

CSE 332 Spring 2023 Midterm

Name: _____

Email address (UWNetID): _____

Instructions:

- The allotted time is 50 minutes.
- Please do not turn the page until the staff says to do so.
- This is a closed-book and closed-notes exam. You are NOT permitted to access electronic devices including calculators.
- Read the directions carefully, especially for problems that require you to show work or provide an explanation.
- When provided, write your answers in the box or on the line provided.
- Unless otherwise noted, every time we ask for an O , Ω , or Θ bound, it must be simplified and tight.
- Unless otherwise noted, assume all data structures are implemented as described in lecture.
- For answers that involve bubbling in a or , make sure to fill in the shape completely.
- If you run out of room on a page, indicate where the answer continues. Try to avoid writing on the very edges of the pages: we scan your exams and edges often get cropped off.
- A formula sheet has been included at the end of the exam.

Advice

- If you feel like you're stuck on a problem, you may want to skip it and come back at the end if you have time.
- Look at the question titles on the cover page to see if you want to start somewhere other than problem 1.
- **Relax and take a few deep breaths. You've got this! :-).**

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Total: 85 points

Q1: Short-answer questions (20 pts)

- For questions asking you about runtime, give a simplified, tight Big-O bound. This means that, for example, $O(5n^2 + 7n + 3)$ (not simplified) or $O(2^n)$ (not tight enough) are unlikely to get points. Unless otherwise specified, all logs are base 2.
- For questions with a mathematical answer, you may leave your answer as an unsimplified formula (e.g. $7 \cdot 10^3$).

We will only grade what is in the provided answer box.

- a. Give the **worst-case** runtime for inserting a value into a binary search tree containing N elements.

$O(\boxed{\quad})$

- b. Give a simplified, tight big-O bound for: $T(N) = T(N/2) + 16$, given $T(1) = 1$

$O(\boxed{\quad})$

- c. Give a simplified, tight big-O bound for: $f(N) = \log^2(N) + 20\log(N^2)$

$O(\boxed{\quad})$

- d. Give the **worst-case** runtime for inserting a value into an AVL tree containing N elements, where the value being inserted is smaller than the current smallest value in the tree.

$O(\boxed{\quad})$

- e. Give a simplified, tight big-O bound for: $f(N) = N \log \log(8^N) + N(\log(N))(\log(N))$

$O(\boxed{\quad})$

Q1: (continued)

(Same instructions as on previous page)

- f. Give the **worst-case** runtime for removing the next element from a CircularArrayFIFOQueue (from P1) containing N elements.

$$O(\boxed{\quad})$$

- g. Give the **worst-case** runtime for deleteMin() from a binary min heap **containing 2^N elements**.

$$O(\boxed{\quad})$$

- h. Suppose a 5-ary heap is stored in an array with the root in cell 0. Give the formula for computing the index of the cell containing the third (aka middle) child of the node in cell x. Give your answer in terms of x.

$$\boxed{\quad}$$

- i. Given M = 5 and L = 13, what is the **minimum** number of leaf nodes in a B-tree with h = 7?

$$\boxed{\quad}$$

- j. Give the **smallest value for c** that could be used **with $n_0 = 1$** in a proof to show that:

$$3n^2 + 5 \text{ is } O(n^2)$$

$$\text{Given } n_0 = 1$$

$$c = \boxed{\quad}$$

Q2: Code Analysis (16 pts)

Describe the worst-case running time for the following pseudocode functions in Big-O notation in terms of the variable n. Your answer **MUST** be tight and simplified. **You do not have to show work or justify your answers for this problem.**

```
a) void summer(int n) {
    for (int i = 0; i*i < n; i++) {
        for (int j = n; j >= 1; j = j/3) {
            System.out.println("Enjoy the sunshine!");
        }
    }
}
```

