

**Definition 1** (non-negative linear combination). *Let  $U \in \mathbb{R}^{d \times p}$ ,  $\mathbf{t} \in \mathbb{R}^p$ ,  $\mathbf{t} \geq \mathbf{0}$ , then  $\sum_{1 \leq j \leq p} t_j U^j$  is called a **non-negative linear combination** of  $U$ .*

**Definition 2** (V-Cone). *Let  $U \in \mathbb{R}^{d \times p}$ . The set of all non-negative linear combinations of  $U$  is denoted  $\text{cone}(U)$ . Such a set is called a **V-Cone**.*

**Definition 3** (convex combination). *Let  $V \in \mathbb{R}^{d \times n}$ ,  $\boldsymbol{\lambda} \in \mathbb{R}^n$ ,  $\boldsymbol{\lambda} \geq \mathbf{0}$ ,  $\sum_{1 \leq j \leq n} \lambda_j = 1$ , then  $\sum_{1 \leq j \leq n} \lambda_j V^j$  is called a **convex combination** of  $V$ . The set of all convex combinations of  $V$  is denoted  $\text{conv}(V)$ .*

**Definition 4** (V-Polyhedron). *Let  $V \in \mathbb{R}^{d \times n}$ ,  $U \in \mathbb{R}^{d \times p}$ . Then the set*

$$\{\mathbf{x} + \mathbf{y} \mid \mathbf{x} \in \text{cone}(U), \mathbf{y} \in \text{conv}(V)\}$$

*is called a V-Polyhedron.*

**Definition 5** (H-Polyhedron). *Let  $A \in \mathbb{R}^{m \times d}$ ,  $\mathbf{b} \in \mathbb{R}^m$ . Then the set*

$$\left\{ \mathbf{x} \in \mathbb{R}^d \mid A\mathbf{x} \leq \mathbf{b} \right\}$$

*is called an H-Polyhedron.*

**Definition 6** (H-Cone). *Let  $A \in \mathbb{R}^{m \times d}$ . Then the set*

$$\left\{ \mathbf{x} \in \mathbb{R}^d \mid A\mathbf{x} \leq \mathbf{0} \right\}$$

*is called an H-Cone.*

The following theorem is the basic result to be proved in this thesis, which states that V-Polyhedra and H-Polyhedra are two different representations of the same objects.

**Theorem 1** (Minkowski-Weyl Theorem). *Every V-Polyhedron is an H-Polyhedron, and every H-Polyhedron is a V-Polyhedron.*

The proof proceeds by first showing that V-Cones are representable as H-Cones, and H-Cones are representable as V-Cones. Then it is shown that the case of polyhedra can be reduced to cones.

**Theorem 2** (Minkowski-Weyl Theorem for Cones). *Every V-Cone is an H-Cone, and every H-Cone is a V-Cone.*

## 1 Every V-Cone is an H-Cone

**Definition 7** (Coordinate Projection). *Let  $I$  be the identity matrix. Then the matrix  $I'$  formed by deleting some rows from  $I$  is called a **coordinate-projection**.*

The proof rests on the following two propositions:

(V1) Every V-Cone is a coordinate-projections of an H-Cone.

(V2) Every coordinate-projection of an H-Cone is an H-Cone.

*Proof.* Given (V1) and (V2), the proof follows simply. Given a V-Cone, we use (V1), to get a description involving coordinate-projection of an H-Cone. Then we can apply (V2) in order to get an H-Cone.  $\square$

*Proof of (V1).* We prove that every V-Cone is a coordinate-projection of an H-Cone, by giving an explicit formula. Let  $U \in \mathbb{R}^{d \times p}$ , and observe that

$$\text{cone}(U) = \{U\mathbf{t} \mid \mathbf{t} \in \mathbb{R}^p, \mathbf{t} \geq \mathbf{0}\} = \{\mathbf{x} \in \mathbb{R}^d \mid (\exists \mathbf{t} \in \mathbb{R}^p) \mathbf{x} = U\mathbf{t}, \mathbf{t} \geq \mathbf{0}\}$$

