

What can be Computed in a Simple Chaotic Way?

Emanuele Natale

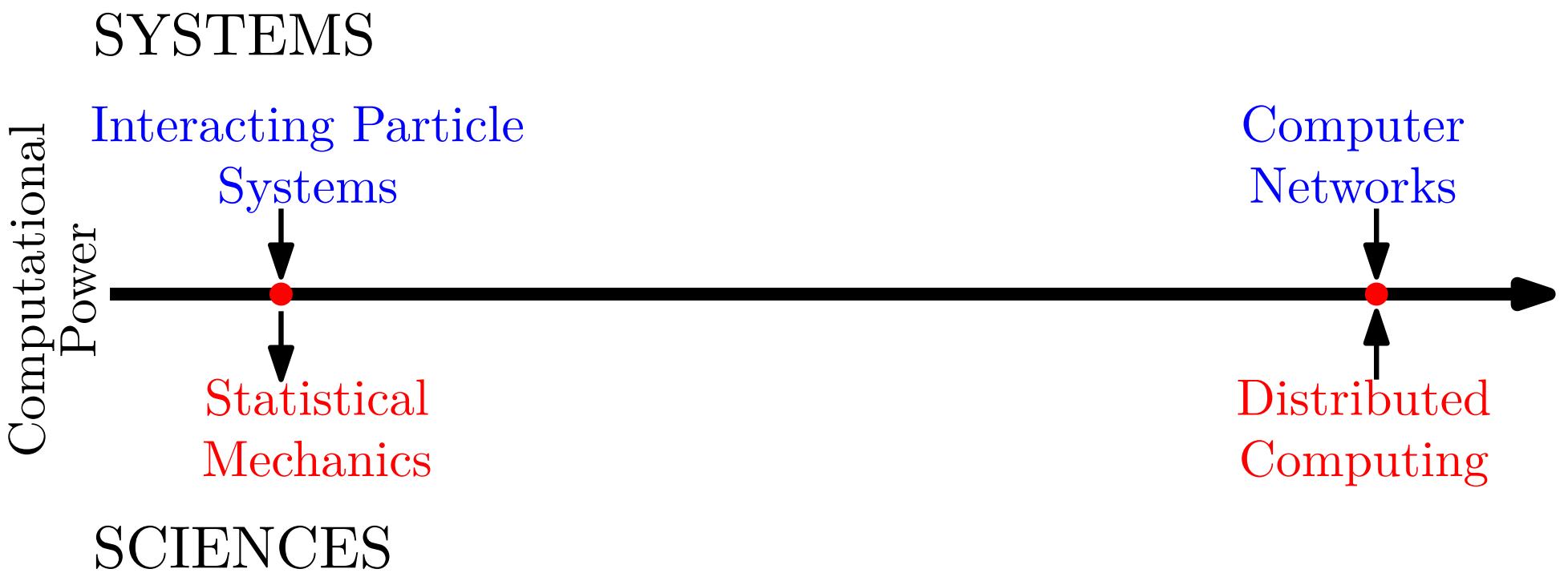


ICTCS 2017

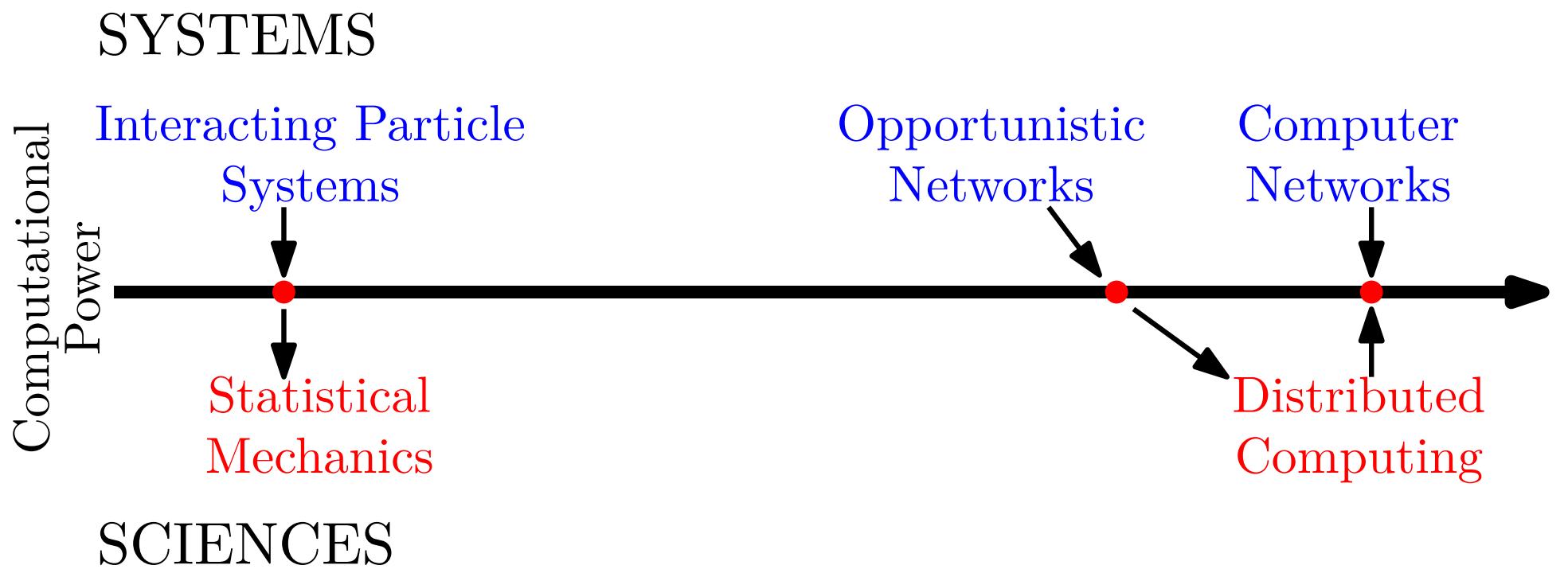
18th Italian Conference on Theoretical Computer Science
26-28 September, Naples, Italy



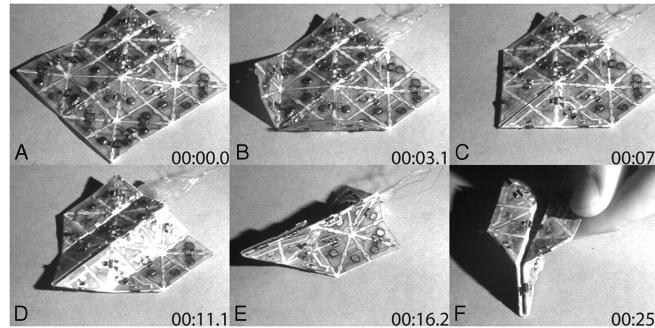
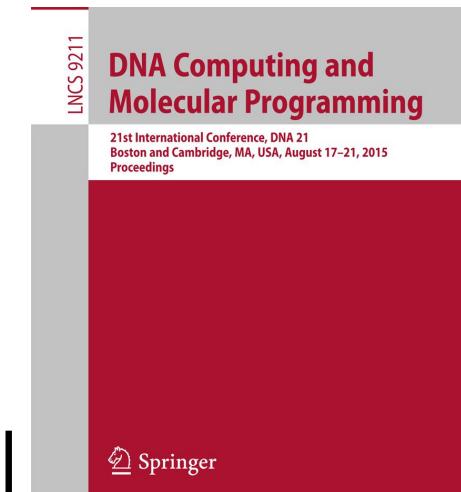
What can *Simple* Systems do?



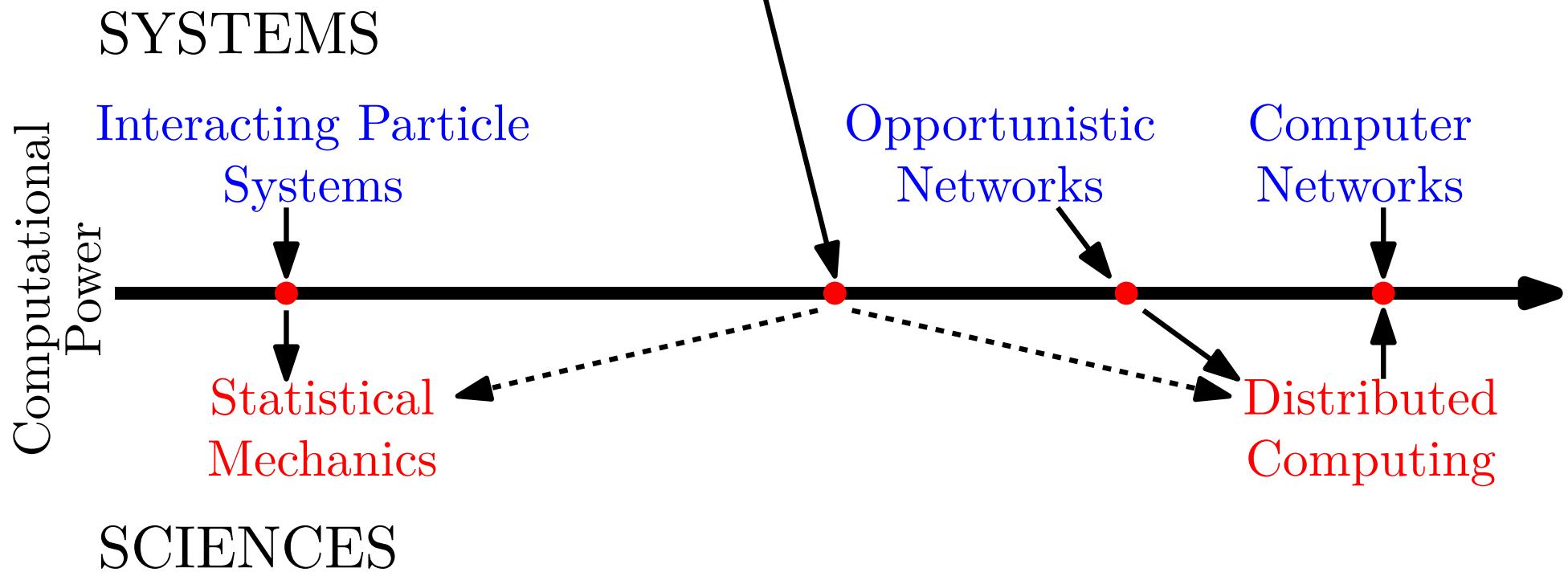
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DNA/Molecular Computing, Programmable Matter, Swarms of Simple Robots



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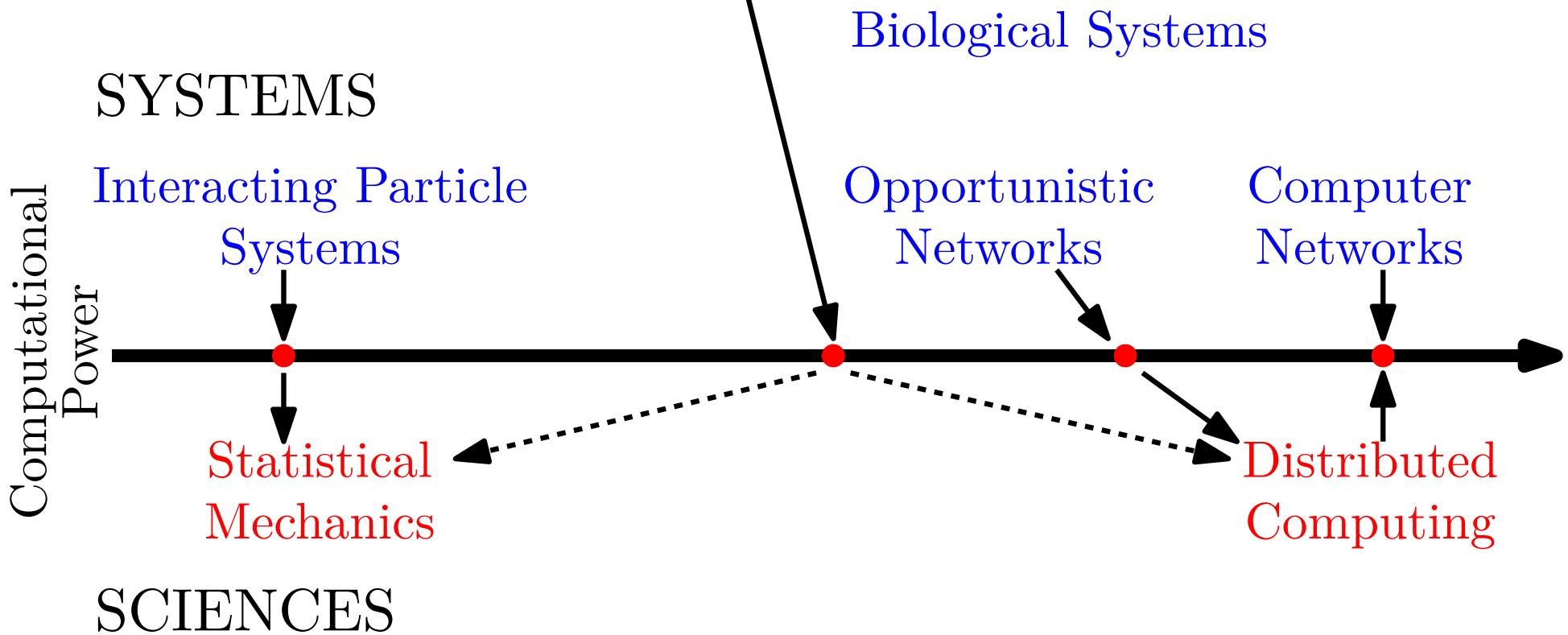


