

# Homework: Matrix functions

July 12, 2023

This is an assignment for the course ‘Numerical Linear Algebra’. The goal of this assignment is to explore matrix functions and associated Krylov methods.

**Please note that the questions are deliberately open-ended! Unlike in the exercise sessions, you are encouraged to make your own analysis, provide your own results and construct your own report in your own way. The assignment is individual.** Don’t be discouraged if not everything works in the end, but show your work and report what you tried and whether you know why something did not work.

**Your report should be a standalone text, should read pleasantly and should not refer to this assignment. Take special care of structure and visual presentation. You are allowed to copy text from this assignment document.**

## 1 Theory

In this section we will be looking at how matrix functions can naturally be defined and equivalences between definitions.

### 1.1 Matrix-valued polynomials, rational functions and power series

Say we have a polynomial  $p(t) := \sum_{k=0}^n c_k t^k$ . Then for any matrix  $A \in \mathbb{C}^{n \times n}$  it is natural to define  $p(A) := \sum_{k=0}^n c_k A^k$ , where we adopt the convention that  $A^0 = I$ . Similarly, for any rational function  $f(t) := \frac{p(t)}{q(t)}$  we want to define  $f(A) := q(A)^{-1}p(A)$ . **When is this undefined? Investigate  $q(A)^{-1}p(A)$**

vs.  $p(A)q(A)^{-1}$ . What do you conclude? Prove this. How do you evaluate  $f(A)$  whenever  $f$  is given by a power series or a (finite-order) Laurent series?

## 1.2 Spectrum-based definition

Supposing the matrix  $A$  is diagonalizable, i.e.  $A = P\Lambda P^{-1}$ , there is an obvious definition for  $f(A)$ , for any  $f$ . **Give this definition. Motivate it in the case  $f$  is a polynomial, rational function or given by a power series. How restrictive is the assumption that  $A$  is diagonalizable? What important matrices are included? Which aren't?** We will mention here that it is possible to extend the above to using the *Jordan canonical form*, but you do not have to do that here.<sup>1</sup>

## 1.3 Interpolation-based definition

Shockingly, in this section you will show that for any  $A \in \mathbb{C}^{n \times n}$  and any sufficiently differentiable  $f$ , we can find a polynomial  $p$  such that  $f(A) = p(A)$ . It all starts by observing that only the values that are actually important for matrix *polynomials* are precisely in the spectrum of  $A$ . First we briefly recall that every matrix  $A \in \mathbb{C}^{n \times n}$  with spectrum  $\{\lambda_1, \dots, \lambda_k\}$  has a *minimal polynomial*  $\phi_A$  given by

$$\phi_A(t) := \prod_{i=1}^k (t - \lambda_i)^{n_i} \quad (1)$$

which is the unique monic minimal degree ( $\deg(\phi_A) = n_1 + \dots + n_k \leq n$ ) polynomial such that  $\phi_A(A) = 0$ . We can use it to show theorem 1.

**Theorem 1** *Let  $A \in \mathbb{C}^{n \times n}$  be a matrix with eigenvalues  $\{\lambda_1, \dots, \lambda_k\}$ , and minimal polynomial given by equation 1. Then for any two polynomials  $p_1, p_2$  we have that  $p_1(A) = p_2(A)$  if and only if*

$$\forall j \in \{1, \dots, k\} : \forall i \in 0, \dots, n_i - 1 : p_1^{(i)}(\lambda_j) = p_2^{(i)}(\lambda_j).$$

**Prove this. Hint: if a polynomial  $p$  satisfies  $p(A) = 0$  then the minimal polynomial of  $A$  divides  $p$ . To what easier statement does theorem 1 reduce in case  $A$  has simple spectrum? Use theorem 1 to motivate the definition 1 of a matrix function.**

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<sup>1</sup>Though you of course can if you want to.

**Definition 1** Let  $A \in \mathbb{C}^{n \times n}$  be a matrix with minimal polynomial given as in equation 1, and let  $f$  be a function that is at least  $\max_k \{n_k - 1\}$  times differentiable. Say  $p$  is its  $(n_1, \dots, n_k)$ -Hermite interpolant i.e. the polynomial satisfying

$$\forall j \in \{1, \dots, k\} : \forall i \in \{0, \dots, n_k - 1\} : p^{(i)}(\lambda_j) = f^{(i)}(\lambda_j)$$

of minimal degree. Then we define  $f(A) = p(A)$ .

**Comment on this definition.** Does it match with the earlier definitions? Is it more general? Is it useful, computable,...?

