# ON DOUBLE-NEGATION SHEAVES IN THE COPRESHEAF TOPOS ON A COFREE COMONOID

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This note fixes a polynomial functor  $p \in \mathbf{Poly}$  and studies the copresheaf topos  $\mathscr{C}_p := \mathscr{F}_P\text{-}\mathbf{Set}$ ; we fix  $\Omega \in \mathscr{C}_p$  the subobject classifier (given by the copresheaf of sieves on  $\mathscr{F}_p^{\mathrm{op}}$ ).

### 1. Presheaves

1.1.  $\mathscr{C}_p$  as a topos of dynamical systems. A presheaf  $F \in \mathscr{F}_p$  consists of the following data (subject to no restrictions):

- (1) For each p-tree  $i \in \mathscr{F}_p$ , a set F(i).
- (2) For each morphism  $f: i \to j$  in  $\mathscr{F}_p$  corresponding with an inclusion of a subtree at height one, a function  $F(f): F(i) \to F(j)$ .

This corresponds with a "dynamical system" with "positions with future knowledge" corresponding with p-trees, states in each position corresponding with each set F(i), and "transition functions" between the states in these positions where one moves to a possible "next state" according to the futures predicted by the p-tree structure. I will henceforth refer to a presheaf  $F \in \mathscr{F}_p$  as a p-dynamical system.

1.2. **Subsystems.** Note that a morphism of p-dynamical systems  $F \to G$  is precisely a map  $F(i) \to G(i)$  from the i-states of F to the i-states of G for each  $i \in \mathscr{F}_p$ , compatible with the transition functions in F and G. A morphism is monic iff each of the constituent morphisms are monic; that is, a subobject corresponds with a subset of the states which is preserved under the transition functions.

As with any topos, the set of subobjects on a p-dynamical system  $\operatorname{Sub}(F)$  comes equipped with a Heyting algebra stricture; the join and meet are given by union and intersection, and the implication  $A \Longrightarrow B$  is given by **FILL IN THE RELEVANT CHARACTERIZATION**. In particular, the negation  $\neg A = (A \Longrightarrow 0)$  is given by the largest dynamical system contained in the set-theoretic complement of A, and hence  $\neg \neq A$  is given by the largest dynamical system whose complement is the same as A.

As with any topos,  $\neg\neg:\Omega\to\Omega$  is a Lawvere-Tierney topology on  $\mathscr{C}_p$ . We will seek to characterize the sheaves with respect to this topology, for which it will be useful to name the following dynamical (sub)systems.

# Example 1.1:

Let  $s \in F(i)$  be a state in a dynamical system. Define the system

$$U_s(j) = \begin{cases} 1 & \exists f : i \to j, \\ 0 & \text{otherwise.} \end{cases}$$

with transition functions are canonically defined. There is a unique monic  $U_s \rightarrow F$  sending  $U_s(i)$  to s. Lemma 1.2. Any dynamical system F is generated by the associated subsystems  $U_s$ ; that is,

$$F = \bigcup_{s \in \coprod_i F(i)} U_s.$$

It is perhaps troublesome that  $U_s$  is defined in reference to a particular state of a dynamical system, as the systems themselves depend only on the position that the state resides in. We can simply define  $U_i$  for position i by the above description; these are contained in the following example.

<sup>&</sup>lt;sup>1</sup>This sentence makes no sense.

## Example 1.3:

The all-ones system  $1_p$  has  $1_p(i) = \{i\}$  for all  $i \in \mathscr{F}_p$ , with transition functions given by the unique endomorphism of 1. Note that there are unique monics  $U_i \rightarrowtail V_i \rightarrowtail 1_p$ .

Consider mentioning that they generate all subobjects under union—they always do, but this time it's particularly simple.

# 2. Sheaves

2.1. **Density.** The following lemma is well known:

**Lemma 2.1.** In a presheaf topos, a subobject  $A \subset C$  is  $\neg \neg$ -dense iff all nonzero subobjects  $0 \subsetneq B \subset C$  intersect A.

This powers the following (easy) proposition, which shows that dense subsystems correspond intuitively with attainable win conditions:

**Proposition 2.2.** A subsystem  $F \subset G$  is dense iff there is no state  $s \in F(i)$  and sequence of morphisms  $i = i_0 \xrightarrow{f_0} i_1 \xrightarrow{f_1} i_2 \xrightarrow{f_2} \cdots$  such that  $f_n(s) \notin G$  for all  $n \in \mathbb{N}$ .