We will collect  $\mathbf{t}$  and  $\mathbf{x}$  on the left side of the inequality, treating  $\mathbf{t}$  as a variable and expressing its constraints as linear inequalities, then project away the coordinates corresponding to  $\mathbf{t}$ . The following expression takes one step:

$$\mathbf{t} \geq \mathbf{0} \Leftrightarrow -I\mathbf{t} \leq \mathbf{0} \quad (1)$$

And using the equality:  $a = 0 \Leftrightarrow a \leq 0 \wedge -a \leq 0$ , and block matrix notation, we take the second step.

$$\mathbf{x} = U\mathbf{t} \Leftrightarrow \mathbf{x} - U\mathbf{t} = \mathbf{0} \Leftrightarrow \begin{pmatrix} I & -U \\ -I & U \end{pmatrix} \begin{pmatrix} \mathbf{x} \\ \mathbf{t} \end{pmatrix} \leq \mathbf{0} \quad (2)$$

Comparing (1) and (2), we define a new matrix  $A' \in \mathbb{R}^{(p+2d) \times (d+p)}$ :

$$A' = \begin{pmatrix} \mathbf{0} & -I \\ I & -U \\ -I & U \end{pmatrix}$$

then we can rewrite  $\text{cone}(U)$ :

$$\text{cone}(U) = \left\{ \mathbf{x} \in \mathbb{R}^d \mid A' \begin{pmatrix} \mathbf{x} \\ \mathbf{t} \end{pmatrix} \leq \mathbf{0} \right\}$$

Let  $\Pi \in \{0, 1\}^{d \times (d+p)}$  be the identity matrix in  $\mathbb{R}^{(d+p) \times (d+p)}$ , but with the last  $p$ -rows deleted. Then  $\Pi$  is a coordinate projection, and the above expression can be written:

$$\text{cone}(U) = \Pi \left( \{\mathbf{y} \in \mathbb{R}^{d+p} \mid A'\mathbf{y} \leq \mathbf{0}\} \right) \quad (3)$$

This is a coordinate projection of an H-Cone, and (V1) is shown.  $\square$

To prove (V2), we use two separate propositions.

**Proposition 1.** Let  $B \in \mathbb{R}^{m' \times (d+p)}$ ,  $B'$  be  $B$  with the last  $p$  columns deleted, and  $\Pi$  the identity matrix with the last  $p$  rows deleted (i.e.  $B' = \Pi B$ ). Furthermore, suppose that the last  $p$  columns of  $B$  are  $\mathbf{0}$ . Then

$$\Pi \left( \{\mathbf{y} \in \mathbb{R}^{d+p} \mid B\mathbf{y} \leq \mathbf{0}\} \right) = \{\mathbf{x} \in \mathbb{R}^d \mid B'\mathbf{x} \leq \mathbf{0}\}$$

*Proof.* Recall that  $B\mathbf{y} \leq \mathbf{0}$  means that  $(\forall i) \langle B_i, \mathbf{y} \rangle \leq 0$ . By the way we've defined  $B$ , any row  $B_i$  of  $B$  can be written  $(B'_i, \mathbf{0})$ , with  $\mathbf{0} \in \mathbb{R}^p$ . Rewriting  $\mathbf{y} \in \mathbb{R}^{d+p}$  as  $(\mathbf{x}, \mathbf{w})$  with  $\mathbf{x} \in \mathbb{R}^d, \mathbf{w} \in \mathbb{R}^p$ , so that  $\mathbf{x} = \Pi(\mathbf{y})$ . Then

$$\langle B_i, \mathbf{y} \rangle = \langle (B'_i, \mathbf{0}), (\mathbf{x}, \mathbf{w}) \rangle = \langle B'_i, \mathbf{x} \rangle = \langle B'_i, \Pi(\mathbf{y}) \rangle$$

It follows that

$$\langle B_i, \mathbf{y} \rangle \leq 0 \Leftrightarrow \langle B'_i, \Pi(\mathbf{y}) \rangle \leq 0$$

Since  $B_i$  is an arbitrary row of  $B$ , the proposition is shown.  $\square$

In order to use the above proposition, we need a matrix with  $\mathbf{0}$  columns. The next proposition shows us how to do so, one column at a time.

