

Problem 1

Collaborators: None

Solver: Nathan Leung

Sources Used: None

a) Proof: $\mathbb{Q}(x, y, s, t, u, v)$

> Define a pulvinizer state machine as described in the problem.

Testing predicate \mathbb{Q} on the start state:

$$\mathbb{Q}(a, b, 1, 0, 0, 1)$$

$$P_1(a, b) : \gcd(a, b) = \gcd(a, b) \quad \checkmark$$

$$P_2(a, 1, 0) : (1)a + (0)b = a \quad \checkmark$$

$$P_3(b, 0, 1) : (0)a + (1)b = b \quad \checkmark$$

Overall, \mathbb{Q} holds for the start state!> Assume $\mathbb{Q}(x, y, s, t, u, v)$ holds true, and initiate a transition to $(x', y', s', t', u', v') = (y, r, u, v, s-qu, t-qv)$

$$P_1(x', y') = P_1(y, r) = \gcd$$

$$\gcd(y, r) = \gcd(y, x \bmod y) = \gcd(x, y), \quad \gcd(x, y) = \gcd(a, b)$$

from the assumption, so by transitive property of equality

 P_1 holds true!

$$P_2(x', s', t') = P_2(y, u, v) \text{ which is } P_2(y, u, v) \text{ for the old state exactly, so } P_2 \text{ holds true!}$$

$$P_3(y', u', v') = P_3(r, s-qu, t-qv)$$

$$(s-qu)a + (t-qv)b = sa + tb - q(ua + vb). \text{ By } P_2(x, s', t),$$

$$sa + tb = x, \text{ and by } P_3(y, u, v), ua + vb = y. \text{ Thus}$$

$$x - qy = (s-qu)a + (t-qv)b = r, = y + r$$

 P_3 holds true!> P_1, P_2, P_3 of the newly formed state are all satisfied, thus satisfying \mathbb{Q} for a state after transitions. Thus, \mathbb{Q} is proven to be a preserved predicate \square

b) Proof:

- > Define a pushover state machine as described in the problem. By the first step of proof a), Q holds true for the initial state. Additionally by the conclusion of proof a) we conclude that Q is a preserved predicate.
 - > By the invariant principle, Q holds true when the state machine terminates. In the final state $y = 0$, so
$$P_1(x, 0) := \gcd(x, 0) = \gcd(a, b)$$
$$P_2(x, s, t) := sa + tb = x$$
.
 - > $\gcd(a, b) = \gcd(x, 0) = x = sa + tb$. Thus we conclude s, t in the final state of the pushover satisfies Bézout's identity.
- c) The pushover for variables x, y follow the exact same transitions as the Euclidean algorithm state machine $(x, y) \rightarrow (y, x \bmod y)$, so the pushover machine should terminate after at most the same number of transitions.

problem 2 -

Collaborators : None

Sources Used : None Solver : Nathan Leung

1106 divides 18062 ^{natural}, since $18062 = (11)(1642)$. However, 2025 is not divisible by 11 since $2025 \equiv 1 \pmod{11}$. If a and b are arbitrary natural numbers, 11 still divides 18062b since $18062b = (11)(1642)(b)$ and from $2025 \equiv 1 \pmod{11}$ we get $2025^a \equiv 1^a \pmod{11}$, $2025^a \equiv 1 \pmod{11}$, $2025^a + 1 \equiv 2 \pmod{11}$, so $2025^a + 1$ is not divisible by 11, and thus cannot be equal to 18062b ! \square

Problem 5 -

Collaborators : None

Solver : Nathan Leung

Sources Used : www.wolframalpha.com

a) $n = (13139465087838462013)(16257701292567269201)$

via factor n from wolframalpha

b) $d \equiv e^{-1} \pmod{(p-1)(q-1)}$ where p and q are the 2 primes calculated in a

$$\equiv 172797418847983865496766528110570298307 \pmod{\phi(n)}$$

via wolframalpha

c) $m \equiv \hat{m}^d \pmod{n}$

$$\equiv 577345663350077735577145663 \pmod{n}$$

via wolframalpha