Schools of fish
[Sumpter et al. '08]

Insects colonies
[Franks et al. '02]



Flocks of birds
[Ben-Shahar et al. '10]



Examples of Natural Algorithms



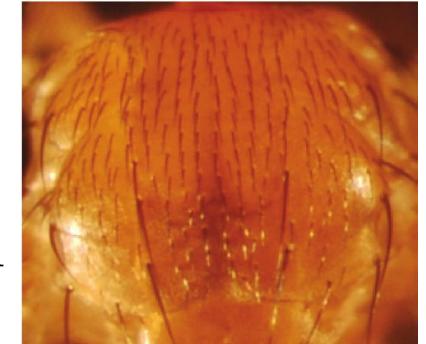
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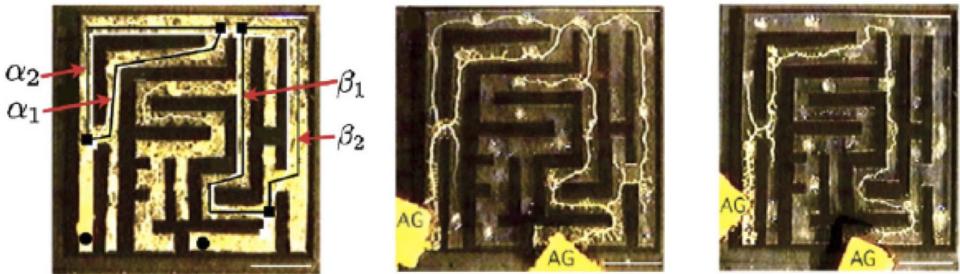
How are sensory
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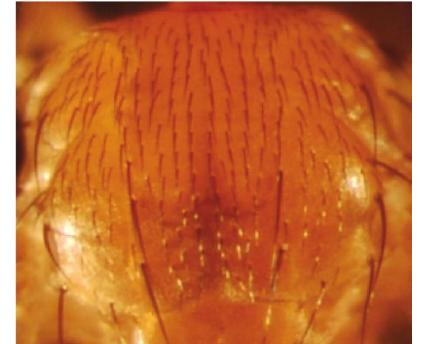


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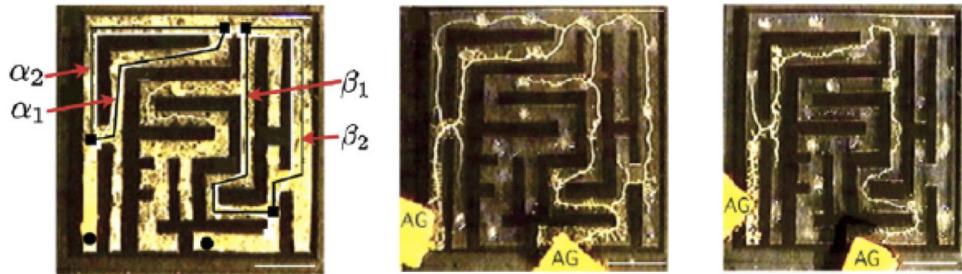
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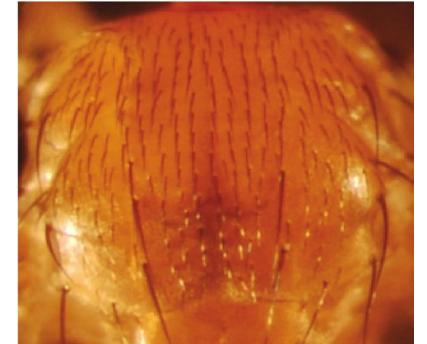


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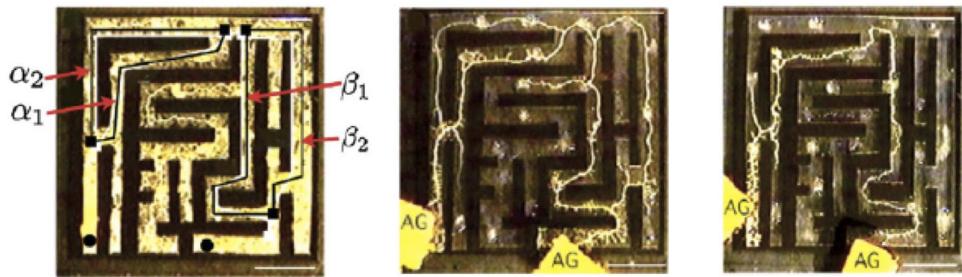
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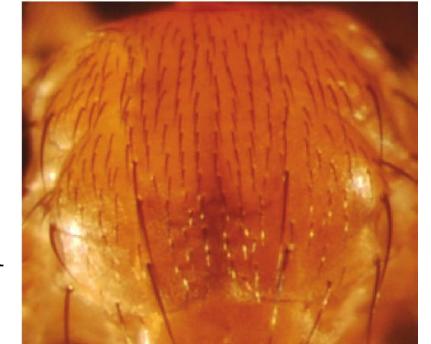


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How ants perform collective navigation [FHBGKKF '16]

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Unstructured Communication Models

Animal communication:

- Chaotic
- Anonymous
- Parsimonious
- Uni-directional
(Passive/Active)
- Noisy

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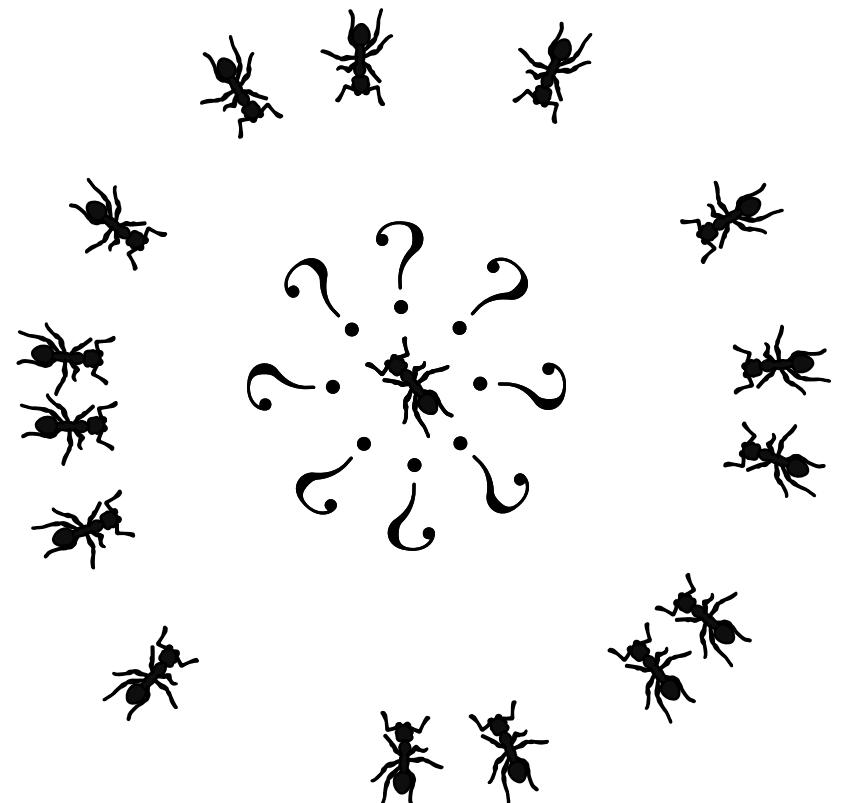
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$PULL(h, \ell)$ model

[Demers '88]: at each round each agent can *observe* h other agents chosen independently and uniformly at random, and *shows* ℓ *bits* to her observers.