*Proof.* THIS NEEDS TO BE FIXED, AND IS ACTUALLY NOT TRUE. IT NEEDS TO INCORPORATE ALL FUTURES. SEE THE EXAMPLE OF THE  $y^2+1$  TREE WHERE THE ROOT DIRECTIONS POINT TOWARDS THE CONSTANT TREE AND ITSELF; YOU CAN NEGLECT THE ENTIRE LEFT BRANCH IN A DENSE SUBSYSTEM. Suppose there is such a state s and sequence I; then, the subobject  $U_{s,I} \subset G$  is nonzero and does not intersect F, so it is not dense.

Now suppose that no such sequence exists, and suppose  $B \subset G$  is a nonzero subobject containing a state  $s \in B(i)$ . Then, picking some sequence  $I = i_0 \to i_1 \to \cdots$ , we have  $B \cap F \supset U_{s,I} \cap F \neq 0$ , so F is dense.

The following dense monics will be useful.

## Example 2.3:

Let F be a p-dynamical system. Suppose  $s \in F(i)$  is a state such that there is no endomorphism  $f: i \to i$  sending F(f)(s) = s. Then, we may define the *strict futures of s* by

$$\widetilde{U}_s(j) := \{ t \in U_s(j) \mid t \neq s \}.$$

It quickly follows from Proposition 2.2 that the monic  $U_{f_n(s),I_{>n}} \rightarrow U_{s,I}$  is dense.

One characterization of the closure of  $F \subset G$  is the largest closed subobject of G containing F as a dense subobject. Using this and Proposition 2.2, we have the following lemma.

**Lemma 2.4.** Let  $F \subset G$  be a p-dynamical subsystem. Then, we may compute the closure of F as CHAR-ACTERIZATION GOES HERE.

$$\neg \neg F$$

2.2. Characterizing sheaves when p is linear. For now I'll cover a special case.

**Proposition 2.5.** Suppose p = Ay + B is a linear polynomial. Then, a p-dynamical system F is a  $\neg \neq$ -sheaf iff every morphism  $f: i \to j$  in  $\mathscr{F}_p$  induces a bijective transition function F(f).

We'll heavily use the following lemma.

**Lemma 2.6.** Let  $f: i \to j$  be a morphism in  $\mathscr{F}_p$  for p a linear polynomial.

- (i) Suppose F(f) fails to be injective for some  $f: i \to j$ . Then, there is some morphism  $h: i \to k$  and two distinct elements distinct elements  $s, t \in F(i)$  such that  $F(g)(s) \neq s$  for all non-identity endomorphisms  $g: i \to i$  and F(f)(s) = F(f)(t).
- (ii) Suppose F(f) fails to be surjective for some  $f: j \to i$ . Then, there is some  $h: k \to i$  and element  $s \in F(i) \operatorname{im} h$  such that  $F(g)(s) \neq s$  for all non-identity endomorphisms  $g: i \to i$ .

<sup>&</sup>lt;sup>2</sup>Is this true?

*Proof.* The lemma is immediate if there are no non-identity endomorphisms of i, so suppose that i has such an endomorphism g. Since p is linear, we have  $\operatorname{End}(i) = \{g^n\}_{n \in \mathbb{N}}$ .

Suppose contrapositively that all  $s \in F(i)$  have some  $n_s$  where  $F(g)^{n_s}(s) = s$ . Then,  $s \in \operatorname{im} F(g)^{n_s} \subset \operatorname{im} F(g)$ , so F(g) is surjective. Conversely, if s, t have F(g)(s) = F(g)(t), then we have

$$s = F(g)^{n_s n_t}(s) = F(g)^{n_s n_t}(t) = t$$

so F(g) is bijective, as desired. This implies that  $F(g^n)$  is bijective for each n.

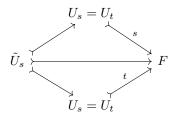
For condition (i), note that there is some composition  $i \xrightarrow{f} j \to i$  expressing f as a prefix of a bijection; hence f is injective.

Now, consider some  $F(f): F(i) \to F(j)$  failing to be injective, so that  $F(g^n)$  fails to be injective for each n, and let  $S:=\{s\in F(i)\mid g^n(s)\neq s\ \forall n\geq 1\}$ . Note that  $F(g)(F(i)-S)\subset F(i)-S$  and  $F(g)|_{F(i)-S}$  is injective (by the above argument) Hence there must be some  $s\in S$  and  $t\in F(i)$  satisfying F(g)(s)=F(g)(t), as desired.