$O(\boxed{\quad})$

```
b) void happy(int n) {
    if (n % 3 == 0) {
        for (int i = 456; i >= 0; i -= 5) {
            for (int j = i; j >= 0; j -= 2) {
                System.out.println("Yipee!");
            }
        }
    } else {
        int p = 0;
        for (int i = 1; i < n; i *= 2) {
            p += 1;
        }
        for (int j = 1; j < p; j *= 2) {
            System.out.println("Smile!");
        }
    }
}
```

$O(\boxed{\quad})$

Q2: (continued)

```
c) AVLTree sunny(int n, BinaryMinHeap input) {  
    // The size of input is n  
    AVLTree output = new AVLTree();  
    for (int i = 0; i < n; i++) {  
        int val = input.deleteMin();  
        output.insert(val);  
    }  
    return output;  
}
```

$O(\boxed{\quad})$

```
d) void funny(int n) {  
    for (int i = 0; i < n; i++) {  
        for (int j = i; j >= 0; j--) {  
            int z = 3 * n;  
            for (int k = 0; k < z; k += 3) {  
                System.out.println("Ha ha ha!");  
            }  
        }  
    }  
}
```

$O(\boxed{\quad})$

Q3: O, Ω, and Θ (9 pts)

For each of the following statements, indicate whether it is always true, sometimes true, or never true. You do not need to include an explanation. Assume that the domain and codomain of all functions in this problem are natural numbers(1, 2, 3 ...).

- a) If $f_a(n)$ is the worst case runtime for insert in a Binary Search Tree and $g_a(n)$ is the worst case runtime for insert in an AVL tree, then $f_a(n)$ is $\Omega(g_a(n))$ (Big-Omega)

Always

Never

Sometimes

- b) $f_b(n)$ is $O(\log((f_b(n))^2))$

Always

Never

Sometimes

- c) If $f_c(n)$ is the worst case runtime of merging three binary min heaps, each with n elements, and $g_c(n)$ is the worst case runtime of inserting $3n$ elements into a binary min heap, then $f_c(n)$ is $O(g_c(n))$

Always

Never

Sometimes

Q4: Write a Recurrence (8 pts)

Give a base case and a recurrence for the runtime of the following function. Use variables appropriately for constants (e.g. c_1 , c_2 , etc.) in your recurrence (you do not need to attempt to count the exact number of operations). **YOU DO NOT NEED TO SOLVE** this recurrence.

```
int MI6(int n) {
    if (n <= 7) {
        for (int i = 0; i < n; i++) {
            System.out.println("The name's Bond, James Bond");
        }
        return n + 1;
    } else {
        if (n % 1 == 1) {
            return MI6(n + 1);
        } else {
            int x = n;
            while (x != 0) {
                System.out.println("Double O-" + x);
                x -= 1;
            }
            for (int i = 1; i < n; i *= 3) {
                System.out.println("Shaken not stirred");
            }
            return MI6(n / 2) + MI6(n - 3) + (n * n);
        }
    }
}
```

$$T(n) = \underline{\hspace{10cm}} \text{ For } n \leq 7$$

$$T(n) = \underline{\hspace{10cm}} \text{ For } n > 7$$

Yipee!!!! YOU DO NOT NEED TO SOLVE this recurrence...

Q5: Solve a Recurrence (10 pts)

Suppose that the running time of an algorithm satisfies the recurrence given below. Find the closed form for T(N). **You may assume that N is a power of 3.** Your answer should **not** be in Big-Oh notation – show the relevant exact constants and bases of logarithms in your answer (e.g. do NOT use “c₁, c₂” in your answer).

Your final answer must **NOT** have any summation symbols or recursion - you may find the list of summations and logarithm identities on the last page of the exam to be useful.

You must show your work and put your final answer on the line below to receive any credit. (e.g. show your general formula to receive partial credit).