**Proposition 2.** *Let  $B \in \mathbb{R}^{m_1 \times (d+p)}$ ,  $1 \leq k \leq p$ , and  $\mathbf{x} = \sum_{i \neq k} x_i \mathbf{e}_i$ . Then there exists a matrix  $B' \in \mathbb{R}^{m_2 \times (d+p)}$  with the following properties:*

1. *Every row of  $B'$  is a positive linear combination of rows of  $B$ .*
2.  *$m_2$  is finite.*
3. *The  $k$ -th column of  $B'$  is  $\mathbf{0}$ .*
4.  *$(\exists t) B(\mathbf{x} + t\mathbf{e}_k) \leq \mathbf{0} \Leftrightarrow B'\mathbf{x} \leq \mathbf{0}$*

*Proof.* Partition the rows of  $B$  as follows:

$$\begin{aligned} P &= i \mid B_i^k > 0 \\ N &= j \mid B_j^k < 0 \\ Z &= l \mid B_l^k = 0 \end{aligned}$$

Then let  $B'$  be a matrix with rows of the following forms:

$$\begin{aligned} C_l &= B_l & \mid l \in Z \\ C_{ij} &= B_i^k B_j - B_j^k B_i & \mid i \in P, j \in N \end{aligned}$$

1 and 2 are clear. 3 can be seen from:

$$\begin{aligned} \langle C_l, \mathbf{e}_k \rangle &= 0 \\ \langle C_{ij}, \mathbf{e}_k \rangle &= \langle B_i^k B_j - B_j^k B_i, \mathbf{e}_k \rangle = B_i^k B_j^k - B_j^k B_i^k = 0 \end{aligned} \tag{4}$$

The right direction of 4 is shown in the following calculations. Because  $B_l^k = 0$ :

$$\langle B_l, \mathbf{x} + t\mathbf{e}_k \rangle = \langle B_l, \mathbf{x} \rangle + tB_l^k = \langle B_l, \mathbf{x} \rangle = \langle C_l, \mathbf{x} \rangle$$

So:

$$\langle B_l, \mathbf{x} + t\mathbf{e}_k \rangle \leq 0 \Rightarrow \langle C_l, \mathbf{x} \rangle \leq 0$$

For rows indexed by  $P, N$ , we observe (13), and have:

$$\langle B_i^k B_j - B_j^k B_i, \mathbf{x} + t\mathbf{e}_k \rangle = \langle B_i^k B_j - B_j^k B_i, \mathbf{x} \rangle$$

Now, we use property 1:

$$\langle B_i, \mathbf{x} + t\mathbf{e}_k \rangle \leq 0, \langle B_j, \mathbf{x} + t\mathbf{e}_k \rangle \leq 0 \Rightarrow \langle B_i^k B_j - B_j^k B_i, \mathbf{x} + t\mathbf{e}_k \rangle \leq 0$$

Therefore

$$\langle B_i^k B_j - B_j^k B_i, \mathbf{x} \rangle \leq 0$$

Now suppose that  $B'\mathbf{x} \leq \mathbf{0}$ . The task is to find a  $t$  so that  $B\mathbf{x} \leq \mathbf{0}$ . Looking at (13), any choice of  $t$  we make will be okay for rows indexed by  $Z$ . So the task is to find a  $t$  so that the inequality holds for rows indexed by  $P$  and  $N$ . Observe

$$\begin{aligned} \forall i \in P, \forall j \in N \quad & \langle B_i^k B_j - B_j^k B_i, \mathbf{x} \rangle \leq 0 \Leftrightarrow \\ \forall i \in P, \forall j \in N \quad & \langle B_i^k B_j, \mathbf{x} \rangle \leq \langle B_j^k B_i, \mathbf{x} \rangle \Leftrightarrow \\ \forall i \in P, \forall j \in N \quad & \langle B_j/B_j^k, \mathbf{x} \rangle \geq \langle B_i/B_i^k, \mathbf{x} \rangle \Leftrightarrow \\ & \min_{j \in N} \langle B_j/B_j^k, \mathbf{x} \rangle \geq \max_{i \in P} \langle B_i/B_i^k, \mathbf{x} \rangle \end{aligned}$$

