



Notebook - Maratona de Programação

Lenhadoras de Segtree

Contents

1 Math	2	3 Template	7
1.1 Ceil	2	3.1 Template	7
1.2 To Decimal	2	3.2 Template Clean	7
1.3 Subsets	2	4 Strings	8
1.4 Matrix Exponentiation	2	4.1 Hash	8
1.5 Crt	3	4.2 Kmp	8
1.6 Binary To Decimal	3	4.3 Generate All Permutations	8
1.7 Fast Exponentiation	3	4.4 Generate All Sequences Length K	8
1.8 Linear Diophantine Equation	3	4.5 Suffix Array	8
1.9 Function Root	4	4.6 Lcs	10
1.10 Sieve Of Eratosthenes	4	4.7 Trie	10
1.11 Horner Algorithm	4	4.8 Z-function	10
1.12 Multiplicative Inverse	4	5 Misc	10
1.13 Representation Arbitrary Base	5	5.1 Split	10
1.14 Set Operations	5	5.2 Int128	11
1.15 Divisors	5	6 Graphs	11
1.16 Check If Bit Is On	5	6.1 Centroid Find	11
1.17 Prime Factors	5	6.2 Bipartite	11
2 DP	5	6.3 Prim	11
2.1 Knapsack With Index	5	6.4 Eulerian Undirected	12
2.2 Substr Palindrome	6	6.5 Ford Fulkerson Edmonds Karp	13
2.3 Edit Distance	6	6.6 Hld Edge	13
2.4 Knapsack	6	6.7 Floyd Warshall	14
2.5 Digits	6	6.8 Lca	14
2.6 Coins	6	6.9 Kuhn	15
2.7 Minimum Coin Change	7	6.10 Eulerian Directed	15
2.8 Kadane	7	6.11 Bellman Ford	16
		6.12 Dinic	17
		6.13 2sat	18
		6.14 Find Cycle	20

6.15	Cycle Path Recovery	20
6.16	Blossom	21
6.17	Centroid Decomposition	21
6.18	Tarjan Bridge	22
6.19	Hld Vertex	22
6.20	Small To Large	24
6.21	Tree Diameter	24
6.22	Dijkstra	24
6.23	Kruskall	25
6.24	Hungarian	25
6.25	Negative Cycle	26
7	Geometry	26
7.1	Shoelace Boundary	26
7.2	Inside Polygon	27
7.3	Closest Pair Points	27
7.4	2d	27
8	Algorithms	30
8.1	Lis	30
8.2	Delta-encoding	30
8.3	Subsets	30
8.4	Binary Search Last True	30
8.5	Ternary Search	30
8.6	Binary Search First True	30
8.7	Biggest K	31
9	Data Structures	31
9.1	Sparse Table	31
9.2	Mergesort Tree Vector	32
9.3	Ordered Set	32
9.4	Priority Queue	33
9.5	Dsu	33
9.6	Two Sets	33
9.7	Psum2d	34
9.8	Dynamic Implicit Sparse	34
9.9	Segtree2d	35
9.10	Minimum And Amount	36
9.11	Lazy Addition To Segment	37
9.12	Segment With Maximum Sum	38
9.13	Range Query Point Update	39
9.14	Lazy Assignment To Segment	40
9.15	Lazy Dynamic Implicit Sparse	41
9.16	Persistent	42
9.17	Sparse Table2d	43

1 Math

1.1 Ceil

```
1 long long division_ceil(long long a, long long b) {
2     return 1 + ((a - 1) / b); // if a != 0
3 }
```

1.2 To Decimal

```
1 const string digits { "0123456789
2     ABCDEFGHIJKLMNOPQRSTUVWXYZ" };
3
4 long long to_decimal(const string& rep, long long
5     base) {
6     long long n = 0;
7
8     for (auto c : rep) {
9         // if the number can't be represented in this
10        base
11        if (c > digits[base - 1]) return -1;
12        n *= base;
13        n += digits.find(c);
14    }
15
16    return n;
17 }
```

1.3 Subsets

```
1 void subsets(vector<int>& nums){
2     int n = nums.size();
3     int powSize = 1 << n;
4
5     for(int counter = 0; counter < powSize; counter++){
6         for(int j = 0; j < n; j++) {
7             if((counter & (1LL << j)) != 0) {
8                 cout << nums[j] << ' ';
9             }
10            cout << '\n';
11        }
12    }
13 }
```

1.4 Matrix Exponentiation

```
1 // Description:
2 // Calculate the nth term of a linear recursion
3
4 // Example Fibonacci:
5 // Given a linear recurrence, for example fibonacci
6 // F(n) = n, x <= 1
7 // F(n) = F(n - 1) + F(n - 2), x > 1
8
9 // The recurrence has two terms, so we can build a
10 // matrix 2 x 1 so that
11 // n + 1 = transition * n
12 // (2 x 1) = (2 x 2) * (2 x 1)
13 // F(n)      = a b * F(n - 1)
14 // F(n - 1)   c d   F(n - 2)
15
16 // Another Example:
17 // Given a grid 3 x n, you want to color it using 3
18 // distinct colors so that
19 // no adjacent place has the same color. In how many
20 // different ways can you do that?
21 // There are 6 ways for the first column to be
22 // colored using 3 distinct colors
23 // ans 6 ways using 2 equal colors and 1 distinct one
```

```
21
22 // Adding another column, there are:
23 // 3 ways to go from 2 equal to 2 equal
24 // 2 ways to go from 2 equal to 3 distinct
25 // 2 ways to go from 3 distinct to 2 equal
26 // 2 ways to go from 3 distinct to 3 distinct
27
28 // So we start with matrix 6 6 and multiply it by the
29 // transition 3 2 and get 18 12
30 //
31 //          2 2          6 6          12 12
32 // the we can exponentiate this matrix to find the
33 // nth column
34
35 // Problem:
36 // https://cses.fi/problemset/task/1722/
37
38 // Complexity:
39 // O(log n)
40
41 // How to use:
42 // vector<vector<ll>> v = {{1, 1}, {1, 0}};
43 // Matriz transition = Matriz(v);
44 // cout << fexp(transition, n)[0][1] << '\n';
45
46 using ll = long long;
47
48 const int MOD = 1e9+7;
49
50 struct Matriz{
51     vector<vector<ll>> mat;
52     int rows, columns;
53
54     vector<ll> operator[](int i){
55         return mat[i];
56     }
57
58     Matriz(vector<vector<ll>>& matriz){
59         mat = matriz;
60         rows = mat.size();
61         columns = mat[0].size();
62     }
63
64     Matriz(int row, int column, bool identity=false){
65         rows = row; columns = column;
66         mat.assign(rows, vector<ll>(columns, 0));
67         if(identity) {
68             for(int i = 0; i < min(rows, columns); i
69             ++){
70                 mat[i][i] = 1;
71             }
72         }
73     }
74
75     Matriz operator * (Matriz a) {
76         assert(columns == a.rows);
77         vector<vector<ll>> resp(rows, vector<ll>(a.
78         columns, 0));
79
80         for(int i = 0; i < rows; i++){
81             for(int j = 0; j < a.columns; j++){
82                 for(int k = 0; k < a.rows; k++){
83                     resp[i][j] = (resp[i][j] + (mat[i
84                     ][k] * 1LL * a[k][j]) % MOD) % MOD;
85                 }
86             }
87         }
88         return Matriz(resp);
89     }
90
91     Matriz operator + (Matriz a) {
92         assert(rows == a.rows && columns == a.columns
93         );
```

```

87     vector<vector<ll>> resp(rows, vector<ll>(
columns,0));
88     for(int i = 0; i < rows; i++){
89         for(int j = 0; j < columns; j++){
90             resp[i][j] = (resp[i][j] + mat[i][j]
+ a[i][j]) % MOD;
91         }
92     }
93     return Matriz(resp);
94 }
95 };
96
97 Matriz fexp(Matriz base, ll exponent){
98     Matriz result = Matriz(base.rows, base.rows, 1);
99     while(exponent > 0){
100         if(exponent & 1LL) result = result * base;
101         base = base * base;
102         exponent = exponent >> 1;
103     }
104     return result;
105 }

```

1.5 Crt

```

1 ll crt(const vector<pair<ll, ll>> &vet){
2     ll ans = 0, lcm = 1;
3     ll a, b, g, x, y;
4     for(const auto &p : vet) {
5         tie(a, b) = p;
6         tie(g, x, y) = gcd(lcm, b);
7         if((a - ans) % g != 0) return -1; // no
solution
8         ans = ans + x * ((a - ans) / g) % (b / g) *
lcm;
9         lcm = lcm * (b / g);
10        ans = (ans % lcm + lcm) % lcm;
11    }
12    return ans;
13 }

```

1.6 Binary To Decimal

```

1 int binary_to_decimal(long long n) {
2     int dec = 0, i = 0, rem;
3
4     while (n!=0) {
5         rem = n % 10;
6         n /= 10;
7         dec += rem * pow(2, i);
8         ++i;
9     }
10
11    return dec;
12 }
13
14 long long decimal_to_binary(int n) {
15     long long bin = 0;
16     int rem, i = 1;
17
18     while (n!=0) {
19         rem = n % 2;
20         n /= 2;
21         bin += rem * i;
22         i *= 10;
23     }
24
25    return bin;
26 }

```

1.7 Fast Exponentiation

```

1 ll fexp(ll b, ll e, ll mod) {
2     ll res = 1;
3     b %= mod;
4     while(e){
5         if(e & 1LL)
6             res = (res * b) % mod;
7         e = e >> 1LL;
8         b = (b * b) % mod;
9     }
10    return res;
11 }

```

1.8 Linear Diophantine Equation

```

1 // int a, b, c, x1, x2, y1, y2; cin >> a >> b >> c >>
x1 >> x2 >> y1 >> y2;
2 // int ans = -1;
3 // if (a == 0 && b == 0) {
4 //     if (c != 0) ans = 0;
5 //     else ans = (x2 - x1 + 1) * (y2 - y1 + 1);
6 // }
7 // else if (a == 0) {
8 //     if (c % b == 0 && y1 <= c / b && y2 >= c / b)
ans = (x2 - x1 + 1);
9 //     else ans = 0;
10 // }
11 // else if (b == 0) {
12 //     if (c % a == 0 && x1 <= c / a && x2 >= c / a)
ans = (y2 - y1 + 1);
13 //     else ans = 0;
14 // }
15
16 // Careful when a or b are negative or zero
17
18 // if (ans == -1) ans = find_all_solutions(a, b, c,
x1, x2, y1, y2);
19 // cout << ans << '\n';
20
21 // Problems:
22 // https://www.spoj.com/problems/CEQU/
23 // http://codeforces.com/problemsets/acmsguru/problem
/99999/106
24
25 // consider trivial case a or b is 0
26 int gcd(int a, int b, int& x, int& y) {
27     if (b == 0) {
28         x = 1;
29         y = 0;
30         return a;
31     }
32     int x1, y1;
33     int d = gcd(b, a % b, x1, y1);
34     x = y1;
35     y = x1 - y1 * (a / b);
36     return d;
37 }
38
39 // x and y are one solution and g is the gcd, all
passed as reference
40 // minx <= x <= maxx miny <= y <= maxy
41 bool find_any_solution(int a, int b, int c, int &x0,
int &y0, int &g) {
42     g = gcd(abs(a), abs(b), x0, y0);
43     if (c % g) {
44         return false;
45     }
46
47     x0 *= c / g;
48     y0 *= c / g;
49     if (a < 0) x0 = -x0;
50     if (b < 0) y0 = -y0;
51     return true;
52 }

```

```

53
54 void shift_solution(int & x, int & y, int a, int b,
    int cnt) {
55     x += cnt * b;
56     y -= cnt * a;
57 }
58
59 // return number of solutions in the interval
60 int find_all_solutions(int a, int b, int c, int minx,
    int maxx, int miny, int maxy) {
61     int x, y, g;
62     if (!find_any_solution(a, b, c, x, y, g))
63         return 0;
64     a /= g;
65     b /= g;
66
67     int sign_a = a > 0 ? +1 : -1;
68     int sign_b = b > 0 ? +1 : -1;
69
70     shift_solution(x, y, a, b, (minx - x) / b);
71     if (x < minx)
72         shift_solution(x, y, a, b, sign_b);
73     if (x > maxx)
74         return 0;
75     int lx1 = x;
76
77     shift_solution(x, y, a, b, (maxx - x) / b);
78     if (x > maxx)
79         shift_solution(x, y, a, b, -sign_b);
80     int rx1 = x;
81
82     shift_solution(x, y, a, b, -(miny - y) / a);
83     if (y < miny)
84         shift_solution(x, y, a, b, -sign_a);
85     if (y > maxy)
86         return 0;
87     int lx2 = x;
88
89     shift_solution(x, y, a, b, -(maxy - y) / a);
90     if (y > maxy)
91         shift_solution(x, y, a, b, sign_a);
92     int rx2 = x;
93
94     if (lx2 > rx2)
95         swap(lx2, rx2);
96     int lx = max(lx1, lx2);
97     int rx = min(rx1, rx2);
98
99     if (lx > rx)
100         return 0;
101     return (rx - lx) / abs(b) + 1;
102 }

```

1.9 Function Root

```

1 const ld EPS1 = 1e-9; // iteration precision error
2 const ld EPS2 = 1e-4; // output precision error
3
4 ld f(ld x) {
5     // exp(-x) == e^(-x)
6     return p * exp(-x) + q * sin(x) + r * cos(x) + s *
    tan(x) + t * x * x + u;
7 }
8
9 ld root(ld a, ld b) {
10     while (b - a >= EPS1) {
11         ld c = (a + b) / 2.0;
12         ld y = f(c);
13
14         if (y < 0) b = c;
15         else a = c;
16     }
17 }

```

```

18     return (a + b) / 2;
19 }
20
21 int main() {
22     ld ans = root(0, 1);
23     if (abs(f(ans)) <= EPS2) cout << fixed <<
        setprecision(4) << ans << '\n';
24     else cout << "No solution\n";
25
26     return 0;
27 }

```

1.10 Sieve Of Eratosthenes

```

1 vector<bool> is_prime(MAX, true);
2 vector<int> primes;
3
4 void sieve() {
5     is_prime[0] = is_prime[1] = false;
6     for (int i = 2; i < MAX; i++) {
7         if (is_prime[i]) {
8             primes.push_back(i);
9
10            for (int j = i + i; j < MAX; j += i)
11                is_prime[j] = false;
12        }
13    }
14 }

```

1.11 Horner Algorithm

```

1 // Description:
2 // Evaluates y = f(x)
3
4 // Problem:
5 // https://onlinejudge.org/index.php?option=
    com_onlinejudge&Itemid=8&page=show_problem&
    problem=439
6
7 // Complexity:
8 // O(n)
9
10 using polynomial = std::vector<int>;
11
12 polynomial p {6, -5, 2}; // p(x) = x^2 - 5x + 6;
13
14 int degree(const polynomial& p) {
15     return p.size() - 1;
16 }
17
18 int evaluate(const polynomial& p, int x) {
19     int y = 0, N = degree(p);
20
21     for (int i = N; i >= 0; --i) {
22         y *= x;
23         y += p[i];
24     }
25
26     return y;
27 }

```

1.12 Multiplicative Inverse

```

1 ll extend_euclid(ll a, ll b, ll &x, ll &y) {
2     if (a == 0)
3     {
4         x = 0; y = 1;
5         return b;
6     }
7     ll x1, y1;
8     ll d = extend_euclid(b%a, a, x1, y1);
9     x = y1 - (b / a) * x1;

```

```

10     y = x1;
11     return d;
12 }
13
14 // gcd(a, m) = 1 para existir solucao
15 // ax + my = 1, ou a*x = 1 (mod m)
16 ll inv_gcd(ll a, ll m) { // com gcd
17     ll x, y;
18     extend_euclid(a, m, x, y);
19     return ((x % m) + m) % m;
20 }
21
22 ll inv(ll a, ll phim) { // com phi(m), se m for primo
23     entao phi(m) = p-1
24     ll e = phim-1;
25     return fexp(a, e, MOD);
26 }

```

1.13 Representation Arbitrary Base

```

1 const string digits { "0123456789
2     ABCDEFGHIJKLMNOPQRSTUVWXYZ" };
3
4 string representation(int n, int b) {
5     string rep;
6
7     do {
8         rep.push_back(digits[n % b]);
9         n /= b;
10    } while (n);
11
12    reverse(rep.begin(), rep.end());
13
14    return rep;
15 }

```

1.14 Set Operations

```

1 // Complexity;
2 // O(n * m) being n and m the sizes of the two sets
3 // 2*(count1+count2)-1 (where countX is the distance
4 // between firstX and lastX):
5
6 vector<int> res;
7 set_union(s1.begin(), s1.end(), s2.begin(), s2.end(),
8     inserter(res, res.begin()));
9 set_intersection(s1.begin(), s1.end(), s2.begin(), s2
10     .end(), inserter(res, res.begin()));
11 // present in the first set, but not in the second
12 set_difference(s1.begin(), s1.end(), s2.begin(), s2
13     .end(), inserter(res, res.begin()));
14 // present in one of the sets, but not in the other
15 set_symmetric_difference(s1.begin(), s1.end(), s2
16     .begin(), s2.end(), inserter(res, res.begin()));

```

1.15 Divisors

```

1 vector<long long> all_divisors(long long n) {
2     vector<long long> ans;
3     for(long long a = 1; a*a <= n; a++){
4         if(n % a == 0) {
5             long long b = n / a;
6             ans.push_back(a);
7             if(a != b) ans.push_back(b);
8         }
9     }
10    sort(ans.begin(), ans.end());
11    return ans;
12 }

```

1.16 Check If Bit Is On

```

1 // msb de 0 é undefined
2 #define msb(n) (32 - __builtin_clz(n))
3 // #define msb(n) (64 - __builtin_clzll(n) )
4 // popcount
5 // turn bit off
6
7 bool bit_on(int n, int bit) {
8     if(1 & (n >> bit)) return true;
9     else return false;
10 }

```

1.17 Prime Factors

```

1 vector<pair<long long, int>> fatora(long long n) {
2     vector<pair<long long, int>> ans;
3     for(long long p = 2; p*p <= n; p++) {
4         if(n % p == 0) {
5             int expoente = 0;
6             while(n % p == 0) {
7                 n /= p;
8                 expoente++;
9             }
10            ans.emplace_back(p, expoente);
11        }
12    }
13    if(n > 1) ans.emplace_back(n, 1);
14    return ans;
15 }

```

2 DP

2.1 Knapsack With Index

```

1 void knapsack(int W, int wt[], int val[], int n) {
2     int i, w;
3     int K[n + 1][W + 1];
4
5     for (i = 0; i <= n; i++) {
6         for (w = 0; w <= W; w++) {
7             if (i == 0 || w == 0)
8                 K[i][w] = 0;
9             else if (wt[i - 1] <= w)
10                K[i][w] = max(val[i - 1] +
11                    K[i - 1][w - wt[i - 1]], K[i -
12                1][w]);
13            else
14                K[i][w] = K[i - 1][w];
15        }
16    }
17
18    int res = K[n][W];
19    cout << res << endl;
20
21    w = W;
22    for (i = n; i > 0 && res > 0; i--) {
23        if (res == K[i - 1][w])
24            continue;
25        else {
26            cout << " " << wt[i - 1] ;
27            res = res - val[i - 1];
28            w = w - wt[i - 1];
29        }
30    }
31
32    int main()
33    {
34        int val[] = { 60, 100, 120 };
35        int wt[] = { 10, 20, 30 };
36        int W = 50;
37        int n = sizeof(val) / sizeof(val[0]);

```

```

38     knapsack(W, wt, val, n);
39
40
41     return 0;
42 }

```

2.2 Substr Palindrome

```

1 // êvoc deve informar se a substring de S formada
  pelos elementos entre os índices i e j
2 // é um palindromo ou não.
3
4 char s[MAX];
5 int calculado[MAX][MAX]; // iniciado com false, ou 0
6 int tabela[MAX][MAX];
7
8 int is_palin(int i, int j){
9     if(calculado[i][j]){
10         return tabela[i][j];
11     }
12     if(i == j) return true;
13     if(i + 1 == j) return s[i] == s[j];
14
15     int ans = false;
16     if(s[i] == s[j]){
17         if(is_palin(i+1, j-1)){
18             ans = true;
19         }
20     }
21     calculado[i][j] = true;
22     tabela[i][j] = ans;
23     return ans;
24 }

```

2.3 Edit Distance

```

1 // Description:
2 // Minimum number of operations required to transform
  a string into another
3 // Operations allowed: add character, remove
  character, replace character
4
5 // Parameters:
6 // str1 - string to be transformed into str2
7 // str2 - string that str1 will be transformed into
8 // m - size of str1
9 // n - size of str2
10
11 // Problem:
12 // https://cses.fi/problemset/task/1639
13
14 // Complexity:
15 // O(m x n)
16
17 // How to use:
18 // memset(dp, -1, sizeof(dp));
19 // string a, b;
20 // edit_distance(a, b, (int)a.size(), (int)b.size());
21
22 // Notes:
23 // Size of dp matriz is m x n
24
25 int dp[MAX][MAX];
26
27 int edit_distance(string &str1, string &str2, int m,
  int n) {
28     if (m == 0) return n;
29     if (n == 0) return m;
30
31     if (dp[m][n] != -1) return dp[m][n];
32
33     if (str1[m - 1] == str2[n - 1]) return dp[m][n] =
  edit_distance(str1, str2, m - 1, n - 1);