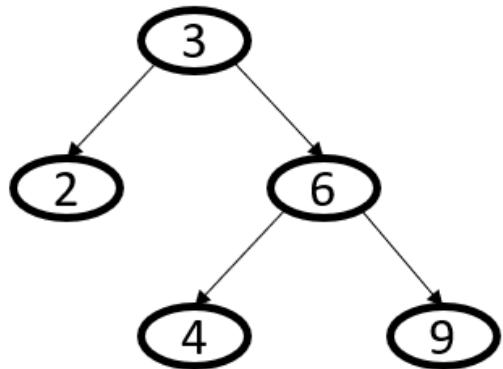
$$T(1) = 5$$

$$T(N) = 3T\left(\frac{N}{3}\right) + N \quad \text{for integers } N > 1$$

Closed Form: _____

Q6: AVL (7 pts)

- a) (3 pts) Draw the AVL tree that results after inserting 7 into this AVL tree. Be sure to draw your final tree in the box below AND **indicate the total number of single and double rotations required for this insertion.**

**Final Tree:**

A large rectangular box intended for drawing the final AVL tree resulting from the insertion of value 7.

Total # of Single Rotations required:

A rectangular box for indicating the total number of single rotations required.

Total # of Double Rotations required:

A rectangular box for indicating the total number of double rotations required.

Q6: (continued)

b) (4 pts) Give the **minimum** and **maximum height** for an AVL tree containing 13 nodes.

Minimum Height: _____

Maximum Height: _____

How many more **nodes** are needed to increase the **minimum** height by 1? _____

How many more **nodes** are needed to increase the **maximum** height by 1? _____

Q7: Heaps (7 pts)

- a. (4 pts) Given a **binary max heap** represented by this array [10, 9, 8, 7, 6, 5, 4, 2, 1], show the resulting heap (in array representation) after inserting the following elements in this order: {11, 33}

Show the final array contents (11 values total, starting from index 0):

0	1	2	3	4	5	6	7	8	9	10

- b. (3 pts) Given the following values: {10, 2, 9, 11, 3, 99, 37, 1}, give the order of insertions into a **binary max heap** that results in the smallest number of `percolateUp` operations.

Order of insertion (8 values total):

_____ , _____ , _____ , _____ , _____ , _____ , _____ , _____

Q8: B-Trees (8 pts)

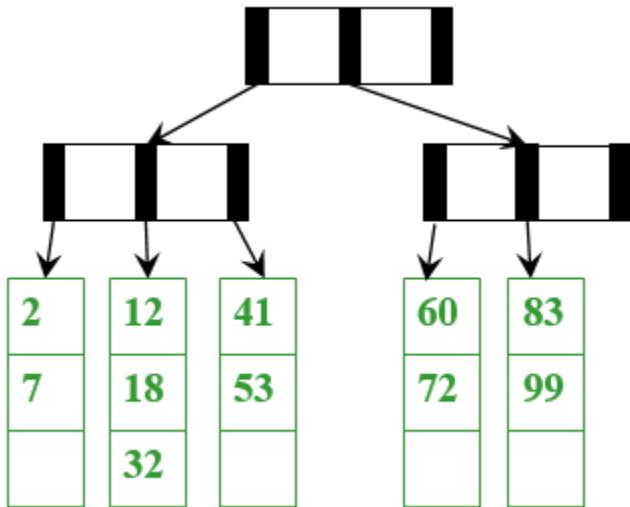
- a) (4 pts) Your friend proposes a modification to the B-tree implementation discussed in lecture: use a "jump search" instead of binary search on internal nodes. You do not need to know anything about the implementation of "jump search", except that the worst case run-time is $O(\sqrt{n})$. Everything else about the find operation on our B-trees remains as we discussed in lecture.

Give a tight big-O running time of the **worst case** of the find operation described above. Keep all factors of M, L, and N in your answer.

Worst case runtime of find: $O(\underline{\hspace{2cm}})$

Q8: (continued)

- b) (1 pt) In the B-tree shown below, **write in the values for the interior nodes:**



- c) (3 pts) Starting with the tree above, draw the tree resulting after inserting 37 into the B-tree. Please circle your final tree.

This is a blank page! Enjoy!