Note that the third inequality changes directions because  $B_j^k < 0$ . Now we choose  $t$  to lie in this last interval, and show that we can use it to satisfy all of the constraints given by  $B$ . So, we have a  $t$  such that

$$\min_{j \in N} \langle B_j/B_j^k, \mathbf{x} \rangle \geq t \geq \max_{i \in P} \langle B_i/B_i^k, \mathbf{x} \rangle$$

In particular,

$$\begin{aligned} (\forall j \in N) \quad & \langle B_j/B_j^k, \mathbf{x} \rangle \geq t \Rightarrow \\ (\forall j \in N) \quad & \langle B_j, \mathbf{x} \rangle - B_j^k t \leq 0 \end{aligned}$$

Again, the inequality changes directions because  $B_j^k < 0$ . Now consider a row  $B_j$  from  $B$ :

$$\langle B_j, \mathbf{x} - t\mathbf{e}_k \rangle = B_j \mathbf{x} - B_j^k t \leq 0$$

Similarly,

$$\begin{aligned} (\forall i \in P) \quad & t \geq B_i/B_i^k \mathbf{x} \Rightarrow \\ (\forall i \in P) \quad & 0 \geq B_i \mathbf{x} - B_i^k t \end{aligned}$$

Now consider a row  $B_i$  from  $B$ :

$$\langle B_i, \mathbf{x} - t\mathbf{e}_k \rangle = B_i \mathbf{x} - B_i^k t \leq 0$$

So, we've demonstrated that  $\mathbf{x} - t\mathbf{e}_k$  satisfies all the constraints from  $B$ , and the left implication is shown. So 4 holds.  $\square$

Now to prove:

(V2) Every coordinate-projection of an H-Cone is an H-Cone.

*proof of (V2).* Here we prove the case that the coordinate projection is onto the first  $d$  of  $d+p$  coordinates. Let  $\{\mathbf{y} \in \mathbb{R}^{d+p} : A'\mathbf{y} \leq \mathbf{0}\}$  be the H-Cone we need to project, and  $\Pi$  the coordinate-projection we need to apply (the identity matrix with the last  $p$  columns deleted). For each  $1 \leq k \leq p$  we can use proposition 2 in an incremental manner, starting with  $A'$ .

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let  $B_0 := A'$ 
for  $1 \leq k \leq p$ 
  let  $B_k :=$  result of proposition 2 applied to  $B_{k-1}, \mathbf{e}_{d+k}$ 
endfor
return  $B_p$ 

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Consider the resulting  $B$ . Property 2 holds throughout, so  $B$  is finite. After each iteration, property 3 holds for  $k$ , so the  $k$ -th column is  $\mathbf{0}$ . Since each iteration only results from non-negative combinations of the result of the previous iteration (property 1), once a column is  $\mathbf{0}$  it remains so. Therefore, at the end of the process, the last  $p$  columns of  $B$  are all  $\mathbf{0}$ . Then, by proposition 1, we can apply  $\Pi$  to  $B$  by simply deleting the last  $p$  columns of  $B$ . Denote this resulting matrix  $A$ . We still need to check:

$$A'\mathbf{y} \leq \mathbf{0} \Leftrightarrow A(\Pi(\mathbf{y})) \leq \mathbf{0} \quad (5)$$

$$(\exists t_1) \dots (\exists t_p) A'(\mathbf{x} + t_1 \mathbf{e}_{d+1} + \dots + t_p \mathbf{e}_{d+p}) \leq \mathbf{0} \Leftrightarrow A\mathbf{x} \leq \mathbf{0} \quad (6)$$

Then, using (5) and (6), it is easy to see that:

$$\Pi \{\mathbf{y} \in \mathbb{R}^{d+p} \mid A'\mathbf{y} \leq \mathbf{0}\} = \{\mathbf{x} \in \mathbb{R}^d \mid A\mathbf{x} \leq \mathbf{0}\} \quad (7)$$