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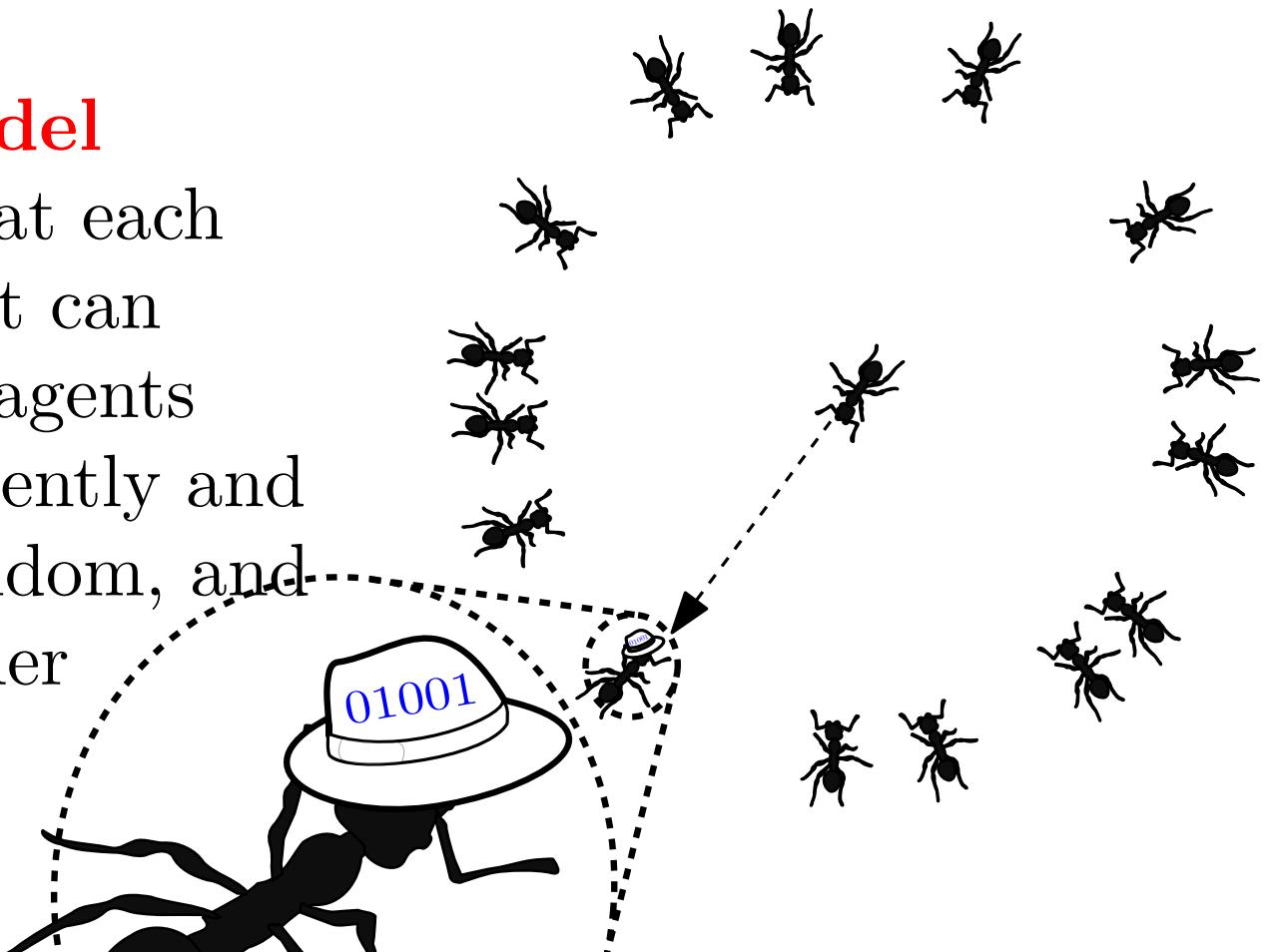
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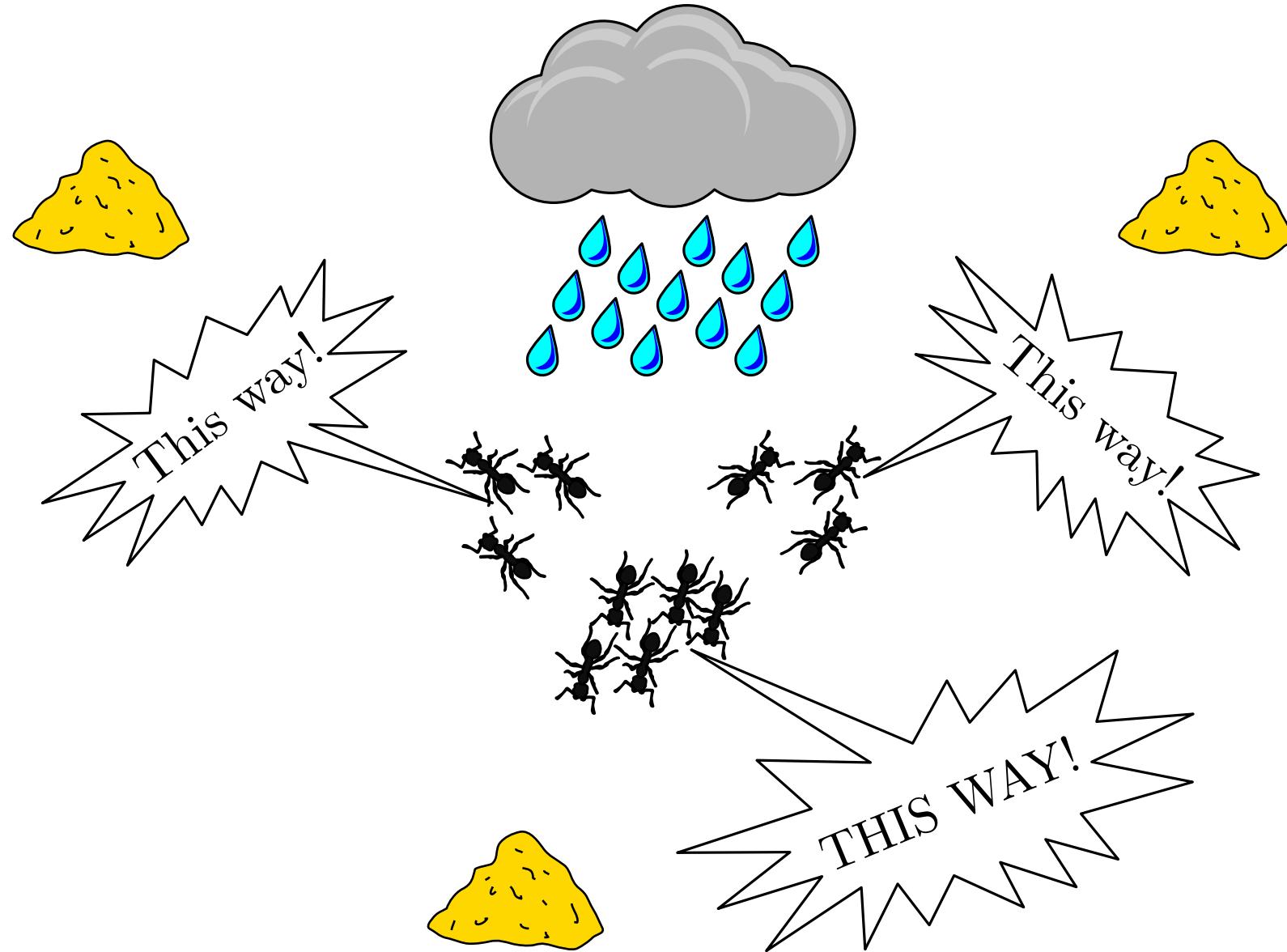
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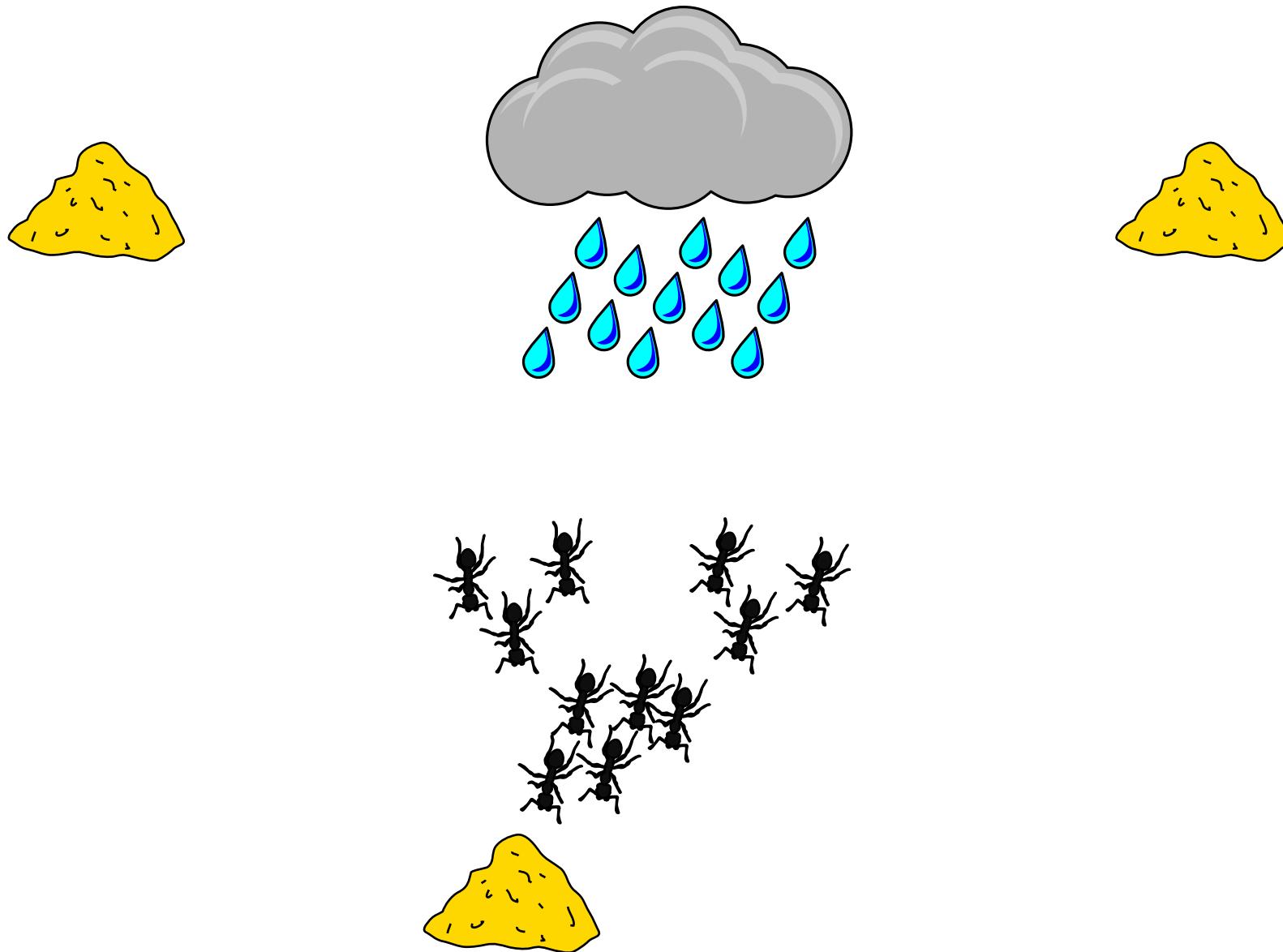
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Natural Algorithms for Plurality Consensus



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Dynamics

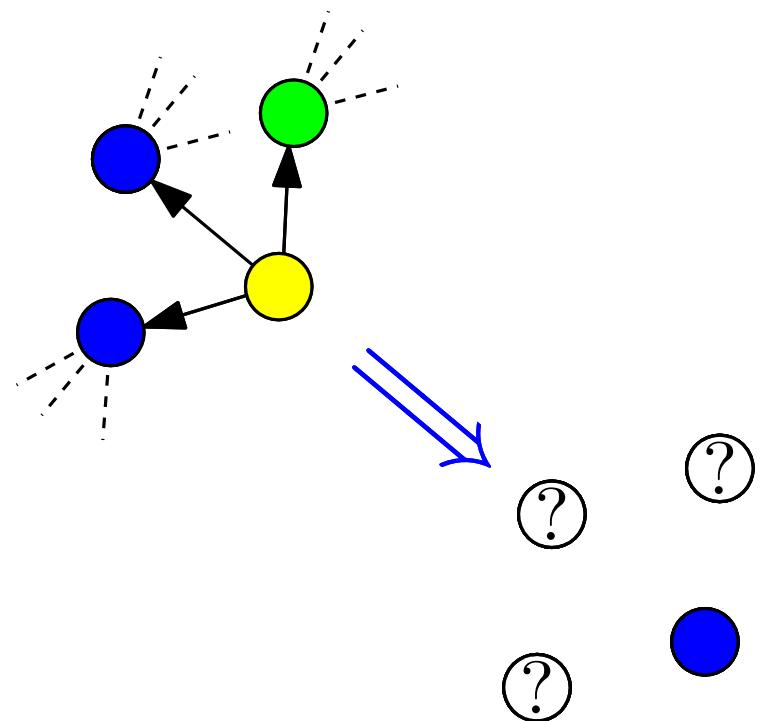
(informal) *Very simple* distributed algorithms: For every graph $G = (V, E)$, agent $u \in V$ and round $t \in \mathbb{N}$, states are updated according to fixed rule $f(\sigma(u), \sigma(S))$ of current state $\sigma(u)$ and symmetric function of states $\sigma(S)$ of a random sample S of neighbors.

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Two examples:

- 3-Majority dynamics

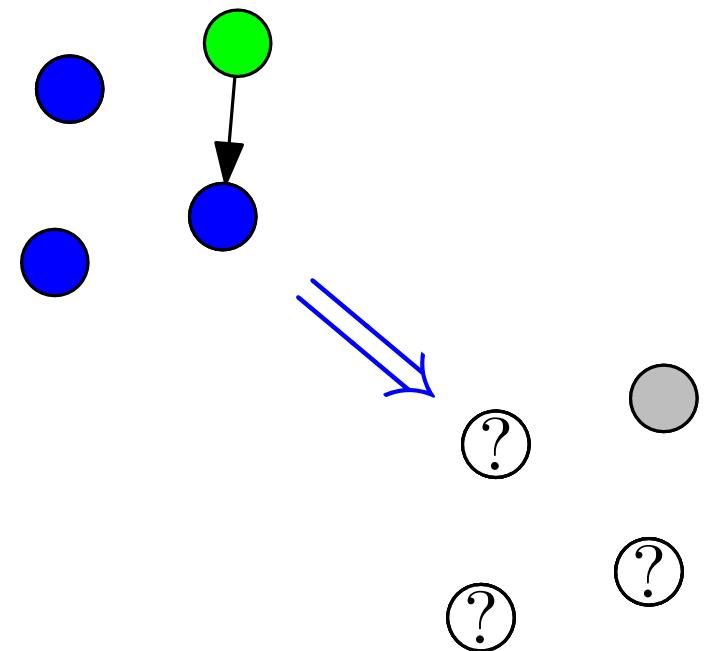


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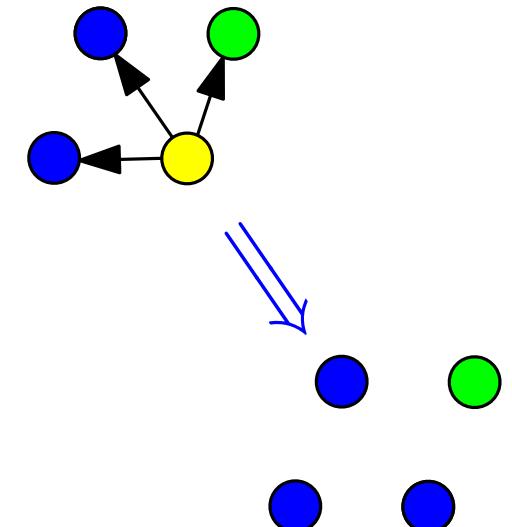
Two examples:

- 3-Majority dynamics
- Undecided-state dynamics



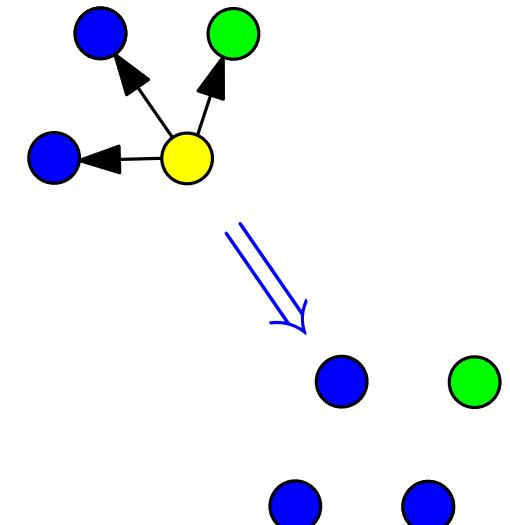
The Power of Dynamics: Plurality Consensus

3-Majority dynamics [SPAA '14, SODA '16]. If plurality has **bias** $\mathcal{O}(\sqrt{kn \log n})$, converges to it in $\mathcal{O}(k \log n)$ rounds w.h.p., even against $o(\sqrt{n/k})$ -bounded adversary. Without bias, converges in $\text{poly}(k)$. h -majority converges in $\Omega(k/h^2)$.

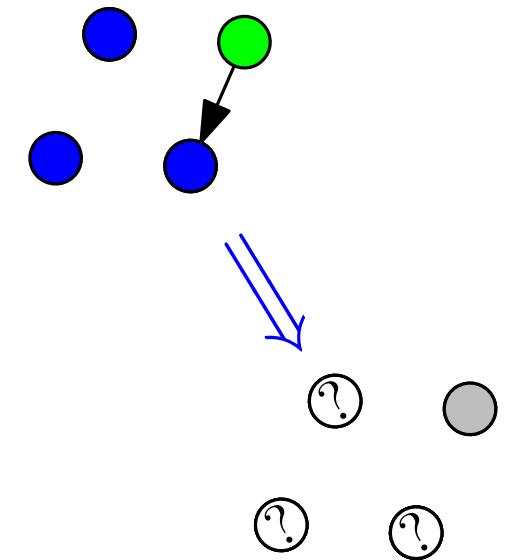


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Undecided-State dynamics [SODA '15]. If majority/second-majority ($c_{maj}/c_{2^{nd} maj}$) is at least $1 + \epsilon$, system converges to plurality within $\tilde{\Theta}(\text{md}(\mathbf{c}))$ rounds w.h.p.