## 1.4 The matrix-vector product $f(A)\mathbf{b}$

Have a look at the matrix in the file `Matrix1` provided in the assignment. Call this matrix  $A$ . **Study its structure. Calculate  $\exp(A)$ . Study this matrix too. What do you notice? Use this to motivate the following claim:**

**Effort should be made to avoid the explicit computation of  $f(A)$ , if only  $f(A)\mathbf{b}$  is needed.**

**Can you give some examples from the course where only (a small amount of) matrix-vector products are needed?**

In this section we will work towards a way to evaluate  $f(A)\mathbf{b}$  in an intuitive way.

Firstly, we will define the notion of  $\phi_{A,\mathbf{b}}$ , i.e. the minimal polynomial of  $A$  with respect to the vector  $\mathbf{b}$ . This is simply the polynomial

$$\phi_{A,\mathbf{b}}(t) := \prod_{i=1}^k (t - \lambda_i)^{m_k} \quad (2)$$

of minimal degree such that  $\phi_{A,\mathbf{b}}(A)\mathbf{b} = \mathbf{0}$ . Here  $\lambda_1, \dots, \lambda_k$  are again the eigenvalues of  $A$ . **Give a non-trivial example of a pair  $(A, \mathbf{b})$  such that  $\phi_{A,\mathbf{b}} \neq \phi_A$ . Also give one where  $\phi_{A,\mathbf{b}} = \phi_A$ .** We are now ready to prove theorem 2

**Theorem 2** Let  $f$  be a sufficiently differentiable function that has no singularities on the spectrum of a given matrix  $A \in \mathbb{C}^{n \times n}$ . Then, with  $p$  the unique Hermite interpolating polynomial of  $A$  w.r.t.  $\mathbf{b}$  i.e.

$$\forall j \in \{1, \dots, k\} : \forall i \in \{0, \dots, m_k - 1\} : p^{(i)}(\lambda_j) = f^{(i)}(\lambda_j)$$

we have that  $f(A)\mathbf{b} = p(A)\mathbf{b}$ .

**Prove this.** Review the Arnoldi method and also show that theorem 2 implies that

$$f(A)\mathbf{b} = \|\mathbf{b}\|_2 Q_m H_m \mathbf{e}_1$$

where  $m := m_1 + \dots + m_k = \text{degr}(\phi_{A,\mathbf{b}})$ ,  $\mathbf{e}_1$  is the first unit vector in  $\mathbb{R}^n$  and  $Q_m, H_m$  are respectively the orthogonal matrix and the Hessenberg matrix of the Arnoldi process after  $m$  iterations.

## 2 Algorithms

In this section we will outline the algorithms you are expected to implement. We distinguish two cases: the ‘dense case’, where  $f(A)$  is to be computed, and the case only  $f(A)\mathbf{b}$  is needed for some  $\mathbf{b}$ .

### 2.1 The dense case

All the ingredients needed for the computation of  $f(A)$  have now been outlined. You have been provided a function `hess_and_phi` that computes the Hessenberg reduction of a given matrix and its associated minimal polynomial. **Study this routine. Briefly explain what it does. What is the format of input and output? How do you evaluate the output?** The reason we chose to do it like this is because Matlab’s built-in `minpoly` does not give accurate results in feasible time.<sup>2</sup> You have also been provided a Matlab routine called ‘`hermite`’ that computes the Hermite interpolant of a given minimal polynomial, using Newton divided differences. **You do not need to discuss this routine, just understand how to use it.**

**Write a Matlab routine called `matrix_function` that takes as inputs a function `f`, and a matrix `A` and outputs  $f(A)$ , as defined by definition 1. You can use `hess_and_phi` and `hermite`. Test your routine on the matrices in ‘`test1.mat`’ using the script ‘`test1.m`’.**

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<sup>2</sup>More involved techniques for this problem exist, but this is outside of the scope of this assignment.

## 2.2 $f(A)\mathbf{b}$

In this section you are asked to implement a routine that *approximates*  $f(A)\mathbf{b}$ . The way this is done is by truncating the Arnoldi process at some predefined dimension  $k$ , thus resulting in an approximation of  $f(A)\mathbf{b}$  in the Krylov space  $\mathcal{K}_k(A, \mathbf{b})$ . **Write a routine called ‘fAb’ that outputs an approximation of  $f(A)\mathbf{b}$  and takes as input a function  $f$ , a matrix  $A \in \mathbb{C}^{n \times n}$ , a vector  $\mathbf{b} \in \mathbb{C}^n$  and an integer  $k$ , representing the maximal dimension of the resulting Krylov space. You can use Arnoldi with iterated Gram-Schmidt. Investigate the convergence of your method for the functions and matrix provided in test2.mat. Discuss your findings. Compare the computational efficiency of this method as opposed to explicitly forming  $f(A)$ . You do not have to use your earlier implementation based on the minimal polynomial; rather than using the theoretical approach above, you can evaluate  $f(H_m)$  directly in Matlab. This will result in better comparisons.**

## 3 Applications

In this section two typical applications of matrix functions are discussed and (simple) model problems are provided for you to solve. It should be stressed that real-world problems quickly become too difficult for these easy standard methods to solve and thus more complex routines are needed. This however falls beyond the scope of this assignment.

