

# Exercise 1: SOLUTIONS

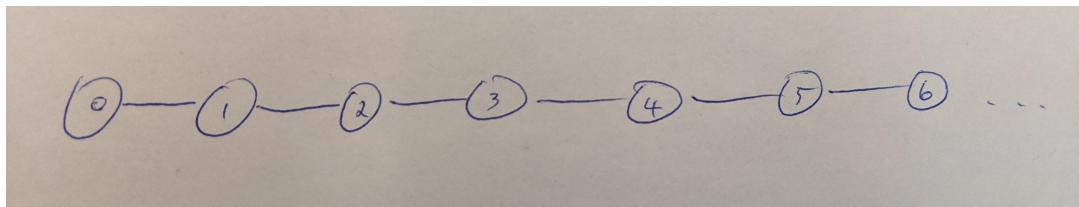
Response by: **SOLUTIONS**

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## Exercise 1-4: Uncountable Sets

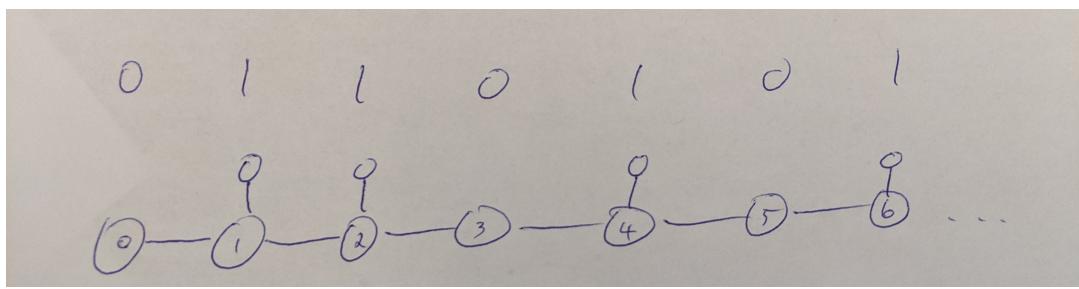
To show that the number of graphs (with a potentially infinite number of vertices) is uncountable, we will provide an injective (one-to-one) mapping from the infinite binary strings (denote with  $\{0, 1\}^\infty$ ) to the infinite graphs (denote with  $G_\infty$ ).

To define this one-to-one mapping, the idea is to have a graph "gadget" to use as the basis to represent a string. We will consider a graph with an infinite number of vertices arranged in a "chain" like this:



In this case we add labels to each vertex for convenience. We will say that vertex  $v_0$  is the (unique) node with degree 1 (i.e. only has one edge connected to it). The vertex adjacent to  $v_0$  we will call  $v_1$ , the vertex adjacent to  $v_1$  that is not  $v_0$  we will call  $v_2$ , and in general will label the vertices adjacent to  $v_i$  as  $v_{i+1}$  and  $v_{i-1}$  (with  $v_{i-1}$  as that vertex that is closer to  $v_0$ ). In this way, we have defined an ordering of the vertices.

We can then map a binary string  $b$  a graph like the one above by adding additional vertices for each 1 that appears in  $b$ . In particular, if bit  $b_i$  in  $b$  is a 1, we will add an additional vertex to the graph that has an edge only to  $v_i$ . The following image illustrates an example of this mapping.



**The mapping is one-to-one:** To show that the mapping is one-to-one, we must show that two different strings map to different graphs. In this case, we know that if string  $b' \neq b''$ , then they must differ by at least one bit. Let's say that they differ at bit  $i$ , and without loss of generality we can say that  $b' = 1$  and  $b'' = 0$ . In this case, if  $b'$  maps to graph  $g'$  and  $b''$  maps to graph  $g''$ , we know that  $g' \neq g''$  because vertex  $i$  in  $g'$  will have degree 3 whereas vertex  $i$  in graph  $g''$  will have degree 2.