CS 111 ASSIGNMENT 4

due February 19, 2024

Problem 1: Give an asymptotic estimate, using the Θ -notation, of the number of letters printed by the algorithms given below. Give a complete justification for your answer, by providing an appropriate recurrence equation and its solution.

```
(a) algorithm PrintAs(n)
                                                           (b) algorithm PrintBs(n)
          if n \leq 1 then
                                                                     if n \geq 4 then
                                                                          for j \leftarrow 1 to n^2
               print("A")
                                                                               do print("B")
          else
               for j \leftarrow 1 to n^3
                                                                          for i \leftarrow 1 to 6 do
                   do print("A")
                                                                               PrintBs(\lfloor n/4 \rfloor)
               for i \leftarrow 1 to 5 do
                                                                          for i \leftarrow 1 to 10 do
                   PrintAs(|n/2|)
                                                                               PrintBs(\lceil n/4 \rceil)
(c) algorithm PrintCs(n)
                                                          (d) algorithm PrintDs(n)
          if n \leq 2 then
                                                                     if n \geq 5 then
               print("C")
                                                                          print("D")
                                                                          print("D")
          else
               for j \leftarrow 1 to n
                                                                        if (x \equiv 0 \pmod{2}) then
                   do print("C")
                                                                              PrintDs(\lfloor n/5 \rfloor)
               PrintCs(|n/3|)
                                                                              PrintDs(\lceil n/5 \rceil)
               PrintCs(|n/3|)
                                                                               x \leftarrow x + 3
               PrintCs(|n/3|)
                                                                          else
                                                                               PrintDs(\lceil n/5 \rceil)
               PrintCs(|n/3|)
                                                                              PrintDs(\lfloor n/5 \rfloor)
                                                                               x \leftarrow 5x + 3
```

In part (d), variable x is a global variable initialized to 1.

The solutions for all the problems will use the master theorem, so I will just state it here. Let $a \ge 1, b > 1, c > 0$ and $d \ge 0$. If T(n) satisfies the recurrence $T(n) = aT(n/b) + cn^d$.

$$T(x) = \begin{cases} \Theta(n^{log_b(a)}) & a > b^d \\ \Theta(n^d log(n)) & a = b^d \\ \Theta(n^d) & a < b^d \end{cases}$$

Solution 1:

- (a) This algorithm can be rewritten as the recurrence: $A(n) = n^3 + 5A(\lfloor n/2 \rfloor)$. Applying the master theorem, with a = 5, b = 2, c = 1, d = 3. Since $5 < 2^3 = 8, T(n) = \Theta(n^3)$.
- (b) This algorithm can be rewritten as the recurrence: $B(n) = n^2 + 6B(\lfloor n/4 \rfloor) + 10B(\lceil n/4 \rceil) = 16B(n/4) + n^2 \implies a = 16, b = 4, c = 1, d = 2$. From the master theorem, since $16 = 4^2$, $B(n) = \Theta(n^2 \log(n))$.
- (c) This algorithm can be rewritten as the recurrence: $C(n) = 4C(\lfloor n/3 \rfloor) + n$. Applying the master theorem, with $a = 4, b = 3, c = 1, d = 1 \implies a > b^1 \implies C(n) = \Theta(n^{\log_3 4})$.
- (d) This algorithm can be rewritten as the recurrence: D(n) = 2 + 2D(n/5). The reason for the 2D(n/5) is because there even though there are four recursive calls in the code, only two are executed in any given iteration. Applying the master theorem, with $a = 2, b = 5, c = 2, d = 0 \implies a < b^d$, we have $B(n) = \Theta(n^{\log_5(2)})$.

Problem 2: We have three sets A, B, C with the following properties:

(a)
$$|B| = 2|A|, |C| = 3|A|,$$

(b)
$$|A \cap B| = 18$$
, $|A \cap C| = 20$, $|B \cap C| = 24$,

(c)
$$|A \cap B \cap C| = 11$$
,

(d)
$$|A \cup B \cup C| = 129$$
.

Use the inclusion-exclusion principle to determine the number of elements in A. Show your work.

Solution 2: According to the inclusion exclusion principle, we have:

$$|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|.$$

We can substitute the values from (b), (c), (d) into the principle:

 $129 = |A| + |B| + |C| - 18 - 20 - 24 + 11 \implies |A| + |B| + |C| = 180$. To solve this, we want to put everything in terms of a single variable. From (a), we know the relationships between the three variables. I will choose to express everything in terms of |A|. Plugging this in, we get

$$|A| + 2|A| + 3|A| = 180 \implies |A| = 30.$$

From (a), we can find the rest of the values: |B| = 60, |C| = 90.

Problem 3: A company, Nice Inc., will award 45 fellowships to high-achieving UCR students from four different majors: computer science, biology, political science and history. They decided to give fellowship awards to at least 8 students majoring in computer science and at most 8 biology majors. The number of political science and history majors should be between 5 and 12 students each. How many possible lists of awardees are there? You need to give a complete derivation for the final answer, using the method developed in class. (Brute force listing of all lists will not be accepted.)

Solution 3: Let's say that the number of CS awards is C, biology is B, political science is P, history is H. We can express the conditions as:

$$C \ge 8$$

 $B \le 8$
 $5 \le P, Q, \le 12 \implies 0 \le P', Q' \le 7$
 $C + B + P + H = 45$

We can substitute B' = B - 5 and P' = P - 5 into the sum. Also, since C has no restrictions on the upper bound, we can kind of ignore it for now. We can now write: B + P' + H' = 27.

We can also let the number of partitions that fit these conditions be the total number of partitions (S()) minus the number of invalid partitions:

$$S(B \le 8 \land P' \le 7 \land H' \le 7)$$

= $S() - S(B \ge 9 \lor P' \ge 8 \lor H' \ge 8).$

We know that the total number of partitions, $S() = {m+k-1 \choose k-1}, m=27, k=4 \implies S() = {30 \choose 3} = 4060.$

We can also calculate the number of invalid partitions using the inclusion-exclusion principle. Also that the number of partitions with conditions is $\binom{m-A+k-1}{k-1}$, where A is the sum of the lower bounds.

$$\begin{split} S(B &\geq 9 \lor P' \geq 8 \lor H' \geq 8) = \\ S(B &\geq 9) + S(P' \geq 8) + S(H' \geq 8) \\ -S(B &\geq 9 \land P' \geq 8) - S(B \geq 9 \land P' \geq 8) - S(B \geq 8 \land P' \geq 8) \\ +S(B &\geq 9 \land P' \geq 8 \land H' \geq 8) \end{split}$$

$$= {27-9+4-1 \choose 4-1} + 2{27-8+4-1 \choose 4-1}$$

$$- 2{27-(9+8)+4-1 \choose 4-1} - {27-(8+8)+4-1 \choose 4-1} - {27-(9+8+8)+4-1 \choose 4-1}$$

$$= 1330 + 1540 + 1540 - 286 - 286 - 364 + 10 = 3484.$$

Subtracting the total with the number of invalid partitions, we get 4060 - 3484 = 576.

Academic integrity declaration. I did this homework by myself. I got help from office hours, though.