```

```

34     return dp[m][n] = 1 + min({edit_distance(str1,
  str2, m, n - 1), edit_distance(str1, str2, m - 1,
  n), edit_distance(str1, str2, m - 1, n - 1)});
35 }

```

2.4 Knapsack

```

1 int val[MAXN], peso[MAXN], dp[MAXN][MAXS];
2
3 int knapsack(int n, int m){ // n Objetos | Peso max
4     for(int i=0; i<=n; i++){
5         for(int j=0; j<=m; j++){
6             if(i==0 || j==0)
7                 dp[i][j] = 0;
8             else if(peso[i-1]<=j)
9                 dp[i][j] = max(val[i-1]+dp[i-1][j-
  peso[i-1]], dp[i-1][j]);
10            else
11                dp[i][j] = dp[i-1][j];
12        }
13    }
14    return dp[n][m];
15 }

```

2.5 Digits

```

1 // achar a quantidade de numeros menores que R que
  possuem no maximo 3 digitos nao nulos
2 // a ideia eh utilizar da ordem lexicografica para
  checar isso pois se temos por exemplo
3 // o numero 8500, a gente sabe que se pegarmos o
  numero 7... qualquer digito depois do 7
4 // sera necessariamente menor q 8500
5
6 string r;
7 int tab[20][2][5];
8
9 // i - digito de R
10 // menor - ja pegou um numero menor que um digito de
  R
11 // qt - quantidade de digitos nao nulos
12 int dp(int i, bool menor, int qt){
13     if(qt > 3) return 0;
14     if(i >= r.size()) return 1;
15     if(tab[i][menor][qt] != -1) return tab[i][menor][
  qt];
16
17     int dr = r[i]-'0';
18     int res = 0;
19
20     for(int d = 0; d <= 9; d++) {
21         int dnn = qt + (d > 0);
22         if(menor == true) {
23             res += dp(i+1, true, dnn);
24         }
25         else if(d < dr) {
26             res += dp(i+1, true, dnn);
27         }
28         else if(d == dr) {
29             res += dp(i+1, false, dnn);
30         }
31     }
32
33     return tab[i][menor][qt] = res;
34 }

```

2.6 Coins

```

1 int tb[1005];
2 int n;
3 vector<int> moedas;
4
5 int dp(int i){

```

```

6     if(i >= n)
7         return 0;
8     if(tb[i] != -1)
9         return tb[i];
10
11     tb[i] = max(dp(i+1), dp(i+2) + moedas[i]);
12     return tb[i];
13 }
14
15 int main(){
16     memset(tb, -1, sizeof(tb));
17 }

```

2.7 Minimum Coin Change

```

1 int n;
2 vector<int> valores;
3
4 int tabela[1005];
5
6 int dp(int k){
7     if(k == 0){
8         return 0;
9     }
10    if(tabela[k] != -1)
11        return tabela[k];
12    int melhor = 1e9;
13    for(int i = 0; i < n; i++){
14        if(valores[i] <= k)
15            melhor = min(melhor, 1 + dp(k - valores[i]));
16    }
17    return tabela[k] = melhor;
18 }

```

2.8 Kadane

```

1 // achar uma subsequencia continua no array que a
2 // soma seja a maior possivel
3 // nesse caso vc precisa multiplicar exatamente 1
4 // elemento da subsequencia
5 // e achar a maior soma com isso
6
7 int n, x, arr[MAX], tab[MAX][2]; // tab[maior
8 // resposta no intervalo][foi multiplicado ou ãno]
9
10 int dp(int i, bool mult) {
11     if (i == n-1) {
12         if (!mult) return arr[n-1]*x;
13         return arr[n-1];
14     }
15     if (tab[i][mult] != -1) return tab[i][mult];
16
17     int res;
18
19     if (mult) {
20         res = max(arr[i], arr[i] + dp(i+1, 1));
21     }
22     else {
23         res = max({
24             arr[i]*x,
25             arr[i]*x + dp(i+1, 1),
26             arr[i] + dp(i+1, 0)
27         });
28     }
29
30     return tab[i][mult] = res;
31 }
32
33 int main() {
34     memset(tab, -1, sizeof(tab));
35 }

```

```

34 int ans = -oo;
35 for (int i = 0; i < n; i++) {
36     ans = max(ans, dp(i, 0));
37 }
38
39 return 0;
40 }
41
42
43
44 int ans = a[0], ans_l = 0, ans_r = 0;
45 int sum = 0, minus_pos = -1;
46
47 for (int r = 0; r < n; ++r) {
48     sum += a[r];
49     if (sum > ans) {
50         ans = sum;
51         ans_l = minus_pos + 1;
52         ans_r = r;
53     }
54     if (sum < 0) {
55         sum = 0;
56         minus_pos = r;
57     }
58 }

```

3 Template

3.1 Template

```

1 #include <bits/stdc++.h>
2 using namespace std;
3
4 #define int long long
5 #define optimize std::ios::sync_with_stdio(false);
6 cin.tie(NULL);
7
8 #define vi vector<int>
9 #define ll long long
10 #define pb push_back
11 #define mp make_pair
12 #define ff first
13 #define ss second
14 #define pii pair<int, int>
15 #define MOD 1000000007
16 #define sqr(x) ((x) * (x))
17 #define all(x) (x).begin(), (x).end()
18 #define FOR(i, j, n) for (int i = j; i < n; i++)
19 #define qle(i, n) (i == n ? "\n" : " ")
20 #define endl "\n"
21 const int oo = 1e9;
22 const int MAX = 1e6;
23
24 int32_t main(){ optimize;
25
26     return 0;
27 }

```

3.2 Template Clean

```

1 // Notes:
2 // Compile and execute
3 // g++ teste.cpp -o teste -std=c++17
4 // ./teste < teste.txt
5
6 // Print with precision
7 // cout << fixed << setprecision(12) << value << endl
8 // ;
9 // File as input and output
10 // freopen("input.txt", "r", stdin);
11 // freopen("output.txt", "w", stdout);

```



```

12
13 #include <bits/stdc++.h>
14 using namespace std;
15
16 #define pb push_back
17 #define mp make_pair
18 #define mt make_tuple
19 #define ff first
20 #define ss second
21 #define ld long double
22 #define ll long long
23 #define int long long
24 #define pii pair<int, int>
25 #define tii tuple<int, int, int>
26
27 int main() {
28     ios::sync_with_stdio(false);
29     cin.tie(NULL);
30
31
32
33     return 0;
34 }

```

4 Strings

4.1 Hash

```

1 // Description:
2 // Turns a string into a integer.
3 // If the hash is different then the strings are
  different.
4 // If the hash is the same the strings may be
  different.
5
6 // Problem:
7 // https://codeforces.com/gym/104518/problem/I
8
9 // Complexity:
10 // O(n) to calculate the hash
11 // O(1) to query
12
13 // Notes:
14 // Primes 1000000007, 1000041323, 100663319,
  201326611, 1000015553, 1000028537
15
16 struct Hash {
17     const ll P = 31;
18     int n; string s;
19     vector<ll> h, hi, p;
20     Hash() {}
21     Hash(string s): s(s), n(s.size()), h(n), hi(n), p
  (n) {
22         for (int i=0;i<n;i++) p[i] = (i ? P*p[i-1]:1)
  % MOD;
23         for (int i=0;i<n;i++)
24             h[i] = (s[i] + (i ? h[i-1]:0) * P) % MOD;
25         for (int i=n-1;i>=0;i--)
26             hi[i] = (s[i] + (i+1<n ? hi[i+1]:0) * P)
  % MOD;
27     }
28     int query(int l, int r) {
29         ll hash = (h[r] - (l ? h[l-1]*p[r-l+1]:MOD :
  0));
30         return hash < 0 ? hash + MOD : hash;
31     }
32     int query_inv(int l, int r) {
33         ll hash = (hi[l] - (r+1 < n ? hi[r+1]*p[r-l
  +1] % MOD : 0));
34         return hash < 0 ? hash + MOD : hash;
35     }
36 };

```

4.2 Kmp

```

1 vector<int> prefix_function(string s) {
2     int n = (int)s.length();
3     vector<int> pi(n);
4     for (int i = 1; i < n; i++) {
5         int j = pi[i-1];
6         while (j > 0 && s[i] != s[j])
7             j = pi[j-1];
8         if (s[i] == s[j])
9             j++;
10        pi[i] = j;
11    }
12    return pi;
13 }

```

4.3 Generate All Permutations

```

1 vector<string> generate_permutations(string s) {
2     int n = s.size();
3     vector<string> ans;
4
5     sort(s.begin(), s.end());
6
7     do {
8         ans.push_back(s);
9     } while (next_permutation(s.begin(), s.end()));
10
11     return ans;
12 }

```

4.4 Generate All Sequences Length K

```

1 // gera todas as ípossveis êsequencias usando as letras
  em set (de comprimento n) e que tenham tamanho k
2 // sequence = ""
3 vector<string> generate_sequences(char set[], string
  sequence, int n, int k) {
4     if (k == 0){
5         return { sequence };
6     }
7
8     vector<string> ans;
9     for (int i = 0; i < n; i++) {
10         auto aux = generate_sequences(set, sequence +
  set[i], n, k - 1);
11         ans.insert(ans.end(), aux.begin(), aux.end())
  ;
12         // for (auto e : aux) ans.push_back(e);
13     }
14
15     return ans;
16 }

```

4.5 Suffix Array

```

1 // Description:
2 // Suffix array is an array with the indexes of the
  starting letter of every
3 // suffix in an array sorted in lexicographical order
  .
4
5 // Problem:
6 // https://codeforces.com/edu/course/2/lesson/2/1/
  practice/contest/269100/problem/A
7
8 // Complexity:
9 // O(n log n) with radix sort
10 // O(n log ^ 2 n) with regular sort
11
12 // Notes:
13 // Relevant Problems

```

```

14 // Substring search: Queries to know whether a given
    substring is present in a string
15 // Binary search for the first suffix that is greater
    or equal
16 //  $O(\log n |p|)$  where  $|p|$  is the total size of the
    substrings queried
17 //
18 // Substring size: Queries to know how many times a
    given substring appears in a string
19 // Binary search both for first and last that is
    greater or equal
20 //
21 // Number of different substrings:
22 // A given suffix gives sz new substrings being sz
    the size of the suffix
23 // We can subtract the lcp (longest common prefix) to
    remove substrings
24 // that were already counted.
25 //
26 // Longest common substring between two strings:
27 // We can calculate the suffix array and lcp array of
    the two strings
28 // concatenated with a character greater than $ and
    smaller than A (like '&')
29 // The answer will be the lcp between two consecutive
    suffixes that belong to different strings
30 // (index at suffix array <= size of the first array)
31
32 void radix_sort(vector<pair<pair<int, int>, int>>& a)
    {
33     int n = a.size();
34     vector<pair<pair<int, int>, int>> ans(n);
35
36     vector<int> count(n);
37
38     for (int i = 0; i < n; i++) {
39         count[a[i].first.second]++;
40     }
41
42     vector<int> p(n);
43
44     p[0] = 0;
45     for (int i = 1; i < n; i++) {
46         p[i] = p[i - 1] + count[i - 1];
47     }
48
49     for (int i = 0; i < n; i++) {
50         ans[p[a[i].first.second]++] = a[i];
51     }
52
53     a = ans;
54
55     count.assign(n, 0);
56
57     for (int i = 0; i < n; i++) {
58         count[a[i].first.first]++;
59     }
60
61     p.assign(n, 0);
62
63     p[0] = 0;
64     for (int i = 1; i < n; i++) {
65         p[i] = p[i - 1] + count[i - 1];
66     }
67
68     for (int i = 0; i < n; i++) {
69         ans[p[a[i].first.first]++] = a[i];
70     }
71
72     a = ans;
73 }
74
75 vector<int> p, c;
76
77 vector<int> suffix_array(string s) {
78     int n = s.size();
79     vector<pair<char, int>> a(n);
80     p.assign(n, 0);
81     c.assign(n, 0);
82
83     for (int i = 0; i < n; i++) {
84         a[i] = mp(s[i], i);
85     }
86
87     sort(a.begin(), a.end());
88
89     for (int i = 0; i < n; i++) {
90         p[i] = a[i].second;
91     }
92
93     c[p[0]] = 0;
94     for (int i = 1; i < n; i++) {
95         if (a[i].first == a[i - 1].first) c[p[i]] = c[p[i
            - 1]];
96         else c[p[i]] = c[p[i - 1]] + 1;
97     }
98
99     int k = 0;
100     while ((1 << k) < n) {
101         vector<pair<pair<int, int>, int>> a(n);
102         for (int i = 0; i < n; i++) {
103             a[i] = mp(mp(c[i], c[(i + (1 << k)) % n]), i);
104         }
105
106         radix_sort(a);
107
108         for (int i = 0; i < n; i++) {
109             p[i] = a[i].second;
110         }
111
112         c[p[0]] = 0;
113         for (int i = 1; i < n; i++) {
114             if (a[i].first == a[i - 1].first) c[p[i]] = c[p
                [i - 1]];
115             else c[p[i]] = c[p[i - 1]] + 1;
116         }
117
118         k++;
119     }
120
121     /* for (int i = 0; i < n; i++) {
122         for (int j = p[i]; j < n; j++) {
123             cout << s[j];
124         }
125         cout << '\n';
126     } */
127
128     return p;
129 }
130
131 // the first suffix will always be $ the (n - 1)th
    character in the string
132 vector<int> lcp_array(string s) {
133     int n = s.size();
134     vector<int> ans(n);
135     // minimum lcp
136     int k = 0;
137     for (int i = 0; i < n - 1; i++) {
138         // indice in the suffix array p of suffix
            starting in i
139         int pi = c[i];
140         // start index of the previous suffix in suffix
            array
141         int j = p[pi - 1];
142         while (s[i + k] == s[j + k]) k++;
143         ans[pi] = k;

```

```

144     k = max(k - 1, 0);
145 }
146
147 return ans;
148 }

```

4.6 Lcs

```

1 // Description:
2 // Finds the longest common subsequence between two
  string
3
4 // Problem:
5 // https://codeforces.com/gym/103134/problem/B
6
7 // Complexity:
8 // O(mn) where m and n are the length of the strings
9
10 string lcsAlgo(string s1, string s2, int m, int n) {
11     int LCS_table[m + 1][n + 1];
12
13     for (int i = 0; i <= m; i++) {
14         for (int j = 0; j <= n; j++) {
15             if (i == 0 || j == 0)
16                 LCS_table[i][j] = 0;
17             else if (s1[i - 1] == s2[j - 1])
18                 LCS_table[i][j] = LCS_table[i - 1][j - 1] +
19                 1;
20             else
21                 LCS_table[i][j] = max(LCS_table[i - 1][j],
22                 LCS_table[i][j - 1]);
23         }
24     }
25
26     int index = LCS_table[m][n];
27     char lcsAlgo[index + 1];
28     lcsAlgo[index] = '\0';
29
30     int i = m, j = n;
31     while (i > 0 && j > 0) {
32         if (s1[i - 1] == s2[j - 1]) {
33             lcsAlgo[index - 1] = s1[i - 1];
34             i--;
35             j--;
36             index--;
37         }
38         else if (LCS_table[i - 1][j] > LCS_table[i][j - 1])
39             i--;
40         else
41             j--;
42     }
43
44     return lcsAlgo;
45 }

```

4.7 Trie

```

1 const int K = 26;
2
3 struct Vertex {
4     int next[K];
5     bool output = false;
6     int p = -1;
7     char pch;
8     int link = -1;
9     int go[K];
10
11     Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
12         fill(begin(next), end(next), -1);
13         fill(begin(go), end(go), -1);

```

```

14     }
15 };
16
17 vector<Vertex> t(1);
18
19 void add_string(string const& s) {
20     int v = 0;
21     for (char ch : s) {
22         int c = ch - 'a';
23         if (t[v].next[c] == -1) {
24             t[v].next[c] = t.size();
25             t.emplace_back(v, ch);
26         }
27         v = t[v].next[c];
28     }
29     t[v].output = true;
30 }
31
32 int go(int v, char ch);
33
34 int get_link(int v) {
35     if (t[v].link == -1) {
36         if (v == 0 || t[v].p == 0)
37             t[v].link = 0;
38         else
39             t[v].link = go(get_link(t[v].p), t[v].pch
40 );
41     }
42     return t[v].link;
43 }
44
45 int go(int v, char ch) {
46     int c = ch - 'a';
47     if (t[v].go[c] == -1) {
48         if (t[v].next[c] != -1)
49             t[v].go[c] = t[v].next[c];
50         else
51             t[v].go[c] = v == 0 ? 0 : go(get_link(v),
52             ch);
53     }
54     return t[v].go[c];
55 }

```

4.8 Z-function

```

1 vector<int> z_function(string s) {
2     int n = (int) s.length();
3     vector<int> z(n);
4     for (int i = 1, l = 0, r = 0; i < n; ++i) {
5         if (i <= r)
6             z[i] = min(r - i + 1, z[i - l]);
7         while (i + z[i] < n && s[z[i]] == s[i + z[i]
8         ])
9             ++z[i];
10         if (i + z[i] - 1 > r)
11             l = i, r = i + z[i] - 1;
12     }
13     return z;
14 }

```

5 Misc

5.1 Split

```

1 vector<string> split(string txt, char key = ' '){
2     vector<string> ans;
3
4     string palTemp = "";
5     for(int i = 0; i < txt.size(); i++){
6
7         if(txt[i] == key){

```

```

8         if(palTemp.size() > 0){
9             ans.push_back(palTemp);
10            palTemp = "";
11        }
12    } else{
13        palTemp += txt[i];
14    }
15 }
16
17 if(palTemp.size() > 0)
18     ans.push_back(palTemp);
19
20 return ans;
21 }
22 }

```

5.2 Int128

```

1 __int128 read() {
2     __int128 x = 0, f = 1;
3     char ch = getchar();
4     while (ch < '0' || ch > '9') {
5         if (ch == '-') f = -1;
6         ch = getchar();
7     }
8     while (ch >= '0' && ch <= '9') {
9         x = x * 10 + ch - '0';
10        ch = getchar();
11    }
12    return x * f;
13 }
14 void print(__int128 x) {
15     if (x < 0) {
16         putchar('-');
17         x = -x;
18     }
19     if (x > 9) print(x / 10);
20     putchar(x % 10 + '0');
21 }

```

6 Graphs

6.1 Centroid Find

```

1 // Description:
2 // Indexed at zero
3 // Find a centroid, that is a node such that when it
4 // is appointed the root of the tree,
5 // each subtree has at most floor(n/2) nodes.
6 // Problem:
7 // https://cses.fi/problemset/task/2079/
8
9 // Complexity:
10 // O(n)
11
12 // How to use:
13 // get_subtree_size(0);
14 // cout << get_centroid(0) + 1 << endl;
15
16 int n;
17 vector<int> adj[MAX];
18 int subtree_size[MAX];
19
20 int get_subtree_size(int node, int par = -1) {
21     int &res = subtree_size[node];
22     res = 1;
23     for (int i : adj[node]) {
24         if (i == par) continue;
25         res += get_subtree_size(i, node);
26     }

```

```

27     return res;
28 }
29
30 int get_centroid(int node, int par = -1) {
31     for (int i : adj[node]) {
32         if (i == par) continue;
33
34         if (subtree_size[i] * 2 > n) { return
35             get_centroid(i, node); }
36     }
37     return node;
38 }
39
40 int main() {
41     cin >> n;
42     for (int i = 0; i < n - 1; i++) {
43         int u, v; cin >> u >> v;
44         u--; v--;
45         adj[u].push_back(v);
46         adj[v].push_back(u);
47     }
48     get_subtree_size(0);
49     cout << get_centroid(0) + 1 << endl;
50 }

```

6.2 Bipartite

```

1 const int NONE = 0, BLUE = 1, RED = 2;
2 vector<vector<int>> graph(100005);
3 vector<bool> visited(100005);
4 int color[100005];
5
6 bool bfs(int s = 1){
7
8     queue<int> q;
9     q.push(s);
10    color[s] = BLUE;
11
12    while (not q.empty()){
13        auto u = q.front(); q.pop();
14
15        for (auto v : graph[u]){
16            if (color[v] == NONE){
17                color[v] = 3 - color[u];
18                q.push(v);
19            }
20            else if (color[v] == color[u]){
21                return false;
22            }
23        }
24    }
25
26    return true;
27 }
28
29 bool is_bipartite(int n){
30
31     for (int i = 1; i <= n; i++)
32         if (color[i] == NONE and not bfs(i))
33             return false;
34
35     return true;
36 }

