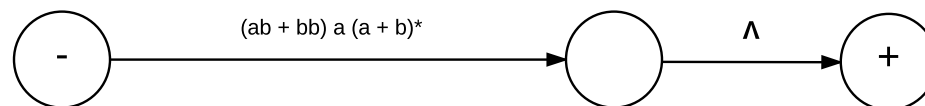
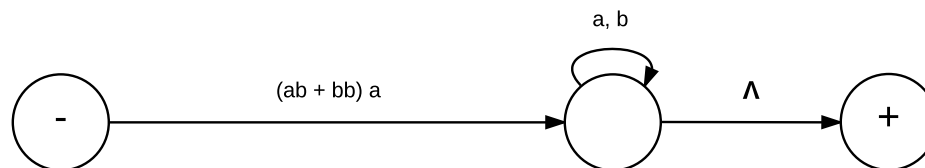
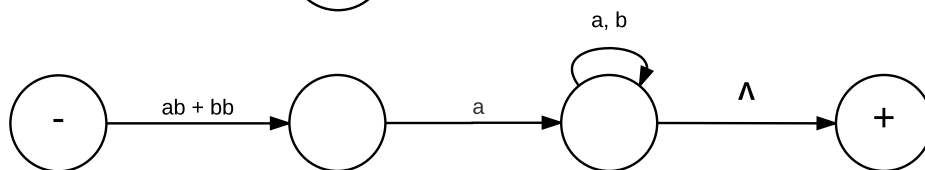
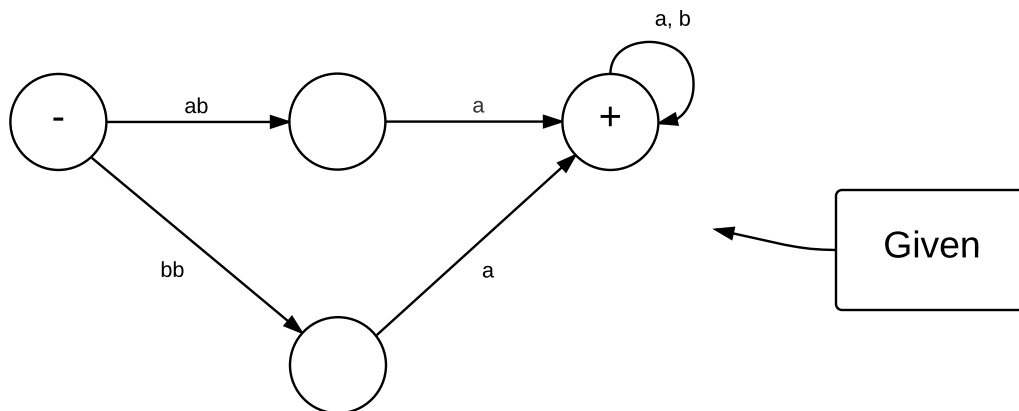


P01. Kleene's Theorem Part 2

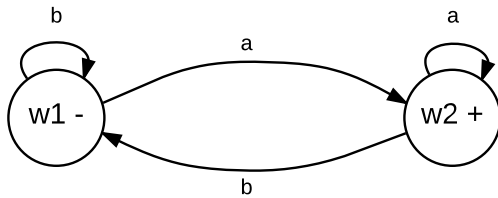
Convert the following TG into a regular expression.



Answer: $(ab + bb) a (a + b)^*$

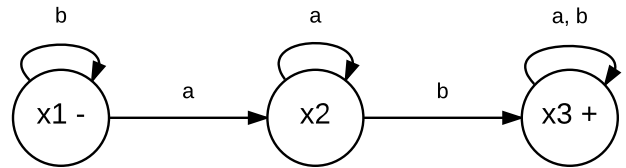
P02. Kleene's Theorem Part 3

Using the algorithm of Kleene's theorem, construct a FA for the union language: $FA1 + FA2$,
Where $FA1$ and $FA2$ are given below:



FA1

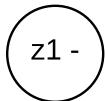
all strings ending in a



FA2

Must have at least 1 a and have at least 1 b after it

(1) Start State: **(z1)** = (w1) or (x1) (-)

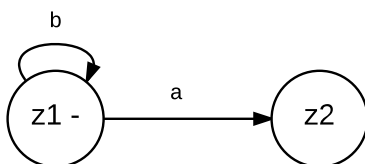


(2) At state **(z1)**

if reading a, $FA1 \rightarrow (w2)$ and $FA2 \rightarrow (x2)$, so **(z2)** = (w2) or (x2)

if reading b, $FA1 \rightarrow (w1)$ and $FA2 \rightarrow (x1)$, so **(z1)** = (w1) or (x1) (not new)

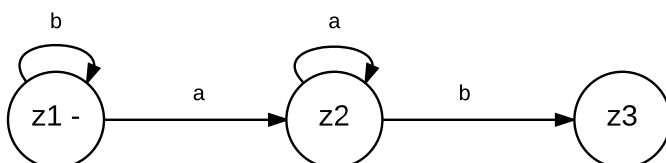
- stays at the same state, both loop when accepting a b



(3) at state **(z2)**

if reading a, $FA1 \rightarrow (w2)$ and $FA2 \rightarrow (x2)$, so **(z2)** = (w2) or (x2) (not new)

if reading b, $FA1 \rightarrow (w1)$ and $FA2 \rightarrow (x3)$, so **(z3)** = (w1) or (x3)

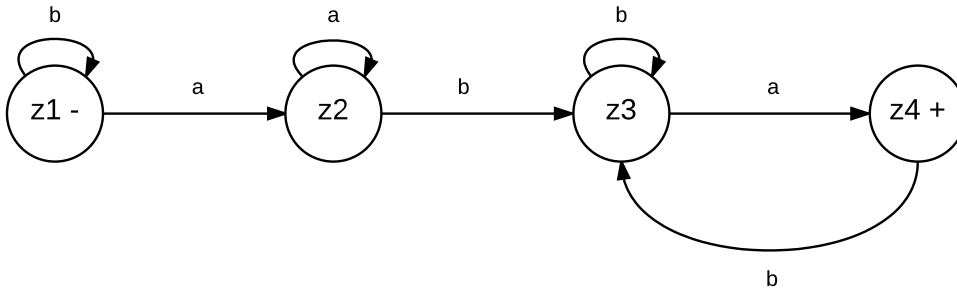


(4) at state (**z3**)

if reading a, FA1->(w2) and FA2->(x3), so (**z4**) = (w2) and (x3) (+)

- is different because these are both now end states

if reading b, FA1->(w1) and FA2->(x3), so back to (**z3**)



(5) at state (**z4**)

if reading a, FA1->(w2) and FA2->(x3), so (**z4**) (+) (not new)

if reading b, FA1->(w1) and FA2->(x3), so (**z3**) (not new)

NO NEW STATES