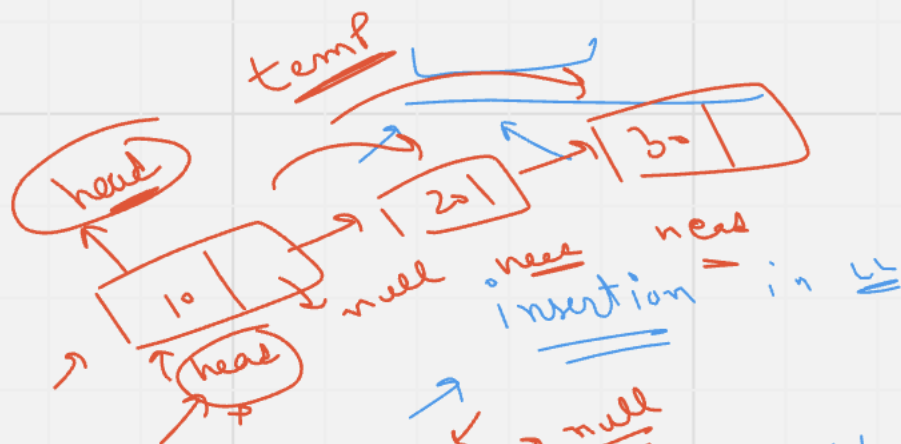




$$1 + 3 \rightarrow 4$$







Node head → null
 Node temp → null

if (head == null)

head = new Node(10);
 temp = head



else {

temp.next = new Node(20);
 temp = temp.next

40