```

6.3 Prim

```

1 int n;
2 vector<vector<int>> adj; // adjacency matrix of graph
3 const int INF = 1000000000; // weight INF means there
4     is no edge

```

```

5 struct Edge {
6     int w = INF, to = -1;
7 };
8
9 void prim() {
10     int total_weight = 0;
11     vector<bool> selected(n, false);
12     vector<Edge> min_e(n);
13     min_e[0].w = 0;
14
15     for (int i=0; i<n; ++i) {
16         int v = -1;
17         for (int j = 0; j < n; ++j) {
18             if (!selected[j] && (v == -1 || min_e[j].
19 w < min_e[v].w))
20                 v = j;
21
22             if (min_e[v].w == INF) {
23                 cout << "No MST!" << endl;
24                 exit(0);
25             }
26
27             selected[v] = true;
28             total_weight += min_e[v].w;
29             if (min_e[v].to != -1)
30                 cout << v << " " << min_e[v].to << endl;
31
32             for (int to = 0; to < n; ++to) {
33                 if (adj[v][to] < min_e[to].w)
34                     min_e[to] = {adj[v][to], v};
35             }
36         }
37
38         cout << total_weight << endl;
39 }

```

6.4 Eulerian Undirected

```

1 // Description:
2 // Hierholzer's Algorithm
3 // An Eulerian path is a path that passes through
4 // every edge exactly once.
5 // An Eulerian circuit is an Eulerian path that
6 // starts and ends on the same node.
7
8 // An Eulerian path exists in an undirected graph if
9 // the degree of every node is even (not counting
10 // self-edges)
11 // except for possibly exactly two nodes that have
12 // and odd degree (start and end nodes).
13 // An Eulerian circuit exists in an undirected graph
14 // if the degree of every node is even.
15
16 // The graph has to be connected (except for isolated
17 // nodes which are allowed because there
18 // are no edges connected to them).
19
20 // Problem:
21 // https://cses.fi/problemset/task/1691
22
23 // Complexity:
24 //  $O(E \cdot \log(E))$  where E is the number of edges
25
26 // How to use
27 // Check whether the path exists before trying to
28 // find it
29 // Find the root - any node that has at least 1
30 // outgoing edge
31 // (if the problem requires that you start from a
32 // node v, the root will be the node v)
33 // Count the degree;
34 //

```

```

25 // for (int i = 0; i < m; i++) {
26 //     int a, b; cin >> a >> b;
27 //     adj[a].pb(b); adj[b].pb(a);
28 //     root = a;
29 //     degree[a]++; degree[b]++;
30 // }
31
32 // Notes
33 // If you want to find a path start and ending nodes
34 // v and u
35 // if ((is_eulerian(n, root, start, end) != 1) || (
36 //     start != v) || (end != u)) cout << "IMPOSSIBLE\n"
37
38 // It can be speed up to work on  $O(E)$  on average by
39 // using unordered_set instead of set
40
41 // It works when there are self loops, but not when
42 // there are multiple edges
43 // If the graph has multiple edges, add more notes to
44 // simulate the edges
45 // e.g
46 // 1 2
47 // 1 2
48 // 1 2
49 // becomes
50 // 3 4
51 // 4 1
52 // 1 2
53
54 vector<bool> visited;
55 vector<int> degree;
56 vector<vector<int>> adj;
57
58 void dfs(int v) {
59     visited[v] = true;
60     for (auto u : adj[v]) {
61         if (!visited[u]) dfs(u);
62     }
63 }
64
65 int is_eulerian(int n, int root, int& start, int& end
66 ) {
67     start = -1, end = -1;
68     if (n == 1) return 2; // only one node
69     visited.assign(n + 1, false);
70     dfs(root);
71
72     for (int i = 1; i <= n; i++) {
73         if (!visited[i] && degree[i] > 0) return 0;
74     }
75
76     for (int i = 1; i <= n; i++) {
77         if (start == -1 && degree[i] % 2 == 1) start = i;
78         else if (end == -1 && degree[i] % 2 == 1) end = i;
79         else if (degree[i] % 2 == 1) return 0;
80     }
81
82     if (start == -1 && end == -1) {start = root; end =
83         root; return 2;} // has eulerian circuit and path
84     if (start != -1 && end != -1) return 1; // has
85         eulerian path
86     return 0; // no eulerian path nor circuit
87 }
88
89 vector<int> path;
90 vector<set<int>> mark;
91
92 void dfs_path(int v) {
93     visited[v] = true;
94
95     while (degree[v] != 0) {
96         degree[v]--;
97     }
98 }

```

```

89     int u = adj[v][degree[v]];
90     if (mark[v].find(u) != mark[v].end()) continue;
91     mark[v].insert(u);
92     mark[u].insert(v);
93     int next_edge = adj[v][degree[v]];
94     dfs_path(next_edge);
95 }
96 path.pb(v);
97 }
98
99 void find_path(int n, int start) {
100     path.clear();
101     mark.resize(n + 1);
102     visited.assign(n + 1, false);
103     dfs_path(start);
104 }

```

6.5 Ford Fulkerson Edmonds Karp

```

1 // Description:
2 // Obtains the maximum possible flow rate given a
  network. A network is a graph with a single
  source vertex and a single sink vertex in which
  each edge has a capacity
3
4 // Complexity:
5 //  $O(V * E^2)$  where V is the number of vertex and E
  is the number of edges
6
7 int n;
8 vector<vector<int>> capacity;
9 vector<vector<int>> adj;
10
11 int bfs(int s, int t, vector<int>& parent) {
12     fill(parent.begin(), parent.end(), -1);
13     parent[s] = -2;
14     queue<pair<int, int>> q;
15     q.push({s, INF});
16
17     while (!q.empty()) {
18         int cur = q.front().first;
19         int flow = q.front().second;
20         q.pop();
21
22         for (int next : adj[cur]) {
23             if (parent[next] == -1 && capacity[cur][
next]) {
24                 parent[next] = cur;
25                 int new_flow = min(flow, capacity[cur
][next]);
26                 if (next == t)
27                     return new_flow;
28                 q.push({next, new_flow});
29             }
30         }
31     }
32
33     return 0;
34 }
35
36 int maxflow(int s, int t) {
37     int flow = 0;
38     vector<int> parent(n);
39     int new_flow;
40
41     while (new_flow = bfs(s, t, parent)) {
42         flow += new_flow;
43         int cur = t;
44         while (cur != s) {
45             int prev = parent[cur];
46             capacity[prev][cur] -= new_flow;
47             capacity[cur][prev] += new_flow;
48             cur = prev;

```

```

49         }
50     }
51
52     return flow;
53 }

```

6.6 Hld Edge

```

1 // Description:
2 // Make queries and updates between two vertexes on a
  tree
3
4 // Problem:
5 // https://www.spoj.com/problems/QTREE/
6
7 // Complexity:
8 //  $O(\log^2 n)$  for both query and update
9
10 // How to use:
11 // HLD hld = HLD(n + 1, adj)
12
13 // Notes
14 // Change the root of the tree on the constructor if
  it's different from 1
15 // Use together with Segtree
16
17 struct HLD {
18     vector<int> parent;
19     vector<int> pos;
20     vector<int> head;
21     vector<int> subtree_size;
22     vector<int> level;
23     vector<int> heavy_child;
24     vector<ftype> subtree_weight;
25     vector<ftype> path_weight;
26     vector<vector<int>> adj;
27     vector<int> at;
28     Segtree seg = Segtree(0);
29     int cpos;
30     int n;
31     int root;
32
33     HLD() {}
34
35     HLD(int n, vector<vector<int>>& adj, int root = 1)
      : adj(adj), n(n), root(root) {
36         seg = Segtree(n);
37         cpos = 0;
38         at.assign(n, 0);
39         parent.assign(n, 0);
40         pos.assign(n, 0);
41         head.assign(n, 0);
42         subtree_size.assign(n, 1);
43         level.assign(n, 0);
44         heavy_child.assign(n, -1);
45         parent[root] = -1;
46         dfs(root, -1);
47         decompose(root, -1);
48     }
49
50     void dfs(int v, int p) {
51         parent[v] = p;
52         if (p != -1) level[v] = level[p] + 1;
53         for (auto u : adj[v]) {
54             if (u != p) {
55                 dfs(u, v);
56                 subtree_size[v] += subtree_size[u];
57                 if (heavy_child[v] == -1 || subtree_size[u] >
subtree_size[heavy_child[v]]) heavy_child[v] = u
58             }
59         }
60     }

```

```

61 void decompose(int v, int chead) {
62     // start a new path
63     if (chead == -1) chead = v;
64
65     // consecutive ids in the hld path
66     at[cpos] = v;
67     pos[v] = cpos++;
68     head[v] = chead;
69
70     // if not a leaf
71     if (heavy_child[v] != -1) decompose(heavy_child[v], chead);
72
73     // light child
74     for (auto u : adj[v]){
75         // start new path
76         if (u != parent[v] && u != heavy_child[v])
77             decompose(u, -1);
78     }
79 }
80
81 ll query_path(int a, int b) {
82     if (a == b) return 0;
83     if (pos[a] < pos[b]) swap(a, b);
84
85     if (head[a] == head[b]) return seg.query(pos[b] + 1, pos[a]);
86     return seg.f(seg.query(pos[head[a]], pos[a]),
87         query_path(parent[head[a]], b));
88 }
89
90 ftype query_subtree(int a) {
91     if (subtree_size[a] == 1) return 0;
92     return seg.query(pos[a] + 1, pos[a] + subtree_size[a] - 1);
93 }
94
95 void update_path(int a, int b, int x) {
96     if (a == b) return;
97     if (pos[a] < pos[b]) swap(a, b);
98
99     if (head[a] == head[b]) return (void)seg.update(
100         pos[b] + 1, pos[a], x);
101     seg.update(pos[head[a]], pos[a], x); update_path(
102         parent[head[a]], b, x);
103 }
104
105 void update_subtree(int a, int val) {
106     if (subtree_size[a] == 1) return;
107     seg.update(pos[a] + 1, pos[a] + subtree_size[a] - 1, val);
108 }
109
110 // vertex
111 void update(int a, int val) {
112     seg.update(pos[a], pos[a], val);
113 }
114
115 // edge
116 void update(int a, int b, int val) {
117     if (parent[a] == b) swap(a, b);
118     update(b, val);
119 }
120
121 int lca(int a, int b) {
122     if (pos[a] < pos[b]) swap(a, b);
123     return head[a] == head[b] ? b : lca(parent[head[a]], b);
124 }

```

6.7 Floyd Warshall

```

1 #include <bits/stdc++.h>
2
3 using namespace std;
4 using ll = long long;
5
6 const int MAX = 507;
7 const long long INF = 0x3f3f3f3f3f3f3fLL;
8
9 ll dist[MAX][MAX];
10 int n;
11
12 void floyd_warshall() {
13     for (int i = 0; i < n; i++) {
14         for (int j = 0; j < n; j++) {
15             if (i == j) dist[i][j] = 0;
16             else if (!dist[i][j]) dist[i][j] = INF;
17         }
18     }
19
20     for (int k = 0; k < n; k++) {
21         for (int i = 0; i < n; i++) {
22             for (int j = 0; j < n; j++) {
23                 // trata o caso no qual o grafo tem
24                 // arestas com peso negativo
25                 if (dist[i][k] < INF && dist[k][j] <
26                     INF){
27                     dist[i][j] = min(dist[i][j], dist
28                         [i][k] + dist[k][j]);
29                 }
30             }
31         }
32     }
33 }

```

6.8 Lca

```

1 // Description:
2 // Find the lowest common ancestor between two nodes
3 // in a tree
4
5 // Problem:
6 // https://cses.fi/problemset/task/1135
7
8 // Complexity:
9 // O(log n)
10
11 // How to use:
12 // preprocess();
13 // lca(a, b);
14
15 // Notes
16 // To calculate the distance between two nodes use
17 // the following formula
18 // level_peso[a] + level_peso[b] - 2*level_peso[lca(a, b)]
19
20 const int MAX = 2e5+10;
21 const int BITS = 30;
22
23 vector<pii> adj[MAX];
24 vector<bool> visited(MAX);
25
26 int up[MAX][BITS + 1];
27 int level[MAX];
28 int level_peso[MAX];
29
30 void find_level() {
31     queue<pii> q;
32     q.push(mp(1, 0));

```

```

32     visited[1] = true;
33
34     while (!q.empty()) {
35         auto [v, depth] = q.front();
36         q.pop();
37         level[v] = depth;
38
39         for (auto [u, d] : adj[v]) {
40             if (!visited[u]) {
41                 visited[u] = true;
42                 up[u][0] = v;
43                 q.push(mp(u, depth + 1));
44             }
45         }
46     }
47 }
48
49 void find_level_peso() {
50     queue<pii> q;
51
52     q.push(mp(1, 0));
53     visited[1] = true;
54
55     while (!q.empty()) {
56         auto [v, depth] = q.front();
57         q.pop();
58         level_peso[v] = depth;
59
60         for (auto [u, d] : adj[v]) {
61             if (!visited[u]) {
62                 visited[u] = true;
63                 up[u][0] = v;
64                 q.push(mp(u, depth + d));
65             }
66         }
67     }
68 }
69
70 int lca(int a, int b) {
71     // get the nodes to the same level
72     int mn = min(level[a], level[b]);
73
74     for (int j = 0; j <= BITS; j++) {
75         if (a != -1 && ((level[a] - mn) & (1 << j))) a
76         = up[a][j];
77         if (b != -1 && ((level[b] - mn) & (1 << j))) b
78         = up[b][j];
79     }
80
81     // special case
82     if (a == b) return a;
83
84     // binary search
85     for (int j = BITS; j >= 0; j--) {
86         if (up[a][j] != up[b][j]) {
87             a = up[a][j];
88             b = up[b][j];
89         }
90     }
91     return up[a][0];
92 }
93
94 void preprocess() {
95     visited = vector<bool>(MAX, false);
96     find_level();
97     visited = vector<bool>(MAX, false);
98     find_level_peso();
99
100     for (int j = 1; j <= BITS; j++) {
101         for (int i = 1; i <= n; i++) {
102             if (up[i][j - 1] != -1) up[i][j] = up[up[i][j - 1]][j - 1];
103         }
104     }

```

```

102     }
103 }

```

6.9 Kuhn

```

1 // Description
2 // Matching algorithm for unweighted bipartite graph
3 //
4 // Problem:
5 // https://codeforces.com/gym/104252/problem/H
6 //
7 // Complexity:
8 //  $O(V * E)$  in which  $V$  is the number of vertexes and
9 //  $E$  is the number of edges
10 //
11 // Notes:
12 // Indexed at zero
13
14 int n, k;
15 // adjacency list
16 vector<vector<int>> g;
17 vector<int> mt;
18 vector<bool> used;
19
20 bool try_kuhn(int v) {
21     if (used[v])
22         return false;
23     used[v] = true;
24     for (int to : g[v]) {
25         if (mt[to] == -1 || try_kuhn(mt[to])) {
26             mt[to] = v;
27             return true;
28         }
29     }
30     return false;
31 }
32
33 int main() {
34     // ... reading the graph g ...
35
36     mt.assign(k, -1);
37     vector<bool> used1(n, false);
38     for (int v = 0; v < n; ++v) {
39         for (int to : g[v]) {
40             if (mt[to] == -1) {
41                 mt[to] = v;
42                 used1[v] = true;
43                 break;
44             }
45         }
46     }
47     for (int v = 0; v < n; ++v) {
48         if (used1[v])
49             continue;
50         used.assign(n, false);
51         try_kuhn(v);
52     }
53
54     for (int i = 0; i < k; ++i)
55         if (mt[i] != -1)
56             printf("%d %d\n", mt[i] + 1, i + 1);

```

6.10 Eulerian Directed

```

1 // Description:
2 // Hierholzer's Algorithm
3 // An Eulerian path is a path that passes through
4 // every edge exactly once.
5 // An Eulerian circuit is an Eulerian path that
6 // starts and ends on the same node.

```



```

5
6 // An Eulerian path exists in an directed graph if
7 // the indegree and outdegree is equal
8 // for every node (not counting self-edges)
9 // except for possibly exactly one node that have
10 // outdegree - indegree = 1
11 // and one node that has indegree - outdegree = 1 (
12 // start and end nodes).
13 // An Eulerian circuit exists in an directed graph if
14 // the indegree and outdegree is equal for every
15 // node.
16
17 // The graph has to be conected (except for isolated
18 // nodes which are allowed because there
19 // are no edges connected to them).
20
21 // Problem:
22 // https://cses.fi/problemset/task/1693
23
24 // Complexity:
25 // O(E) where E is the number of edges
26
27 // How to use
28 // Check whether the path exists before trying to
29 // find it
30 // Find the root - any node that has at least 1
31 // outgoing edge
32 // (if the problem requires that you start from a
33 // node v, the root will be the node v)
34 // Count the degree;
35 //
36 // for (int i = 0; i < m; i++) {
37 //   int a, b; cin >> a >> b;
38 //   adj[a].pb(b);
39 //   root = a;
40 //   outdegree[a]++; indegree[b]++;
41 // }
42
43 // Notes
44 // It works when there are self loops, but not when
45 // there are multiple edges
46
47 vector<bool> visited;
48 vector<int> outdegree, indegree;
49 vector<vector<int>> adj, undir;
50
51 void dfs(int v) {
52     visited[v] = true;
53     for (auto u : undir[v]) {
54         if (!visited[u]) dfs(u);
55     }
56 }
57
58 int is_eulerian(int n, int root, int &start, int& end
59 ) {
60     start = -1, end = -1;
61     if (n == 1) return 2; // only one node
62     visited.assign(n + 1, false);
63     dfs(root);
64
65     for (int i = 1; i <= n; i++) {
66         if (!visited[i] && (i == n || i == 1 || outdegree
67         [i] + indegree[i] > 0)) return 0;
68     }
69
70     // start => node with indegree - outdegree = 1
71     // end => node with outdegree - indegree = 1
72     for (int i = 1; i <= n; i++) {
73         if (start == -1 && indegree[i] - outdegree[i] ==
74         1) start = i;
75         else if (end == -1 && outdegree[i] - indegree[i]
76         == 1) end = i;
77         else if (indegree[i] != outdegree[i]) return 0;
78     }
79
80     if (start == -1 && end == -1) {start = root; end =
81     root; return 2;} // has eulerian circuit and path
82     if (start != -1 && end != -1) {swap(start, end);
83     return 1;} // has eulerian path
84     return 0; // no eulerian path nor circuit
85 }
86
87 vector<int> path;
88
89 void dfs_path(int v) {
90     visited[v] = true;
91
92     while (outdegree[v] != 0) {
93         int u = adj[v][--outdegree[v]];
94         int next_edge = adj[v][outdegree[v]];
95         dfs_path(next_edge);
96     }
97     path.pb(v);
98 }
99
100 void find_path(int n, int start) {
101     path.clear();
102     visited.assign(n + 1, false);
103     dfs_path(start);
104     reverse(path.begin(), path.end());
105 }

```

6.11 Bellman Ford

```

1 // Description:
2 // Finds the shortest path from a vertex v to any
3 // other vertex
4
5 // Problem:
6 // https://cses.fi/problemset/task/1673
7
8 // Complexity:
9 // O(n * m)
10
11 struct Edge {
12     int a, b, cost;
13     Edge(int a, int b, int cost) : a(a), b(b), cost(
14     cost) {}
15 };
16
17 int n, m;
18 vector<Edge> edges;
19 const int INF = 1e9+10;
20
21 void bellman_ford(int v, int t) {
22     vector<int> d(n + 1, INF);
23     d[v] = 0;
24     vector<int> p(n + 1, -1);
25
26     for (;;) {
27         bool any = false;
28         for (Edge e : edges) {
29             if (d[e.a] >= INF) continue;
30             if (d[e.b] > d[e.a] + e.cost) {
31                 d[e.b] = d[e.a] + e.cost;
32                 p[e.b] = e.a;
33                 any = true;
34             }
35         }
36         if (!any) break;
37     }
38
39     if (d[t] == INF)
40         cout << "No path from " << v << " to " << t << ".
41         ";
42     else {

```

```

40     vector<int> path;
41     for (int cur = t; cur != -1; cur = p[cur]) {
42         path.push_back(cur);
43     }
44     reverse(path.begin(), path.end());
45
46     cout << "Path from " << v << " to " << t << ": ";
47     for (int u : path) {
48         cout << u << ' ';
49     }
50 }
51 }