A Global Measure of Bias

$$\text{md}(\mathbf{c}^{(0)}) := \sum_{i=1}^k \left(\frac{c_i^{(0)}}{c_{maj}^{(0)}} \right)^2 = 1 + \mathcal{D} \left(\begin{array}{c} \text{Bar chart showing } c_i^{(0)} \text{ values for } i=1, \dots, k \\ \text{Bar chart showing } c_i^{(t)} \text{ values for } i=1, \dots, k \end{array} \right)$$

$$1 \leq \text{md} \left(\begin{array}{c} \text{Bar chart showing } c_i^{(t)} \text{ values for } i=1, \dots, k \\ \text{Majority bar } c_{maj}^{(t)} \text{ highlighted with a dashed box and red bracket} \end{array} \right) \ll \text{md} \left(\begin{array}{c} \text{Bar chart showing } c_i^{(t)} \text{ values for } i=1, \dots, k \\ \text{Second-majority bar } c_{2^{nd}maj}^{(t)} \text{ highlighted with a dashed box and red bracket} \end{array} \right) \leq k$$

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Applications: Broadcast Problem

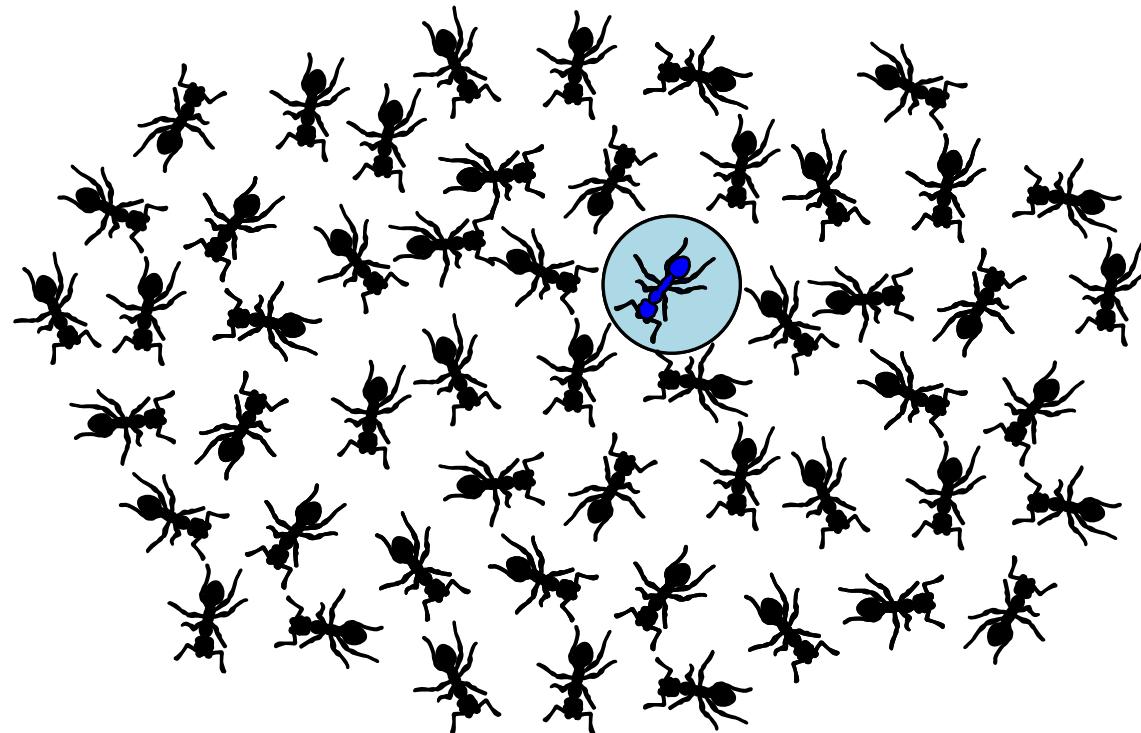


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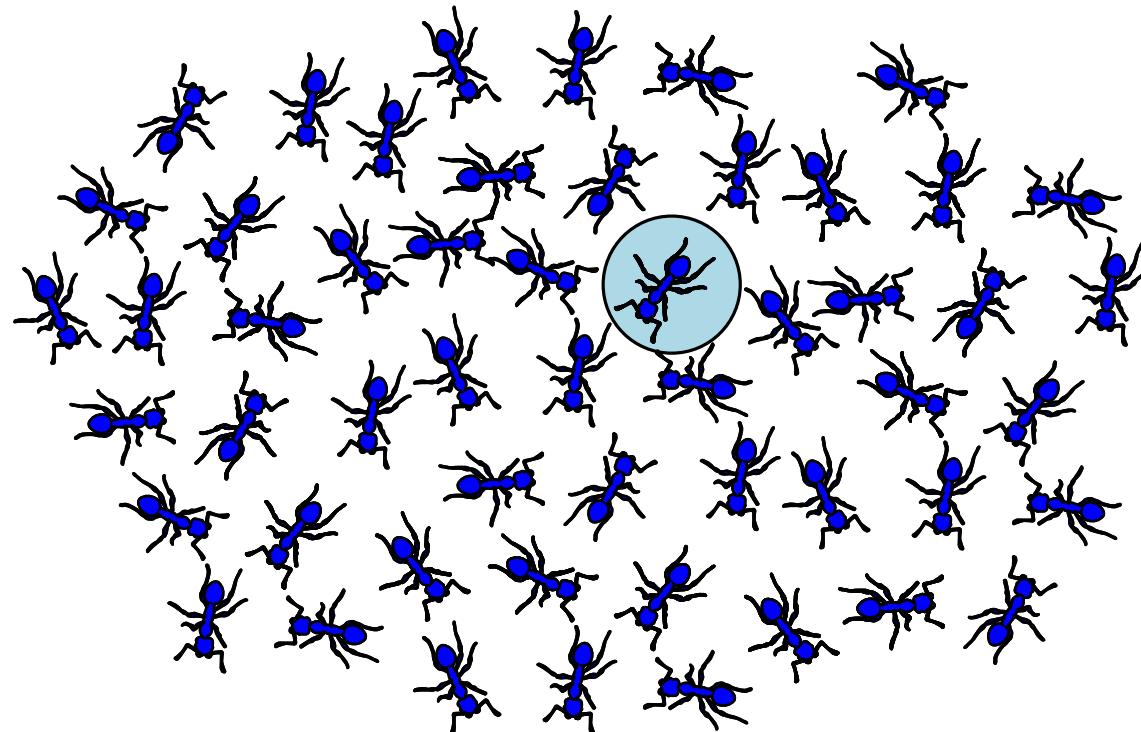
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Sources' bits (and other agents' states) may change in response to *external environment*.



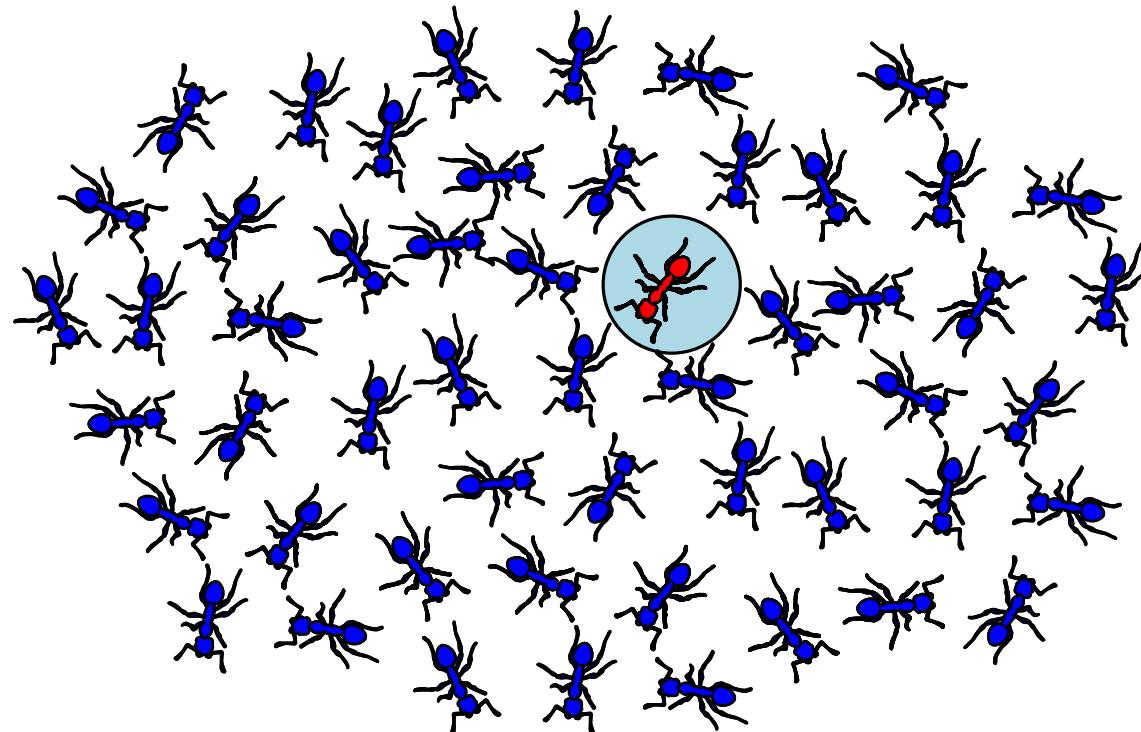
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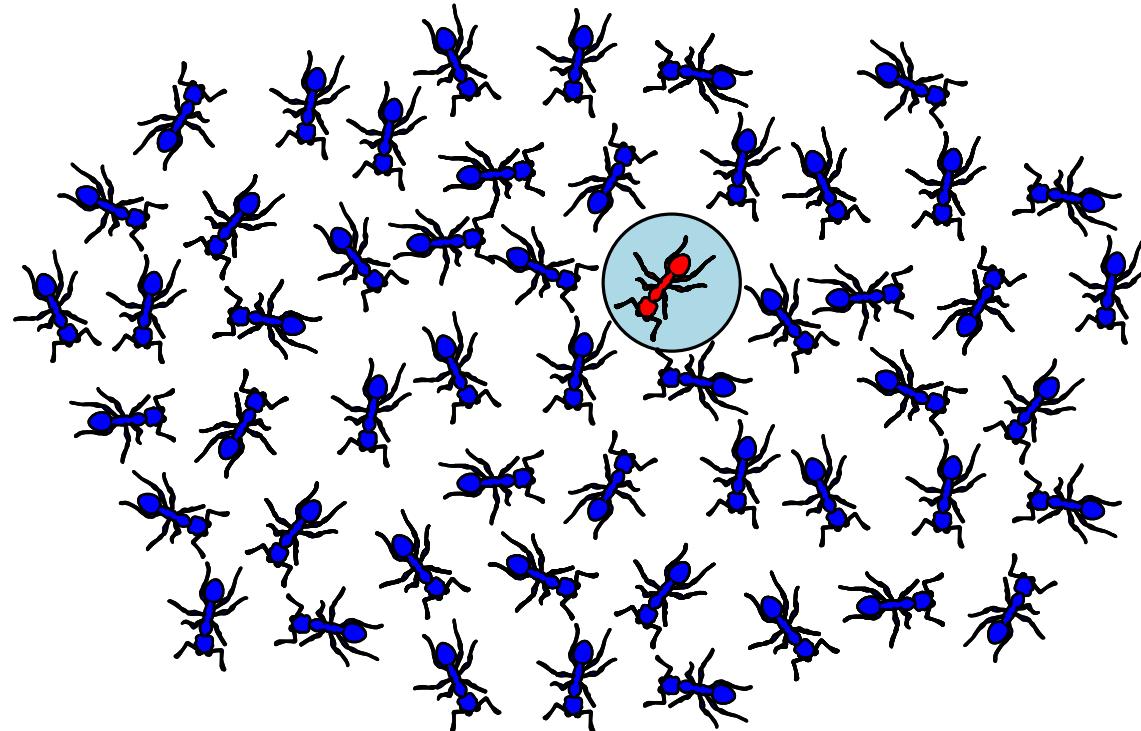
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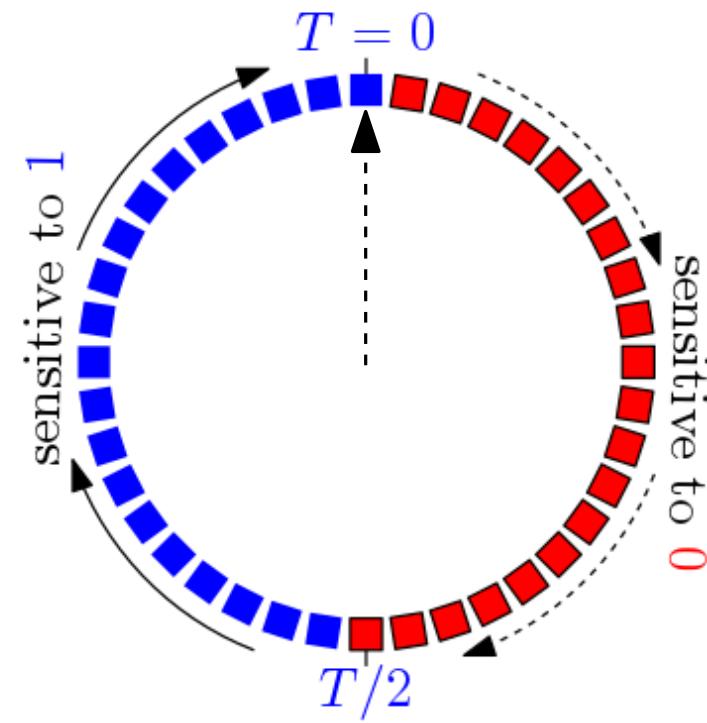
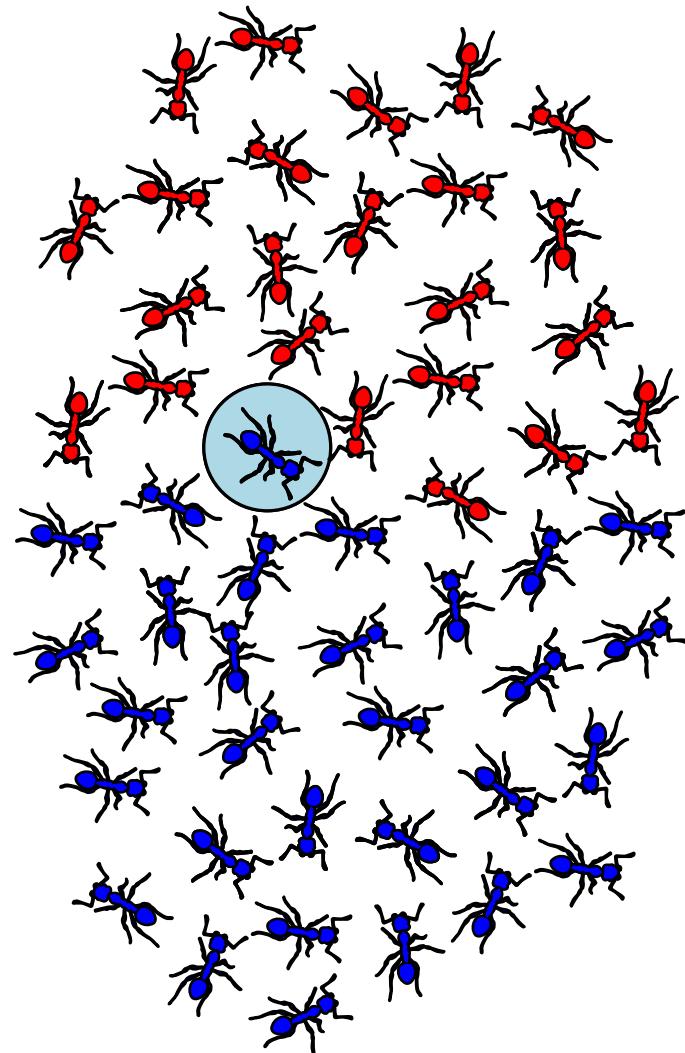
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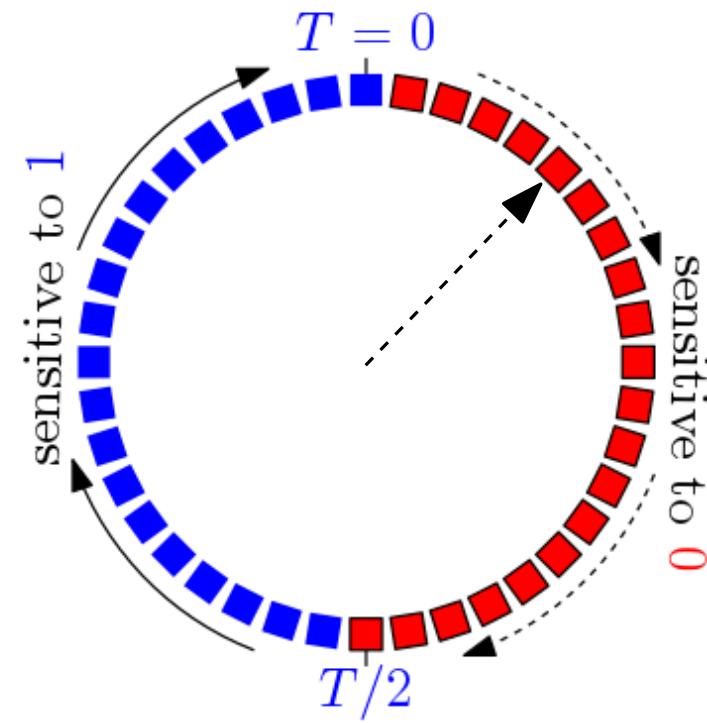
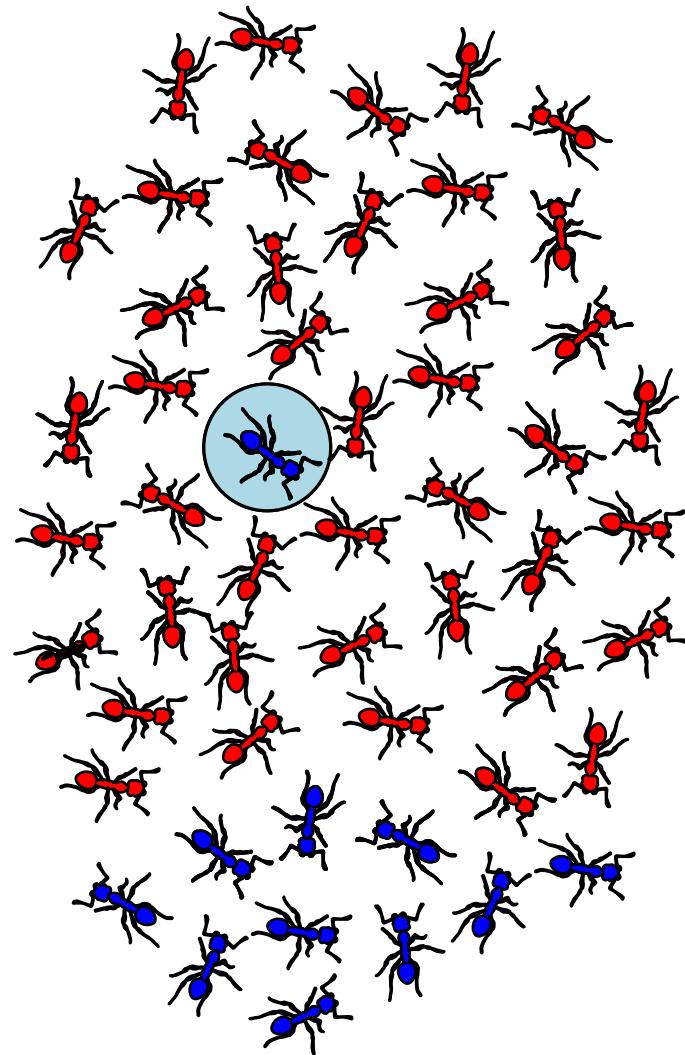


