

# P01. Regular Language

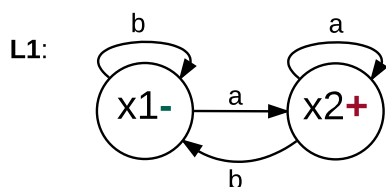
For the following pair of regular expressions, find the Finite Automaton that defines  $L1 \cap L2$  :

L1:  $(a+b)^*a$

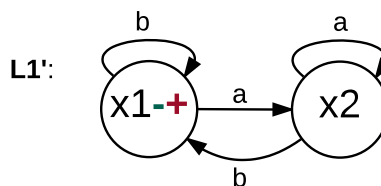
L2:  $(a+b)^*aa(a+b)^*$

**Must find the compliment of both languages**

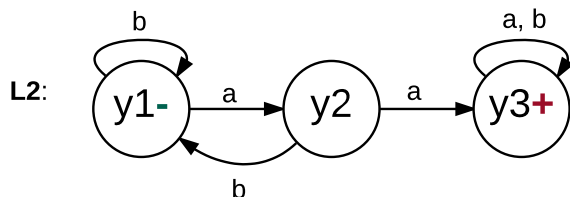
Note: L1 accepts all strings that end in a



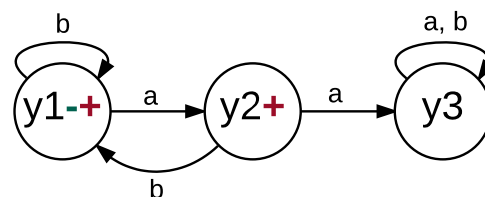
Find Compliment



Note: L2 accepts all strings that contain 2 as



Find Compliment L2':



**Must find  $L1' + L2'$**

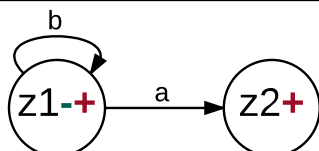
(1) Start State:  $(z1) = (x1) \text{ or } (y1) (-)(+)$



(2) At state  $(z1)$  which is  $(x1)$  or  $(y1)$

if reading a,  $L1' \rightarrow (x2)$  and  $L2' \rightarrow (y2)$ , so  $(z2) = (x2) \text{ or } (y2) (+)$

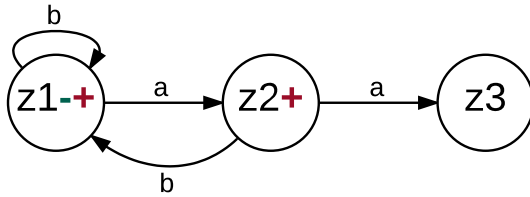
if reading b,  $L1' \rightarrow (x1)$  and  $L2' \rightarrow (y1)$ , so  $(z1) = (x1) \text{ or } (y1) \text{ (not new) } (+)$



(3) At state (**z2**) which is (x2) or (y2)

if reading a,  $L1' \rightarrow (x2)$  and  $L2' \rightarrow (y3)$ , so (**z3**) = (x2) or (y3)

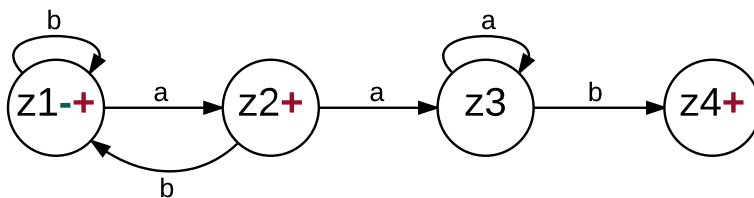
if reading b,  $L1' \rightarrow (x1)$  and  $L2' \rightarrow (y1)$ , so (**z1**) = (x1) or (y1) (not new) (+)



(4) At state (**z3**) which is (x2) or (y3)

if reading a,  $L1' \rightarrow (x2)$  and  $L2' \rightarrow (y3)$ , so (**z3**) = (x2) or (y3) (not new)

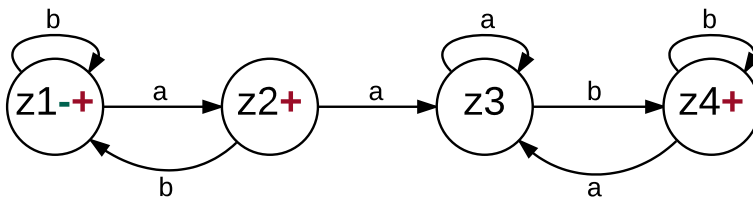
if reading b,  $L1' \rightarrow (x1)$  and  $L2' \rightarrow (y3)$ , so (**z4**) = (x1) or (y3) (+)



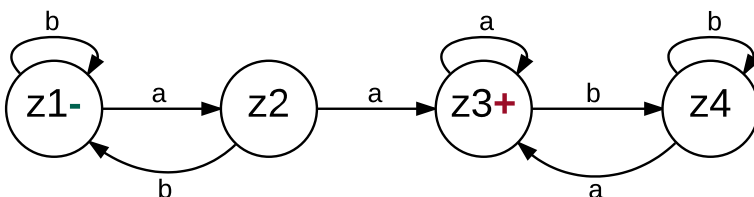
(4) At state (**z4**) which is (x1) or (y3)

if reading a,  $L1' \rightarrow (x2)$  and  $L2' \rightarrow (y3)$ , so (**z3**) = (x2) or (y3) (not new)

if reading b,  $L1' \rightarrow (x1)$  and  $L2' \rightarrow (y3)$ , so (**z4**) = (x1) or (y3) (not new)



**Must find the compliment of this union of  $L1' + L2'$**



## P02. Non-regular Language

Use the pumping lemma to show that the following language is non-regular:

$$a^{(n)} b^{(2n)}, n > 1$$

1. Need to find the substrings  $x$ ,  $y$  and  $z$  such that  $y$  can be pumped
2. This language will have some  $a$ s then twice that number of  $b$ s

Lets try a few out:

$\underline{x}$      $\underline{y}$      $\underline{z}$   
 $a$      $aa$      $bbbbbb$

If  $y$  is pumped there will be too many  $a$ s

$\underline{x}$      $\underline{y}$      $\underline{z}$   
 $aaa$      $bbbb$      $bb$

If  $y$  is pumped there will be too many  $b$ s

$\underline{x}$      $\underline{y}$      $\underline{z}$   
 $aa$      $abb$      $bbbb$

If  $y$  is pumped the string will be out of order, we will get  $abbabb$

Cannot find a substring to be pumped that satisfies the properties of the language, so this language is non regular