```

6.12 Dinic

```

1 // Description:
2 // Obtains the maximum possible flow rate given a
  // network. A network is a graph with a single
  // source vertex and a single sink vertex in which
  // each edge has a capacity
3
4 // Problem:
5 // https://codeforces.com/gym/103708/problem/J
6
7 // Complexity:
8 //  $O(V^2 * E)$  where  $V$  is the number of vertex and  $E$ 
  // is the number of edges
9
10 // Unit network
11 // A unit network is a network in which for any
  // vertex except source and sink either incoming or
  // outgoing edge is unique and has unit capacity (
  // matching problem).
12 // Complexity on unit networks:  $O(E * \sqrt{V})$ 
13
14 // Unity capacity networks
15 // A more generic settings when all edges have unit
  // capacities, but the number of incoming and
  // outgoing edges is unbounded
16 // Complexity on unity capacity networks:  $O(E * \sqrt{E})$ 
17
18 // How to use:
19 // Dinic dinic = Dinic(num_vertex, source, sink);
20 // dinic.add_edge(vertex1, vertex2, capacity);
21 // cout << dinic.max_flow() << '\n';
22
23 #include <bits/stdc++.h>
24
25 #define pb push_back
26 #define mp make_pair
27 #define pii pair<int, int>
28 #define ff first
29 #define ss second
30 #define ll long long
31
32 using namespace std;
33
34 const ll INF = 1e18+10;
35
36 struct Edge {
37     int from;
38     int to;
39     ll capacity;
40     ll flow;
41     Edge* residual;
42
43     Edge() {}
44
45     Edge(int from, int to, ll capacity) : from(from),
46     to(to), capacity(capacity) {
47         flow = 0;
48     }
49 }

```

```

48
49 ll get_capacity() {
50     return capacity - flow;
51 }
52
53 ll get_flow() {
54     return flow;
55 }
56
57 void augment(ll bottleneck) {
58     flow += bottleneck;
59     residual->flow -= bottleneck;
60 }
61
62 void reverse(ll bottleneck) {
63     flow -= bottleneck;
64     residual->flow += bottleneck;
65 }
66
67 bool operator<(const Edge& e) const {
68     return true;
69 }
70 };
71
72 struct Dinic {
73     int source;
74     int sink;
75     int nodes;
76     ll flow;
77     vector<vector<Edge*>> adj;
78     vector<int> level;
79     vector<int> next;
80     vector<int> reach;
81     vector<bool> visited;
82     vector<vector<int>> path;
83
84     Dinic(int source, int sink, int nodes) : source(
85     source), sink(sink), nodes(nodes) {
86         adj.resize(nodes + 1);
87     }
88
89     void add_edge(int from, int to, ll capacity) {
90         Edge* e1 = new Edge(from, to, capacity);
91         Edge* e2 = new Edge(to, from, 0);
92         // Edge* e2 = new Edge(to, from, capacity);
93         e1->residual = e2;
94         e2->residual = e1;
95         adj[from].pb(e1);
96         adj[to].pb(e2);
97     }
98
99     bool bfs() {
100         level.assign(nodes + 1, -1);
101         queue<int> q;
102         q.push(source);
103         level[source] = 0;
104
105         while (!q.empty()) {
106             int node = q.front();
107             q.pop();
108
109             for (auto e : adj[node]) {
110                 if (level[e->to] == -1 && e->
111                 get_capacity() > 0) {
112                     level[e->to] = level[e->from] +
113                     1;
114                     q.push(e->to);
115                 }
116             }
117         }
118
119         return level[sink] != -1;
120     }
121 }

```

```

118 ll dfs(int v, ll flow) {
119     if (v == sink)
120         return flow;
121
122     int sz = adj[v].size();
123     for (int i = next[v]; i < sz; i++) {
124         Edge* e = adj[v][i];
125         if (level[e->to] == level[e->from] + 1 &&
126             e->get_capacity() > 0) {
127             ll bottleneck = dfs(e->to, min(flow,
128             e->get_capacity()));
129             if (bottleneck > 0) {
130                 e->augment(bottleneck);
131                 return bottleneck;
132             }
133             next[v] = i + 1;
134         }
135     }
136     return 0;
137 }
138
139 ll max_flow() {
140     flow = 0;
141     while(bfs()) {
142         next.assign(nodes + 1, 0);
143         ll sent = -1;
144         while (sent != 0) {
145             sent = dfs(source, INF);
146             flow += sent;
147         }
148     }
149     return flow;
150 }
151
152 void reachable(int v) {
153     visited[v] = true;
154
155     for (auto e : adj[v]) {
156         if (!visited[e->to] && e->get_capacity()
157 > 0) {
158             reach.pb(e->to);
159             visited[e->to] = true;
160             reachable(e->to);
161         }
162     }
163 }
164
165 void print_min_cut() {
166     reach.clear();
167     visited.assign(nodes + 1, false);
168     reach.pb(source);
169     reachable(source);
170
171     for (auto v : reach) {
172         for (auto e : adj[v]) {
173             if (!visited[e->to] && e->
174 get_capacity() == 0) {
175                 cout << e->from << ' ' << e->to
176 << '\n';
177             }
178         }
179     }
180
181 ll build_path(int v, int id, ll flow) {
182     visited[v] = true;
183     if (v == sink) {
184         return flow;
185     }
186     for (auto e : adj[v]) {
187         if (!visited[e->to] && e->get_flow() > 0)
188         {
189             visited[e->to] = true;
190             ll bottleneck = build_path(e->to, id,
191 min(flow, e->get_flow()));
192             if (bottleneck > 0) {
193                 path[id].pb(e->to);
194                 e->reverse(bottleneck);
195                 return bottleneck;
196             }
197         }
198     }
199     return 0;
200 }
201
202 void print_flow_path() {
203     path.clear();
204     ll sent = -1;
205     int id = -1;
206     while (sent != 0) {
207         visited.assign(nodes + 1, false);
208         path.pb(vector<int>{});
209         sent = build_path(source, ++id, INF);
210         path[id].pb(source);
211     }
212     path.pop_back();
213
214     for (int i = 0; i < id; i++) {
215         cout << path[i].size() << '\n';
216         reverse(path[i].begin(), path[i].end());
217         for (auto e : path[i]) {
218             cout << e << ' ';
219         }
220         cout << '\n';
221     }
222 }
223
224 int main() {
225     ios::sync_with_stdio(false);
226     cin.tie(NULL);
227
228     int n, m; cin >> n >> m;
229
230     Dinic dinic = Dinic(1, n, n);
231
232     for (int i = 1; i <= m; i++) {
233         int v, u; cin >> v >> u;
234         dinic.add_edge(v, u, 1);
235     }
236
237     cout << dinic.max_flow() << '\n';
238     // dinic.print_min_cut();
239     // dinic.print_flow_path();
240
241     return 0;
242 }

```

6.13 2sat

```

1 // Description:
2 // Solves expression of the type (a v b) ^ (c v d) ^
  (e v f)
3
4 // Problem:
5 // https://cses.fi/problemset/task/1684
6
7 // Complexity:
8 // O(n + m) where n is the number of variables and m
  is the number of clauses
9

```

```

10 #include <bits/stdc++.h>
11 #define pb push_back
12 #define mp make_pair
13 #define pii pair<int, int>
14 #define ff first
15 #define ss second
16
17 using namespace std;
18
19 struct SAT {
20     int nodes;
21     int curr = 0;
22     int component = 0;
23     vector<vector<int>> adj;
24     vector<vector<int>> rev;
25     vector<vector<int>> condensed;
26     vector<pii> departure;
27     vector<bool> visited;
28     vector<int> scc;
29     vector<int> order;
30
31     // 1 to nodes
32     // nodes + 1 to 2 * nodes
33     SAT(int nodes) : nodes(nodes) {
34         adj.resize(2 * nodes + 1);
35         rev.resize(2 * nodes + 1);
36         visited.resize(2 * nodes + 1);
37         scc.resize(2 * nodes + 1);
38     }
39
40     void add_imp(int a, int b) {
41         adj[a].pb(b);
42         rev[b].pb(a);
43     }
44
45     int get_not(int a) {
46         if (a > nodes) return a - nodes;
47         return a + nodes;
48     }
49
50     void add_or(int a, int b) {
51         add_imp(get_not(a), b);
52         add_imp(get_not(b), a);
53     }
54
55     void add_nor(int a, int b) {
56         add_or(get_not(a), get_not(b));
57     }
58
59     void add_and(int a, int b) {
60         add_or(get_not(a), b);
61         add_or(a, get_not(b));
62         add_or(a, b);
63     }
64
65     void add_nand(int a, int b) {
66         add_or(get_not(a), b);
67         add_or(a, get_not(b));
68         add_or(get_not(a), get_not(b));
69     }
70
71     void add_xor(int a, int b) {
72         add_or(a, b);
73         add_or(get_not(a), get_not(b));
74     }
75
76     void add_xnor(int a, int b) {
77         add_or(get_not(a), b);
78         add_or(a, get_not(b));
79     }
80
81     void departure_time(int v) {
82         visited[v] = true;

```

```

83
84         for (auto u : adj[v]) {
85             if (!visited[u]) departure_time(u);
86         }
87
88         departure.pb(mp(++curr, v));
89     }
90
91     void find_component(int v, int component) {
92         scc[v] = component;
93         visited[v] = true;
94
95         for (auto u : rev[v]) {
96             if (!visited[u]) find_component(u,
97                 component);
98         }
99     }
100
101     void topological_order(int v) {
102         visited[v] = true;
103
104         for (auto u : condensed[v]) {
105             if (!visited[u]) topological_order(u);
106         }
107
108         order.pb(v);
109     }
110
111     bool is_possible() {
112         component = 0;
113         for (int i = 1; i <= 2 * nodes; i++) {
114             if (!visited[i]) departure_time(i);
115         }
116
117         sort(departure.begin(), departure.end(),
118             greater<pii>());
119
120         visited.assign(2 * nodes + 1, false);
121
122         for (auto [_ , node] : departure) {
123             if (!visited[node]) find_component(node,
124                 ++component);
125         }
126
127         for (int i = 1; i <= nodes; i++) {
128             if (scc[i] == scc[i + nodes]) return
129                 false;
130         }
131
132         return true;
133     }
134
135     int find_value(int e, vector<int> &ans) {
136         if (e > nodes && ans[e - nodes] != 2) return
137             !ans[e - nodes];
138         if (e <= nodes && ans[e + nodes] != 2) return
139             !ans[e + nodes];
140         return 0;
141     }
142
143     vector<int> find_ans() {
144         condensed.resize(component + 1);
145
146         for (int i = 1; i <= 2 * nodes; i++) {
147             for (auto u : adj[i]) {
148                 if (scc[i] != scc[u]) condensed[scc[i]
149                     ].pb(scc[u]);
150             }
151         }
152
153         visited.assign(component + 1, false);
154
155         for (int i = 1; i <= component; i++) {

```

```

149         if (!visited[i]) topological_order(i);
150     }
151
152     reverse(order.begin(), order.end());
153
154     // 0 - false
155     // 1 - true
156     // 2 - no value yet
157     vector<int> ans(2 * nodes + 1, 2);
158
159     vector<vector<int>> belong(component + 1);
160
161     for (int i = 1; i <= 2 * nodes; i++) {
162         belong[scc[i]].pb(i);
163     }
164
165     for (auto p : order) {
166         for (auto e : belong[p]) {
167             ans[e] = find_value(e, ans);
168         }
169     }
170
171     return ans;
172 }
173 };
174
175 int main() {
176     ios::sync_with_stdio(false);
177     cin.tie(NULL);
178
179     int n, m; cin >> n >> m;
180
181     SAT sat = SAT(m);
182
183     for (int i = 0; i < n; i++) {
184         char op1, op2; int a, b; cin >> op1 >> a >>
185         op2 >> b;
186         if (op1 == '+' && op2 == '+') sat.add_or(a, b);
187         if (op1 == '-' && op2 == '-') sat.add_or(sat.get_not(a), sat.get_not(b));
188         if (op1 == '+' && op2 == '-') sat.add_or(a, sat.get_not(b));
189         if (op1 == '-' && op2 == '+') sat.add_or(sat.get_not(a), b);
190     }
191
192     if (!sat.is_possible()) cout << "IMPOSSIBLE\n";
193     else {
194         vector<int> ans = sat.find_ans();
195         for (int i = 1; i <= m; i++) {
196             cout << (ans[i] == 1 ? '+' : '-') << ' ';
197         }
198         cout << '\n';
199     }
200
201     return 0;
202 }

```

6.14 Find Cycle

```

1  bitset<MAX> visited;
2  vector<int> path;
3  vector<int> adj[MAX];
4
5  bool dfs(int u, int p){
6
7      if (visited[u]) return false;
8
9      path.pb(u);
10     visited[u] = true;
11
12     for (auto v : adj[u]){

```

```

13         if (visited[v] and u != v and p != v){
14             path.pb(v); return true;
15         }
16
17         if (dfs(v, u)) return true;
18     }
19
20     path.pop_back();
21     return false;
22 }
23
24 bool has_cycle(int N){
25
26     visited.reset();
27
28     for (int u = 1; u <= N; ++u){
29         path.clear();
30         if (not visited[u] and dfs(u,-1))
31             return true;
32     }
33
34     return false;
35 }
36 }

```

6.15 Cycle Path Recovery

```

1  int n;
2  vector<vector<int>> adj;
3  vector<char> color;
4  vector<int> parent;
5  int cycle_start, cycle_end;
6
7  bool dfs(int v) {
8      color[v] = 1;
9      for (int u : adj[v]) {
10         if (color[u] == 0) {
11             parent[u] = v;
12             if (dfs(u))
13                 return true;
14         } else if (color[u] == 1) {
15             cycle_end = v;
16             cycle_start = u;
17             return true;
18         }
19     }
20     color[v] = 2;
21     return false;
22 }
23
24 void find_cycle() {
25     color.assign(n, 0);
26     parent.assign(n, -1);
27     cycle_start = -1;
28
29     for (int v = 0; v < n; v++) {
30         if (color[v] == 0 && dfs(v))
31             break;
32     }
33
34     if (cycle_start == -1) {
35         cout << "Acyclic" << endl;
36     } else {
37         vector<int> cycle;
38         cycle.push_back(cycle_start);
39         for (int v = cycle_end; v != cycle_start; v =
40             parent[v])
41             cycle.push_back(v);
42         cycle.push_back(cycle_start);
43         reverse(cycle.begin(), cycle.end());
44
45         cout << "Cycle found: ";
46         for (int v : cycle)

```

```

46         cout << v << " ";
47     cout << endl;
48 }
49 }

```

6.16 Blossom

```

1 // Description:
2 // Matching algorithm for general graphs (non-
   bipartite)
3
4 // Problem:
5 // https://acm.timus.ru/problem.aspx?space=1&num=1099
6
7 // Complexity:
8 // O (n ^3)
9
10 // vector<pii> Blossom(vector<vector<int>>& graph) {
11 vector<int> Blossom(vector<vector<int>>& graph) {
12     int n = graph.size(), timer = -1;
13     vector<int> mate(n, -1), label(n), parent(n),
14         orig(n), aux(n, -1), q;
15     auto lca = [&](int x, int y) {
16         for (timer++; ; swap(x, y)) {
17             if (x == -1) continue;
18             if (aux[x] == timer) return x;
19             aux[x] = timer;
20             x = (mate[x] == -1 ? -1 : orig[parent[mate[x]
21 ]]);
22         }
23     };
24     auto blossom = [&](int v, int w, int a) {
25         while (orig[v] != a) {
26             parent[v] = w; w = mate[v];
27             if (label[w] == 1) label[w] = 0, q.push_back(w)
28         };
29         orig[v] = orig[w] = a; v = parent[w];
30     };
31     auto augment = [&](int v) {
32         while (v != -1) {
33             int pv = parent[v], nv = mate[pv];
34             mate[v] = pv; mate[pv] = v; v = nv;
35         }
36     };
37     auto bfs = [&](int root) {
38         fill(label.begin(), label.end(), -1);
39         iota(orig.begin(), orig.end(), 0);
40         q.clear();
41         label[root] = 0; q.push_back(root);
42         for (int i = 0; i < (int)q.size(); ++i) {
43             int v = q[i];
44             for (auto x : graph[v]) {
45                 if (label[x] == -1) {
46                     label[x] = 1; parent[x] = v;
47                     if (mate[x] == -1)
48                         return augment(x), 1;
49                     label[mate[x]] = 0; q.push_back(mate[x]);
50                 } else if (label[x] == 0 && orig[v] != orig[x]
51 ) {
52                     int a = lca(orig[v], orig[x]);
53                     blossom(x, v, a); blossom(v, x, a);
54                 }
55             }
56         }
57         return 0;
58     };
59     // Time halves if you start with (any) maximal
   matching.
60     for (int i = 0; i < n; i++)
61         if (mate[i] == -1)
62             bfs(i);
63     return mate;

```

```

62
63     /*
64     vector<bool> used(n, false);
65     vector<pii> ans;
66     for (int i = 0; i < n; i++) {
67         if (matching[i] == -1 || used[i]) continue;
68         used[i] = true;
69         used[matching[i]] = true;
70         ans.emplace_back(i, matching[i]);
71     }
72     return ans;
73     */
74 }

```

6.17 Centroid Decomposition

```

1 int n;
2 vector<set<int>> adj;
3 vector<char> ans;
4
5 vector<bool> removed;
6
7 vector<int> subtree_size;
8
9 int dfs(int u, int p = 0) {
10     subtree_size[u] = 1;
11
12     for(int v : adj[u]) {
13         if(v != p && !removed[v]) {
14             subtree_size[u] += dfs(v, u);
15         }
16     }
17
18     return subtree_size[u];
19 }
20
21 int get_centroid(int u, int sz, int p = 0) {
22     for(int v : adj[u]) {
23         if(v != p && !removed[v]) {
24             if(subtree_size[v]*2 > sz) {
25                 return get_centroid(v, sz, u);
26             }
27         }
28     }
29
30     return u;
31 }
32
33 char get_next(char c) {
34     if (c != 'Z') return c + 1;
35     return '$';
36 }
37
38 bool flag = true;
39
40 void solve(int node, char c) {
41     int center = get_centroid(node, dfs(node));
42     ans[center] = c;
43     removed[center] = true;
44
45     for (auto u : adj[center]) {
46         if (!removed[u]) {
47             char next = get_next(c);
48             if (next == '$') {
49                 flag = false;
50                 return;
51             }
52             solve(u, next);
53         }
54     }
55 }
56
57 int32_t main(){

```

```

58 ios::sync_with_stdio(false);
59 cin.tie(NULL);
60
61 cin >> n;
62 adj.resize(n + 1);
63 ans.resize(n + 1);
64 removed.resize(n + 1);
65 subtree_size.resize(n + 1);
66
67 for (int i = 1; i <= n - 1; i++) {
68     int u, v; cin >> u >> v;
69     adj[u].insert(v);
70     adj[v].insert(u);
71 }
72
73 solve(1, 'A');
74
75 if (!flag) cout << "Impossible!\n";
76 else {
77     for (int i = 1; i <= n; i++) {
78         cout << ans[i] << ' ';
79     }
80     cout << '\n';
81 }
82
83 return 0;
84 }

```

6.18 Tarjan Bridge

```

1 // Description:
2 // Find a bridge in a connected undirected graph
3 // A bridge is an edge so that if you remove that
4 // edge the graph is no longer connected
5
6 // Problem:
7 // https://cses.fi/problemset/task/2177/
8
9 // Complexity:
10 // O(V + E) where V is the number of vertices and E
11 // is the number of edges
12
13 int n;
14 vector<vector<int>>> adj;
15
16 vector<bool> visited;
17 vector<int> tin, low;
18 int timer;
19
20 void dfs(int v, int p) {
21     visited[v] = true;
22     tin[v] = low[v] = timer++;
23     for (int to : adj[v]) {
24         if (to == p) continue;
25         if (visited[to]) {
26             low[v] = min(low[v], tin[to]);
27         } else {
28             dfs(to, v);
29             low[v] = min(low[v], low[to]);
30             if (low[to] > tin[v]) {
31                 IS_BRIDGE(v, to);
32             }
33         }
34     }
35 }
36
37 void find_bridges() {
38     timer = 0;
39     visited.assign(n, false);
40     tin.assign(n, -1);
41     low.assign(n, -1);
42     for (int i = 0; i < n; ++i) {
43         if (!visited[i]) {
44             dfs(i, -1);
45         }
46     }
47 }

```

```

42         dfs(i, -1);
43     }
44 }

```

6.19 Hld Vertex

```

1 // Description:
2 // Make queries and updates between two vertexes on a
3 // tree
4 // Query path - query path (a, b) inclusive
5 // Update path - update path (a, b) inclusive
6 // Query subtree - query subtree of a
7 // Update subtree - update subtree of a
8 // Update - update vertex or edge
9 // Lca - get lowest common ancestor of a and b
10 // Search - perform a binary search to find the last
11 // node with a certain property
12 // on the path from a to the root
13
14 // Problem:
15 // https://codeforces.com/gym/101908/problem/L
16
17 // Complexity:
18 // O(log^2 n) for both query and update
19
20 // How to use:
21 // HLD hld = HLD(n + 1, adj)
22
23 // Notes
24 // Change the root of the tree on the constructor if
25 // it's different from 1
26 // Use together with Segtree
27
28 typedef long long ftype;
29
30 struct HLD {
31     vector<int> parent;
32     vector<int> pos;
33     vector<int> head;
34     vector<int> subtree_size;
35     vector<int> level;
36     vector<int> heavy_child;
37     vector<ftype> subtree_weight;
38     vector<ftype> path_weight;
39     vector<vector<int>>> adj;
40     vector<int> at;
41     Segtree seg = Segtree(0);
42     int cpos;
43     int n;
44     int root;
45     vector<vector<int>>> up;
46
47     HLD() {}
48
49     HLD(int n, vector<vector<int>>>& adj, int root = 1)
50         : adj(adj), n(n), root(root) {
51         seg = Segtree(n);
52         cpos = 0;
53         at.resize(n);
54         parent.resize(n);
55         pos.resize(n);
56         head.resize(n);
57         subtree_size.assign(n, 1);
58         level.assign(n, 0);
59         heavy_child.assign(n, -1);
60         parent[root] = -1;
61         dfs(root, -1);
62         decompose(root, -1);
63     }
64
65     void dfs(int v, int p) {
66         parent[v] = p;
67         if (p != -1) level[v] = level[p] + 1;
68     }