More generally, system is initialized in *arbitrary state* (self-stabilization).

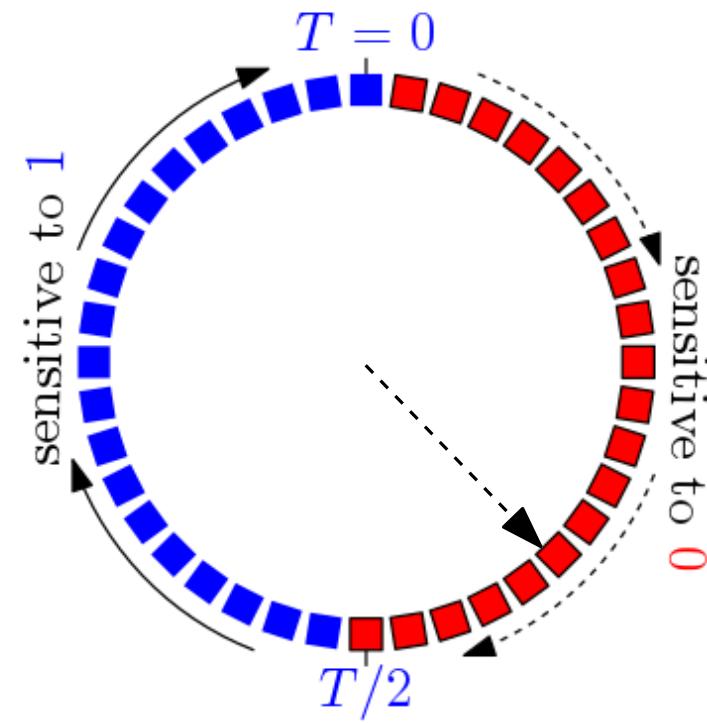
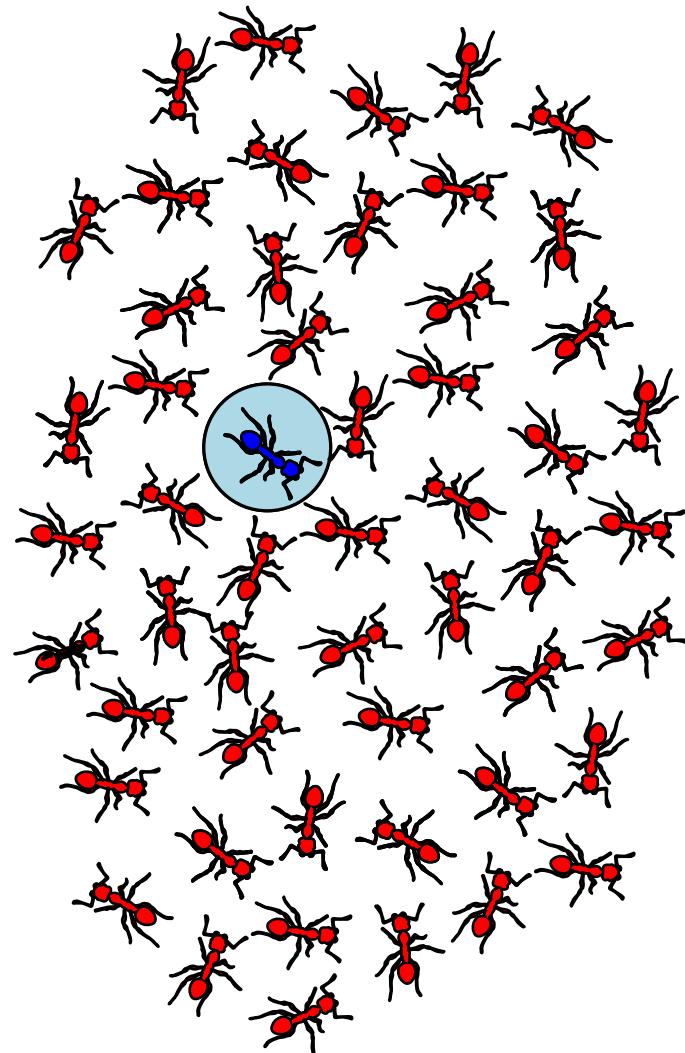
(Self-Stab.) Broadcast vs Synchronization