### 3.1 The matrix exponential

Consider the simple system of ODEs

$$\frac{d\mathbf{x}}{dt} = A\mathbf{x},$$

with initial condition  $\mathbf{x}(0) = \mathbf{x}_0 \in \mathbb{R}^n$ . Then we know the solution to be given by  $\mathbf{x}(t) = e^{At}\mathbf{x}_0$ . However, for all but the stablest systems, this is not a good method, due to issues such as stability and stiffness<sup>3</sup>

Here we consider for instance the 2D convection-diffusion equation for the

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<sup>3</sup>we will not investigate the issues of stiffness here.

flow  $\mathbf{u}(x, y)$ :

$$\frac{d\mathbf{u}}{dt} = \epsilon \Delta \mathbf{u} + \alpha \cdot \nabla \mathbf{u}$$

with Dirichlet boundary conditions and  $\epsilon \in \mathbb{R}_0^+$  and  $\alpha \in \mathbf{R}^2$ . Simple time-stepping methods are known to be unstable at large time-steps, and our exponential scheme suffers from similar problems, i.e.  $t$  cannot be taken too large. **Study the MatLab file `convection_diffusion.m`. Explain how the ODE is formed. What discretization is used? Add your `fAb` routine to the file. Compare your results with the original routine. Be thorough.**

### 3.2 The sign function

In control theory we are often interested in the eigenvalues  $\lambda$  of system matrices with  $\operatorname{Re}(\lambda) > 0$ , since they correspond to unstable poles. In the design of controllers it is therefore interesting to have an efficient way to count the number of eigenvalues of a matrix in the right half-plane  $\operatorname{Re}(z) > 0$ . Here we will build such a method.

**Theorem 3** *Let  $A \in \mathbb{C}^{n \times n}$  be a matrix with  $k_-$  eigenvalues in the left plane,  $k_+$  eigenvalues in the right plane and none on the imaginary axis, counting multiplicity. Let  $\operatorname{sgn} : \mathbb{C} \mapsto \{1, -1\}$  be defined by*

$$\operatorname{sgn}(z) = \begin{cases} 1 & \operatorname{Re}(z) \geq 0 \\ -1 & \operatorname{Re}(z) < 0. \end{cases}$$

*Then  $\operatorname{trace}(\operatorname{sgn}(A)) = k_+ - k_-$ .*

**Prove this.** With this theorem in mind, interpret algorithm 1 (where we use MatLab notation for some of the operations). The function ‘mode’ here refers to the most frequently observed value. **Be thorough.**

**Implement a computationally efficient version of algorithm 1.** Apply your algorithm to the system `control.mat`. First plot the eigenvalues of the system using `eigs`. What do you observe? Do you expect Arnoldi to converge in few iterations? How would you remedy this? Hint: look at chapter 4 of [1] for inspiration. Compare the computed number of eigenvalues in the right half-plane to the actual number. Discuss the advantages and disadvantages of your scheme.

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**Algorithm 1:** Computing  $k_+(A)$ 

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**Data:**  $A \in \mathbb{C}^{n \times n}$ ,  $d \in \mathbb{N}$ ,  $N \in \mathbb{N}$ **Result:**  $k_+ \in \mathbb{R}$ 

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1  $\mathbf{q} \leftarrow \mathbf{0} \in \mathbb{R}^N$ ;  
2 for  $i = 1$  to  $N$  do  
3    $\mathbf{u} \leftarrow \text{randn}(n, 1)$ ;  
4    $\hat{\mathbf{u}} \leftarrow \frac{\mathbf{u}}{\|\mathbf{u}\|}$ ;  
5    $[H, Q] \leftarrow \text{arnoldi}(A, \hat{\mathbf{u}}, k)$ ;  
6    $q(i) \leftarrow \text{trace}(\text{sign}(H))$ ;  
7 end  
8  $q := \text{mode}(\mathbf{q})$ ;  
9  $q \leftarrow (q + k)/2$ 
```

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## 4 Quotation and questions

Points are awarded based on the correctness of your results and the quality of your code (40%), the insight demonstrated in your report (40%) and the organization and presentation of your report (20%). Don't hesitate to email me at [simon.dirckx@kuleuven.be](mailto:simon.dirckx@kuleuven.be) if anything is unclear or if you suspect something is wrong. If you are stuck on something you can always ask for a hint. This will be taken into account in the evaluation but will not severely impact your grade.

## References

- [1] Saad, Y., *Numerical Methods for Large Eigenvalue Problems*, Society for Industrial and Applied Mathematics, 2011