```

```

64     for (auto u : adj[v]) {
65         if (u != p) {
66             dfs(u, v);
67             subtree_size[v] += subtree_size[u];
68             if (heavy_child[v] == -1 || subtree_size[u] > subtree_size[heavy_child[v]]) heavy_child[v] = u;
69         }
70     }
71 }
72
73 void decompose(int v, int head) {
74     // start a new path
75     if (head == -1) head = v;
76
77     // consecutive ids in the hld path
78     at[cpos] = v;
79     pos[v] = cpos++;
80     head[v] = head;
81
82     // if not a leaf
83     if (heavy_child[v] != -1) decompose(heavy_child[v], head);
84
85     // light child
86     for (auto u : adj[v]){
87         // start new path
88         if (u != parent[v] && u != heavy_child[v])
89             decompose(u, -1);
90     }
91
92     ftype query_path(int a, int b) {
93         if(pos[a] < pos[b]) swap(a, b);
94
95         if(head[a] == head[b]) return seg.query(pos[b], pos[a]);
96         return seg.f(seg.query(pos[head[a]], pos[a]), query_path(parent[head[a]], b));
97     }
98
99     // iterative
100     /*ftype query_path(int a, int b) {
101         ftype ans = 0;
102
103         while (head[a] != head[b]) {
104             if (level[head[a]] > level[head[b]]) swap(a, b);
105             ans = seg.merge(ans, seg.query(pos[head[b]], pos[b]));
106             b = parent[head[b]];
107         }
108
109         if (level[a] > level[b]) swap(a, b);
110         ans = seg.merge(ans, seg.query(pos[a], pos[b]));
111         return ans;
112     }*/
113
114     ftype query_subtree(int a) {
115         return seg.query(pos[a], pos[a] + subtree_size[a] - 1);
116     }
117
118     void update_path(int a, int b, int x) {
119         if(pos[a] < pos[b]) swap(a, b);
120
121         if(head[a] == head[b]) return (void)seg.update(pos[b], pos[a], x);
122         seg.update(pos[head[a]], pos[a], x); update_path(parent[head[a]], b, x);
123     }
124
125     void update_subtree(int a, int val) {
126         seg.update(pos[a], pos[a] + subtree_size[a] - 1, val);
127     }
128
129     void update(int a, int val) {
130         seg.update(pos[a], pos[a], val);
131     }
132
133     //edge
134     void update(int a, int b, int val) {
135         if (level[a] > level[b]) swap(a, b);
136         update(b, val);
137     }
138
139     int lca(int a, int b) {
140         if(pos[a] < pos[b]) swap(a, b);
141         return head[a] == head[b] ? b : lca(parent[head[a]], b);
142     }
143
144     void search(int a) {
145         a = parent[a];
146         if (a == -1) return;
147         if (seg.query(pos[head[a]], pos[head[a]] + subtree_size[head[a]] - 1) + pos[a] - pos[head[a]] + 1 == subtree_size[head[a]]) {
148             seg.update(pos[head[a]], pos[a], 1);
149             return search(parent[head[a]]);
150         }
151         int l = pos[head[a]], r = pos[a] + 1;
152         while (l < r) {
153             int m = (l + r) / 2;
154             if (seg.query(m, m + subtree_size[at[m]] - 1) + pos[a] - m + 1 == subtree_size[at[m]]) {
155                 r = m;
156             }
157             else l = m + 1;
158         }
159         seg.update(l, pos[a], 1);
160     }
161
162     /* k-th ancestor of x
163     int x, k; cin >> x >> k;
164
165     for (int b = 0; b <= BITS; b++) {
166         if (x != -1 && (k & (1 << b))) {
167             x = up[x][b];
168         }
169     }
170
171     cout << x << '\n';
172     */
173     void preprocess() {
174         up.assign(n + 1, vector<int>(31, -1));
175
176         for (int i = 1; i < n; i++) {
177             up[i][0] = parent[i];
178         }
179
180         for (int i = 1; i < n; i++) {
181             for (int j = 1; j <= 30; j++) {
182                 if (up[i][j - 1] != -1) up[i][j] = up[up[i][j - 1]][j - 1];
183             }
184         }
185     }
186
187     int getKth(int p, int q, int k){
188         int a = lca(p, q), d;
189
190         if( a == p ){
191             d = level[q] - level[p] + 1;
192             swap(p, q);

```



```

193     k = d - k + 1 ;
194 }
195 else if( a == q ) ;
196 else {
197     if( k > level[p] - level[a] + 1 ) {
198         d = level[p] + level[q] - 2 * level[a] +
199         1 ;
200         k = d - k + 1 ;
201         swap(p,q);
202     }
203     else ;
204 }
205 int lg ; for( lg = 1 ; (1 << lg) <= level[p] ; ++
206 lg ); lg--;
207 k--;
208 for( int i = lg ; i >= 0 ; i-- ){
209     if( (1 << i) <= k ){
210         p = up[p][i];
211         k -= ( 1 << i );
212     }
213 }
214 return p;
215 }
216 };

```

6.20 Small To Large

```

1 // Problem:
2 // https://codeforces.com/contest/600/problem/E
3
4 void process_colors(int curr, int parent) {
5
6     for (int n : adj[curr]) {
7         if (n != parent) {
8             process_colors(n, curr);
9
10            if (colors[curr].size() < colors[n].size
11            ()) {
12                sum_num[curr] = sum_num[n];
13                vmax[curr] = vmax[n];
14                swap(colors[curr], colors[n]);
15            }
16            for (auto [item,vzs] : colors[n]) {
17                if(colors[curr][item]+vzs > vmax[curr
18                ]){
19                    vmax[curr] = colors[curr][item] +
20                    vzs;
21                    sum_num[curr] = item;
22                }
23                else if(colors[curr][item]+vzs ==
24                vmax[curr]){
25                    sum_num[curr] += item;
26                }
27            }
28            colors[curr][item] += vzs;
29        }
30    }
31 }
32
33 int32_t main() {
34
35     int n; cin >> n;
36
37     for (int i = 1; i <= n; i++) {
38         int a; cin >> a;
39         colors[i][a] = 1;
40         vmax[i] = 1;
41         sum_num[i] = a;
42     }

```

```

43
44     for (int i = 1; i < n; i++) {
45         int a, b; cin >> a >> b;
46
47         adj[a].push_back(b);
48         adj[b].push_back(a);
49     }
50
51     process_colors(1, 0);
52
53     for (int i = 1; i <= n; i++) {
54         cout << sum_num[i] << (i < n ? " " : "\n");
55     }
56
57     return 0;
58 }
59
60

```

6.21 Tree Diameter

```

1 #include<bits/stdc++.h>
2
3 using namespace std;
4
5 const int MAX = 3e5+17;
6
7 vector<int> adj[MAX];
8 bool visited[MAX];
9
10 int max_depth = 0, max_node = 1;
11
12 void dfs (int v, int depth) {
13     visited[v] = true;
14
15     if (depth > max_depth) {
16         max_depth = depth;
17         max_node = v;
18     }
19
20     for (auto u : adj[v]) {
21         if (!visited[u]) dfs(u, depth + 1);
22     }
23 }
24
25 int tree_diameter() {
26     dfs(1, 0);
27     max_depth = 0;
28     for (int i = 0; i < MAX; i++) visited[i] = false;
29     dfs(max_node, 0);
30     return max_depth;
31 }

```

6.22 Dijkstra

```

1 const int MAX = 2e5+7;
2 const int INF = 1000000000;
3 vector<vector<pair<int, int>>> adj(MAX);
4
5 void dijkstra(int s, vector<int> & d, vector<int> & p
6 ) {
7     int n = adj.size();
8     d.assign(n, INF);
9     p.assign(n, -1);
10
11     d[s] = 0;
12     set<pair<int, int>> q;
13     q.insert({0, s});
14     while (!q.empty()) {
15         int v = q.begin()->second;
16         q.erase(q.begin());

```

```

17     for (auto edge : adj[v]) {
18         int to = edge.first;
19         int len = edge.second;
20
21         if (d[v] + len < d[to]) {
22             q.erase({d[to], to});
23             d[to] = d[v] + len;
24             p[to] = v;
25             q.insert({d[to], to});
26         }
27     }
28 }
29
30 vector<int> restore_path(int s, int t) {
31     vector<int> path;
32
33     for (int v = t; v != s; v = p[v])
34         path.push_back(v);
35     path.push_back(s);
36
37     reverse(path.begin(), path.end());
38     return path;
39 }
40
41 int adj[MAX][MAX];
42 int dist[MAX];
43 int minDistance(int dist[], bool sptSet[], int V) {
44     int min = INT_MAX, min_index;
45
46     for (int v = 0; v < V; v++)
47         if (sptSet[v] == false && dist[v] <= min)
48             min = dist[v], min_index = v;
49
50     return min_index;
51 }
52
53 void dijkstra(int src, int V) {
54     bool sptSet[V];
55     for (int i = 0; i < V; i++)
56         dist[i] = INT_MAX, sptSet[i] = false;
57
58     dist[src] = 0;
59
60     for (int count = 0; count < V - 1; count++) {
61         int u = minDistance(dist, sptSet, V);
62
63         sptSet[u] = true;
64
65         for (int v = 0; v < V; v++)
66             if (!sptSet[v] && adj[u][v]
67                 && dist[u] != INT_MAX
68                 && dist[u] + adj[u][v] < dist[v])
69                 dist[v] = dist[u] + adj[u][v];
70     }
71 }
72
73 }
74

```

6.23 Kruskal

```

1 struct DSU {
2     int n;
3     vector<int> link, sizes;
4
5     DSU(int n) {
6         this->n = n;
7         link.assign(n+1, 0);
8         sizes.assign(n+1, 1);
9
10        for (int i = 0; i <= n; i++)
11            link[i] = i;
12    }

```

```

13
14 int find(int x) {
15     while (x != link[x])
16         x = link[x];
17
18     return x;
19 }
20
21 bool same(int a, int b) {
22     return find(a) == find(b);
23 }
24
25 void unite(int a, int b) {
26     a = find(a);
27     b = find(b);
28
29     if (a == b) return;
30
31     if (sizes[a] < sizes[b])
32         swap(a, b);
33
34     sizes[a] += sizes[b];
35     link[b] = a;
36 }
37 };
38
39 struct Edge {
40     int u, v;
41     long long weight;
42
43     Edge() {}
44
45     Edge(int u, int v, long long weight) : u(u), v(v),
46         weight(weight) {}
47
48     bool operator<(const Edge& other) const {
49         return weight < other.weight;
50     }
51
52     bool operator>(const Edge& other) const {
53         return weight > other.weight;
54     }
55 };
56
57 vector<Edge> kruskal(vector<Edge> edges, int n) {
58     vector<Edge> result; // arestas da MST
59     long long cost = 0;
60
61     sort(edges.begin(), edges.end());
62
63     DSU dsu(n);
64
65     for (auto e : edges) {
66         if (!dsu.same(e.u, e.v)) {
67             cost += e.weight;
68             result.push_back(e);
69             dsu.unite(e.u, e.v);
70         }
71     }
72
73     return result;
74 }

```

6.24 Hungarian

```

1 // Description:
2 // A matching algorithm for weighted bipartite graphs
3 // that returns
4 // a perfect match
5
6 // Problem:
7 // https://codeforces.com/gym/103640/problem/H

```

```

8 // Complexity:
9 //  $O(V^3)$  in which  $V$  is the number of vertices
10
11 // Notes:
12 // Indexed at 1
13
14 //  $n$  is the number of items on the right side and  $m$ 
15 // the number of items
16 // on the left side of the graph
17 // Returns minimum assignment cost and which items
18 // were matched
19
20 pair<int, vector<pii>> hungarian(int n, int m, vector
21 <vector<int>> A) {
22     vector<int> u (n+1), v (m+1), p (m+1), way (m+1);
23     for (int i=1; i<=n; ++i) {
24         p[0] = i;
25         int j0 = 0;
26         vector<int> minv (m+1, INF);
27         vector<char> used (m+1, false);
28         do {
29             used[j0] = true;
30             int i0 = p[j0], delta = INF, j1;
31             for (int j=1; j<=m; ++j)
32                 if (!used[j]) {
33                     int cur = A[i0][j]-u[i0]-v[j];
34                     if (cur < minv[j])
35                         minv[j] = cur, way[j] = j0;
36                 }
37             for (int j=0; j<=m; ++j)
38                 if (used[j])
39                     u[p[j]] += delta, v[j] -= delta;
40             else
41                 minv[j] -= delta;
42             j0 = j1;
43         } while (p[j0] != 0);
44         do {
45             int j1 = way[j0];
46             p[j0] = p[j1];
47             j0 = j1;
48         } while (j0);
49     }
50
51     vector<pair<int, int>> result;
52     for (int i = 1; i <= m; ++i){
53         result.push_back(make_pair(p[i], i));
54     }
55
56     int C = -v[0];
57     return mp(C, result);
58 }

```

6.25 Negative Cycle

```

1 // Description
2 // Detects any cycle in which the sum of edge weights
3 // is negative.
4 // Alternatively, we can detect whether there is a
5 // negative cycle
6 // starting from a specific vertex.
7
8 // Problem:
9 // https://cses.fi/problemset/task/1197
10
11 // Complexity:
12 //  $O(n * m)$ 
13
14 // Notes

```

```

13 // In order to consider only the negative cycles
14 // located on the path from a to b,
15 // Reverse the graph, run a dfs from node b and mark
16 // the visited nodes
17 // Consider only the edges that connect to visited
18 // nodes when running bellman-ford
19 // on the normal graph
20
21 struct Edge {
22     int a, b, cost;
23     Edge(int a, int b, int cost) : a(a), b(b), cost(
24         cost) {}
25 };
26
27 int n, m;
28 vector<Edge> edges;
29 const int INF = 1e9+10;
30
31 void negative_cycle() {
32     // uncomment to find negative cycle starting from a
33     // vertex v
34     // vector<int> d(n + 1, INF);
35     // d[v] = 0;
36     vector<int> d(n + 1, 0);
37     vector<int> p(n + 1, -1);
38     int x;
39     // uncomment to find all negative cycles
40     // // set<int> s;
41     for (int i = 1; i <= n; ++i) {
42         x = -1;
43         for (Edge e : edges) {
44             // if (d[e.a] >= INF) continue;
45             if (d[e.b] > d[e.a] + e.cost) {
46                 // d[e.b] = max(-INF, d[e.a] + e.cost);
47                 d[e.b] = d[e.a] + e.cost;
48                 p[e.b] = e.a;
49                 x = e.b;
50                 // // s.insert(e.b);
51             }
52         }
53     }
54
55     if (x == -1)
56         cout << "NO\n";
57     else {
58         // // int y = all nodes in set s
59         int y = x;
60         for (int i = 1; i <= n; ++i) {
61             y = p[y];
62         }
63
64         vector<int> path;
65         for (int cur = y; cur = p[cur]) {
66             path.push_back(cur);
67             if (cur == y && path.size() > 1) break;
68         }
69         reverse(path.begin(), path.end());
70
71         cout << "YES\n";
72         for (int u : path)
73             cout << u << ' ';
74         cout << '\n';
75     }
76 }

```

7 Geometry

7.1 Shoelace Boundary

```

1 // Description
2 // Shoelace formula finds the area of a polygon

```

```

3 // Boundary points return the number of integer
  points on the edges of a polygon
4 // not counting the vertexes
5
6 // Problem
7 // https://codeforces.com/gym/101873/problem/G
8
9 // Complexity
10 // O(n)
11
12 // before dividing by two
13 int shoelace(vector<point> & points) {
14     int n = points.size();
15     vector<point> v(n + 2);
16
17     for (int i = 1; i <= n; i++) {
18         v[i] = points[i - 1];
19     }
20     v[n + 1] = points[0];
21
22     int sum = 0;
23     for (int i = 1; i <= n; i++) {
24         sum += (v[i].x * v[i + 1].y - v[i + 1].x * v[i].y);
25     }
26
27     sum = abs(sum);
28     return sum;
29 }
30
31 int boundary_points(vector<point> & points) {
32     int n = points.size();
33     vector<point> v(n + 2);
34
35     for (int i = 1; i <= n; i++) {
36         v[i] = points[i - 1];
37     }
38     v[n + 1] = points[0];
39
40     int ans = 0;
41     for (int i = 1; i <= n; i++) {
42         if (v[i].x == v[i + 1].x) ans += abs(v[i].y -
43         v[i + 1].y) - 1;
44         else if (v[i].y == v[i + 1].y) ans += abs(v[i].x -
45         v[i + 1].x) - 1;
46         else ans += gcd(abs(v[i].x - v[i + 1].x), abs
47         (v[i].y - v[i + 1].y)) - 1;
48     }
49     return points.size() + ans;
50 }

```

7.2 Inside Polygon

```

1 // Description
2 // Checks if a given point is inside, outside or on
  the boundary of a polygon
3
4 // Problem
5 // https://cses.fi/problemset/task/2192/
6
7 // Complexity
8 // O(n)
9
10 int inside(vp &p, point pp){
11     // 1 - inside / 0 - boundary / -1 - outside
12     int n = p.size();
13     for(int i=0;i<n;i++){
14         int j = (i+1)%n;
15         if(line({p[i], p[j]}).inside_seg(pp))
16             return 0; // boundary
17     }
18     int inter = 0;
19     for(int i=0;i<n;i++){

```

```

20         int j = (i+1)%n;
21         if(p[i].x <= pp.x and pp.x < p[j].x and ccw(p
22         [i], p[j], pp)==1)
23             inter++; // up
24         else if(p[j].x <= pp.x and pp.x < p[i].x and
25         ccw(p[i], p[j], pp)==-1)
26             inter++; // down
27     }
28     if(inter%2==0) return -1; // outside
29     else return 1; // inside

```

7.3 Closest Pair Points

```

1 // Description
2 // Find the squared distance between the closest two
  points among n points
3 // Also finds which pair of points is closest (could
  be more than one)
4
5 // Problem
6 // https://cses.fi/problemset/task/2194/
7
8 // Complexity
9 // O(n log n)
10
11 ll closest_pair_points(vp &vet){
12     pair<point, point> ans;
13     int n = vet.size();
14     sort(vet.begin(), vet.end());
15     set<point> s;
16
17     ll best_dist = LLONG_MAX;
18     int j=0;
19     for(int i=0;i<n;i++){
20         ll d = ceil(sqrt(best_dist));
21         while(j<n and vet[i].x-vet[j].x >= d){
22             s.erase(point(vet[j].y, vet[j].x));
23             j++;
24         }
25
26         auto it1 = s.lower_bound({vet[i].y - d, vet[i]
27         }.x});
28         auto it2 = s.upper_bound({vet[i].y + d, vet[i]
29         }.x});
30
31         for(auto it=it1; it!=it2; it++){
32             ll dx = vet[i].x - it->y;
33             ll dy = vet[i].y - it->x;
34
35             if(best_dist > dx*dx + dy*dy){
36                 best_dist = dx*dx + dy*dy;
37                 // closest pair points
38                 ans = mp(vet[i], point(it->y, it->x))
39             }
40         }
41         s.insert(point(vet[i].y, vet[i].x));
42     }
43     // best distance squared
44     return best_dist;
45 }

```

7.4 2d

```

1 #define vp vector<point>
2 #define ld long double
3 const ld EPS = 1e-6;
4 const ld PI = acos(-1);