The key observation of this verification utilizes property 4 of proposition 2:

$$(\exists t) B(\mathbf{x} + t \mathbf{e}_k) \leq \mathbf{0} \Leftrightarrow B'\mathbf{x} \leq \mathbf{0}$$

In what follows, let  $\mathbf{x} = \sum_{1 \leq j \leq d} x_j \mathbf{e}_j$ . The above property is applied sequentially to the sets  $B_k$  as follows:

$$\begin{array}{lll}
(\exists t_p)(\exists t_{p-1}) \dots (\exists t_1) & B_0(\mathbf{x} + t_1 \mathbf{e}_p + t_2 \mathbf{e}_{p-1} + \dots + t_p \mathbf{e}_d) \leq \mathbf{0} & \Leftrightarrow \\
(\exists t_p) \dots (\exists t_2) & B_1(\mathbf{x} + t_2 \mathbf{e}_{d+2} + \dots + t_p \mathbf{e}_{d+p}) \leq \mathbf{0} & \Leftrightarrow \\
& \vdots & \vdots \\
(\exists t_p) & B_{p-1}(\mathbf{x} + t_p \mathbf{e}_{d+p}) \leq \mathbf{0} & \Leftrightarrow \\
& B_p \mathbf{x} \leq \mathbf{0} & \Leftrightarrow
\end{array}$$

Because  $A' = B_0$ , and  $A$  is  $B_p$  with the last  $p$  columns deleted, (5) and (6) hold, therefore (7) holds, and the proof of (V2) is complete, and we've shown that a coordinate projection of an H-Cone is again an H-Cone.  $\square$

With (V1) and (V2) proven, we are now certain that any V-Cone is also an H-Cone.

## 2 Every H-Cone is a V-Cone

**Definition 8** (Coordinate Hyperplane). *A set of the form*

$$\{\mathbf{x} \in \mathbb{R}^{d+m} \mid \langle \mathbf{x}, \mathbf{e}_k \rangle = 0\} = \{\mathbf{x} \in \mathbb{R}^{d+m} \mid x_k = 0\}$$

*is called a coordinate-hyperplane.*

We will use coordinate-hyperplanes in the following way. We consider a V-Cone intersected with some coordinate hyperplanes, and write it in the following way:

$$\left\{ \mathbf{x} \in \mathbb{R}^d \mid (\exists \mathbf{t} \geq 0) \begin{pmatrix} \mathbf{x} \\ \mathbf{0} \end{pmatrix} = U' \mathbf{t} \right\} \quad (8)$$

If we suppose that  $U' \subset \mathbb{R}^{d+m}$ , and  $\Pi$  is the identity matrix with the last  $m$  rows deleted, then this is just a convenient way of writing:

$$\Pi(\text{cone}(U') \cap \{x_{d+1} = 0\} \cap \cdots \cap \{x_{d+m} = 0\}) \quad (9)$$

The proof rests on the following three propositions:

*H1* Every H-Cone is a coordinate-projection of a V-Cone intersected with some coordinate hyperplanes.

*H2* Every V-Cone intersected with a coordinate-hyperplane is a V-Cone

*H3* Every coordinate-projection of a V-Cone is an V-Cone.

*Proof.* Given *H1*, *H2*, and *H3*, the proof follows simply. Given an H-Cone, we use *H1* to get a description involving the coordinate-projection of a V-Cone intersected with some coordinate-hyperplanes. We apply *H2* as many times as necessary to eliminate the intersections, then we can apply *H3* in order to get a V-Cone.  $\square$

*Proof of H1.* Let  $A \in \mathbb{R}^{m \times d}$ , we now show that the H-Cone

$$\{\mathbf{x} \in \mathbb{R}^d \mid A\mathbf{x} \leq \mathbf{0}\}$$

can be written as the projection of a V-Cone intersected with some hyperplanes. Define  $U'$ :

$$U' = \left\{ \begin{pmatrix} \mathbf{e}_j \\ A^j \end{pmatrix}, \begin{pmatrix} -\mathbf{e}_j \\ -A^j \end{pmatrix}, \begin{pmatrix} \mathbf{0} \\ \mathbf{e}_i \end{pmatrix}, 1 \leq j \leq d, 1 \leq i \leq m \right\}$$

