System Calls

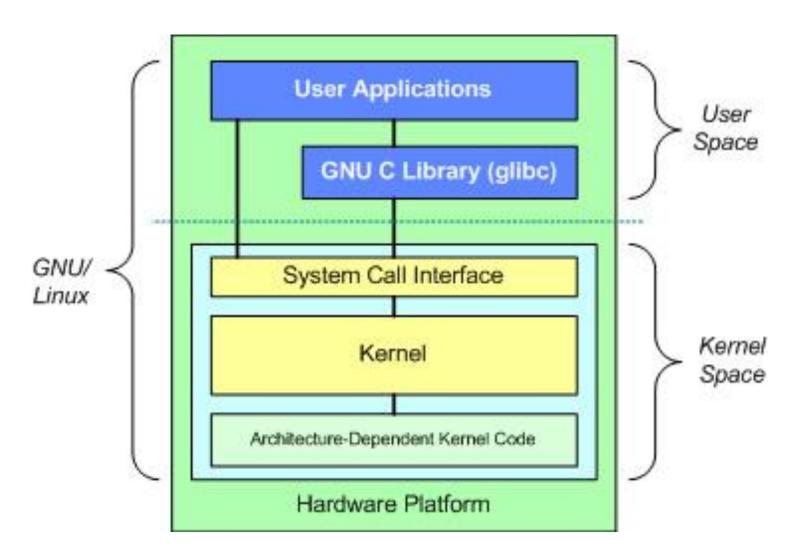
Using Tanenbaum's Modern Operating Systems (3rd edition)

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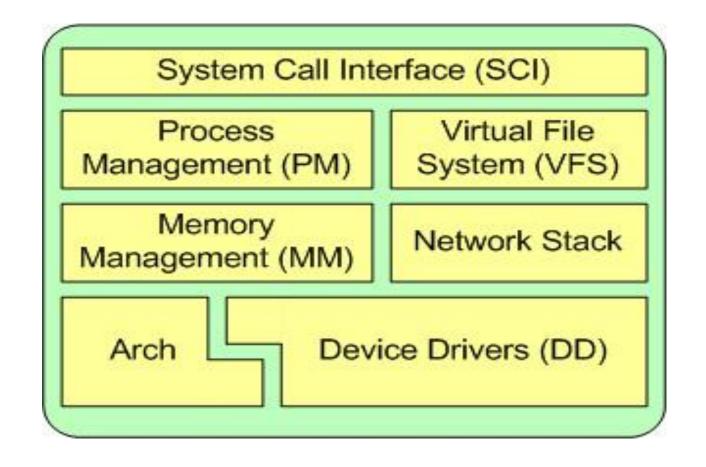
System Call

- What is a system call?
 - API to the OS
 - Requests a service from an operating system's kernel
- Two CPU modes :
 - User mode:
 - Limits the address space of the program.
 - Prevents the application from directly using devices.
 - Kernel mode:
 - All is allowed.
- POSIX defines about 100 system calls:
 - open, read, write, close, wait, execve, fork, exit, kill

Linux OS Architecture



Linux OS Architecture (II)



glibc

- The C standard library provides macros, type definitions, and functions:
 - Basic manipulation:
 - String handling.
 - Mathematical computations.
 - System calls:
 - I/O processing.
 - Memory allocation.
 - Other OS services.

glibc (II)

- By calling a system-call:
 - Sets the parameters and TRAP instruction, which:
 - Switches into kernel mode transfers control to OS
 - Jumps to a single fixed location

• "fork" and "execve" are glibc functions that in turn call the "fork" and "execve" system-calls.

System Calls (I)

- Process control:
 - load
 - execute
 - create process
 - terminate process
 - get/set process attributes
 - wait for time, wait event, signal event
 - allocate, free memory

System Calls (II)

- File management:
 - create file, delete file
 - open, close
 - read, write, reposition
 - get/set file attributes
- Device management:
 - attach or detach devices
 - read, write, reposition
 - get/set device attributes

System Calls (III)

• Information:

- get/set time or date
- get/set system data
- get/set process, file, or device attributes

Communications:

- create, delete communications connection
- send, receive messages
- transfer status information
- attach or detach remote devices

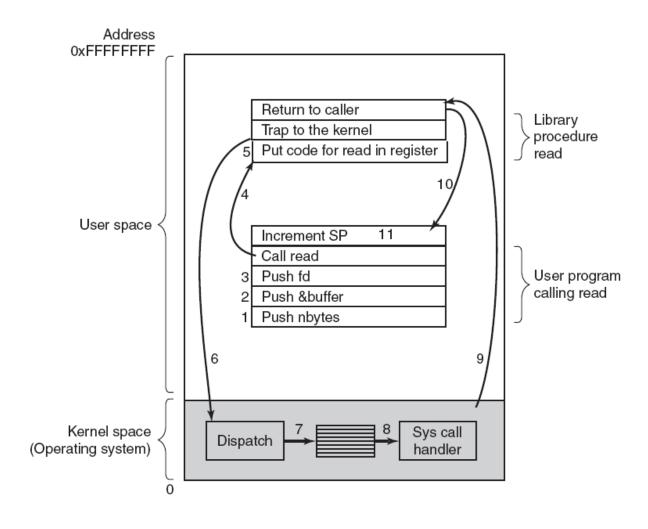
System Call Flow

C program:

```
...
count = read(fd,buffer,nbyts);
...
```

- counts is the number of read bytes.
- if failed count=-1 and errno is set.
- The library function read is called
- The system call number is put into a register
- A TRAP instruction is executed to switch to kernel mode
- The kernel starts at a fixed address
- The kernel executes the system call handler that is a function pointer assigned to the specific system call number
- The handler is executed and control is returned to the user-space library at the instruction after the TRAP

System Call (11 steps)



Blocking System Call

Example:

- The user calls read() from the keyboard but nothing has been typed yet.
- What happens?
 - The system call will run the kernel code for blocking the process until a keyboard interrupt arrives
 - In the meantime, another process is taken from the Ready Queue to be executed

POSIX System Calls (I)

Process management

Call	Description
pid = fork()	Create a child process identical to the parent
pid = waitpid(pid, &statloc, options)	Wait for a child to terminate
s = execve(name, argv, environp)	Replace a process' core image
exit(status)	Terminate process execution and return status

File management

Description
Open a file for reading, writing, or both
Close an open file
Read data from a file into a buffer
Write data from a buffer into a file
Move the file pointer
Get a file's status information

POSIX System Calls (II)

Directory and file system management

Call	Description
s = mkdir(name, mode)	Create a new directory
s = rmdir(name)	Remove an empty directory
s = link(name1, name2)	Create a new entry, name2, pointing to name1
s = unlink(name)	Remove a directory entry
s = mount(special, name, flag)	Mount a file system
s = umount(special)	Unmount a file system

Miscellaneous

Call	Description
s = chdir(dirname)	Change the working directory
s = chmod(name, mode)	Change a file's protection bits
s = kill(pid, signal)	Send a signal to a process
seconds = time(&seconds)	Get the elapsed time since Jan. 1, 1970

Homework + Interview Questions

- What is "system call"?
- Explain in detail how system call is executed
- What is TRAP?