```

```

5
6 // typedef ll cod;
7 // bool eq(cod a, cod b){ return (a==b); }
8 typedef ld cod;
9 bool eq(cod a, cod b){ return abs(a - b) <= EPS; }
10
11 struct point{
12     cod x, y;
13     int id;
14     point(cod x=0, cod y=0): x(x), y(y){}
15
16     point operator+(const point &o) const{ return {x+o.x, y+o.y}; }
17     point operator-(const point &o) const{ return {x-o.x, y-o.y}; }
18     point operator*(cod t) const{ return {x*t, y*t}; }
19     point operator/(cod t) const{ return {x/t, y/t}; }
20     cod operator*(const point &o) const{ return x * o.x + y * o.y; }
21     cod operator^(const point &o) const{ return x * o.y - y * o.x; }
22     bool operator<(const point &o) const{
23         return (eq(x, o.x) ? y < o.y : x < o.x);
24     }
25     bool operator==(const point &o) const{
26         return eq(x, o.x) and eq(y, o.y);
27     }
28     friend ostream& operator<<(ostream& os, point p) {
29         return os << "(" << p.x << "," << p.y << ")"; }
30 };
31
32 int ccw(point a, point b, point e){ // -1=dir; 0=
33     collinear; 1=esq;
34     cod tmp = (b-a) ^ (e-a); // vector from a to b
35     return (tmp > EPS) - (tmp < -EPS);
36 }
37 ld norm(point a){ // Modulo
38     return sqrt(a * a);
39 }
40 cod norm2(point a){
41     return a * a;
42 }
43 bool nulo(point a){
44     return (eq(a.x, 0) and eq(a.y, 0));
45 }
46 point rotccw(point p, ld a){
47     // a = PI*a/180; // graus
48     return point((p.x*cos(a)-p.y*sin(a)), (p.y*cos(a)+p.x*sin(a)));
49 }
50 point rot90cw(point a) { return point(a.y, -a.x); }
51 point rot90ccw(point a) { return point(-a.y, a.x); }
52
53 ld proj(point a, point b){ // a sobre b
54     return a*b/norm(b);
55 }
56 ld angle(point a, point b){ // em radianos
57     ld ang = a*b / norm(a) / norm(b);
58     return acos(max(min(ang, (ld)1), (ld)-1));
59 }
60 ld angle_vec(point v){
61     // return 180/PI*atan2(v.x, v.y); // graus
62     return atan2(v.x, v.y);
63 }
64 ld order_angle(point a, point b){ // from a to b ccw
65     (a in front of b)
66     ld aux = angle(a,b)*180/PI;
67     return ((a^b)<=0 ? aux:360-aux);
68 }
69 bool angle_less(point a1, point b1, point a2, point
b2){ // ang(a1,b1) <= ang(a2,b2)
70     point p1((a1*b1), abs((a1^b1)));
71     point p2((a2*b2), abs((a2^b2)));
72     return (p1^p2) <= 0;
73 }
74 ld area(vp &p){ // (points sorted)
75     ld ret = 0;
76     for(int i=2;i<(int)p.size();i++)
77         ret += (p[i]-p[0])^(p[i-1]-p[0]);
78     return abs(ret/2);
79 }
80 ld areaT(point &a, point &b, point &c){
81     return abs((b-a)^(c-a))/2.0;
82 }
83
84 point center(vp &A){
85     point c = point();
86     int len = A.size();
87     for(int i=0;i<len;i++)
88         c=c+A[i];
89     return c/len;
90 }
91
92 point forca_mod(point p, ld m){
93     ld cm = norm(p);
94     if(cm<EPS) return point();
95     return point(p.x*m/cm,p.y*m/cm);
96 }
97
98 ld param(point a, point b, point v){
99     // v = t*(b-a) + a // return t;
100     // assert(line(a, b).inside_seg(v));
101     return ((v-a) * (b-a)) / ((b-a) * (b-a));
102 }
103
104 bool simetric(vp &a){ //ordered
105     int n = a.size();
106     point c = center(a);
107     if(n&1) return false;
108     for(int i=0;i<n/2;i++)
109         if(ccw(a[i], a[i+n/2], c) != 0)
110             return false;
111     return true;
112 }
113
114 point mirror(point m1, point m2, point p){
115     // mirror point p around segment m1m2
116     point seg = m2-m1;
117     ld t0 = ((p-m1)*seg) / (seg*seg);
118     point ort = m1 + seg*t0;
119     point pm = ort-(p-ort);
120     return pm;
121 }
122
123 ///////////////
124 // Line //
125 ///////////////
126
127 struct line{
128     point p1, p2;
129     cod a, b, c; // ax+by+c = 0;
130     // y-y1 = ((y2-y1)/(x2-x1))(x-x1)
131     line(point p1=0, point p2=0): p1(p1), p2(p2){
132         a = p1.y - p2.y;
133         b = p2.x - p1.x;
134         c = p1 ^ p2;
135     }
136
137     line(cod a=0, cod b=0, cod c=0): a(a), b(b), c(c)
138     {
139         // Gera os pontos p1 p2 dados os coeficientes
140         // isso aqui eh um lixo mas quebra um galho

```

```

140 kkkkkk
141     if(b==0){
142         p1 = point(1, -c/a);
143         p2 = point(0, -c/a);
144     }else{
145         p1 = point(1, (-c-a*1)/b);
146         p2 = point(0, -c/b);
147     }
148 }
149
150 cod eval(point p){
151     return a*p.x+b*p.y+c;
152 }
153 bool inside(point p){
154     return eq(eval(p), 0);
155 }
156 point normal(){
157     return point(a, b);
158 }
159
160 bool inside_seg(point p){
161     return (
162         ((p1-p) ^ (p2-p)) == 0 and
163         ((p1-p) * (p2-p)) <= 0
164     );
165 }
166 };
167
168 // be careful with precision error
169 vp inter_line(line l1, line l2){
170     ld det = l1.a*l2.b - l1.b*l2.a;
171     if(det==0) return {};
172     ld x = (l1.b*l2.c - l1.c*l2.b)/det;
173     ld y = (l1.c*l2.a - l1.a*l2.c)/det;
174     return {point(x, y)};
175 }
176
177 // segments not collinear
178 vp inter_seg(line l1, line l2){
179     vp ans = inter_line(l1, l2);
180     if(ans.empty() or !l1.inside_seg(ans[0]) or !l2.
181         inside_seg(ans[0]))
182         return {};
183     return ans;
184 }
185
186 bool seg_has_inter(line l1, line l2){
187     // if collinear
188     if (l1.inside_seg(l2.p1) || l1.inside_seg(l2.p2)
189     || l2.inside_seg(l1.p1) || l2.inside_seg(l1.p2))
190     return true;
191
192     return ccw(l1.p1, l1.p2, l2.p1) * ccw(l1.p1, l1.
193     p2, l2.p2) < 0 and
194     ccw(l2.p1, l2.p2, l1.p1) * ccw(l2.p1, l2.
195     p2, l1.p2) < 0;
196 }
197
198 ld dist_seg(point p, point a, point b){ // point -
199     seg
200     if((p-a)*(b-a) < EPS) return norm(p-a);
201     if((p-b)*(a-b) < EPS) return norm(p-b);
202     return abs((p-a)^(b-a)) / norm(b-a);
203 }
204
205 ld dist_line(point p, line l){ // point - line
206     return abs(l.eval(p))/sqrt(l.a*l.a + l.b*l.b);
207 }
208
209 line bisector(point a, point b){
210     point d = (b-a)*2;
211     return line(d.x, d.y, a*a - b*b);
212 }

```

```

206
207 line perpendicular(line l, point p){ // passes
208     through p
209     return line(l.b, -l.a, -l.b*p.x + l.a*p.y);
210 }
211
212 // Circle //
213
214 struct circle{
215     point c; cod r;
216     circle() : c(0, 0), r(0){}
217     circle(const point o) : c(o), r(0){}
218     circle(const point a, const point b){
219         c = (a+b)/2;
220         r = norm(a-c);
221     }
222     circle(const point a, const point b, const point
223     cc){
224         assert(ccw(a, b, cc) != 0);
225         c = inter_line(bisector(a, b), bisector(b, cc
226     ))[0];
227         r = norm(a-c);
228     }
229     bool inside(const point &a) const{
230         return norm(a - c) <= r + EPS;
231     }
232 };
233
234 pair<point, point> tangent_points(circle cr, point p)
235 {
236     ld d1 = norm(p-cr.c), theta = asin(cr.r/d1);
237     point p1 = rotccw(cr.c-p, -theta);
238     point p2 = rotccw(cr.c-p, theta);
239     assert(d1 >= cr.r);
240     p1 = p1 * (sqrt(d1*d1-cr.r*cr.r) / d1) + p;
241     p2 = p2 * (sqrt(d1*d1-cr.r*cr.r) / d1) + p;
242     return {p1, p2};
243 }
244
245 circle incircle(point p1, point p2, point p3){
246     ld m1 = norm(p2-p3);
247     ld m2 = norm(p1-p3);
248     ld m3 = norm(p1-p2);
249     point c = (p1*m1 + p2*m2 + p3*m3)*(1/(m1+m2+m3));
250     ld s = 0.5*(m1+m2+m3);
251     ld r = sqrt(s*(s-m1)*(s-m2)*(s-m3)) / s;
252     return circle(c, r);
253 }
254
255 circle circumcircle(point a, point b, point c) {
256     circle ans;
257     point u = point((b-a).y, -(b-a).x);
258     point v = point((c-a).y, -(c-a).x);
259     point n = (c-b)*0.5;
260     ld t = (u^n)/(v^u);
261     ans.c = ((a+c)*0.5) + (v*t);
262     ans.r = norm(ans.c-a);
263     return ans;
264 }
265
266 vp inter_circle_line(circle C, line L){
267     point ab = L.p2 - L.p1, p = L.p1 + ab * ((C.c-L.
268     p1)*(ab) / (ab*ab));
269     ld s = (L.p2-L.p1)^(C.c-L.p1), h2 = C.r*C.r - s*s
270     / (ab*ab);
271     if (h2 < -EPS) return {};
272     if (eq(h2, 0)) return {p};
273     point h = (ab/norm(ab)) * sqrt(h2);
274     return {p - h, p + h};

```

```

273 }
274
275 vp inter_circle(circle C1, circle C2){
276     if(C1.c == C2.c) { assert(C1.r != C2.r); return
        {};}
277     point vec = C2.c - C1.c;
278     ld d2 = vec*vec, sum = C1.r+C2.r, dif = C1.r-C2.r
        ;
279     ld p = (d2 + C1.r*C1.r - C2.r*C2.r)/(d2*2), h2 =
        C1.r*C1.r - p*p*d2;
280     if (sum*sum < d2 or dif*dif > d2) return {};
281     point mid = C1.c + vec*p, per = point(-vec.y, vec
        .x) * sqrt(max((ld)0, h2) / d2);
282     if(eq(per.x, 0) and eq(per.y, 0)) return {mid};
283     return {mid + per, mid - per};
284 }
285
286 // minimum circle cover O(n) amortizado
287 circle min_circle_cover(vp v){
288     random_shuffle(v.begin(), v.end());
289     circle ans;
290     int n = v.size();
291     for(int i=0; i<n; i++){ if(!ans.inside(v[i])){
292         ans = circle(v[i]);
293         for(int j=0; j<i; j++){ if(!ans.inside(v[j])){
294             ans = circle(v[i], v[j]);
295             for(int k=0; k<j; k++){ if(!ans.inside(v[k])){
296                 ans = circle(v[i], v[j], v[k]);
297             }
298         }
299     }
300     return ans;
301 }

```

8 Algorithms

8.1 Lis

```

1 int lis(vector<int> const& a) {
2     int n = a.size();
3     vector<int> d(n, 1);
4     for (int i = 0; i < n; i++) {
5         for (int j = 0; j < i; j++) {
6             if (a[j] < a[i])
7                 d[i] = max(d[i], d[j] + 1);
8         }
9     }
10
11     int ans = d[0];
12     for (int i = 1; i < n; i++) {
13         ans = max(ans, d[i]);
14     }
15     return ans;
16 }

```

8.2 Delta-encoding

```

1 #include <bits/stdc++.h>
2 using namespace std;
3
4 int main(){
5     int n, q;
6     cin >> n >> q;
7     int [n];
8     int delta[n+2];
9
10    while(q--){
11        int l, r, x;
12        cin >> l >> r >> x;
13        delta[l] += x;

```

```

14        delta[r+1] -= x;
15    }
16
17    int curr = 0;
18    for(int i=0; i < n; i++){
19        curr += delta[i];
20        v[i] = curr;
21    }
22
23    for(int i=0; i < n; i++){
24        cout << v[i] << ' ';
25    }
26    cout << '\n';
27
28    return 0;
29 }

```

8.3 Subsets

```

1 void subsets(vector<int>& nums){
2     int n = nums.size();
3     int powSize = 1 << n;
4
5     for(int counter = 0; counter < powSize; counter++){
6         for(int j = 0; j < n; j++){
7             if((counter & (1LL << j)) != 0) {
8                 cout << nums[j] << ' ';
9             }
10        }
11        cout << '\n';
12    }
13 }

```

8.4 Binary Search Last True

```

1 int last_true(int lo, int hi, function<bool(int)> f)
2 {
3     lo--;
4     while (lo < hi) {
5         int mid = lo + (hi - lo + 1) / 2;
6         if (f(mid)) {
7             lo = mid;
8         } else {
9             hi = mid - 1;
10        }
11    }
12    return lo;

```

8.5 Ternary Search

```

1 double ternary_search(double l, double r) {
2     double eps = 1e-9; //set the error
3     limit here
4     while (r - l > eps) {
5         double m1 = l + (r - l) / 3;
6         double m2 = r - (r - l) / 3;
7         double f1 = f(m1); //evaluates the
8         function at m1
9         double f2 = f(m2); //evaluates the
10        function at m2
11        if (f1 < f2)
12            l = m1;
13        else
14            r = m2;
15    }
16    return f(l); //return the
17    maximum of f(x) in [l, r]

```

8.6 Binary Search First True

```

1 int first_true(int lo, int hi, function<bool(int)> f) 54
2 {
3     hi++;
4     while (lo < hi) {
5         int mid = lo + (hi - lo) / 2;
6         if (f(mid)) {
7             hi = mid;
8         } else {
9             lo = mid + 1;
10        }
11    }
12    return lo;
13 }

```

8.7 Biggest K

```

1 // Description: Gets sum of k biggest or k smallest
  elements in an array
2
3 // Problem: https://atcoder.jp/contests/abc306/tasks/
  abc306_e
4
5 // Complexity: O(log n)
6
7 struct SetSum {
8     ll s = 0;
9     multiset<ll> mt;
10    void add(ll x){
11        mt.insert(x);
12        s += x;
13    }
14    int pop(ll x){
15        auto f = mt.find(x);
16        if(f == mt.end()) return 0;
17        mt.erase(f);
18        s -= x;
19        return 1;
20    }
21 };
22
23 struct BigK {
24     int k;
25     SetSum gt, mt;
26     BigK(int _k){
27         k = _k;
28     }
29     void balancear(){
30         while((int)gt.mt.size() < k && (int)mt.mt.
size()){
31             auto p = (prev(mt.mt.end()));
32             gt.add(*p);
33             mt.pop(*p);
34         }
35         while((int)mt.mt.size() && (int)gt.mt.size()
&&
36             *(gt.mt.begin()) < *(prev(mt.mt.end())) ){
37             ll u = *(gt.mt.begin());
38             ll v = *(prev(mt.mt.end()));
39             gt.pop(u); mt.pop(v);
40             gt.add(v); mt.add(u);
41         }
42     }
43     void add(ll x){
44         mt.add(x);
45         balancear();
46     }
47     void rem(ll x){
48         //x = -x;
49         if(mt.pop(x) == 0)
50             gt.pop(x);
51         balancear();
52     }
53 };

```

```

55 int main() {
56     ios::sync_with_stdio(false);
57     cin.tie(NULL);
58
59     int n, k, q; cin >> n >> k >> q;
60
61     BigK big = BigK(k);
62
63     int arr[n] = {};
64
65     while (q--) {
66         int pos, num; cin >> pos >> num;
67         pos--;
68         big.rem(arr[pos]);
69         arr[pos] = num;
70         big.add(arr[pos]);
71
72         cout << big.gt.s << '\n';
73     }
74
75     return 0;
76 }

```

9 Data Structures

9.1 Sparse Table

```

1 // Description:
2 // Data structure to query for minimum and maximum
3
4 // Problem:
5 // https://cses.fi/problemset/task/1647/
6
7 // Complexity:
8 // Build O(n log n)
9 // Query O(1)
10
11 #include <bits/stdc++.h>
12
13 using namespace std;
14
15 const int MAX = 2e5+17;
16 const int INF = 1e9+17;
17
18 struct SparseTable {
19     int n;
20     vector<int> arr;
21     vector<vector<int>> st;
22     vector<int> log_2;
23
24     SparseTable(vector<int>& arr, int& n) : arr(arr), n
(n) {
25         build();
26     }
27
28     void build() {
29         log_2.resize(MAX + 1);
30
31         log_2[1] = 0;
32         for (int i = 2; i <= MAX; i++) {
33             log_2[i] = log_2[i/2] + 1;
34         }
35
36         int K = log_2[n + 1];
37
38         st.resize(MAX, vector<int>(K + 1));
39
40         for (int i = 0; i < MAX; i++) {
41             for (int j = 0; j < K + 1; j++) {
42                 st[i][j] = INF;
43             }
44         }
45     }
46 }

```



```

44     }
45
46     for (int i = 0; i < n; i++) {
47         st[i][0] = arr[i];
48     }
49
50     for (int j = 1; j <= K; j++) {
51         for (int i = 0; i + (1 << j) < MAX; i++) {
52             st[i][j] = min(st[i][j-1], st[i + (1 << (j - 1))][j - 1]);
53         }
54     }
55 }
56
57 int query(int l, int r) {
58     int j = log2[r - l + 1];
59     return min(st[l][j], st[r - (1 << j) + 1][j]);
60 }
61 };

```

9.2 Mergesort Tree Vector

```

1 // Description:
2 // In each node, the tree keeps a sorted list of
  elements in that range.
3 // It can be used to find how many elements are
  greater than x in a given range.
4 // It can also be used to find the position of an
  element if the list was sorted.
5 // query(i, j, k) - how many elements greater than k
  are in the range (i, j)
6
7 // Problem:
8 // https://www.spoj.com/problems/KQUERY
9
10 // Complexity:
11 // O(n log n) for build
12 // O(log ~ 2 n) for query
13
14 struct MergeSortTree {
15     vector<vector<int>> tree;
16     int n;
17
18     MergeSortTree(int n) : n(n) {
19         int sz = 1;
20         while (sz < n) sz *= 2;
21
22         tree.assign(2 * sz, vector<int>());
23     }
24
25     vector<int> merge(vector<int>& a, vector<int>& b) {
26         vector<int> res((int)a.size() + (int)b.size());
27         int it = 0, jt = 0, curr = 0;
28
29         while (it < (int)a.size() && jt < (int)b.size())
30         {
31             if (a[it] <= b[jt]) {
32                 res[curr++] = a[it++];
33             } else {
34                 res[curr++] = b[jt++];
35             }
36         }
37
38         while (it < (int)a.size()) {
39             res[curr++] = a[it++];
40         }
41
42         while (jt < (int)b.size()) {
43             res[curr++] = b[jt++];
44         }
45
46         return res;
47     }

```

```

47
48 void build(int pos, int ini, int fim, vector<int>&
  v) {
49     if (ini == fim) {
50         if (ini < (int)v.size()) {
51             tree[pos].pb(v[ini]);
52         }
53         return;
54     }
55
56     int mid = ini + (fim - ini) / 2;
57
58     build(2 * pos + 1, ini, mid, v);
59     build(2 * pos + 2, mid + 1, fim, v);
60
61     tree[pos] = merge(tree[2 * pos + 1], tree[2 * pos
  + 2]);
62 }
63
64 // how many elements greater than val in vector v
65 int search(vector<int>& v, int val) {
66     auto it = upper_bound(v.begin(), v.end(), val);
67     if (it == v.end()) return 0;
68     return (int)v.size() - (it - v.begin());
69 }
70
71 // how many elements greater than val in the range
  (p, q)
72 int query(int pos, int ini, int fim, int p, int q,
  int val) {
73     if (fim < p || ini > q) {
74         return 0;
75     }
76
77     if (ini >= p && fim <= q) {
78         return search(tree[pos], val);
79     }
80
81     int mid = ini + (fim - ini) / 2;
82     return query(2 * pos + 1, ini, mid, p, q, val) +
  query(2 * pos + 2, mid + 1, fim, p, q, val);
83 }
84
85 void build(vector<int>& v) {
86     build(0, 0, n - 1, v);
87 }
88
89 int query(int p, int q, int val) {
90     return query(0, 0, n - 1, p, q, val);
91 }
92 };

```

9.3 Ordered Set

```

1 // Description:
2 // insert(k) - add element k to the ordered set
3 // erase(k) - remove element k from the ordered set
4 // erase(it) - remove element it points to from the
  ordered set
5 // order_of_key(k) - returns number of elements
  strictly smaller than k
6 // find_by_order(n) - return an iterator pointing to
  the k-th element in the ordered set (counting
  from zero).
7
8 // Problem:
9 // https://cses.fi/problemset/task/2169/
10
11 // Complexity:
12 // O(log n) for all operations
13
14 // How to use:
15 // ordered_set<int> os;

```

```

16 // cout << os.order_of_key(1) << '\n';
17 // cout << os.find_by_order(1) << '\n';
18
19 // Notes
20 // The ordered set only contains different elements
21 // By using less_equal<T> instead of less<T> on using
   ordered_set declaration
22 // The ordered_set becomes an ordered_multiset
23 // So the set can contain elements that are equal
24
25 #include <ext/pb_ds/assoc_container.hpp>
26 #include <ext/pb_ds/tree_policy.hpp>
27
28 using namespace __gnu_pbds;
29 template <typename T>
30 using ordered_set = tree<T, null_type, less<T>,
   rb_tree_tag, tree_order_statistics_node_update>;
31
32 void Erase(ordered_set<int>& a, int x){
33     int r = a.order_of_key(x);
34     auto it = a.find_by_order(r);
35     a.erase(it);
36 }