Let  $F(f): F(j) \to F(i)$  fail to be surjective, so that  $F(g^n)$  fails to be surjective for each n; a composition  $i \to j \xrightarrow{f} i$  proves that the subset  $S := \{s \in F(i) \mid g^n(s) \neq s \ \forall n \geq 1\}$  is nonempty. Recall that  $F(g)(F(i) - S) \subset F(i) - S$ , and note that  $F(g)|_{F(i)-S}$  surjects onto F(i) - S; hence there must be some  $s \in F(i) - S$  which is not in the image of g, as desired.

Proof of Proposition 2.5. First suppose that F has bijective transition functions, and suppose that  $A \to E$  is a dense monic and  $\tilde{\varphi}: A \to F$  is a morphism of p-dynamical systems. Then, for each  $s \in E(i)$ , there is a morphism  $f: i \to j$  such that  $f(s) \in A(j)$ ; we define an extension  $\varphi: E \to F$  by  $\varphi(s) := F(f)^{-1}(\tilde{\varphi}(f(s)))$ . This is the unique choice of  $\varphi(s)$  such that  $F(f)(\varphi(s)) = \varphi(E(f)(s))$  for any f, so it is well defined as the unique extension of  $\tilde{\varphi}$  to E.

Conversely, suppose first that F(f) fails to be injective for some  $f: i \to j$ . Then, choose s, t, h as in Lemma 2.6, and note that  $U_s = U_t$  as p-dynamical systems; however, they are distinct subsystems of F. Further, note that there are factorizations (where the leftmost monic is dense)



which prove that the map  $\operatorname{Hom}(E,F) \to \operatorname{Hom}(A,F)$  is not injective.

Next, suppose that F(f) fails to be surjective for some  $f: j \to i$  and choose some  $s \in i$  as in Lemma 2.6. The system  $U_s$  occurs as a dense subobject of some system U who includes a state in position h; however, the morphism  $U_s \to F$  doesn't extend to U, as desired.

This is left here for archival purposes; this argument works for linear systems with states in trees that are "non-referencial."

2.3. Sheaves as ancient systems which remember history. We say that a p-dynamical system F is ancient if all transition functions  $F(i) \to F(j)$  are epic, and say that F remembers history if  $F(i) \to F(j)$  are monic.

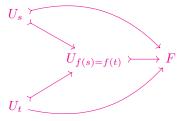
**Proposition 2.7.** A p-dynamical system F is a  $\neg \neg$ -separated presheaf iff it remembers history. A p-dynamical system F is a  $\neg \neg$ -sheaf iff it is ancient and remembers history.

*Proof.* We first verify the statement for separated presheaves. We have to verify that, F remembers history iff for all dense monics  $A \mapsto E$ , the induced map

$$\operatorname{Hom}_{\mathscr{C}_p}(E,F) \to \operatorname{Hom}_{\mathscr{C}_p}(A,F)$$

is monic.

First suppose that F does not remember history, and there are two distinct states  $s, t \in F(i)$  and a morphism  $f: i \to j$  such that f(s) = f(t). Then, since  $U_s = U_t$  as (dense) subobjects of  $U_{f(s)} = U_{f(t)}$ , we have two distinct extensions



so that F is not separated.

Conversely, suppose that F does remember history, and suppose we have two morphisms  $\varphi, \varphi' : E \to F$  which each restrict to A identically. At position j possessing morphism  $f : j \to i$ , this is represented by commuting parallel morphisms

$$E(i) \longrightarrow F(i)$$

$$\uparrow \qquad \qquad \downarrow$$

$$E(j) \Longrightarrow F(j)$$

which must coincide since the morphism  $F(j) \rightarrow F(i)$  is monic.<sup>3</sup>

For  $\neg\neg$ -sheaves, it suffices to prove that the morphism  $\operatorname{Hom}_{\mathscr{C}_p}(E,F) \to \operatorname{Hom}_{\mathscr{C}_p}(A,F)$  is epic iff F is ancient. This will work by a similar argument; if F is not ancient, we may pick a state  $s \in \coprod_i F(i)$  and morphism  $f: j \to i$  such that F(j) is nonempty (say, containing  $t \in F(j)$ ) and such that  $s \notin \operatorname{im} f$ , and consider the embedding  $U_s \to F$ . Then,  $U_s$  is isomorphic as a p-dynamical system to  $U_{f(t)} \subset U_t$ , but there is no morphism  $U_t \to F$  extending the given morphism.

Conversely, if F is ancient, we probably have to assume AoC, and then just choose preimages to extend a morphism  $A \to F$  to E. I'll fill this in fully later.

<sup>&</sup>lt;sup>3</sup>This needs to be corrected later; it isn't necessarily the case that any E(i) is totally subsumed by A. Instead, we have to make some argument using colimits, which I don't particularly want to do..