Then we claim:

$$\{\mathbf{x} \in \mathbb{R}^d \mid A\mathbf{x} \leq \mathbf{0}\} = \left\{ \mathbf{x} \in \mathbb{R}^d \mid (\exists \mathbf{t} \geq 0) \begin{pmatrix} \mathbf{x} \\ \mathbf{0} \end{pmatrix} = U' \mathbf{t} \right\} \quad (10)$$

First, considering (8) and (9), observe that this is a coordinate-projection of a V-Cone intersected with some coordinate-hyperplanes. Next, we note that

$$\begin{pmatrix} \mathbf{x} \\ A\mathbf{x} \end{pmatrix} = \sum_{1 \leq j \leq d} x_j \begin{pmatrix} \mathbf{e}_j \\ A^j \end{pmatrix}$$

We can write this as a sum with all positive coefficients if we split up the  $x_j$  as follows:

$$x_j^+ = \begin{cases} x_j & x_j \geq 0 \\ 0 & x_j < 0 \end{cases} \quad x_j^- = \begin{cases} 0 & x_j \geq 0 \\ -x_j & x_j < 0 \end{cases}$$

Then we have

$$\begin{pmatrix} \mathbf{x} \\ A\mathbf{x} \end{pmatrix} = \sum_{1 \leq j \leq d} x_j^+ \begin{pmatrix} \mathbf{e}_j \\ A^j \end{pmatrix} + \sum_{1 \leq j \leq d} x_j^- \begin{pmatrix} -\mathbf{e}_j \\ -A^j \end{pmatrix} \quad (11)$$

where  $x_j^+, x_j^- \geq 0$ . Also observe that

$$A\mathbf{x} \leq \mathbf{0} \Leftrightarrow (\exists \mathbf{w} \geq \mathbf{0}) \mid A\mathbf{x} + \mathbf{w} = \mathbf{0}$$

This can also be written

$$A\mathbf{x} \leq \mathbf{0} \Leftrightarrow (\exists \mathbf{w} \geq \mathbf{0}) \mid \begin{pmatrix} \mathbf{x} \\ A\mathbf{x} \end{pmatrix} + \begin{pmatrix} \mathbf{0} \\ \mathbf{w} \end{pmatrix} = \begin{pmatrix} \mathbf{x} \\ \mathbf{0} \end{pmatrix} \quad (12)$$

(11) and (12) together show

$$A\mathbf{x} \leq \mathbf{0} \Rightarrow (\exists \mathbf{t} \geq \mathbf{0}) \begin{pmatrix} \mathbf{x} \\ \mathbf{0} \end{pmatrix} = U'\mathbf{t}$$

Conversely, suppose

$$(\exists \mathbf{t} \geq \mathbf{0}) \begin{pmatrix} \mathbf{x} \\ \mathbf{0} \end{pmatrix} = U'\mathbf{t}$$

We would like to show that  $A\mathbf{x} \leq \mathbf{0}$ . Let  $x_j^+, x_j^-, w_i$  take the values of  $\mathbf{t}$  that are coefficients of  $\begin{pmatrix} \mathbf{e}_j \\ A^j \end{pmatrix}$ ,  $\begin{pmatrix} -\mathbf{e}_j \\ -A^j \end{pmatrix}$ , and  $\begin{pmatrix} \mathbf{0} \\ \mathbf{e}_i \end{pmatrix}$  respectively, and denote  $x_j = x_j^+ - x_j^-$ .