```

9.4 Priority Queue

```

1 // Description:
2 // Keeps the largest (by default) element at the top
   of the queue
3
4 // Problem:
5 // https://cses.fi/problemset/task/1164/
6
7 // Complexity:
8 // O(log n) for push and pop
9 // O(1) for looking at the element at the top
10
11 // How to use:
12 // priority_queue<int> pq;
13 // pq.push(1);
14 // pq.top();
15 // pq.pop()
16
17 // Notes
18 // To use the priority queue keeping the smallest
   element at the top
19
20 priority_queue<int, vector<int>, greater<int>> pq;

```

9.5 Dsu

```

1 #include <bits/stdc++.h>
2
3 using namespace std;
4
5 const int MAX = 1e6+17;
6
7 struct DSU {
8     int n;
9     vector<int> link, sizes;
10
11     DSU(int n) {
12         this->n = n;
13         link.assign(n+1, 0);
14         sizes.assign(n+1, 1);
15
16         for (int i = 0; i <= n; i++)
17             link[i] = i;
18     }
19
20     int find(int x) {
21         while (x != link[x])

```

```

22             x = link[x];
23
24         return x;
25     }
26
27     bool same(int a, int b) {
28         return find(a) == find(b);
29     }
30
31     void unite(int a, int b) {
32         a = find(a);
33         b = find(b);
34
35         if (a == b) return;
36
37         if (sizes[a] < sizes[b])
38             swap(a, b);
39
40         sizes[a] += sizes[b];
41         link[b] = a;
42     }
43
44     int size(int x) {
45         return sizes[x];
46     }
47 };
48
49 int main() {
50     ios::sync_with_stdio(false);
51     cin.tie(NULL);
52
53     int cities, roads; cin >> cities >> roads;
54     vector<int> final_roads;
55     int ans = 0;
56     DSU dsu = DSU(cities);
57     for (int i = 0, a, b; i < roads; i++) {
58         cin >> a >> b;
59         dsu.unite(a, b);
60     }
61
62     for (int i = 2; i <= cities; i++) {
63         if (!dsu.same(1, i)) {
64             ans++;
65             final_roads.push_back(i);
66             dsu.unite(1, i);
67         }
68     }
69
70     cout << ans << '\n';
71     for (auto e : final_roads) {
72         cout << "1 " << e << '\n';
73     }
74
75 }

```

9.6 Two Sets

```

1 // Description
2 // The values are divided in two multisets so that
   one of them contain all values that are
3 // smaller than the median and the other one contains
   all values that are greater or equal to the
   median.
4
5 // Problem:
6 // https://atcoder.jp/contests/abc306/tasks/abc306_e
7 // Problem I - Maratona Feminina de çãProgramao da
   Unicamp 2023
8 // https://codeforces.com/group/WYIydkIPyE/contest
   /450037/attachments
9
10 // Complexity:
11 // Add and remove elements - O(log n)

```

```

12 // Return sum of biggest or smallest set or return
    the median - 0(1)
13
14 using ll = long long;
15
16 struct TwoSets {
17     multiset<int> small;
18     multiset<int> big;
19     ll sums = 0;
20     ll sumb = 0;
21     int n = 0;
22
23     int size_small() {
24         return small.size();
25     }
26
27     int size_big() {
28         return big.size();
29     }
30
31     void balance() {
32         while (size_small() > n / 2) {
33             int v = *small.rbegin();
34             small.erase(prev(small.end()));
35             big.insert(v);
36             sums -= v;
37             sumb += v;
38         }
39         while (size_big() > n - n / 2) {
40             int v = *big.begin();
41             big.erase(big.begin());
42             small.insert(v);
43             sumb -= v;
44             sums += v;
45         }
46     }
47
48     void add(int x) {
49         n++;
50         small.insert(x);
51         sums += x;
52         while (!small.empty() && *small.rbegin() > *big.
begin()) {
53             int v = *small.rbegin();
54             small.erase(prev(small.end()));
55             big.insert(v);
56             sums -= v;
57             sumb += v;
58         }
59         balance();
60     }
61
62     bool rem(int x) {
63         n--;
64         auto it1 = small.find(x);
65         auto it2 = big.find(x);
66         bool flag = false;
67         if (it1 != small.end()) {
68             sums -= *it1;
69             small.erase(it1);
70             flag = true;
71         } else if (it2 != big.end()) {
72             sumb -= *it2;
73             big.erase(it2);
74             flag = true;
75         }
76         balance();
77         return flag;
78     }
79
80     ll sum_small() {
81         return sums;
82     }

```

```

83
84     ll sum_big() {
85         return sumb;
86     }
87
88     int median() {
89         return *big.begin();
90     }
91 };

```

9.7 Psum2d

```

1 // Description:
2 // Queries the sum of a rectangle that goes from grid
    [from_row][from_col] to grid[to_row][to_col]
3
4 // Problem:
5 // https://cses.fi/problemset/task/1652/
6
7 // Complexity:
8 // O(n) build
9 // O(1) query
10
11 for (int i = 1; i <= n; i++) {
12     for (int j = 1; j <= n; j++) {
13         psum[i][j] = grid[i][j] + psum[i - 1][j] + psum[i
][j - 1] - psum[i - 1][j - 1];
14     }
15 }
16
17 while (q--) {
18     int from_row, to_row, from_col, to_col;
19     cin >> from_row >> from_col >> to_row >> to_col;
20     cout << psum[to_row][to_col] - psum[from_row - 1][
to_col] -
21     psum[to_row][from_col - 1] + psum[from_row - 1][
from_col - 1] << '\n';
22 }

```

9.8 Dynamic Implicit Sparse

```

1 // Description:
2 // Indexed at one
3
4 // When the indexes of the nodes are too big to be
    stored in an array
5 // and the queries need to be answered online so we
    can't sort the nodes and compress them
6 // we create nodes only when they are needed so there
    'll be (Q*log(MAX)) nodes
7 // where Q is the number of queries and MAX is the
    maximum index a node can assume
8
9 // Query - get sum of elements from range (l, r)
    inclusive
10 // Update - update element at position id to a value
    val
11
12 // Problem:
13 // https://cses.fi/problemset/task/1648
14
15 // Complexity:
16 // O(log n) for both query and update
17
18 // How to use:
19 // MAX is the maximum index a node can assume
20
21 // Segtree seg = Segtree(MAX);
22
23 typedef long long ftype;
24
25 const int MAX = 1e9+17;

```

```

26
27 struct Segtree {
28     vector<ftype> seg, d, e;
29     const ftype NEUTRAL = 0;
30     int n;
31
32     Segtree(int n) {
33         this->n = n;
34         create();
35         create();
36     }
37
38     ftype f(ftype a, ftype b) {
39         return a + b;
40     }
41
42     ftype create() {
43         seg.push_back(0);
44         e.push_back(0);
45         d.push_back(0);
46         return seg.size() - 1;
47     }
48
49     ftype query(int pos, int ini, int fim, int p, int
50 q) {
51     if (q < ini || p > fim) return NEUTRAL;
52     if (pos == 0) return 0;
53     if (p <= ini && fim <= q) return seg[pos];
54     int m = (ini + fim) >> 1;
55     return f(query(e[pos], ini, m, p, q), query(d
56 [pos], m + 1, fim, p, q));
57
58 void update(int pos, int ini, int fim, int id,
59 int val) {
60     if (ini > id || fim < id) {
61         return;
62     }
63
64     if (ini == fim) {
65         seg[pos] = val;
66         return;
67     }
68
69     int m = (ini + fim) >> 1;
70
71     if (id <= m) {
72         if (e[pos] == 0) e[pos] = create();
73         update(e[pos], ini, m, id, val);
74     } else {
75         if (d[pos] == 0) d[pos] = create();
76         update(d[pos], m + 1, fim, id, val);
77     }
78
79     seg[pos] = f(seg[e[pos]], seg[d[pos]]);
80
81     ftype query(int p, int q) {
82         return query(1, 1, n, p, q);
83     }
84
85     void update(int id, int val) {
86         update(1, 1, n, id, val);
87     }
88 };

```

9.9 Segtree2d

```

1 // Description:
2 // Indexed at zero
3 // Given a N x M grid, where i represents the row and
  j the column, perform the following operations

```

```

4 // update(i, j) - update the value of grid[i][j]
5 // query(i1, j1, i2, j2) - return the sum of values
  inside the rectangle
6 // defined by grid[i1][j1] and grid[i2][j2] inclusive
7
8 // Problem:
9 // https://cses.fi/problemset/task/1739/
10
11 // Complexity:
12 // Time complexity:
13 // O(log N * log M) for both query and update
14 // O(N * M) for build
15 // Memory complexity:
16 // 4 * M * N
17
18 // How to use:
19 // Segtree2D seg = Segtree2D(n, m);
20 // vector<vector<int>> v(n, vector<int>(m));
21 // seg.build(v);
22
23 struct Segtree2D {
24     const int MAXN = 1025;
25     const int NEUTRAL = 0;
26     int N, M;
27
28     vector<vector<int>> seg;
29
30     Segtree2D(int N, int M) {
31         this->N = N;
32         this->M = M;
33         seg.assign(4*MAXN, vector<int>(4*MAXN,
34 NEUTRAL));
35     }
36
37     int f(int a, int b) {
38         return max(a, b);
39     }
40
41     void buildY(int noX, int lX, int rX, int noY, int
42 lY, int rY, vector<vector<int>> &v){
43         if(lY == rY){
44             if(lX == rX){
45                 seg[noX][noY] = v[rX][rY];
46             }else{
47                 seg[noX][noY] = f(seg[2*noX+1][noY],
48 seg[2*noX+2][noY]);
49             }
50         }else{
51             int m = (lY+rY)/2;
52
53             buildY(noX, lX, rX, 2*noY+1, lY, m, v);
54             buildY(noX, lX, rX, 2*noY+2, m+1, rY, v);
55
56             seg[noX][noY] = f(seg[noX][2*noY+1], seg[
57 noX][2*noY+2]);
58         }
59     }
60
61     void buildX(int noX, int lX, int rX, vector<
62 vector<int>> &v){
63         if(lX != rX){
64             int m = (lX+rX)/2;
65
66             buildX(2*noX+1, lX, m, v);
67             buildX(2*noX+2, m+1, rX, v);
68         }
69
70         buildY(noX, lX, rX, 0, 0, M - 1, v);
71     }
72
73     void updateY(int noX, int lX, int rX, int noY,
74 int lY, int rY, int y){
75         if(lY == rY){

```

```

70         if(lX == rX){
71             seg[noX][noY] = !seg[noX][noY];
72         }else{
73             seg[noX][noY] = seg[2*noX+1][noY] +
seg[2*noX+2][noY];
74         }
75     }else{
76         int m = (lY+rY)/2;
77
78         if(y <= m){
79             updateY(noX, lX, rX, 2*noY+1,lY, m, y
);
80         }else if(m < y){
81             updateY(noX, lX, rX, 2*noY+2, m+1, rY
, y);
82         }
83
84         seg[noX][noY] = seg[noX][2*noY+1] + seg[
noX][2*noY+2];
85     }
86 }
87
88 void updateX(int noX, int lX, int rX, int x, int
y){
89     int m = (lX+rX)/2;
90
91     if(lX != rX){
92         if(x <= m){
93             updateX(2*noX+1, lX, m, x, y);
94         }else if(m < x){
95             updateX(2*noX+2, m+1, rX, x, y);
96         }
97     }
98
99     updateY(noX, lX, rX, 0, 0, M - 1, y);
100 }
101
102 int queryY(int noX, int noY, int lY, int rY, int
aY, int bY){
103     if(aY <= lY && rY <= bY) return seg[noX][noY
];
104
105     int m = (lY+rY)/2;
106
107     if(bY <= m) return queryY(noX, 2*noY+1, lY, m
, aY, bY);
108     if(m < aY) return queryY(noX, 2*noY+2, m+1,
rY, aY, bY);
109
110     return f(queryY(noX, 2*noY+1, lY, m, aY, bY),
queryY(noX, 2*noY+2, m+1, rY, aY, bY));
111 }
112
113 int queryX(int noX, int lX, int rX, int aX, int
bX, int aY, int bY){
114     if(aX <= lX && rX <= bX) return queryY(noX,
0, 0, M - 1, aY, bY);
115
116     int m = (lX+rX)/2;
117
118     if(bX <= m) return queryX(2*noX+1, lX, m, aX,
bX, aY, bY);
119     if(m < aX) return queryX(2*noX+2, m+1, rX, aX
, bX, aY, bY);
120
121     return f(queryX(2*noX+1, lX, m, aX, bX, aY,
bY), queryX(2*noX+2, m+1, rX, aX, bX, aY, bY));
122 }
123
124 void build(vector<vector<int>> &v) {
125     buildX(0, 0, N - 1, v);
126 }
127
128 int query(int aX, int aY, int bX, int bY) {
129     return queryX(0, 0, N - 1, aX, bX, aY, bY);
130 }
131
132 void update(int x, int y) {
133     updateX(0, 0, N - 1, x, y);
134 }
135 };

```

9.10 Minimum And Amount

```

1 // Description:
2 // Query - get minimum element in a range (l, r)
   inclusive
3 // and also the number of times it appears in that
   range
4 // Update - update element at position id to a value
   val
5
6 // Problem:
7 // https://codeforces.com/edu/course/2/lesson/4/1/
   practice/contest/273169/problem/C
8
9 // Complexity:
10 // O(log n) for both query and update
11
12 // How to use:
13 // Segtree seg = Segtree(n);
14 // seg.build(v);
15
16 #define pii pair<int, int>
17 #define mp make_pair
18 #define ff first
19 #define ss second
20
21 const int INF = 1e9+17;
22
23 typedef pii ftype;
24
25 struct Segtree {
26     vector<ftype> seg;
27     int n;
28     const ftype NEUTRAL = mp(INF, 0);
29
30     Segtree(int n) {
31         int sz = 1;
32         while (sz < n) sz *= 2;
33         this->n = sz;
34
35         seg.assign(2*sz, NEUTRAL);
36     }
37
38     ftype f(ftype a, ftype b) {
39         if (a.ff < b.ff) return a;
40         if (b.ff < a.ff) return b;
41
42         return mp(a.ff, a.ss + b.ss);
43     }
44
45     ftype query(int pos, int ini, int fim, int p, int
q) {
46         if (ini >= p && fim <= q) {
47             return seg[pos];
48         }
49
50         if (q < ini || p > fim) {
51             return NEUTRAL;
52         }
53
54         int e = 2*pos + 1;
55         int d = 2*pos + 2;
56         int m = ini + (fim - ini) / 2;
57

```

```

58     return f(query(e, ini, m, p, q), query(d, m + 1, fim, p, q));
59 }
60
61 void update(int pos, int ini, int fim, int id,
62 int val) {
63     if (ini > id || fim < id) {
64         return;
65     }
66
67     if (ini == id && fim == id) {
68         seg[pos] = mp(val, 1);
69
70         return;
71     }
72
73     int e = 2*pos + 1;
74     int d = 2*pos + 2;
75     int m = ini + (fim - ini) / 2;
76
77     update(e, ini, m, id, val);
78     update(d, m + 1, fim, id, val);
79
80     seg[pos] = f(seg[e], seg[d]);
81 }
82
83 void build(int pos, int ini, int fim, vector<int>
84 &v) {
85     if (ini == fim) {
86         if (ini < (int)v.size()) {
87             seg[pos] = mp(v[ini], 1);
88         }
89         return;
90     }
91
92     int e = 2*pos + 1;
93     int d = 2*pos + 2;
94     int m = ini + (fim - ini) / 2;
95
96     build(e, ini, m, v);
97     build(d, m + 1, fim, v);
98
99     seg[pos] = f(seg[e], seg[d]);
100 }
101
102 ftype query(int p, int q) {
103     return query(0, 0, n - 1, p, q);
104 }
105
106 void update(int id, int val) {
107     update(0, 0, n - 1, id, val);
108 }
109
110 void build(vector<int> &v) {
111     build(0, 0, n - 1, v);
112 }
113
114 void debug() {
115     for (auto e : seg) {
116         cout << e.ff << ' ' << e.ss << '\n';
117     }
118     cout << '\n';
119 }
120 };

```

9.11 Lazy Addition To Segment

```

1 // Description:
2 // Query - get sum of elements from range (l, r)
3 // inclusive
4 // Update - add a value val to elementos from range (l, r) inclusive

```

```

5 // Problem:
6 // https://codeforces.com/edu/course/2/lesson/5/1/
7 // practice/contest/279634/problem/A
8 // Complexity:
9 // O(log n) for both query and update
10
11 // How to use:
12 // Segtree seg = Segtree(n);
13 // seg.build(v);
14
15 // Notes
16 // Change neutral element and f function to perform a
17 // different operation
18
19 const long long INF = 1e18+10;
20
21 typedef long long ftype;
22
23 struct Segtree {
24     vector<ftype> seg;
25     vector<ftype> lazy;
26     int n;
27     const ftype NEUTRAL = 0;
28     const ftype NEUTRAL_LAZY = -1; // change to -INF
29     if there are negative numbers
30
31     Segtree(int n) {
32         int sz = 1;
33         while (sz < n) sz *= 2;
34         this->n = sz;
35
36         seg.assign(2*sz, NEUTRAL);
37         lazy.assign(2*sz, NEUTRAL_LAZY);
38     }
39
40     ftype apply_lazy(ftype a, ftype b, int len) {
41         if (b == NEUTRAL_LAZY) return a;
42         if (a == NEUTRAL_LAZY) return b * len;
43         else return a + b * len;
44     }
45
46     void propagate(int pos, int ini, int fim) {
47         if (ini == fim) {
48             return;
49         }
50
51         int e = 2*pos + 1;
52         int d = 2*pos + 2;
53         int m = ini + (fim - ini) / 2;
54
55         lazy[e] = apply_lazy(lazy[e], lazy[pos], 1);
56         lazy[d] = apply_lazy(lazy[d], lazy[pos], 1);
57
58         seg[e] = apply_lazy(seg[e], lazy[pos], m -
59 ini + 1);
60         seg[d] = apply_lazy(seg[d], lazy[pos], fim -
61 m);
62
63         lazy[pos] = NEUTRAL_LAZY;
64     }
65
66     ftype f(ftype a, ftype b) {
67         return a + b;
68     }
69
70     ftype query(int pos, int ini, int fim, int p, int
71 q) {
72     propagate(pos, ini, fim);
73
74     if (ini >= p && fim <= q) {
75         return seg[pos];
76     }

```

```

72         if (q < ini || p > fim) {
73             return NEUTRAL;
74         }
75     }
76
77     int e = 2*pos + 1;
78     int d = 2*pos + 2;
79     int m = ini + (fim - ini) / 2;
80
81     return f(query(e, ini, m, p, q), query(d, m + 1, fim, p, q));
82 }
83
84 void update(int pos, int ini, int fim, int p, int q, int val) {
85     propagate(pos, ini, fim);
86
87     if (ini > q || fim < p) {
88         return;
89     }
90
91     if (ini >= p && fim <= q) {
92         lazy[pos] = apply_lazy(lazy[pos], val, 1);
93     }
94     seg[pos] = apply_lazy(seg[pos], val, fim - ini + 1);
95     return;
96 }
97
98 int e = 2*pos + 1;
99 int d = 2*pos + 2;
100 int m = ini + (fim - ini) / 2;
101
102 update(e, ini, m, p, q, val);
103 update(d, m + 1, fim, p, q, val);
104
105 seg[pos] = f(seg[e], seg[d]);
106 }
107
108 void build(int pos, int ini, int fim, vector<int> &v) {
109     if (ini == fim) {
110         if (ini < (int)v.size()) {
111             seg[pos] = v[ini];
112         }
113         return;
114     }
115
116     int e = 2*pos + 1;
117     int d = 2*pos + 2;
118     int m = ini + (fim - ini) / 2;
119
120     build(e, ini, m, v);
121     build(d, m + 1, fim, v);
122
123     seg[pos] = f(seg[e], seg[d]);
124 }
125
126 ftype query(int p, int q) {
127     return query(0, 0, n - 1, p, q);
128 }
129
130 void update(int p, int q, int val) {
131     update(0, 0, n - 1, p, q, val);
132 }
133
134 void build(vector<int> &v) {
135     build(0, 0, n - 1, v);
136 }
137
138 void debug() {
139     for (auto e : seg) {

```

```

140         cout << e << ' ';
141     }
142     cout << '\n';
143     for (auto e : lazy) {
144         cout << e << ' ';
145     }
146     cout << '\n';
147     cout << '\n';
148 }
149 }