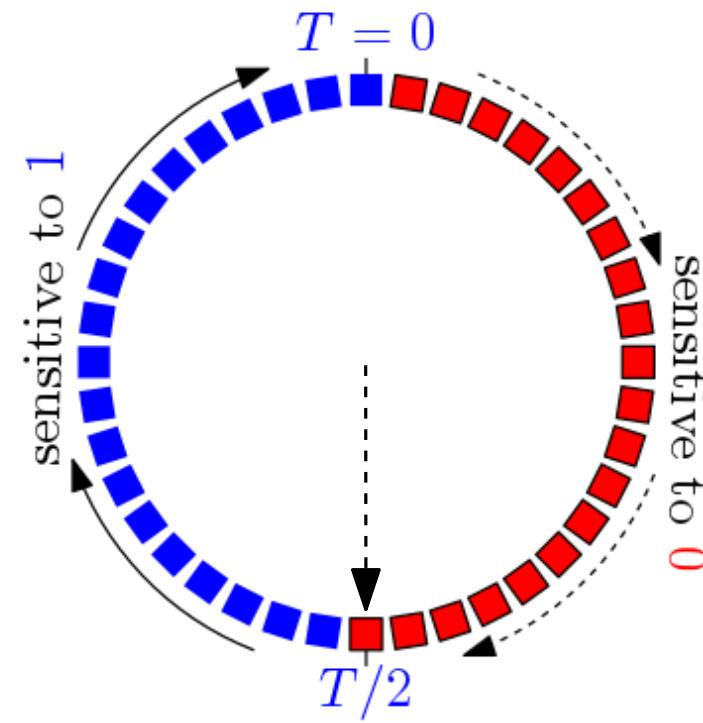
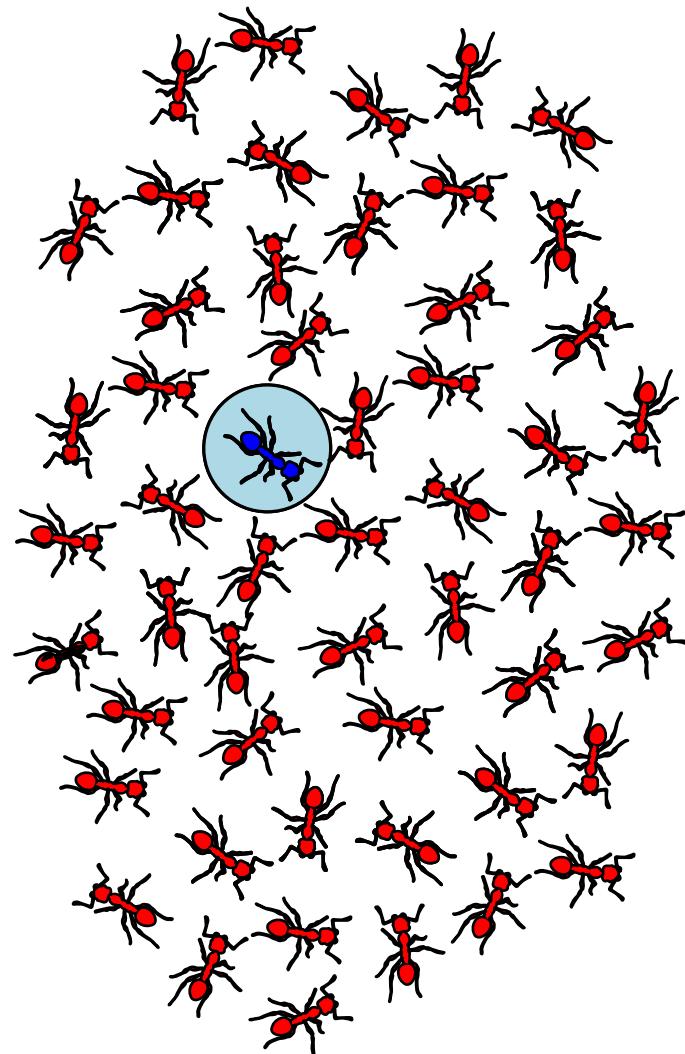
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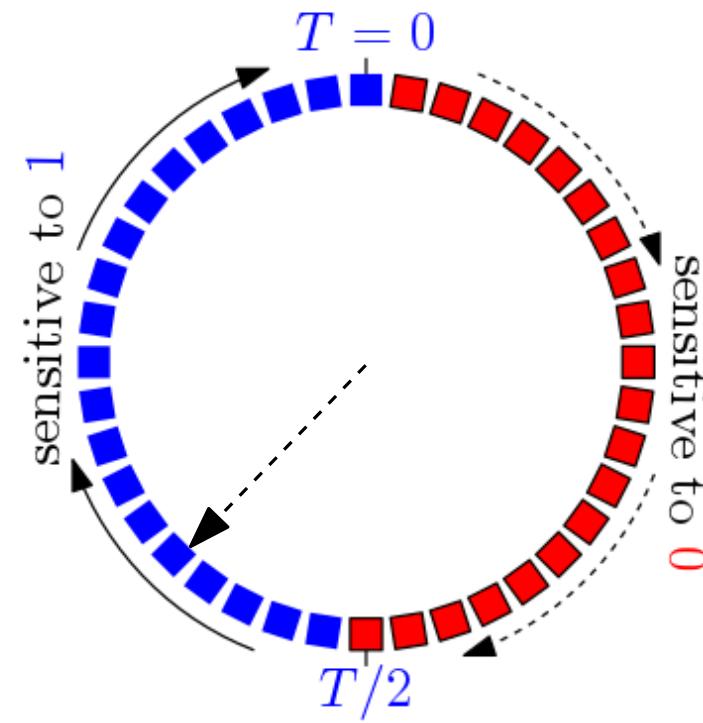
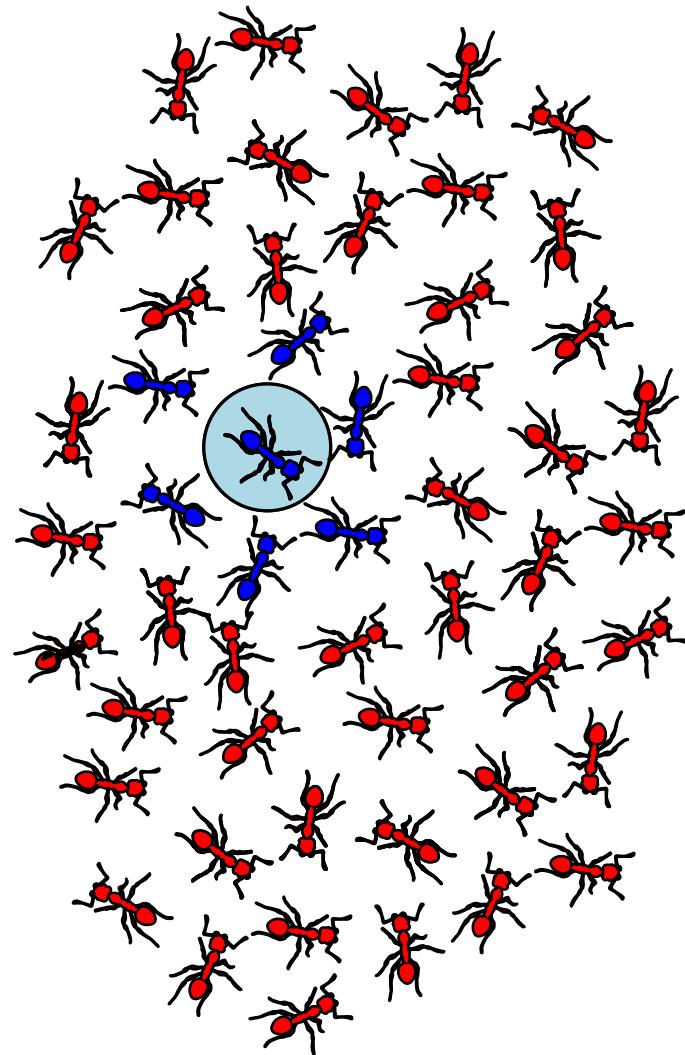
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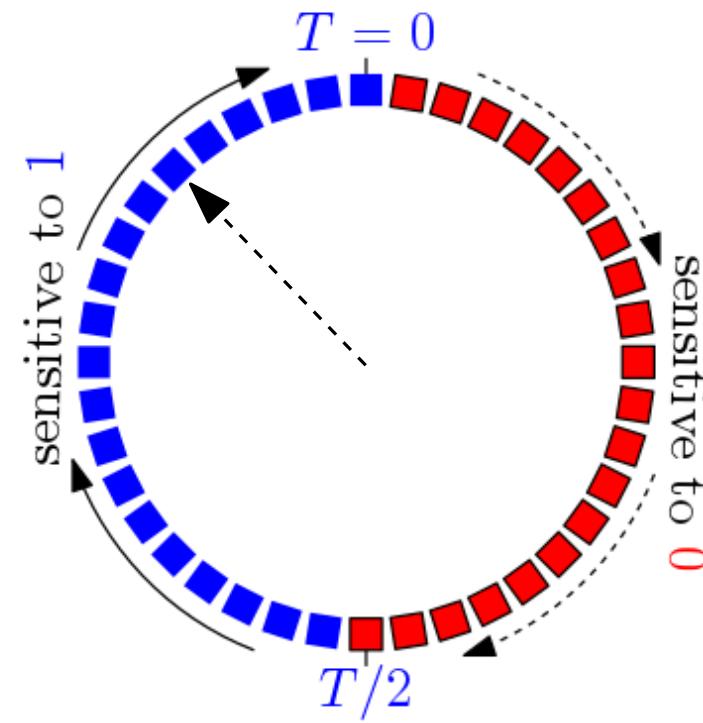
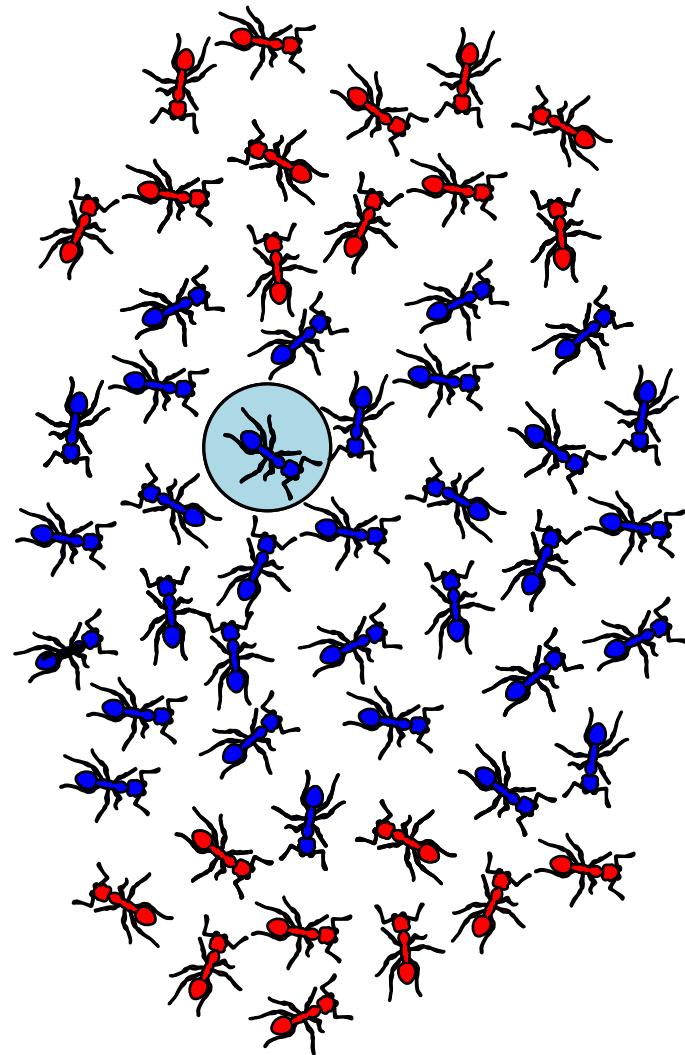
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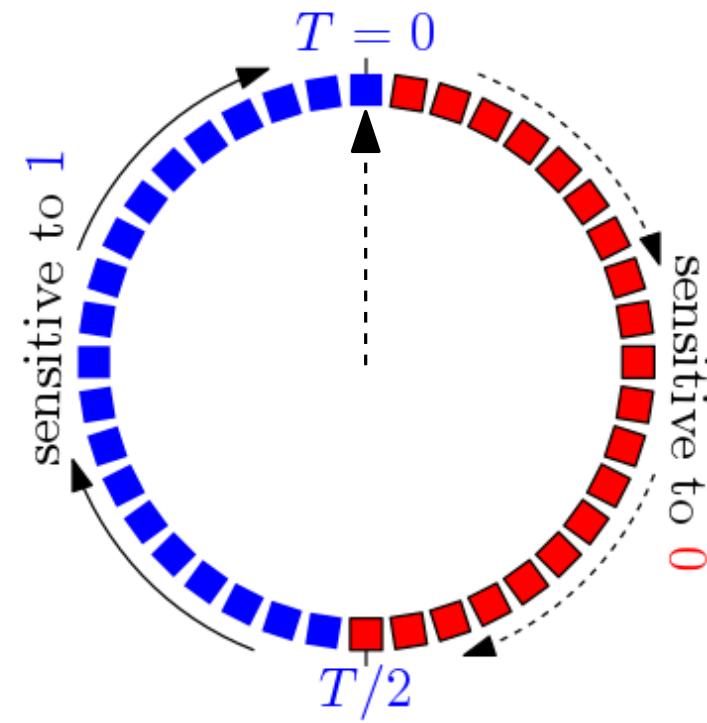
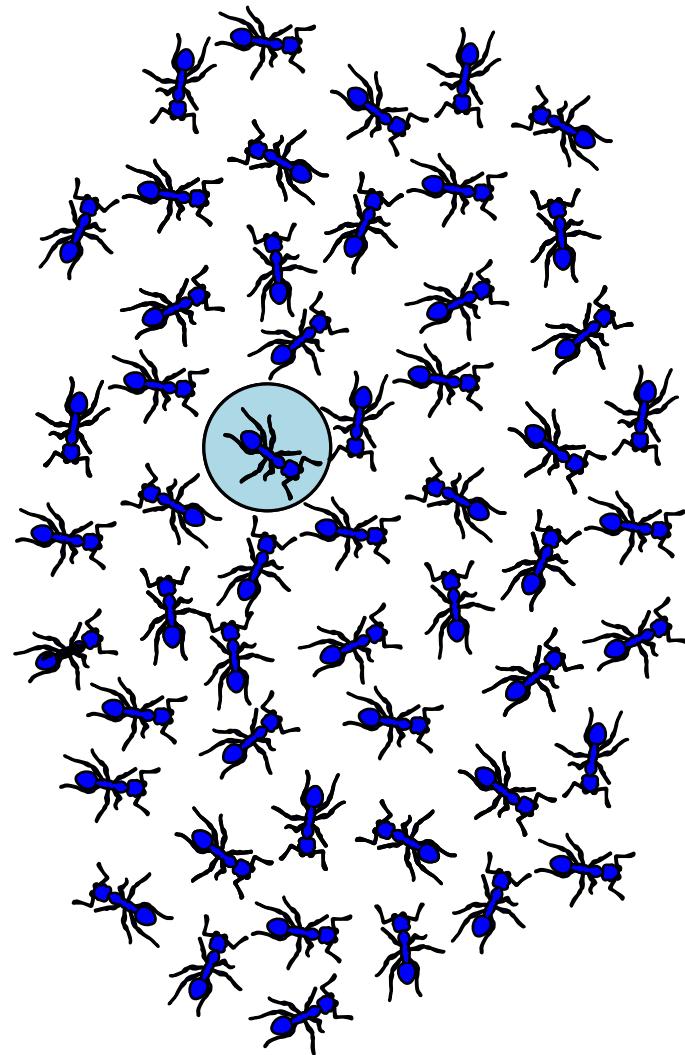
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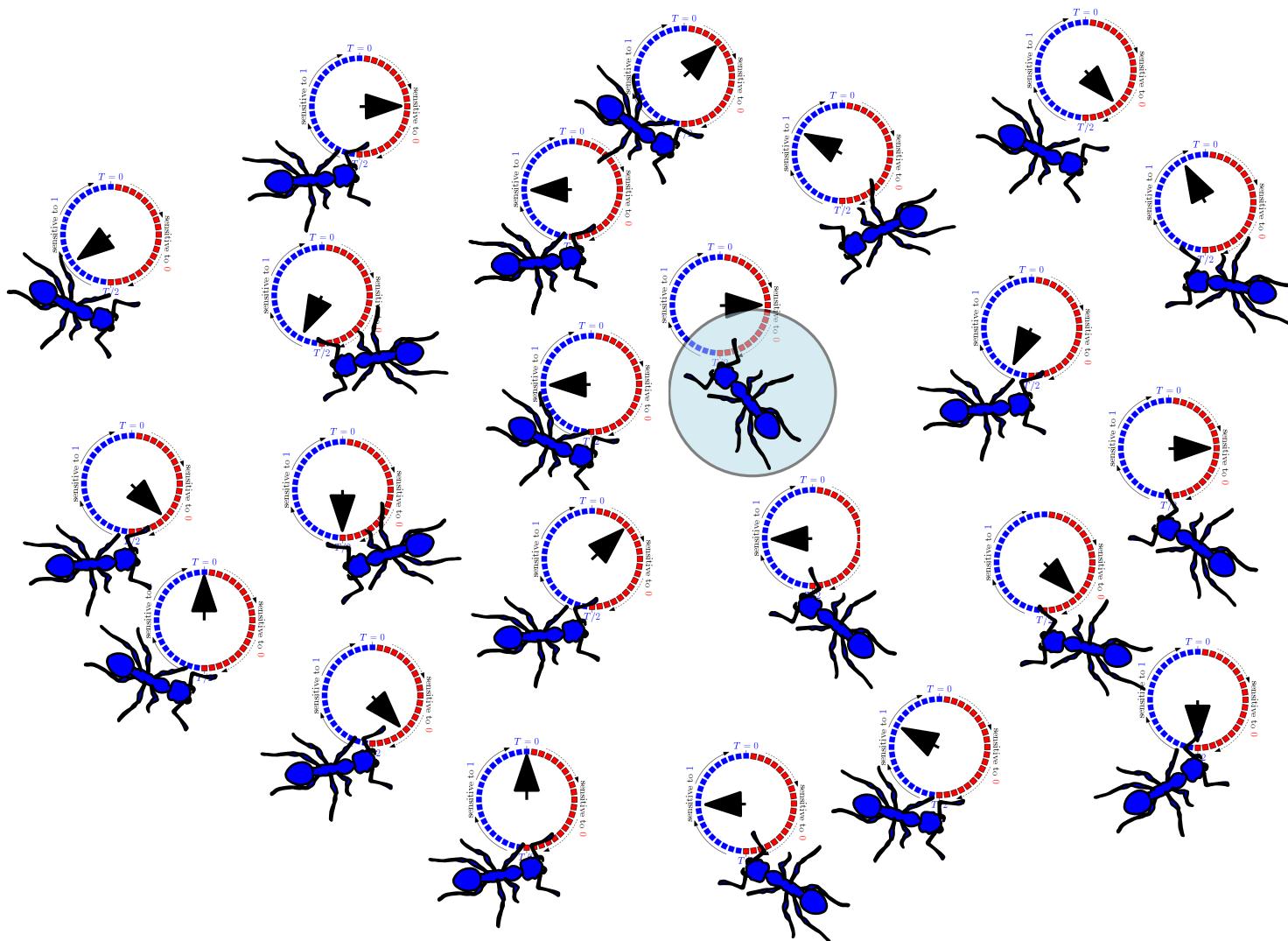


