

P01. CFG = PDA

Convert the given CFG to its CNF (I separated them out to make it easier to convert)

$S \rightarrow XaaX$
 $X \rightarrow aX$
 $X \rightarrow bX$
 $X \rightarrow \Lambda$

[Chomsky Normal Form: every non-terminal \rightarrow double non-terminal (XX) or a single terminal (x)]

First remove any nullable non-terminals

nullable non-terminals: X

remove($X \rightarrow \Lambda$)

Before, X could be replaced by Λ . From these rules that could contain null string:

$S \rightarrow XaaX$
 $X \rightarrow aX$
 $X \rightarrow bX$

We can make new rules:

$S \rightarrow Xaa$
 $S \rightarrow aaX$
 $S \rightarrow aa$
 $X \rightarrow a$
 $X \rightarrow b$

So with no null productions our rules are now:

$S \rightarrow XaaX \mid Xaa \mid aaX \mid aa$
 $X \rightarrow aX \mid bX \mid a \mid b$

There are no Unit Productions (non-terminal \rightarrow one non-terminal), so we can skip this step.

Now we need to replace any rule that does not have the form XX or x by introducing new non-terminals

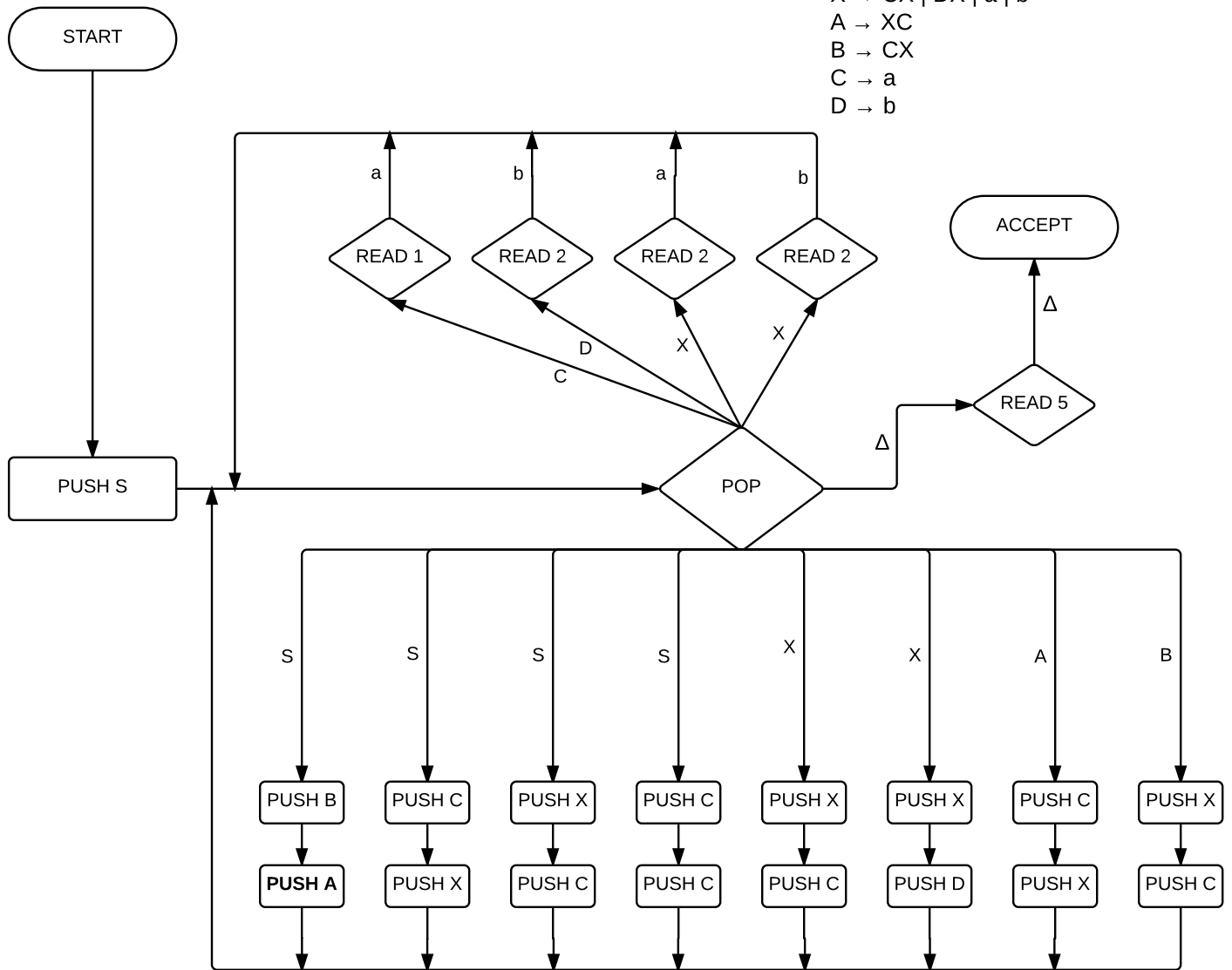
$A \rightarrow XC$
 $B \rightarrow CX$
 $C \rightarrow a$
 $D \rightarrow b$

and now our rules, including the modified versions, are:

$S \rightarrow AB \mid XC \mid CX \mid CC$
 $X \rightarrow CX \mid DX \mid a \mid b$
 $A \rightarrow XC$
 $B \rightarrow CX$
 $C \rightarrow a$
 $D \rightarrow b$

Construct a PDA that accepts the same language generated by the given CFG using the algorithm of Theorem 30:

$S \rightarrow AB \mid XC \mid CX \mid CC$
 $X \rightarrow CX \mid DX \mid a \mid b$
 $A \rightarrow XC$
 $B \rightarrow CX$
 $C \rightarrow a$
 $D \rightarrow b$



P02. CFL

Find a CFG for this language: All words that start with an a or are of the form $a^n b^n$

Just defining it so I can see it in a different way:

$$S \rightarrow a(a + b)^* \mid a^n b^n$$

$$S \rightarrow a \mid aX \mid W$$

$$X \rightarrow aX \mid bX \mid a \mid b$$

$$W \rightarrow aWb \mid ab \mid \Lambda$$