

L6: Even More Func. Programming

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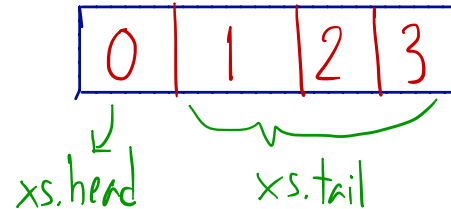
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Let's Pick Up from Tuesday

- We talked about recursion
- We talked about a tuple and a list

Why Is This Code Bad?

- def badMin(xs: List[Int]): Int =
 if (xs.isEmpty) {
 2147483647 // really bad idea but what can i do?
 }
 else if (xs.tail.isEmpty) {
 xs.head
 } else if (xs.head < badMin(xs.tail)) {
 xs.head
 } else {
 badMin(xs.tail)
 }
}
- What if we do val x = badMin(List(1,2,3,...,30))
vs. val x = badMin(List(30,29,...,1))



* Call badMin twice

Better Code

- ```
def betterMin(xs: List[Int]): Int =
 if (xs.isEmpty) {
 2147483647 // really bad idea but what can i do?
 } else if (xs.tail.isEmpty) {
 xs.head
 } else {
 val tailMin = betterMin(xs.tail)
 if (xs.head < tailMin) xs.head else tailMin
 }
```
- This code call the function once, instead of twice
- Still, we have not handled the empty list case
  - Options come to the rescue!

# Options

- Option is a type
  - Option[T]
- Think of it as Option[T] is either
  - None expressing emptiness (Nothing in here)
  - Some(v: T) keeping a value v of type T

# Options – Usage and Examples

- `val x: Option[String] = None`
- `val y: Option[String] = Some("hi")`  
input
- `val z: Option[(Double, String)] = Some((3.14, "Pi"))`
- `val q: Option[List[Double]] = Some(List(3.1, 2.5, 9.0))`
- `val r: Option[List[Double]] = None`

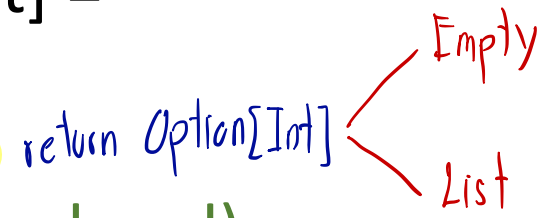
# Options – Accessing Options' Values

- If `t: Option[T]` then
- `t.isEmpty: Boolean` and `t.nonempty: Boolean` indicates whether `t` is empty or non-empty
- If `t` is non-empty and `Some(v: T)` then `t.get` evaluates to `v`
  - Throws an exception otherwise
- To avoid the exception, you can use `t.getOrElse(whenEmpty: T): T`
  - This is similar to `t.get`, but if empty the expression evaluates to `whenEmpty`
- Pattern matching also works with options

```
def addOne(x: Option[Int]): Option[Int] = x match {
 case None => None
 case Some(value) => Some(1+value)
}
```

# Let's Fix BetterMin

- Let's get rid of that one weird case when xs is empty
- ```
def betterMin(xs: List[Int]): Option[Int] =  
  if (xs.isEmpty) None  
  else { val tlAns = betterMin(xs.tail)  
    if (tlAns.nonEmpty && tlAns.get < xs.head)  
      tlAns  
    else  
      Some(xs.head)  
    }
```


- We can also separate the empty and non-empty cases

Let's Fix BetterMin

- ```
def betterMin(xs: List[Int]): Option[Int] =
 if (xs.isEmpty) None else {
 def minNonEmpty(ys: List[Int]): Int =
 if (ys.tail.isEmpty)
 ys.head
 else {
 val tlAns = minNonEmpty(ys.tail)
 if (tlAns < ys.head) tlAns else ys.head
 }
 Some(minNonEmpty(xs))
 }
```

# Tail Recursion

- Let's see how to evaluate `fac(4)` from  
`def fac(n: Int): Int =`  
    `if (n==0) 1 else n * fac(n - 1)`
- See how the expression keeps growing?
- During this time, Scala needs to remember all these values - *Take memories*
- **Q:** Can we rewrite the code so that it does not grow?

# Tail Recursion

- Let's use tail recursion

- ```
def facTail(n: Int) = {  
  def tailFac(n: Int, prod: Int): Int =  
    if (n==0) prod else tailFac(n-1, prod*n)  
  tailFac(n, 1)  
}
```

$f(4) \rightarrow \text{tailFac}(4, 1) \rightarrow \text{if}(4 == 0) \ 1 \ \text{else} \ \text{tailFac}(4-1, 1 \cdot 4)$

$\text{tailFac}(3, 4) \rightarrow \text{if}(3 == 0) \ 4 \ \text{else} \ \text{tailFac}(3-1, 4 \cdot 3)$

$\text{tailFac}(2, 12) \rightarrow \text{if}(2 == 0) \ 12 \ \text{else} \ \text{tailFac}(2-1, 12 \cdot 2)$

$\text{tailFac}(1, 24) \rightarrow \text{if}(1 == 0) \ 24 \ \text{else} \ \text{tailFac}(1-1, 24 \cdot 1)$

$\text{tailFac}(0, 24) \rightarrow \text{if}(0 == 0) \ 24 \ \text{else} \ \text{tailFac}(-1, 0)$

$\rightarrow 24$

Accumulator

Tail Recursion

- The **difference** here is that the last line of the function is a call to a function
 - Not to itself but to a tail call
 - This is call *tail recursive*
- Benefits:
 - **Stack frames can be recycled**
 - Compile to a very nice iterative program with no additional state on each stack call
 - Reduce burden on the compiler
- In the previous example, *prod* is the accumulator
 - Accumulate the answer we have so far instead of waiting for the call to return

More Example

- How can I make a tail recursive out of
def sum(xs: List[Int]): Int =
 if (xs.isEmpty) 0 else xs.head + sum(xs.tail)
- def sum(xs: List[Int]): Int = {
 def tailSum(ys: List[Int], acc: Int): Int =
 if (ys.isEmpty) — Base Case
 acc
 else
 tailSum(ys.tail, acc + ys.head)
 tailSum(xs, 0)
}

Common Things in Tail Recursion

- When we are at the base case, the helper function returns the accumulator
- Accumulator stores the partial computation we have seen so far
- This tail-call is very similar but more general to a while loop
 - Why? Tail call can actually call to other functions

Before We Leave Today

In-class Exercise 6

- Write a function ***def find(xs: List[(Int, String)], key: Int): Option[String]*** that takes in a list of key-value (Int,String)-pairs and returns the string value matching the given integer key. It should return None if nothing matches it.
- Write a function ***def rev(xs: List[Int]): List[Int]*** that takes a list and produces the reverse of the input list. Can you write it as a tail-recursive function?
- Write a function ***def fib(n: Int): Long*** that computes the n-th Fibonacci number in a tail-recursive manner.