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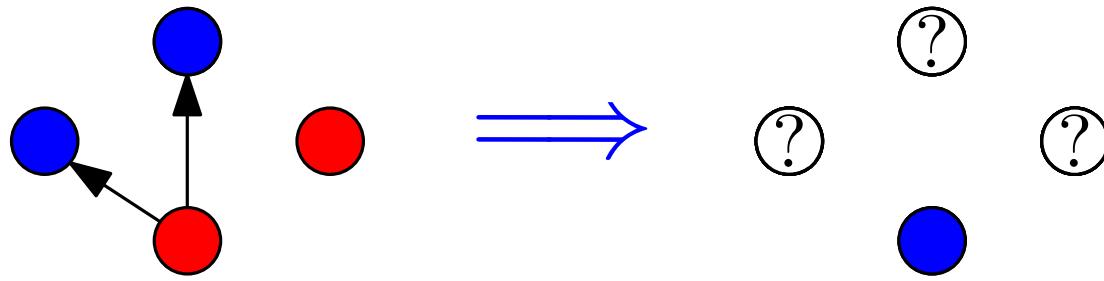


(Self-Stab.) Broadcast vs Synchronization

Self-stabilizing algorithms converge from
any initial configuration

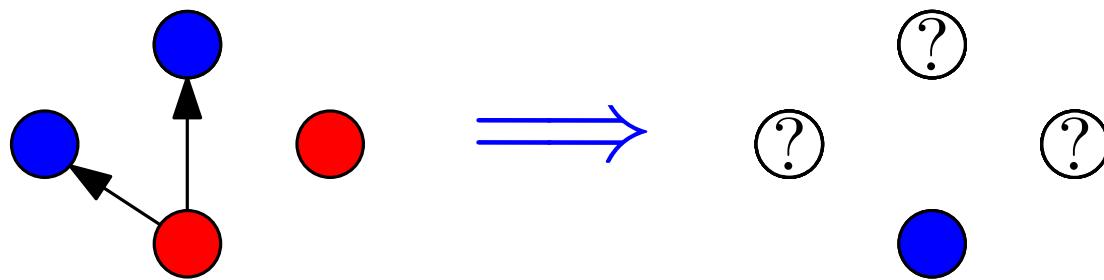


Self-Stabilizing Clock Sync. in the \mathcal{PULL} Model

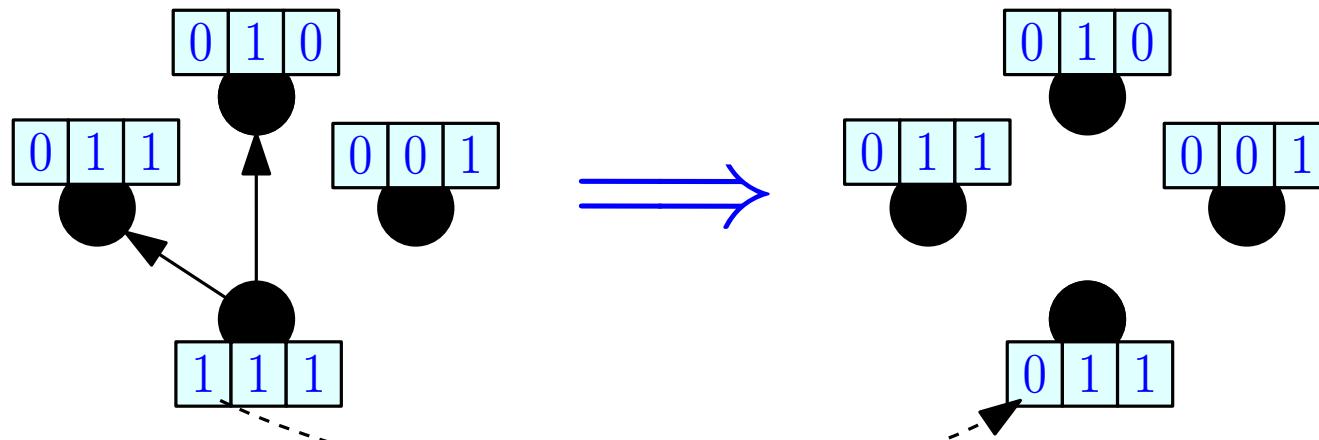


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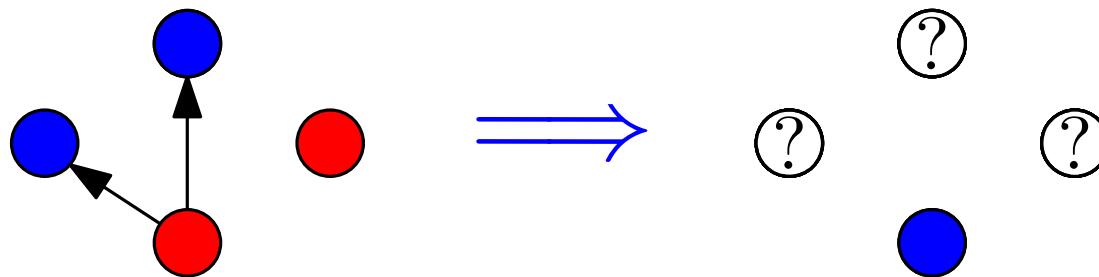
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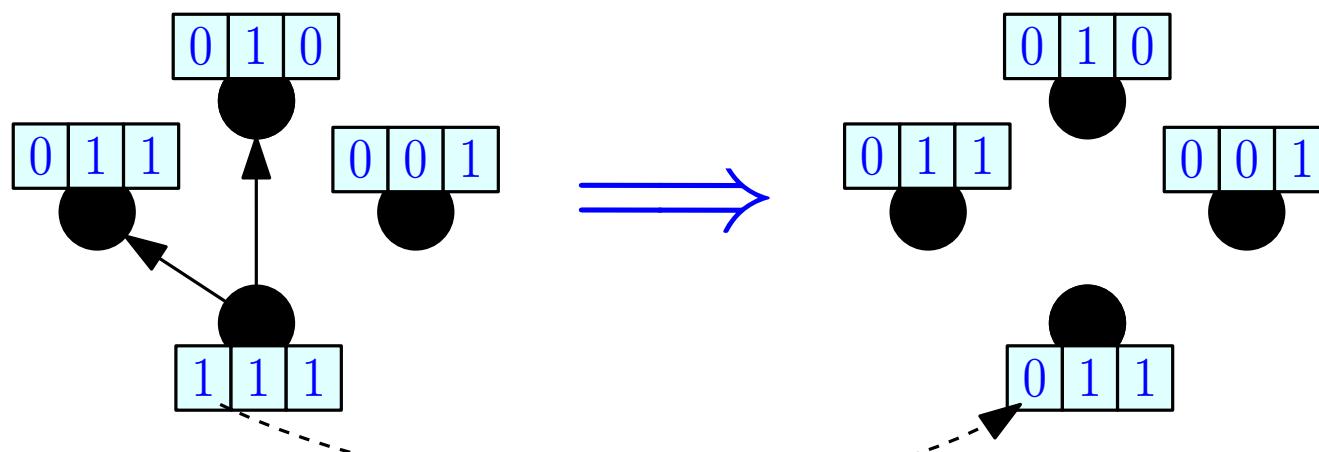
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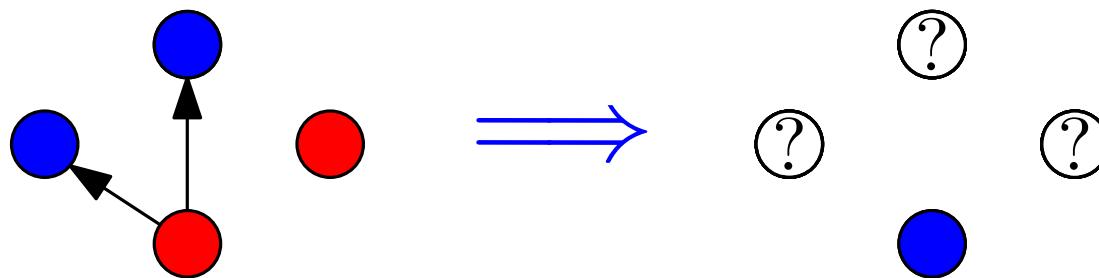


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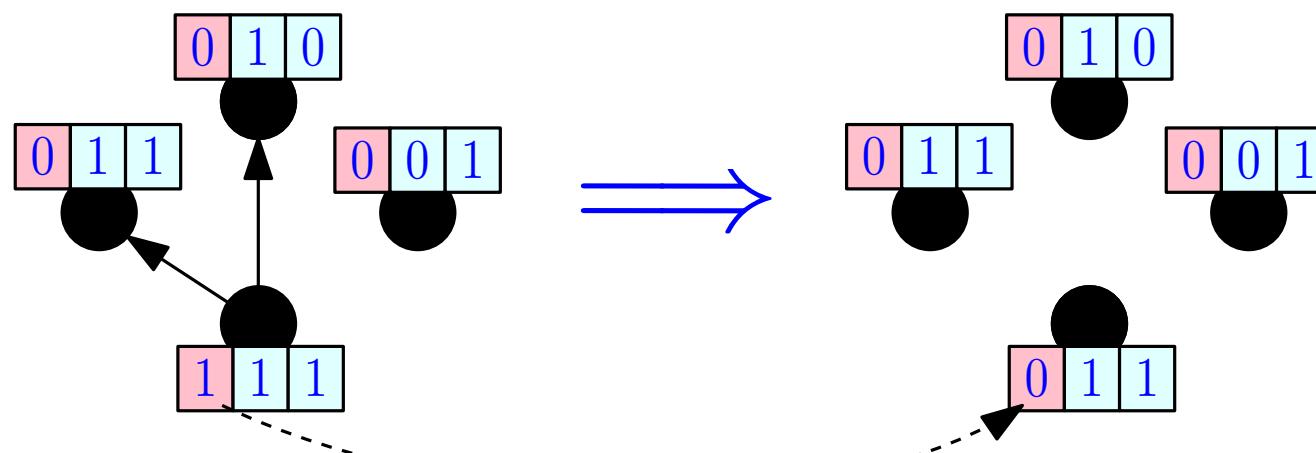


T -clock can be sync. in $\mathcal{O}(\log n \log T)$ rounds w.h.p. using $\log T$ bits. Binary broadcast can be done in **1-bit \mathcal{PULL}** ...