Then we have

$$\begin{aligned} \begin{pmatrix} \mathbf{x} \\ \mathbf{0} \end{pmatrix} &= \sum_{1 \leq j \leq d} x_j^+ \begin{pmatrix} \mathbf{e}_j \\ A^j \end{pmatrix} + \sum_{1 \leq j \leq d} x_j^- \begin{pmatrix} -\mathbf{e}_j \\ -A^j \end{pmatrix} + \sum_{1 \leq i \leq n} w_i \begin{pmatrix} \mathbf{0} \\ \mathbf{e}_i \end{pmatrix} \\ &= \sum_{1 \leq j \leq d} x_j \begin{pmatrix} \mathbf{e}_j \\ A^j \end{pmatrix} + \sum_{1 \leq i \leq n} w_i \begin{pmatrix} \mathbf{0} \\ \mathbf{e}_i \end{pmatrix} \\ &= \begin{pmatrix} \mathbf{x} \\ A\mathbf{x} \end{pmatrix} + \begin{pmatrix} \mathbf{0} \\ \mathbf{w} \end{pmatrix} \end{aligned}$$

where  $\mathbf{w} \geq \mathbf{0}$ . By (12) we have  $A\mathbf{x} \leq \mathbf{0}$ . So (10) holds.  $\square$

The proof of  $H2$  relies upon the following proposition.

**Proposition 3.** *Let  $Y \in \mathbb{R}^{(d+m) \times n_1}$ ,  $1 \leq k \leq m$ , and  $\mathbf{x}$  satisfy  $x_k = 0$ . Then there exists a matrix  $Y' \in \mathbb{R}^{(d+m) \times n_2}$  with the following properties:*

1. *Every column of  $Y'$  is a positive linear combination of rows of  $B$ .*
2.  *$n_2$  is finite.*
3. *The  $k$ -th row of  $Y'$  is  $\mathbf{0}$ .*
4.  *$(\exists \mathbf{t} \geq \mathbf{0}) \mathbf{x} = Y \mathbf{t} \Leftrightarrow (\exists \mathbf{t}' \geq \mathbf{0}) \mathbf{x} = Y' \mathbf{t}'$*

Recall that  $Y^i$  is the  $i$ -th column of  $Y$ , and  $Y_k^i$  is the element of  $Y$  in the  $i$ -th column and  $k$ -th row.

*Proof.* We partition the columns of  $Y$ :

$$\begin{aligned} P &= i \mid Y_k^i > 0 \\ N &= j \mid Y_k^j < 0 \\ Z &= l \mid Y_k^l = 0 \end{aligned}$$

We then define  $Y'$ :

$$Y' = \{Y^l \mid l \in Z\} \cup \{Y_k^i Y^j - Y_k^j Y^i \mid i \in P, j \in N\}$$

1 and 2 are clear. 3 can be seen from:

$$\begin{aligned} \langle Y^l, \mathbf{e}^k \rangle &= 0 \\ \langle Y^{ij}, \mathbf{e}^k \rangle &= \langle Y_k^i Y^j - Y_k^j Y^i, \mathbf{e}^k \rangle = Y_k^i Y_k^j - Y_k^j Y_k^i = 0 \end{aligned} \tag{13}$$

Before moving on to the proof, we first note how to write our vectors.

$$\begin{aligned} Y \mathbf{t} &= \sum_{k \in Z} t_k Y^k + \sum_{i \in P} t_i Y^i + \sum_{j \in N} t_j Y^j \\ Y' \mathbf{t} &= \sum_{k \in Z} t_k Y^k + \sum_{\substack{i \in P \\ j \in N}} t_{ij} (Y_k^i Y^j - Y_k^j Y^i) \end{aligned}$$

Then, to show that the proposition is true, we need only show that, given any  $t_i, t_j \geq 0$  ( $t_{ij} \geq 0$ ), there exists  $t_{ij} \geq 0$  ( $t_i, t_j \geq 0$ ) such that

$$\sum_{i \in P} t_i Y^i + \sum_{j \in N} t_j Y^j = \sum_{\substack{i \in P \\ j \in N}} t_{ij} (Y_k^i Y^j - Y_k^j Y^i) \tag{14}$$



**Proposition 4.** *Suppose that*

$$\sum_{i \in P} t_i Y_{d+1}^i + \sum_{j \in N} t_j Y_{d+1}^j = 0 \quad Y_k^j < 0 < Y_k^i$$