```

9.12 Segment With Maximum Sum

```

1 // Description:
2 // Query - get sum of segment that is maximum among
3 // all segments
4 // E.g
5 // Array: 5 -4 4 3 -5
6 // Maximum segment sum: 8 because 5 + (-4) + 4 + 3 = 8
7 // Update - update element at position id to a value val
8 // Problem:
9 // https://codeforces.com/edu/course/2/lesson/4/2/practice/contest/273278/problem/A
10
11 // Complexity:
12 // O(log n) for both query and update
13
14 // How to use:
15 // Segtree seg = Segtree(n);
16 // seg.build(v);
17
18 // Notes
19 // The maximum segment sum can be a negative number
20 // In that case, taking zero elements is the best choice
21 // So we need to take the maximum between 0 and the query
22 // max(0LL, seg.query(0, n).max_seg)
23
24 using ll = long long;
25
26 typedef ll ftype_node;
27
28 struct Node {
29     ftype_node max_seg;
30     ftype_node pref;
31     ftype_node suf;
32     ftype_node sum;
33
34     Node(ftype_node max_seg, ftype_node pref, ftype_node suf, ftype_node sum) : max_seg(max_seg), pref(pref), suf(suf), sum(sum) {};
35 };
36
37 typedef Node ftype;
38
39 struct Segtree {
40     vector<ftype> seg;
41     int n;
42     const ftype NEUTRAL = Node(0, 0, 0, 0);
43
44     Segtree(int n) {
45         int sz = 1;
46         // potencia de dois mais proxima
47         while (sz < n) sz *= 2;
48         this->n = sz;
49
50         // numero de nos da seg
51         seg.assign(2*sz, NEUTRAL);
52     }

```

```

53 ftype f(ftype a, ftype b) {
54     ftype_node max_seg = max({a.max_seg, b.
55     max_seg, a.suf + b.pref});
56     ftype_node pref = max(a.pref, a.sum + b.pref);
57     ftype_node suf = max(b.suf, b.sum + a.suf);
58     ftype_node sum = a.sum + b.sum;
59
60     return Node(max_seg, pref, suf, sum);
61 }
62
63 ftype query(int pos, int ini, int fim, int p, int
64 q) {
65     if (ini >= p && fim <= q) {
66         return seg[pos];
67     }
68     if (q < ini || p > fim) {
69         return NEUTRAL;
70     }
71
72     int e = 2*pos + 1;
73     int d = 2*pos + 2;
74     int m = ini + (fim - ini) / 2;
75
76     return f(query(e, ini, m, p, q), query(d, m +
77     1, fim, p, q));
78 }
79
80 void update(int pos, int ini, int fim, int id,
81 int val) {
82     if (ini > id || fim < id) {
83         return;
84     }
85     if (ini == id && fim == id) {
86         seg[pos] = Node(val, val, val, val);
87         return;
88     }
89
90     int e = 2*pos + 1;
91     int d = 2*pos + 2;
92     int m = ini + (fim - ini) / 2;
93
94     update(e, ini, m, id, val);
95     update(d, m + 1, fim, id, val);
96
97     seg[pos] = f(seg[e], seg[d]);
98 }
99
100 void build(int pos, int ini, int fim, vector<int>
101 &v) {
102     if (ini == fim) {
103         // se a posição existir no array original
104         // seg tamanho potencia de dois
105         if (ini < (int)v.size()) {
106             seg[pos] = Node(v[ini], v[ini], v[ini]
107             ], v[ini]);
108         }
109         return;
110     }
111
112     int e = 2*pos + 1;
113     int d = 2*pos + 2;
114     int m = ini + (fim - ini) / 2;
115
116     build(e, ini, m, v);
117     build(d, m + 1, fim, v);
118
119     seg[pos] = f(seg[e], seg[d]);
120 }
121
122 ftype query(int p, int q) {
123     return query(0, 0, n - 1, p, q);
124 }
125
126 void update(int id, int val) {
127     update(0, 0, n - 1, id, val);
128 }
129
130 void build(vector<int> &v) {
131     build(0, 0, n - 1, v);
132 }
133
134 void debug() {
135     for (auto e : seg) {
136         cout << e.max_seg << ' ' << e.pref << ' '
137         << e.suf << ' ' << e.sum << '\n';
138     }
139     cout << '\n';
140 }
141
142 };

```

9.13 Range Query Point Update

```

1 // Description:
2 // Indexed at zero
3 // Query - get sum of elements from range (l, r)
4 // inclusive
5 // Update - update element at position id to a value
6 // val
7
8 // Problem:
9 // https://codeforces.com/edu/course/2/lesson/4/1/
10 // practice/contest/273169/problem/B
11
12 // Complexity:
13 // O(log n) for both query and update
14
15 // How to use:
16 // Segtree seg = Segtree(n);
17 // seg.build(v);
18
19 // Notes
20 // Change neutral element and f function to perform a
21 // different operation
22
23 // If you want to change the operations to point
24 // query and range update
25 // Use the same segtree, but perform the following
26 // operations
27 // Query - seg.query(0, id);
28 // Update - seg.update(l, v); seg.update(r + 1, -v);
29
30 typedef long long ftype;
31
32 struct Segtree {
33     vector<ftype> seg;
34     int n;
35     const ftype NEUTRAL = 0;
36
37     Segtree(int n) {
38         int sz = 1;
39         while (sz < n) sz *= 2;
40         this->n = sz;
41         seg.assign(2*sz, NEUTRAL);
42     }
43
44     ftype f(ftype a, ftype b) {
45         return a + b;
46     }
47 }

```



```

43 ftype query(int pos, int ini, int fim, int p, int q) {
44     if (ini >= p && fim <= q) {
45         return seg[pos];
46     }
47
48     if (q < ini || p > fim) {
49         return NEUTRAL;
50     }
51
52     int e = 2*pos + 1;
53     int d = 2*pos + 2;
54     int m = ini + (fim - ini) / 2;
55
56     return f(query(e, ini, m, p, q), query(d, m +
57 1, fim, p, q));
58
59 void update(int pos, int ini, int fim, int id,
60 int val) {
61     if (ini > id || fim < id) {
62         return;
63     }
64
65     if (ini == id && fim == id) {
66         seg[pos] = val;
67
68         return;
69     }
70
71     int e = 2*pos + 1;
72     int d = 2*pos + 2;
73     int m = ini + (fim - ini) / 2;
74
75     update(e, ini, m, id, val);
76     update(d, m + 1, fim, id, val);
77
78     seg[pos] = f(seg[e], seg[d]);
79
80 void build(int pos, int ini, int fim, vector<int>
81 &v) {
82     if (ini == fim) {
83         if (ini < (int)v.size()) {
84             seg[pos] = v[ini];
85         }
86         return;
87     }
88
89     int e = 2*pos + 1;
90     int d = 2*pos + 2;
91     int m = ini + (fim - ini) / 2;
92
93     build(e, ini, m, v);
94     build(d, m + 1, fim, v);
95
96     seg[pos] = f(seg[e], seg[d]);
97
98 ftype query(int p, int q) {
99     return query(0, 0, n - 1, p, q);
100 }
101
102 void update(int id, int val) {
103     update(0, 0, n - 1, id, val);
104 }
105
106 void build(vector<int> &v) {
107     build(0, 0, n - 1, v);
108 }
109
110 void debug() {
111     for (auto e : seg) {

```

```

112         cout << e << ' ';
113     }
114     cout << '\n';
115 }
116 };

```

9.14 Lazy Assignment To Segment

```

1 const long long INF = 1e18+10;
2
3 typedef long long ftype;
4
5 struct Segtree {
6     vector<ftype> seg;
7     vector<ftype> lazy;
8     int n;
9     const ftype NEUTRAL = 0;
10    const ftype NEUTRAL_LAZY = -1; // Change to -INF
    if there are negative numbers
11
12    Segtree(int n) {
13        int sz = 1;
14        // potencia de dois mais proxima
15        while (sz < n) sz *= 2;
16        this->n = sz;
17
18        // numero de nos da seg
19        seg.assign(2*sz, NEUTRAL);
20        lazy.assign(2*sz, NEUTRAL_LAZY);
21    }
22
23    ftype apply_lazy(ftype a, ftype b, int len) {
24        if (b == NEUTRAL_LAZY) return a;
25        if (a == NEUTRAL_LAZY) return b * len;
26        else return b * len;
27    }
28
29    void propagate(int pos, int ini, int fim) {
30        if (ini == fim) {
31            return;
32        }
33
34        int e = 2*pos + 1;
35        int d = 2*pos + 2;
36        int m = ini + (fim - ini) / 2;
37
38        lazy[e] = apply_lazy(lazy[e], lazy[pos], 1);
39        lazy[d] = apply_lazy(lazy[d], lazy[pos], 1);
40
41        seg[e] = apply_lazy(seg[e], lazy[pos], m -
42 ini + 1);
43        seg[d] = apply_lazy(seg[d], lazy[pos], fim -
44 m);
45
46        lazy[pos] = NEUTRAL_LAZY;
47    }
48
49    ftype f(ftype a, ftype b) {
50        return a + b;
51    }
52
53    ftype query(int pos, int ini, int fim, int p, int
54 q) {
55        propagate(pos, ini, fim);
56
57        if (ini >= p && fim <= q) {
58            return seg[pos];
59        }
60
61        if (q < ini || p > fim) {
62            return NEUTRAL;
63        }

```

```

62     int e = 2*pos + 1;
63     int d = 2*pos + 2;
64     int m = ini + (fim - ini) / 2;
65
66     return f(query(e, ini, m, p, q), query(d, m +
67 1, fim, p, q));
68 }

```

```

69 void update(int pos, int ini, int fim, int p, int
70 q, int val) {
71     propagate(pos, ini, fim);
72
73     if (ini > q || fim < p) {
74         return;
75     }
76
77     if (ini >= p && fim <= q) {
78         lazy[pos] = apply_lazy(lazy[pos], val, 1)
79 ;
80         seg[pos] = apply_lazy(seg[pos], val, fim
81 - ini + 1);
82
83         return;
84     }
85
86     int e = 2*pos + 1;
87     int d = 2*pos + 2;
88     int m = ini + (fim - ini) / 2;
89
90     update(e, ini, m, p, q, val);
91     update(d, m + 1, fim, p, q, val);
92
93     seg[pos] = f(seg[e], seg[d]);
94 }
95
96 void build(int pos, int ini, int fim, vector<int>
97 &v) {
98     if (ini == fim) {
99         // se a posição existir no array original
100         // seg tamanho potencia de dois
101         if (ini < (int)v.size()) {
102             seg[pos] = v[ini];
103         }
104         return;
105     }
106
107     int e = 2*pos + 1;
108     int d = 2*pos + 2;
109     int m = ini + (fim - ini) / 2;
110
111     build(e, ini, m, v);
112     build(d, m + 1, fim, v);
113
114     seg[pos] = f(seg[e], seg[d]);
115 }
116
117 ftype query(int p, int q) {
118     return query(0, 0, n - 1, p, q);
119 }
120
121 void update(int p, int q, int val) {
122     update(0, 0, n - 1, p, q, val);
123 }
124
125 void build(vector<int> &v) {
126     build(0, 0, n - 1, v);
127 }
128
129 void debug() {
130     for (auto e : seg) {
131         cout << e << ' ';
132     }
133     cout << '\n';

```

```

130     for (auto e : lazy) {
131         cout << e << ' ';
132     }
133     cout << '\n';
134     cout << '\n';
135 }
136 };

```

9.15 Lazy Dynamic Implicit Sparse

```

1 // Description:
2 // Indexed at one
3
4 // When the indexes of the nodes are too big to be
5 // stored in an array
6 // and the queries need to be answered online so we
7 // can't sort the nodes and compress them
8 // we create nodes only when they are needed so there
9 // 'll be (Q*log(MAX)) nodes
10 // where Q is the number of queries and MAX is the
11 // maximum index a node can assume
12
13 // Query - get sum of elements from range (l, r)
14 // inclusive
15 // Update - update element at position id to a value
16 // val
17
18 // Problem:
19 // https://oj.uz/problem/view/IZh012_apple
20
21 // Complexity:
22 // O(log n) for both query and update
23
24 // How to use:
25 // MAX is the maximum index a node can assume
26 // Create a default null node
27 // Create a node to be the root of the segtree
28
29 // Segtree seg = Segtree(MAX);
30
31 const int MAX = 1e9+10;
32 const long long INF = 1e18+10;
33
34 typedef long long ftype;
35
36 struct Segtree {
37     vector<ftype> seg, d, e, lazy;
38     const ftype NEUTRAL = 0;
39     const ftype NEUTRAL_LAZY = -1; // change to -INF
40     if the elements can be negative
41     int n;
42
43     Segtree(int n) {
44         this->n = n;
45         create();
46         create();
47     }
48
49     ftype apply_lazy(ftype a, ftype b, int len) {
50         if (b == NEUTRAL_LAZY) return a;
51         else return b * len; // change to a + b * len
52         to add to an element instead of updating it
53     }
54
55     void propagate(int pos, int ini, int fim) {
56         if (seg[pos] == 0) return;
57
58         if (ini == fim) {
59             return;
60         }
61
62         int m = (ini + fim) >> 1;

```

```

56     if (e[pos] == 0) e[pos] = create();
57     if (d[pos] == 0) d[pos] = create();
58
59     lazy[e[pos]] = apply_lazy(lazy[e[pos]], lazy[
pos], 1);
60     lazy[d[pos]] = apply_lazy(lazy[d[pos]], lazy[
pos], 1);
61
62     seg[e[pos]] = apply_lazy(seg[e[pos]], lazy[
pos], m - ini + 1);
63     seg[d[pos]] = apply_lazy(seg[d[pos]], lazy[
pos], fim - m);
64
65     lazy[pos] = NEUTRAL_LAZY;
66 }
67
68 ftype f(ftype a, ftype b) {
69     return a + b;
70 }
71
72 ftype create() {
73     seg.push_back(0);
74     e.push_back(0);
75     d.push_back(0);
76     lazy.push_back(-1);
77     return seg.size() - 1;
78 }
79
80 ftype query(int pos, int ini, int fim, int p, int
q) {
81     propagate(pos, ini, fim);
82     if (q < ini || p > fim) return NEUTRAL;
83     if (pos == 0) return 0;
84     if (p <= ini && fim <= q) return seg[pos];
85     int m = (ini + fim) >> 1;
86     return f(query(e[pos], ini, m, p, q), query(d
[pos], m + 1, fim, p, q));
87 }
88
89 void update(int pos, int ini, int fim, int p, int
q, int val) {
90     propagate(pos, ini, fim);
91     if (ini > q || fim < p) {
92         return;
93     }
94
95     if (ini >= p && fim <= q) {
96         lazy[pos] = apply_lazy(lazy[pos], val, 1)
;
97         seg[pos] = apply_lazy(seg[pos], val, fim
- ini + 1);
98
99         return;
100     }
101
102     int m = (ini + fim) >> 1;
103
104     if (e[pos] == 0) e[pos] = create();
105     update(e[pos], ini, m, p, q, val);
106
107     if (d[pos] == 0) d[pos] = create();
108     update(d[pos], m + 1, fim, p, q, val);
109
110     seg[pos] = f(seg[e[pos]], seg[d[pos]]);
111 }
112
113 ftype query(int p, int q) {
114     return query(1, 1, n, p, q);
115 }
116
117 void update(int p, int q, int val) {
118     update(1, 1, n, p, q, val);
119 }

```

```
120 };
```

9.16 Persistent

```

1 // Description:
2 // Persistent segtree allows for you to save the
different versions of the segtree between each
update
3 // Indexed at one
4 // Query - get sum of elements from range (l, r)
inclusive
5 // Update - update element at position id to a value
val
6
7 // Problem:
8 // https://cses.fi/problemset/task/1737/
9
10 // Complexity:
11 // O(log n) for both query and update
12
13 // How to use:
14 // vector<int> raiz(MAX); // vector to store the
roots of each version
15 // Segtree seg = Segtree(INF);
16 // raiz[0] = seg.create(); // null node
17 // curr = 1; // keep track of the last version
18
19 // raiz[k] = seg.update(raiz[k], idx, val); //
updating version k
20 // seg.query(raiz[k], l, r) // querying version k
21 // raiz[++curr] = raiz[k]; // create a new version
based on version k
22
23 const int MAX = 2e5+17;
24 const int INF = 1e9+17;
25
26 typedef long long ftype;
27
28 struct Segtree {
29     vector<ftype> seg, d, e;
30     const ftype NEUTRAL = 0;
31     int n;
32
33     Segtree(int n) {
34         this->n = n;
35     }
36
37     ftype f(ftype a, ftype b) {
38         return a + b;
39     }
40
41     ftype create() {
42         seg.push_back(0);
43         e.push_back(0);
44         d.push_back(0);
45         return seg.size() - 1;
46     }
47
48     ftype query(int pos, int ini, int fim, int p, int
q) {
49         if (q < ini || p > fim) return NEUTRAL;
50         if (pos == 0) return 0;
51         if (p <= ini && fim <= q) return seg[pos];
52         int m = (ini + fim) >> 1;
53         return f(query(e[pos], ini, m, p, q), query(d
[pos], m + 1, fim, p, q));
54     }
55
56     int update(int pos, int ini, int fim, int id, int
val) {
57         int novo = create();
58
59         seg[novo] = seg[pos];

```

```

60     e[novo] = e[pos];
61     d[novo] = d[pos];
62
63     if (ini == fim) {
64         seg[novo] = val;
65         return novo;
66     }
67
68     int m = (ini + fim) >> 1;
69
70     if (id <= m) e[novo] = update(e[novo], ini, m
71 , id, val);
72     else d[novo] = update(d[novo], m + 1, fim, id
73 , val);
74
75     seg[novo] = f(seg[e[novo]], seg[d[novo]]);
76
77     return novo;
78
79 ftype query(int pos, int p, int q) {
80     return query(pos, 1, n, p, q);
81 }
82
83 int update(int pos, int id, int val) {
84     return update(pos, 1, n, id, val);
85 };

```

9.17 Sparse Table2d

```

1 // Description
2 // Minimum queries in a 2D grid
3
4 // Problem:
5 // https://codeforces.com/group/YgJmumGtHD/contest
6 // 103794/problem/D
7
8 // Complexity:
9 // Build O(N * M * log(N) * log(M))
10 // Query O(1)
11 // Memory Complexity: O(N * M * log(N) * log(M))
12
13 const int MAX = 410;
14
15 struct SparseTable2D {
16     vector<vector<int>>> matrix;
17     vector<vector<vector<vector<int>>>> table;
18     int n, m;
19
20     SparseTable2D(vector<vector<int>>& matrix, int n,
21 int m) : matrix(matrix), n(n), m(m) {
22         table.resize(MAX, vector<vector<vector<int>>>(MAX
23 , vector<vector<int>>>(log2(MAX) + 1, vector<int>>(
24 log2(MAX) + 1)))));
25         build();
26     }
27
28     int f(int a, int b) {
29         return max(a, b);
30     }
31
32     void build() {
33         for (int i = 0; i < n; i++) {
34             for (int j = 0; j < m; j++) {

```

```

31         table[i][j][0][0] = matrix[i][j];
32     }
33 }
34
35 for (int k = 1; k <= (int)(log2(n)); k++) {
36     for (int i = 0; i + (1 << k) - 1 < n; i++) {
37         for (int j = 0; j + (1 << k) - 1 < m; j++) {
38             table[i][j][k][0] = f(
39                 table[i][j][k - 1][0],
40                 table[i + (1 << (k - 1))][j][k - 1][0]);
41         }
42     }
43 }
44
45 for (int k = 1; k <= (int)(log2(m)); k++) {
46     for (int i = 0; i < n; i++) {
47         for (int j = 0; j + (1 << k) - 1 < m; j++) {
48             table[i][j][0][k] = f(
49                 table[i][j][0][k - 1],
50                 table[i][j + (1 << (k - 1))][0][k - 1]);
51         }
52     }
53 }
54
55 for (int k = 1; k <= (int)(log2(n)); k++) {
56     for (int l = 1; l <= (int)(log2(m)); l++) {
57         for (int i = 0; i + (1 << k) - 1 < n; i++) {
58             for (int j = 0; j + (1 << l) - 1 < m; j++)
59 {
60                 table[i][j][k][l] = f(
61                     f(
62                         table[i][j][k - 1][l - 1],
63                         table[i + (1 << (k - 1))][j][k - 1][l
64 - 1]),
65                     f(
66                         table[i][j + (1 << (l - 1))][k - 1][l
67 - 1],
68                         table[i + (1 << (k - 1))][j + (1 << (
69 l - 1))][k - 1][l - 1])
70                     );
71                 }
72             }
73         }
74     }
75 }
76
77 int query(int x1, int y1, int x2, int y2) {
78     int k = log2(x2 - x1 + 1);
79     int l = log2(y2 - y1 + 1);
80
81     return f(
82         f(
83             table[x1][y1][k][l],
84             table[x2 - (1 << k) + 1][y1][k][l]
85         ),
86         f(
87             table[x1][y2 - (1 << l) + 1][k][l],
88             table[x2 - (1 << k) + 1][y2 - (1 << l) + 1][k
89 ][l]
90         )
91     );
92 }
93 };

```