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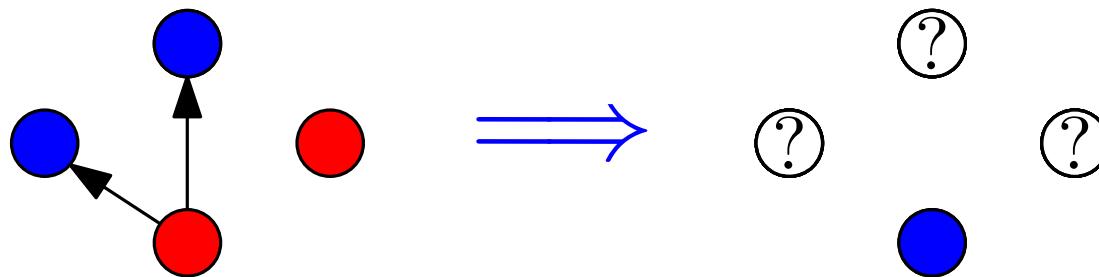


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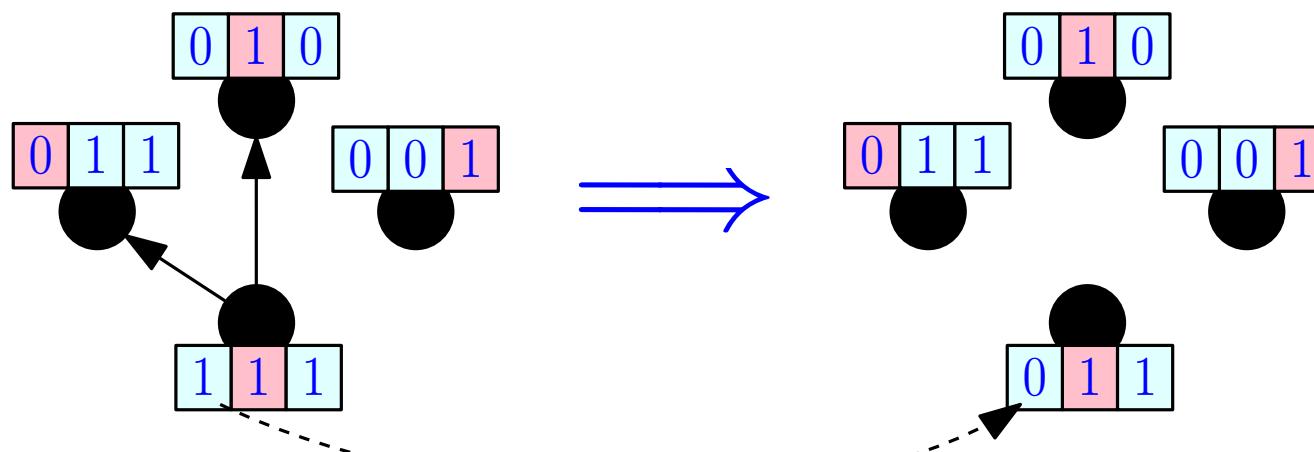


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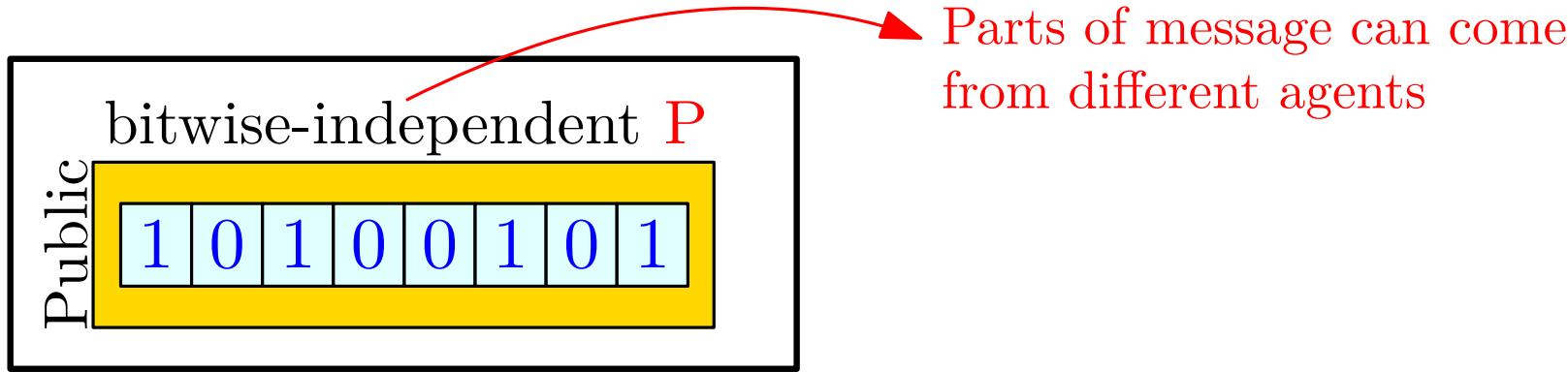


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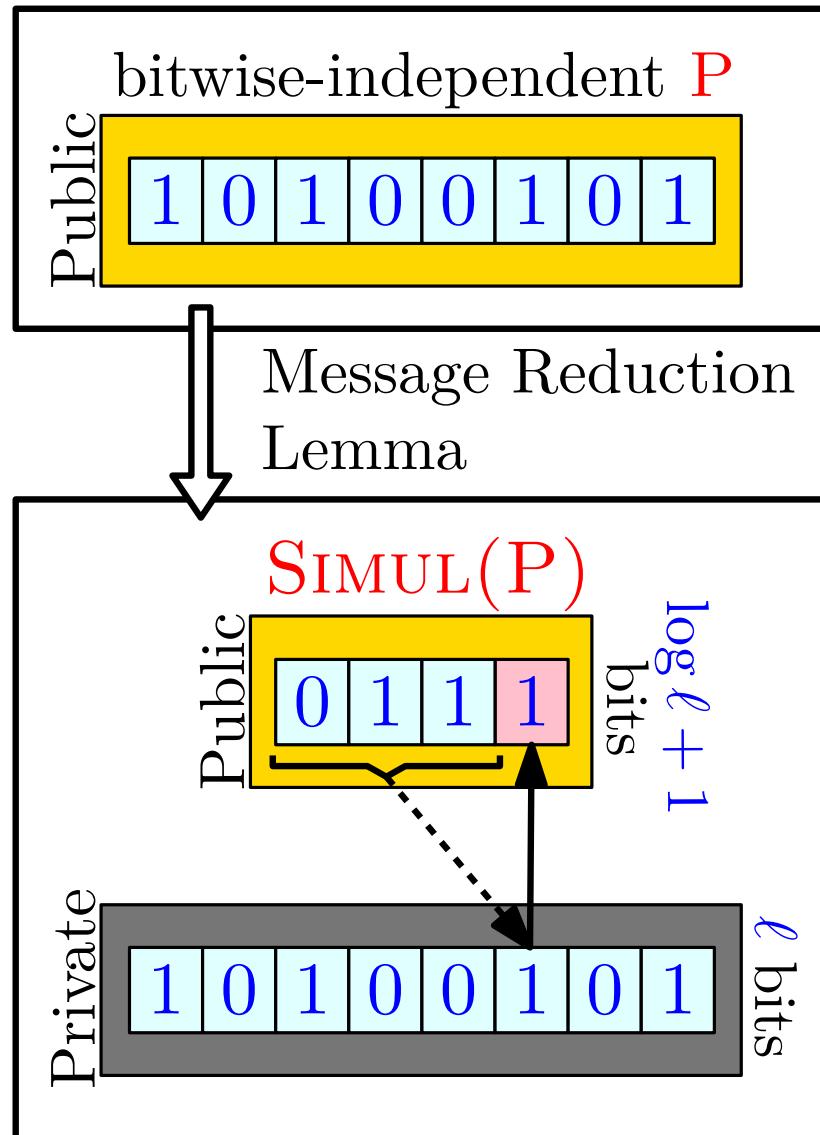


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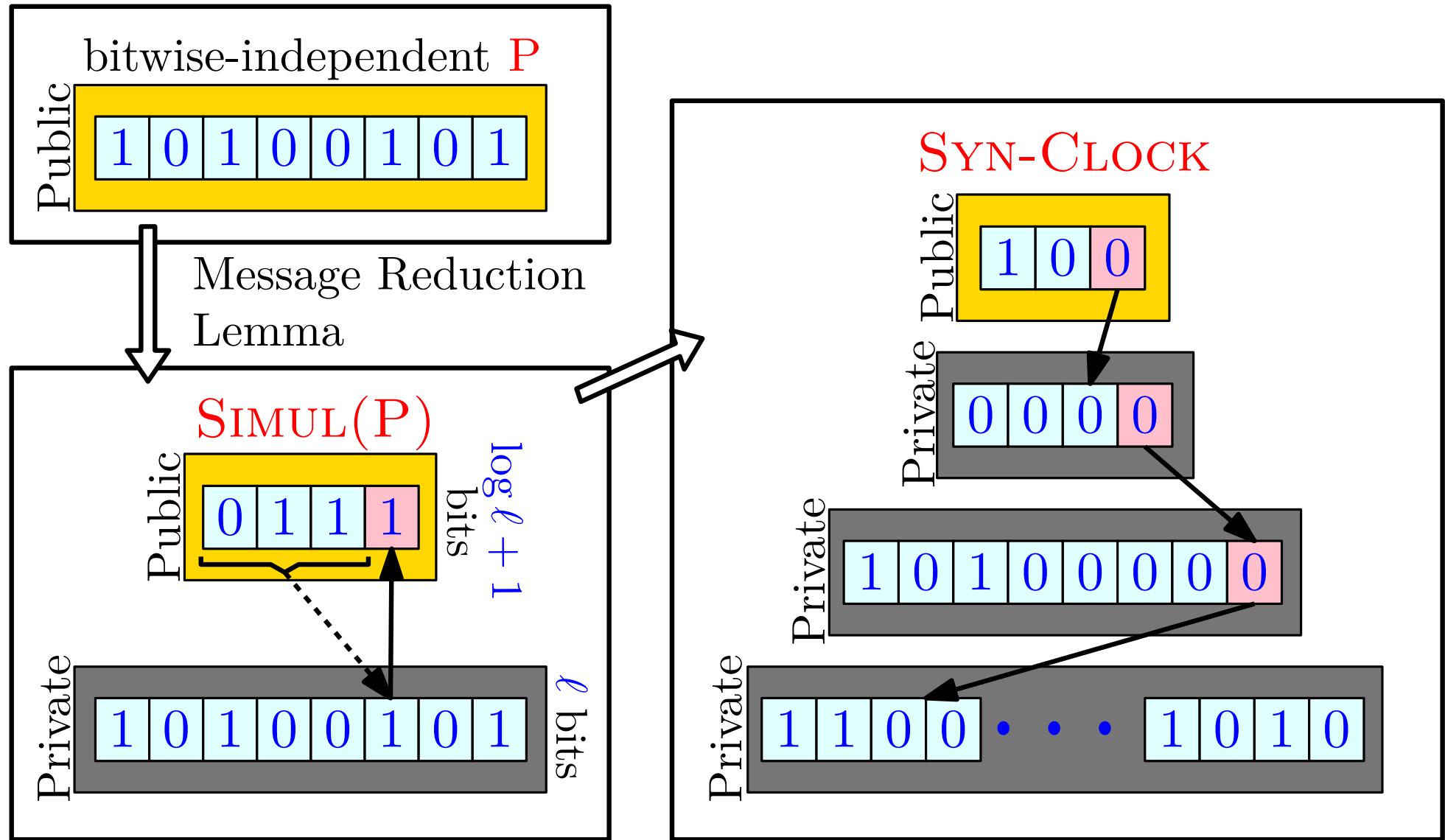
The Message Reduction Lemma



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Results: 3 Bits suffice...

Theorem (Clock Syncronization). SYN-CLOCK is a *self-stabilizing* clock synchronization protocol which synchronizes a clock modulo T in $\tilde{\mathcal{O}}(\log n \log T)$ rounds w.h.p. using **3-bit messages**.

Corollary (Self-stabilizing Majority Broadcast). SYN-PHASE-SPREAD is a *self-stabilizing* Majority Broadcast protocol which converges in $\tilde{\mathcal{O}}(\log n)$ rounds w.h.p using **3-bit messages**, provided majority is supported by $(\frac{1}{2} + \epsilon)$ -fraction of source agents.

Conclusions

Biology demands the study of systems in-between interacting-particle systems and human-made ones.

TCS can analyze **natural algorithms**, helping to understand principles behind the systems' ability to compute in **simple chaotic ways**.