*Then the following holds*

$$\begin{aligned} (t_i, t_j \geq 0) &\Rightarrow (\exists t_{ij} \geq 0) \quad \text{such that (14) holds} \\ (t_{ij} \geq 0) &\Rightarrow (\exists t_i, t_j \geq 0) \text{ such that (14) holds} \end{aligned}$$

*Proof.* First note that if all  $t_i = 0, t_j = 0$ , then choosing  $t_{ij} = 0$  satisfies (14), likewise if all  $t_{ij} = 0$ , then  $t_i = 0, t_j = 0$  satisfies (14). So suppose that some  $t_i \neq 0, t_j \neq 0, t_{ij} \neq 0$ .

The right hand side of (14) can be written

$$\sum_{j \in N} \left( \sum_{i \in P} t_{ij} Y_k^i \right) Y^j + \sum_{i \in P} \left( - \sum_{j \in N} t_{ij} Y_k^j \right) Y^i$$

This means, given  $t_{ij} \geq 0$ , we can choose  $t_j = \sum_{i \in P} t_{ij} Y_k^i$ , and  $t_i = - \sum_{j \in N} t_{ij} Y_k^j$ , both of which are greater than 0.

Now suppose we have been given  $t_i \geq 0, t_j \geq 0$ . First observe:

$$0 = \sum_{i \in P} t_i Y_k^i + \sum_{j \in N} t_j Y_k^j \Rightarrow \sum_{i \in P} t_i Y_k^i = - \sum_{j \in N} t_j Y_k^j$$

Denote the value in this equality as  $\sigma$ , and note that  $\sigma > 0$ . Then

$$\begin{aligned} \sum_{i \in P} t_i Y^i &= \frac{- \sum_{j \in N} t_j Y_k^j}{\sigma} \sum_{i \in P} t_i Y^i = \sum_{\substack{i \in P \\ j \in N}} - \frac{t_i t_j}{\sigma} Y_k^j Y^i \\ \sum_{j \in N} t_j Y^j &= \frac{\sum_{i \in P} t_i Y_k^i}{\sigma} \sum_{j \in N} t_j Y^j = \sum_{\substack{i \in P \\ j \in N}} \frac{t_i t_j}{\sigma} Y_k^i Y^j \end{aligned}$$

Combining these results, we have

$$\sum_{i \in P} t_i Y^i + \sum_{j \in N} t_j Y^j = \sum_{\substack{i \in P \\ j \in N}} \frac{t_i t_j}{\sigma} (Y_k^i Y^j - Y_k^j Y^i)$$

□

Finally, we can conclude that, given  $\mathbf{t} \geq \mathbf{0}$ , if  $Y\mathbf{t}$  has a 0 in the final coordinate, then we can write it as  $Y'\mathbf{t}'$  where  $\mathbf{t}' \geq \mathbf{0}$ , and any non-negative linear combination of vectors from  $Y'$  can be written as a non-negative linear combination of vectors from  $Y$ , and will necessarily have the  $k$ -th coordinate be 0 by property 3. So property 4 holds. □

*Proof of H2.* In proposition 3, the assumption that  $x_k = 0$  in property 4 creates the set  $\text{cone}(Y) \cap \{\mathbf{x} \mid x_k = 0\}$ . This set, by property 4, is  $\text{cone}(Y')$ .  $\square$

*Proof of H3.* We shall prove that the coordinate-projection of a V-Cone is again a V-Cone. Let  $\Pi$  be the relevant projection, then we have:

$$\Pi\{U\mathbf{t} \mid \mathbf{t} \geq \mathbf{0}\} = \{\Pi(U\mathbf{t}) \mid \mathbf{t} \geq \mathbf{0}\} = \{\Pi(U)\mathbf{t} \mid \mathbf{t} \geq \mathbf{0}\}$$

The last equality follows from associativity of matrix multiplication. Therefore,

$$\Pi(\text{cone}(U)) = \text{cone}(\Pi(U))$$

$\square$